

# Distributed tasking problem for track and search

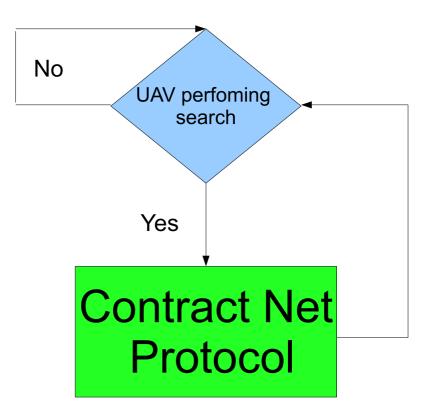
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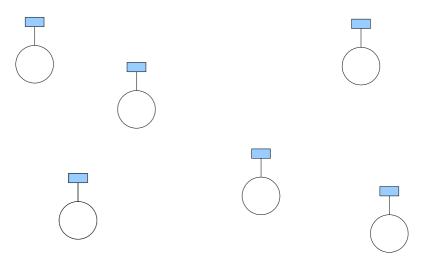
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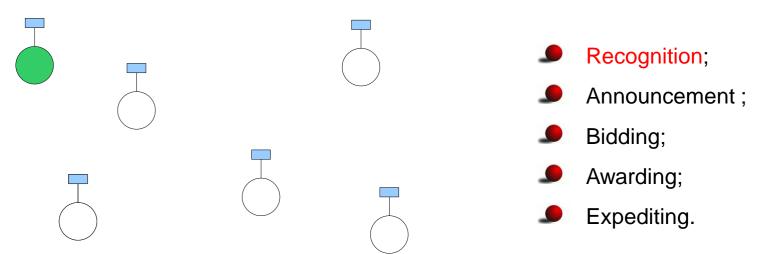
#### Introduction



- Recognition;
- Announcement;
- Bidding;
- Awarding;
- Expediting.



- Recognition;
- Announcement;
- Bidding;
- Awarding;
- Expediting.

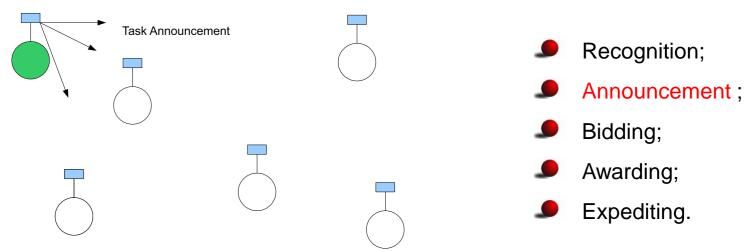


In this stage, an agent recognises it has a problem it wants help with.

Agent has a goal, and either

- realises it cannot achieve the goal in isolation does not have capability;
- realises it would prefer not to achieve the goal in isolation (typically because of solution quality, deadline, etc)

As a result, it needs to involve other agents.

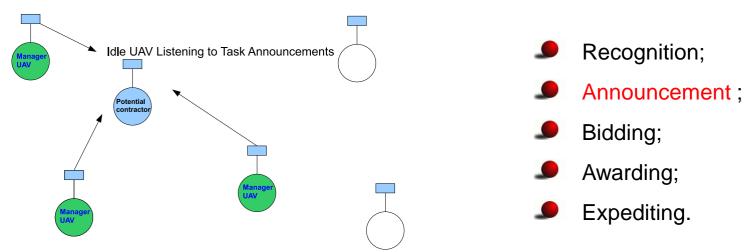


In this stage, the agent with the task sends out an announcement of the task which includes a specification of the task to be achieved.

Specification must encode:

- description of task itself (maybe executable);
- any constraints (e.g., deadlines, quality constraints).
- meta-task information (e.g.,bids must be submitted by...)

The announcement is then broadcast.

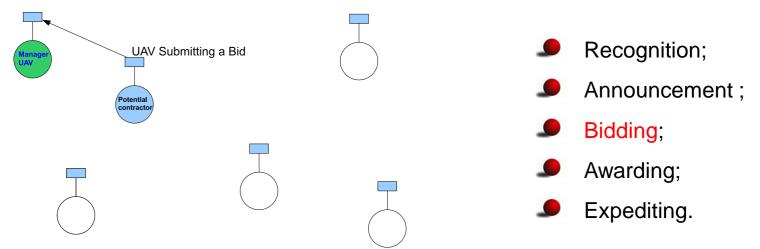


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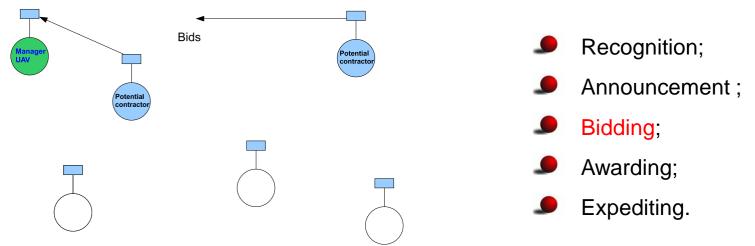


UAVs that receive the announcement decide for themselves whether they wish to bid for the task.

