

Github repo: [https://github.com/915-Petruta-Razvan/LFTC\\_Labs](https://github.com/915-Petruta-Razvan/LFTC_Labs)

## Documentation

### Grammar

Details: class used to represent a grammar read from a file.

Methods:

- PrintNonterminalSymbols(): prints the nonterminals of the grammar
- PrintTerminalSymbols(): prints the terminals of the grammar
- PrintAllProductions(): prints the set of productions of the grammar
- CheckCFG(): checks if a grammar is a context free grammar (the left hand side of a production is a nonterminal and the right hand side is a combination of terminal and nonterminal symbols)
- PrintProductionLHSsForANonterminal(nonterminal: string): prints the productions in which the nonterminal appears in the left hand side
- GetProductionLHSsForANonterminal(nonterminal: string): gets the productions in which the nonterminal appears in the left hand side
- PrintProductionRHSsForANonterminal(nonterminal: string): prints the productions in which the nonterminal appears in the right hand side
- GetProductionRHSsForANonterminal(nonterminal: string): gets the productions in which the nonterminal appears in the right hand side
- PrintStartingSymbol(): prints the starting symbol of the grammar
- InitFromFile(): read the grammar from a file

### Production

Details: this class is used to represent a production having a left hand side and a right hand side.

### LL1Parser

Details: parser implementation for LL(1)

Methods:

- PerformConcatenationOfSizeOne(nonterminals: List<string>, terminal: string):
  - o If a nonterminal can derive  $\epsilon$ , the first symbol of a string derived from the sequence could come from the FIRST set of the next nonterminal.
  - o The process continues until a nonterminal that cannot derive  $\epsilon$  is encountered, or all nonterminals have been checked.
  - o if all nonterminals can derive  $\epsilon$ , and there's a terminal, the terminal is also included in the resulting set (as the entire nonterminal

- sequence can derive  $\epsilon$ , leaving the terminal as the first symbol).
- GenerateFirstDictionary():
  - o Initial Pass: The method first adds terminals or  $\epsilon$  that are directly derivable from each nonterminal.
  - o Iterative Refinement: The method then iteratively refines these sets. This is necessary because the FIRST set of a nonterminal may depend on the FIRST sets of other nonterminals. For example, if a nonterminal A has a production  $A \rightarrow B C$ , then the FIRST set of B (and possibly C, if B can derive  $\epsilon$ ) contributes to the FIRST set of A.
  - o The loop continues until no more changes occur in the FIRST sets
- GenerateFollowDictionary():
  - o Initial Setup: The method starts by adding  $\epsilon$  to the FOLLOW set of the starting symbol, as it's the first symbol in the derivation process.
  - o Handling Productions: The method examines each production where a nonterminal appears on the RHS. Depending on the position and the symbol following the nonterminal, different rules are applied to update the FOLLOW set.
  - o Terminal and Nonterminal Handling: If a terminal follows the nonterminal, it's added directly to the FOLLOW set. If another nonterminal follows, the method adds all terminals from its FIRST set (except  $\epsilon$ ) to the FOLLOW set. If  $\epsilon$  is in the FIRST set of this following nonterminal, the FOLLOW set of the LHS nonterminal is also included.
  - o Iterative Refinement: The method iteratively updates the FOLLOW sets. Since the FOLLOW set of one nonterminal can depend on others, multiple iterations are needed until no further changes occur.

