

The Ford-Fulkerson Algorithm: Unleashing Maximum Flow in Networks

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Introduction:

The world we live in is interconnected by various systems - transportation networks, communication pathways, and supply chains - all of which involve the efficient movement of resources, information, or entities. To optimize these systems, computer scientists and engineers developed the concept of network flow, a mathematical framework that reveals the dynamics of such intricate interconnections. Central to this realm is the Ford-Fulkerson algorithm, an ingenious method that has revolutionized how we calculate maximum flow within networks.

Understanding Network Flow:

Imagine a network as a web of nodes and edges, each representing a point of interest and the pathways connecting them, respectively. This abstraction is used to depict anything from road networks to data transmission routes. Network flow refers to the movement of something - be it goods, data, or even people - along these pathways. The Ford-Fulkerson algorithm, named after Delbert Fulkerson and Lester Ford, focuses on the challenge of finding the maximum flow that can be pushed from a designated source to a sink node while obeying capacity constraints on the edges.

The Algorithm in Action:

The Ford-Fulkerson algorithm is surprisingly intuitive in its approach. It starts by initializing the flow across all edges to zero. Then, in each iteration, it searches for augmenting paths from the source to the sink using a technique called Depth-First Search (DFS). If such a path exists, it finds the minimum capacity along that path and adds it to the flow. The process is repeated until no augmenting path can be found, at which point the algorithm terminates, leaving us with the maximum possible flow.

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Pseudocode:

Let's visualize the algorithm through pseudocode:

```
Algorithm FordFulkerson(G, source, sink):
```

```
    Initialize flow  $f[e]$  for each edge  $e$  in  $G$  to 0 // Initially, no flow
```

```
    while there's an augmenting path  $p$  from source to sink in the residual graph  $G_f$ :
```

```
        Find the minimum capacity  $c_{\min}$  along path  $p$  // Calculate the maximum flow that can be added
```

```
        // Update flow along path  $p$ 
```

```
        for each edge  $e$  in path  $p$ :
```

```
             $f[e] += c_{\min}$  // Increase flow by minimum capacity
```

```
        // Update the residual graph  $G_f$  based on the new flow
```

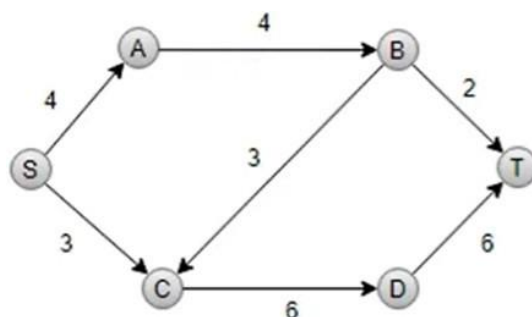
```
        for each edge  $e$  in path  $p$ :
```

```
            Reduce the residual capacity of  $e$  by  $c_{\min}$ 
```

```
            Increase the residual capacity of the reverse edge of  $e$  by  $c_{\min}$ 
```

```
    return the flow  $f$ 
```

Problem Solving: Ford-Fulkerson algorithm to find the maximum flow from node S to T in the weighted directed graph below.



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④ Ford - Fulkerson algorithm:

Edges = water pipes
 weight = max cap of pipe

residual capacity: it is the capacity of the edge after subtracting the flow from the max capacity.

residual graph: A graph with the same vertices & same edges, but we use the residual capacities as capacities.

Path 1: $S \rightarrow A \rightarrow B \rightarrow T$

$\min = 2$

$\text{min} = 2$

$\text{flow}(u, v) = \text{flow}(u, v) - \text{min cap}$

$\text{flow}(v, u) = \text{flow}(v, u) + \text{min cap}$

$\text{flow}(S, A) = 4 - 2 = 2$

$(A, S) = 0 + 2 = 2$

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$A \rightarrow B$
 $flow(A, B) = 4 - 2 = 2$
 $flow(B, A) = 0 + 2 = 2$