#### Factors:

- agent must decide whether it is capable of expediting task;
- agent must determine quality constraints and price information (if relevant).

If they do choose to bid, then they submit a tender.

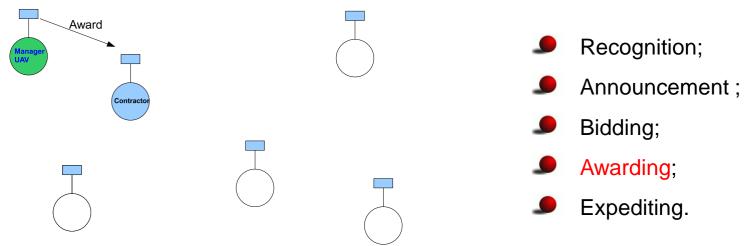


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#### Factors:

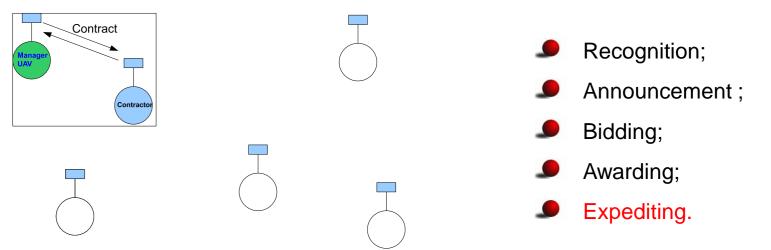
- agent must decide whether it is capable of expediting task;
- agent must determine quality constraints and price information (if relevant).

If they do choose to bid, then they submit a tender.



Agent that sent task announcement must choose between bids and decide who to "award the contract" to.

The result of this process is communicated to agents that submitted a bid.



Agent that sent task announcement must choose between bids and decide who to "award the contract" to.

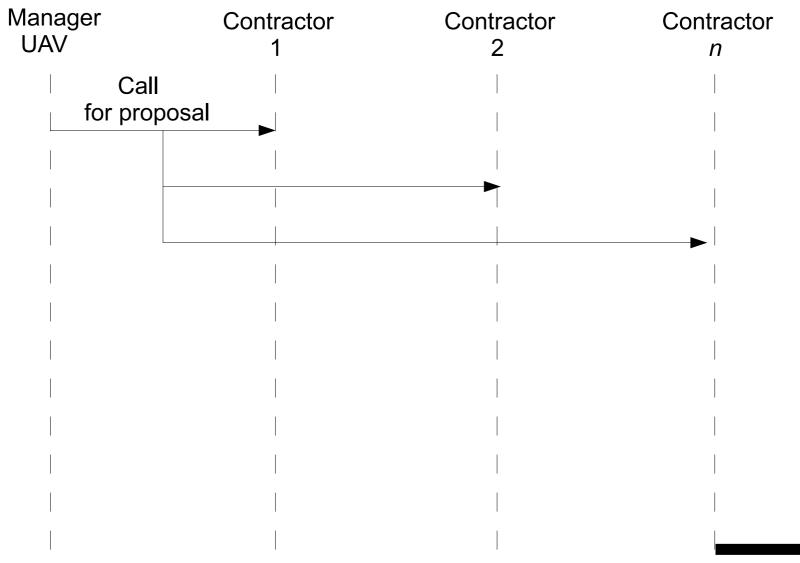
The result of this process is communicated to agents that submitted a bid.