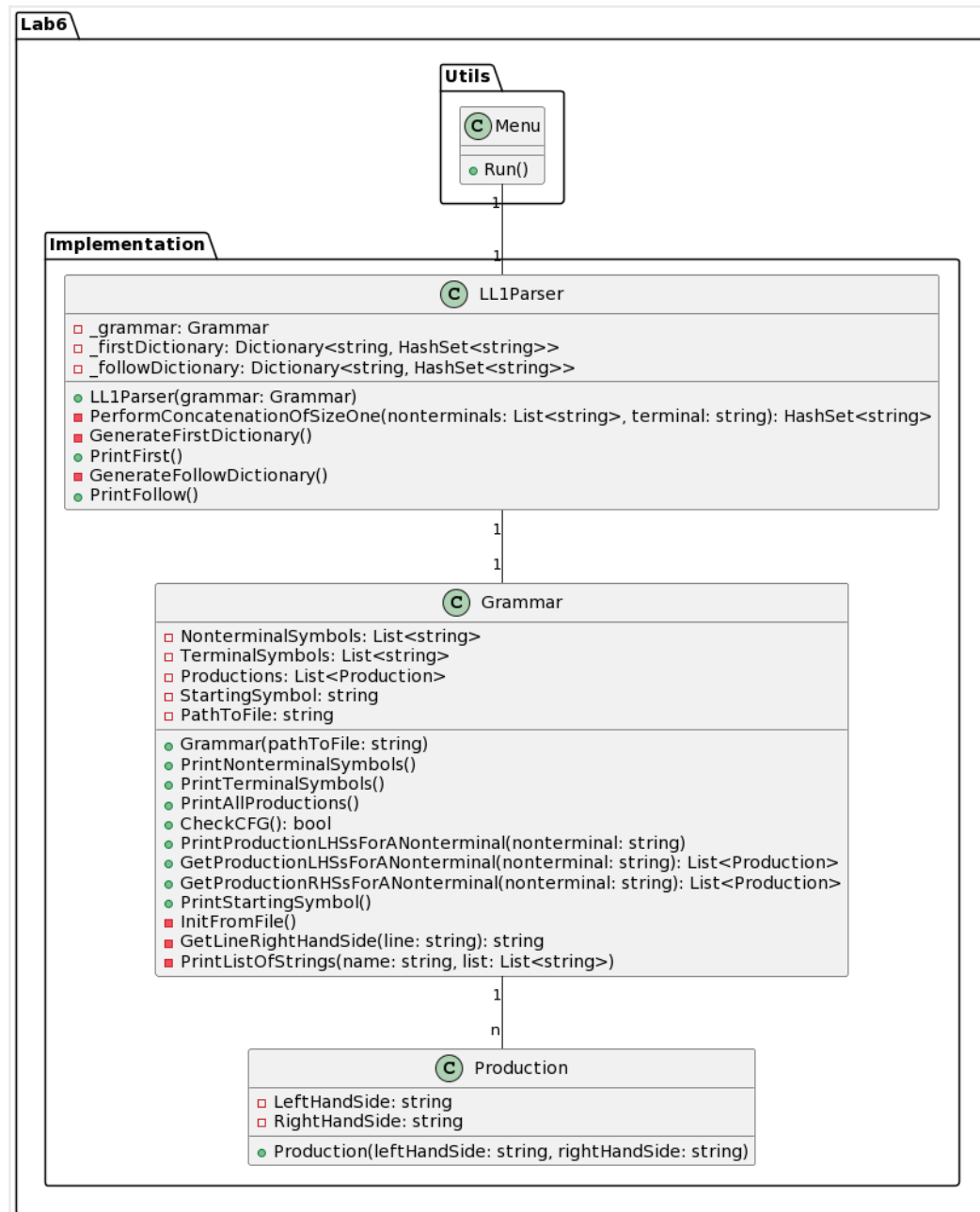
## Menu

Details: this class is used to create a menu for the user to interact with.

Possible options:

- EXIT
- Show nonterminals
- Show terminals
- Show all productions
- Show productions for a given nonterminal (LHS)
- Show productions for a given nonterminal (RHS)
- Show starting symbol
- Is the grammar a context free grammar?
- Show FIRST
- Show FOLLOW

## Class Diagram



Testing the algorithm:

Grammar from seminar 8:

**Ex.:** Given the CFG  $G = (\{S, A, B, C, D\}, \{+, *, a, (, )\}, P, S)$ ,

- $P$  :
- (1)  $S \rightarrow BA$
  - (2)  $A \rightarrow +BA$
  - (3)  $A \rightarrow \varepsilon$
  - (4)  $B \rightarrow DC$
  - (5)  $C \rightarrow *DC$
  - (6)  $C \rightarrow \varepsilon$
  - (7)  $D \rightarrow (S)$
  - (8)  $D \rightarrow a$ ,

$\text{FIRST}(S) = \{ (, a \}$   
 $\text{FIRST}(A) = \{ +, \varepsilon \}$   
 $\text{FIRST}(B) = \{ (, a \}$   
 $\text{FIRST}(C) = \{ *, \varepsilon \}$   
 $\text{FIRST}(D) = \{ (, a \}$

$\text{FOLLOW}(S) = \{ \varepsilon, ) \}$   
 $\text{FOLLOW}(A) = \{ \varepsilon, ) \}$   
 $\text{FOLLOW}(B) = \{ +, \varepsilon, ) \}$   
 $\text{FOLLOW}(C) = \{ +, \varepsilon, ) \}$   
 $\text{FOLLOW}(D) = \{ *, +, \varepsilon, ) \}$

Program execution:

```
0: EXIT
1: Show nonterminal symbols
2: Show terminal symbols
3: Show set of productions
4: Show productions for a given nonterminal (LHS)
5: Show productions for a given nonterminal (RHS)
6: Show starting symbol
7: Is the grammar a context free grammar?
8. Show FIRST
9. Show FOLLOW
> 1
N = {S,A,B,C,D}
> 2
E = {+,* ,a,(,)}
> 3
P = {
    S -> B A
    A -> + B A
    A -> ε
    B -> D C
    C -> * D C
    C -> ε
    D -> ( S )
    D -> a
}
> 6
S = S
```

> 8

S: (, a

A: +,  $\epsilon$

B: (, a

C: \*,  $\epsilon$

D: (, a

> 9

S:  $\epsilon$ , )

A:  $\epsilon$ , )

B: +,  $\epsilon$ , )

C: +,  $\epsilon$ , )

D: \*,  $\epsilon$ , +, )