$B \rightarrow T$
 $flow(B, T) = 2 - 2 = 0$
 $flow(T, B) = 0 + 2 = 2$

\downarrow

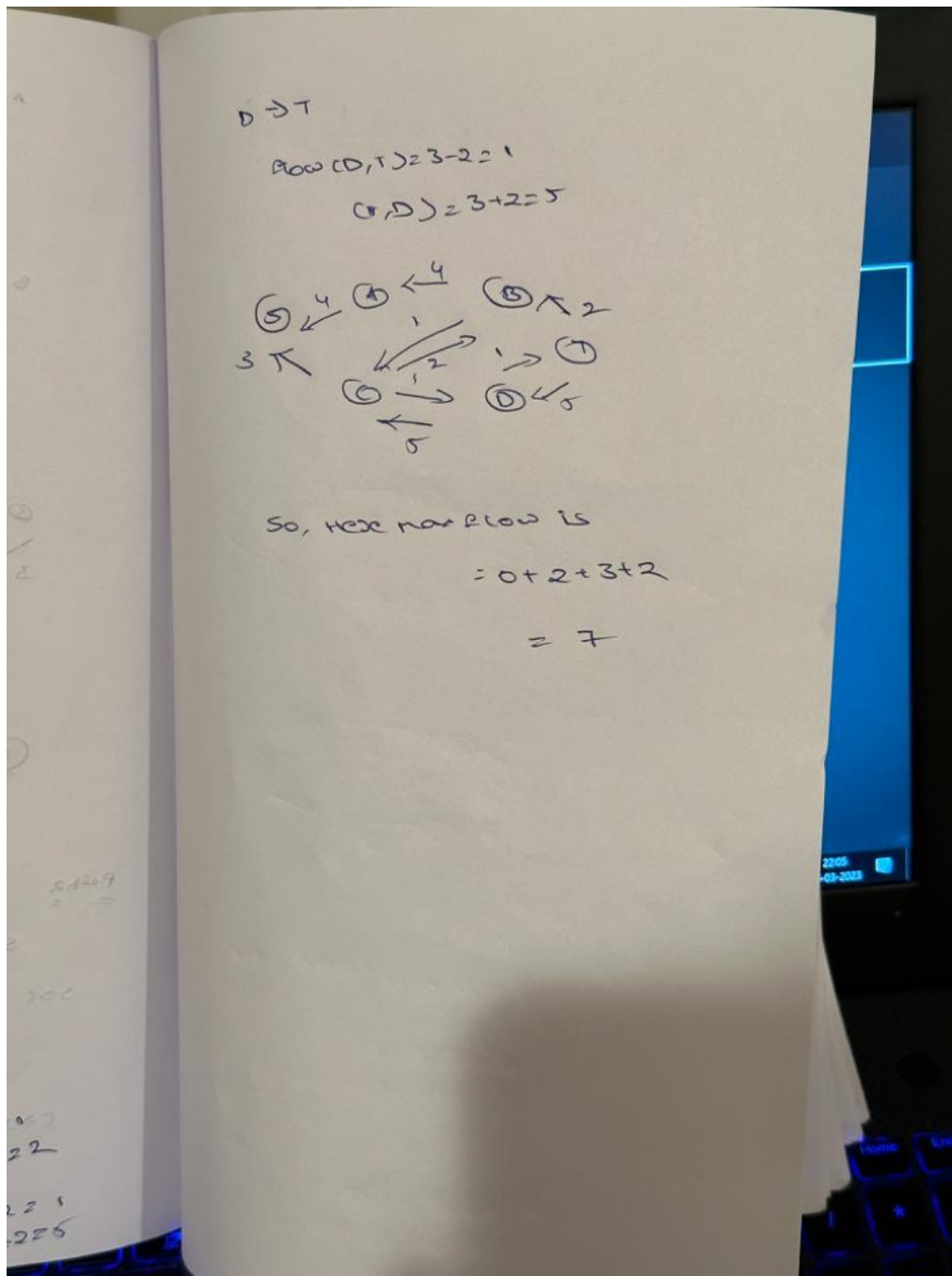
$S \rightarrow C \rightarrow D \rightarrow T$ $min cap = 3$
 $flow(C, S) = 3 - 3 = 0$
 $flow(S, C) = 0 + 3 = 3$
 $C \rightarrow D$
 $flow(C, D) = 6 - 3 = 3$
 $flow(D, C) = 0 + 3 = 3$

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$D \rightarrow T$
 $Flow(C, D) = 6 - 3 = 3$
 $Flow(C, T) = 6 + 3 = 9$

$P = 4$
 $S \rightarrow C \rightarrow D \rightarrow T$
 $S \rightarrow A \rightarrow B \rightarrow C \rightarrow D \rightarrow T$
 $min\{cap = 2$
 $max\{flow = 0 + 2 + 3 + 2$
 $S \rightarrow A$
 $Flow(S, A) = 2 - 2 = 0$
 $CA, D = 2 + 2 = 4$
 $A \rightarrow B$
 $Flow(CA, B) = 2 - 2 = 0$
 $CB, A = 2 + 2 = 4$
 $B \rightarrow C$
 $Flow(CB, C) = 3 - 2 = 1$
 $CC, B = 0 + 2 = 2$
 $C \rightarrow D$
 $Flow(C, D) = 3 - 2 = 1$
 $CD, C = 3 + 2 = 5$

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Conclusion:

In the realm of network flow, the Ford-Fulkerson algorithm stands tall as a cornerstone. Its ingenious approach to calculating maximum flow showcases the elegance of combining graph theory, optimization, and computer science. By understanding how this algorithm optimizes flows, we gain a deeper appreciation for the intricate interconnections that shape our world - a world where the movement of resources, information, and entities flows seamlessly through pathways of networks.

References:

- 1) https://en.wikipedia.org/wiki/Ford%E2%80%93Fulkerson_algorithm
- 2) <https://www.geeksforgeeks.org/ford-fulkerson-algorithm-for-maximum-flow-problem/>

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