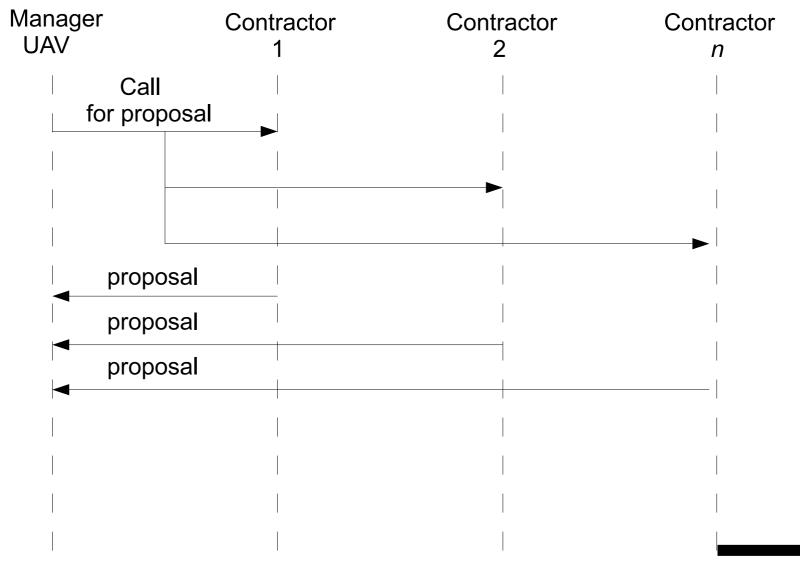
The successful contractor then expedites the task.

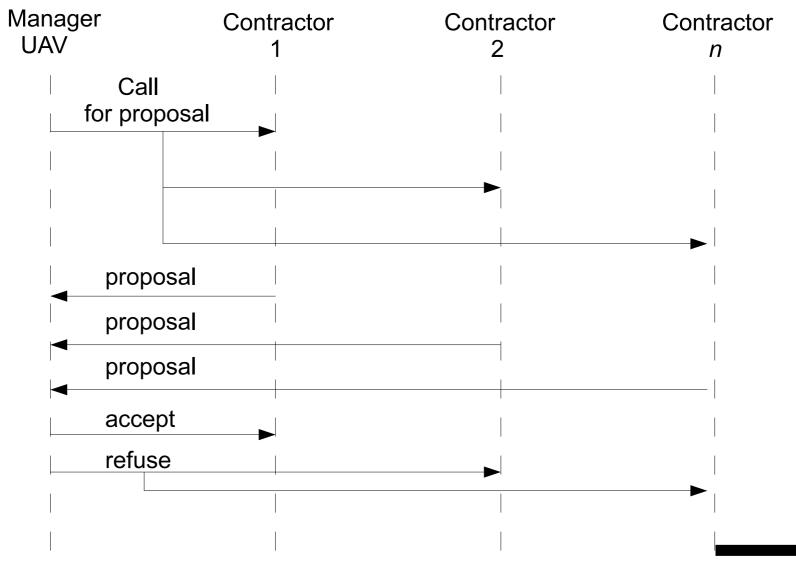
May involve generating further manager-contractor relationships: sub-contracting.

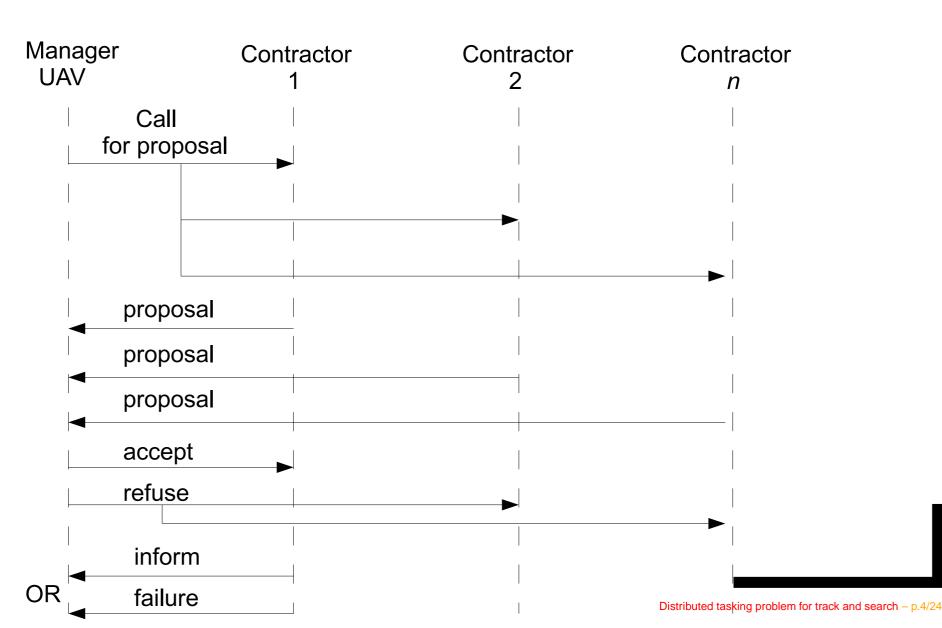
May involve another contract net.

Manager UAV	Contractor 1	Contractor 2	Contractor <i>n</i>
			1





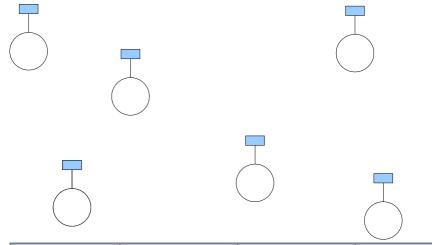




#### **Issues for Implementing Contract Net**

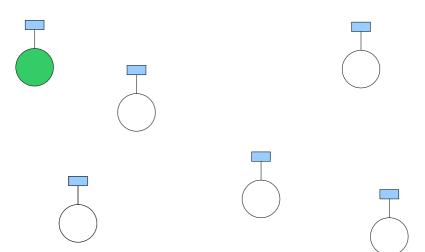
#### How to. . .

- ... specify tasks?
- ... specify quality of service?
- ... decide how to bid?
- ... select between competing offers?
- ... differentiate between offers based on multiple criteria?



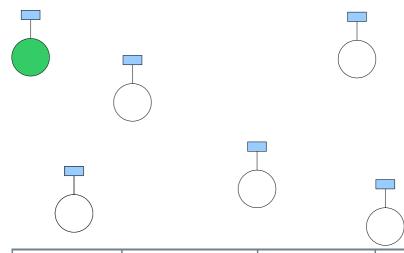
- A bundle of targets,  $\mathbf{b}_i \doteq \{b_{i1},...,b_{i|\mathbf{b}_i|}\}$
- A corresponding path,  $\mathbf{p}_i \doteq \{p_{i1},...,p_{i|\mathbf{p}_i|}\}$
- lacksquare A vector of times  $au_i \doteq \{ au_{i1},..., au_{i| au_i|}\}$

i	$Target_1$	$Target_2$	$Target_k$	$Target_{N_t}$	Values
Bandle					
Path					
Time					



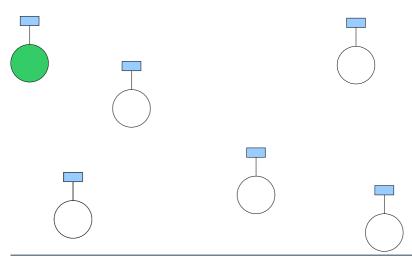
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- lacksquare A vector of times  $au_i \doteq \{ au_{i1},..., au_{i| au_i|}\}$

i	$Target_1$	$Target_2$	$Target_k$	$Target_{N_t}$	Values
Bandle	✓				$b_i = [b_{i1}]$
Path					
Time					

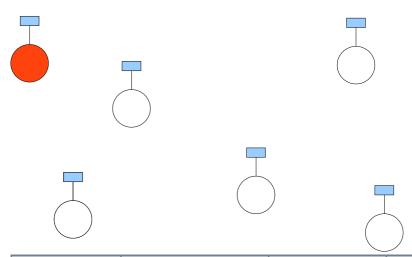


Recognition phase
Type of target (static or moving);

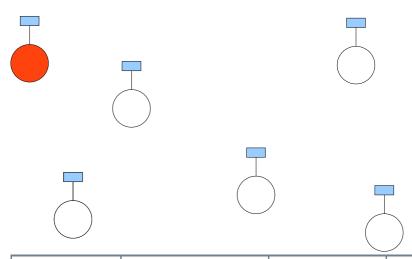
i	$Target_1$	$Target_2$	$Target_k$	$Target_{N_t}$	Values
Bandle	✓				$b_i = [b_{i1}]$
Path					
Time					



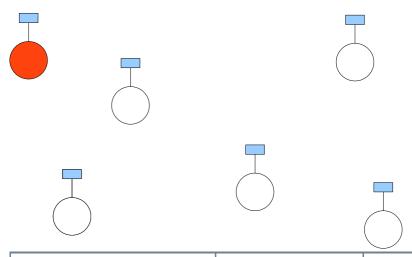
			$\smile$		
i	$Target_1$	$Target_2$	$Target_k$	$Target_{N_t}$	Values
Bandle	✓				$b_i = [b_{i1}]$
Path					
Time					



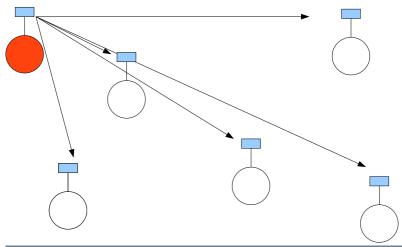
			<u> </u>		
i	$Target_1$	$Target_2$	$Target_k$	$Target_{N_t}$	Values
Bandle	✓				$b_i = [b_{i1}, b_{i2}]$
Path	$p_{i1}$				$p_i = [p_{i1}]$
Time					



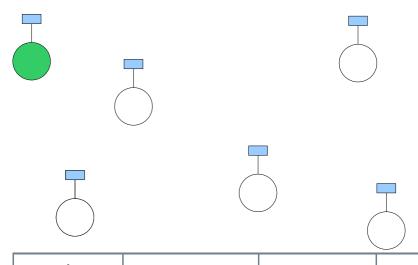
i	$Target_1$	$Target_2$	$Target_k$	$Target_{N_t}$	Values
Bandle	✓				$b_i = [b_{i1}]$
Path	$p_{i1}$				$p_i = [p_{i1}]$
Time	$ au_{i1}$				$\tau_i = [\tau_{i1}]$



i	$Target_1$	$Target_2$	$Target_k$	$Target_{N_t}$	Values
Winning Agent	i				$z_i = [z_{i1}]$
WinningBids	$y_{i1}$				$y_i = [y_{i1}]$



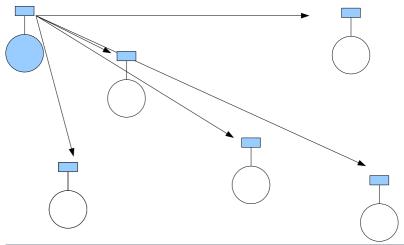
i	$Target_1$	$Target_2$	$Target_k$	$Target_{N_t}$	Values
Bandle	$\checkmark$				$b_i = [b_{i1}]$
Winning Agent	i				$z_i = [z_{i1}]$
WinningBids	$y_{i1}$				$y_i = [y_{i1}]$



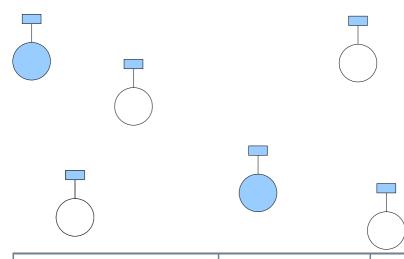
Recognition phase

Type of target (static or moving); Do the same until number of static target  $nst \leq 2$ 

i	$Target_1$	$Target_2$	$Target_k$	$Target_{N_t}$	Values
Bandle	√				$b_i = [b_{i1}]$
Path					
Time					

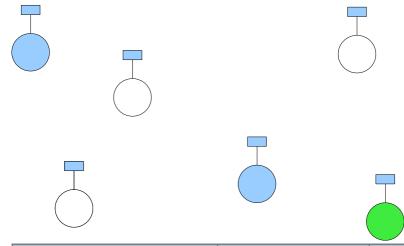


i	$Target_1$	$Target_2$	$Target_k$	$Target_{N_t}$	Values
Bandle	$\checkmark$				$b_i = [b_{i1}]$
Winning Agent	i				$z_i = [z_{i1}]$
WinningBids	$y_{i1}$				$y_i = [y_{i1}]$



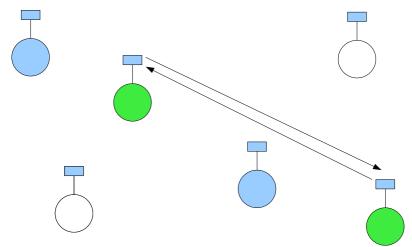
For example two static targets a found then each UAVs have the following information

k	$Target_1$	$Target_2$	$Target_k$	$Target_{N_t}$	Values
Bandle	✓	✓			$b_k = [b_{k1}, b_{k2}]$
Winning Agent	i	j			$z_k = [z_{k1}, z_{k2}]$
WinningBids	$y_{k1}$	$y_{k2}$			$y_k = [y_{k1}, y_{k2}]$



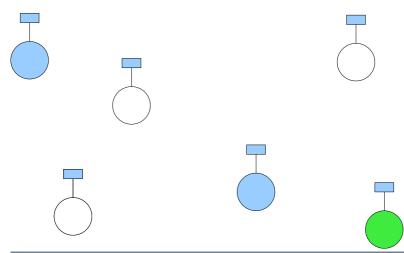
Recognition phase
Type of target (static or moving);

k	$Target_1$	$Target_2$	$Target_k$	$Target_{N_t}$	Values
Bandle	✓	✓	*		$b_k = [b_{k1}, b_{k2}, b_{kk}]$
Winning Agent	i	j			$z_k = [z_{k1}, z_{k2}]$
WinningBids	$y_{k1}$	$y_{k2}$			$y_k = [y_{k1}, y_{k2}]$



Recognition phase
Type of target (static or moving);
Check of existence of another new static target, if exist more then 1, select a manager UAV.

k	$Target_1$	$Target_2$	$Target_k$	$Target_{N_t}$	Values
Bandle	✓	✓	*	*	$[b_{k1}, b_{k2}, b_{kk}, b_{kN_t}]$
Winning Agent	i	j			$z_k = [z_{k1}, z_{k2}]$
WinningBids	$y_{k1}$	$y_{k2}$			$y_k = [y_{k1}, y_{k2}]$

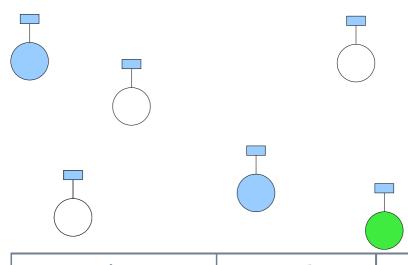


Recognition phase

Then do "Separation procedure": where input are: locations of static targets and

Number of subgroups =Total static target - number of new static targets)

k	$Target_1$	$Target_2$	$Target_k$	$Target_{N_t}$	Values
Bandle	✓	✓	¥	*	$[b_{k1}, b_{k2}, b_{kk}, b_{kN_t}]$
Winning Agent	i	j			$z_k = [z_{k1}, z_{k2}]$
WinningBids	$y_{k1}$	$y_{k2}$			$y_k = [y_{k1}, y_{k2}]$

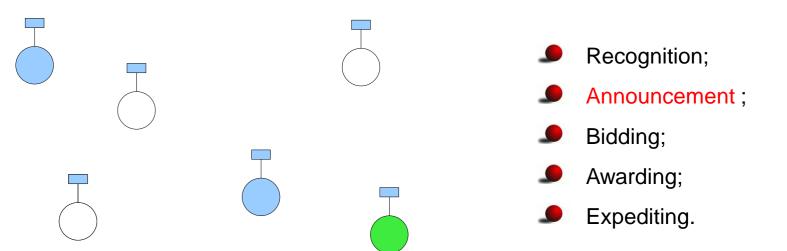


Recognition phase

Then do "Separation procedure": where input are: locations of static targets and

Number of subgroups = current length of bundle |b| - number of new static targets)

		<u> </u>			
k	$Task_1$	$Task_2$			Values
Bandle			The state of the s	*	$[b_{k1}, b_{k2}, b_{kk}, b_{kN_t}]$
Winning Agent					$z_k = [z_{kt1}, z_{kt2}]$
WinningBids	$y_{kt1}$	$y_{kt2}$			$y_k = [y_{kt1}, y_{kt2}]$



k	$Task_1$	$Task_2$			Values
Bandle			The state of the s	The state of the s	$[b_{k1}, b_{k2}, b_{kk}, b_{kN_t}]$
Winning Agent					$z_k = [z_{kt1}, z_{kt2}]$
WinningBids	$y_{kt1}$	$y_{kt2}$			$y_k = [y_{kt1}, y_{kt2}]$

#### **Definition of UAVs and targets**

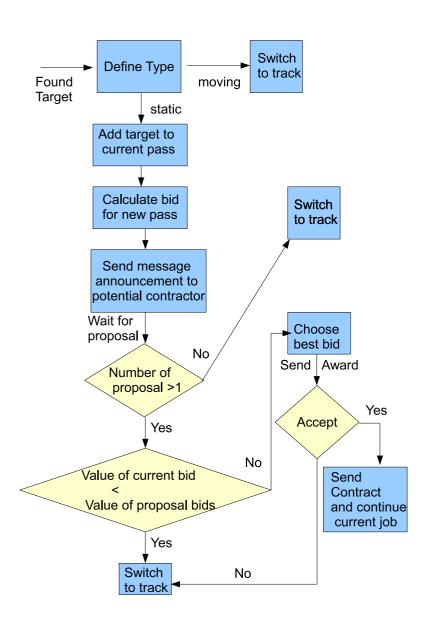
#### Possible target (task) fields:

- id task id;
- type -task type;
- value -task reward;
- start-task start time (sec);
- end task expiry time (sec);
- duration -task default duration (sec);
- x- task position (meters);
- y-task position (meters);
- z-task position (meters).

#### Possible UAVs fields:

- id- agent id;
- type- agent type;
- avail- agent availability (expected time in sec);
- x- agent position (meters);
- y- agent position (meters);
- z- agent position (meters);
- velocity agent cruise velocity (m/s));
- fuel-(agent fuel per meter)).

#### Manager statecharts



#### Manager UAV (Case 1)

Manager $UAV_i$			Values
Bundle			$b_i = []$
Path			$p_i = []$
Time			$ au_i = []$

#### $Performing\ Search$

Manager $UAV_i$			Values
$Winning\ Agent$			$z_i = []$
WinningBids			$y_i = []$

#### Manager UAV (Case 1)

Manager $UAV_i$	$Target_1$		Values
Bundle	¥		$b_i = []$
Path			$p_i = []$
Time			$ au_i = []$

#### $Found\ target$

Manager $UAV_i$			Values
Winning Agent			$z_i = []$
WinningBids			$y_i = []$

Manager $UAV_i$	$Target_1$			Values
Bundle	*			$b_i = []$
Path				$p_i = []$
Time				$ au_i = []$
		mov	ing	

Manager $UAV_i$			Values
$oxed{Winning Agent}$			$z_i = []$
WinningBids			$y_i = []$

Manager $UAV_i$	$Target_1$		Values
Bundle	<b>√</b>		$b_i = []$
Path			$p_i = []$
Time			$ au_i = []$



Switch to track

Manager $UAV_i$	$Target_1$			Values
Bundle	*			$b_i = []$
Path				$p_i = []$
Time				$ au_i = []$
		stat	tic	

 Manager  $UAV_i$  Values 

  $Winning\ Agent$   $z_i = []$  

 WinningBids  $y_i = []$ 

Manager $UAV_i$	$Target_1$		Values
Bundle	¥		$b_i = [b_{i1}]$
Path	$p_{i1}$		$p_i = [p_{i1}]$
Time	$ au_{i1}$		$\tau_i = [\tau_{i1}]$

Calculate arrival time  $\tau_{i1}(p)$  and corresponding bid  $y_{i1}$ 

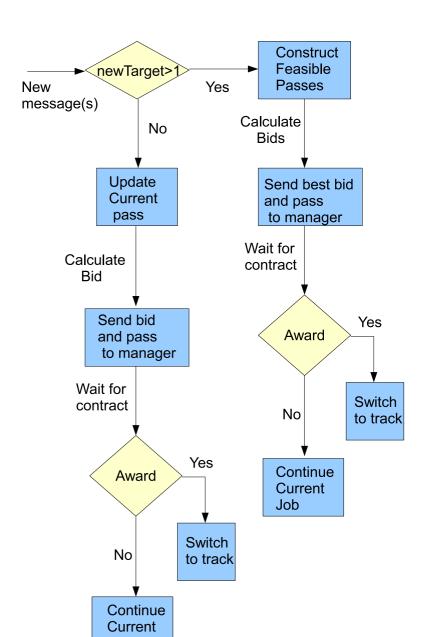
Manager $UAV_i$			Values
Winning Agent			$z_i = []$
WinningBids			$y_i = []$

Manager $UAV_i$	$Target_1$		Values
Bundle	*		$b_i = [b_{i1}]$
Path	$p_{i1}$		$p_i = [p_{i1}]$
Time	$ au_{i1}$		$ au_i = [ au_{i1}]$

Calculate arrival time  $\tau_{i1}(p)$  and corresponding bid  $y_{i1}$ 

Manager $UAV_i$	$Target_1$		Values
Winning Agent	i		$z_i = [z_{i1}]$
WinningBids	$y_{i1}$		$y_i = [y_{i1}]$

#### Potential contractors statecharts



Potential Contractor $UAV_j$			Values
Bundle			$b_j = []$
Path			$p_j = []$
Time			$ au_j = []$

Performing Search

Potential Contractor $UAV_j$	$Target_1$		Values
Bundle	¥		$b_i = []$
Path			$p_i = []$
Time			$ au_i = []$

 $Recieve\ message$ 

Potential Contractor $UAV_j$	$Target_1$		Values
Bundle	¥		$b_j = [b_{j1}]$
Path	$p_{j1}$		$p_j = [p_{j1}]$
Time	$ au_{j1}$		$\tau_j = [\tau_{j1}]$

Calculate arrival time  $\tau_{j1}(p)$  and corresponding bid  $y_{j1}$ 

Potential Contractor $UAV_j$	$Target_1$		Values
$Winning\ Agent$	j		$z_j = [z_{j1}]$
WinningBids	$y_{j1}$		$y_j = [y_{j1}]$

$UAV_k$	$Target_2$	$Target_k$	Values
Bundle	$\checkmark$	<b>√</b>	$b_k = [b_{k2}, b_{kk}]$
Path	$p_{k2}$	$p_{kk}$	$p_k = [p_{kk}, p_{k2}]$
Time	$ au_{k2}$	$ au_{kk}$	$ au_k = [ au_{kk},  au_{k2}]$

Performing Tracking

$UAV_k$	$Target_1$	$Target_2$	$Target_k$	Values
Bundle	¥	$\checkmark$	✓	$b_k = [b_{k2}, b_{kk}]$
Path		$p_{k2}$	$p_{kk}$	$p_k = [p_{kk}, p_{k2}]$
Time		$ au_{k2}$	$ au_{kk}$	$ au_k = [ au_{kk},  au_{k2}]$

 $Recieve\ message$ 

$UAV_k$	$Target_1$	$Target_2$	$Target_k$	Values
Bundle	<b>√</b>	$\checkmark$	✓	$\mathbf{b}_i \leftarrow (\mathbf{b}_i \oplus_{end} 1)$
Path		$p_{k2}$	$p_{kk}$	$p_k = [p_{kk}, p_{k2}]$
Time		$ au_{k2}$	$ au_{kk}$	$ au_k = [ au_{kk},  au_{k2}]$

 $Update\ current\ bundle\ of\ targets$ 

$UAV_k$	$Target_1$	$Target_2$	$Target_k$	Values
Bundle	$\checkmark$	$\checkmark$	✓	$\mathbf{b}_k \leftarrow (\mathbf{b}_k \oplus_{end} 1)$
Path	$p_{k1}$	$p_{k2}$	$p_{kk}$	$\mathbf{p}_k \leftarrow (\mathbf{p}_k \oplus_{n_1} 1)$
Time		$ au_{k2}$	$ au_{kk}$	$\tau_k = [\tau_{kk}, \tau_{k2}]$

Update current pass

$UAV_k$	$Target_1$	$Target_2$	$Target_k$	Values
Bundle	$\checkmark$	✓	✓	$\mathbf{b}_k \leftarrow (\mathbf{b}_k \oplus_{end} 1)$
Path	$p_{k1}$	$p_{k2}$	$p_{kk}$	$\mathbf{p}_k \leftarrow (\mathbf{p}_k \oplus_{n_1^*} 1)$
Time		$ au_{k2}$	$ au_{kk}$	$ au_k = [ au_{kk},  au_{k2}]$

Update current pass

 $\Rightarrow$  And optimal location  $n_1^*$  is then given by  $n_1^* = \max_{n_1} c_1(\tau_{k1}^*(\mathbf{p}_k \oplus_{n_1} 1))$ 

$UAV_k$	$Target_1$	$Target_2$	$Target_k$	Values
Bundle	<b>√</b>	<b>√</b>	$\checkmark$	$\mathbf{b}_k \leftarrow (\mathbf{b}_k \oplus_{end} 1)$
Path	$p_{k1}$	$p_{k2}$	$p_{kk}$	$\mathbf{p}_k \leftarrow (\mathbf{p}_k \oplus_{n_1^*} 1)$
Time	$ au_{k1}$	$ au_{k2}$	$ au_{kk}$	$\tau_k \leftarrow (\tau_k \oplus_{n_1*} \tau_{k1}(\mathbf{p}_k \oplus_{n_1*}$

 $Update\; current\; pass$ 

 $\Longrightarrow$  And optimal location  $n_1^*$  is then given by  $n_1^* = \max_{n_1} c_1(\tau_{k1}^*(\mathbf{p}_k \oplus_{n_1} 1))$ 

$UAV_k$	$Target_1$	$Target_2$	$Target_k$	Values
Bundle	<b>√</b>	✓	✓	$\mathbf{b}_k \leftarrow (\mathbf{b}_k \oplus_{end} 1)$
Path	$p_{k1}$	$p_{k2}$	$p_{kk}$	$\mathbf{p}_k \leftarrow (\mathbf{p}_k \oplus_{n_1^*} 1)$
Time	$ au_{k1}$	$ au_{k2}$	$ au_{kk}$	$\tau_k \leftarrow (\tau_k \oplus_{n_1 *} \tau_{k1}(\mathbf{p}_k \oplus_{n_1 *}$

#### $Update\ current\ pass$

- $\Longrightarrow$  And optimal location  $n_1^*$  is then given by  $n_1^* = \max_{n_1} c_1(\tau_{k1}^*(\mathbf{p}_k \oplus_{n_1} 1))$
- $\Longrightarrow$  Then the final score for new task j (which is include  $|b_k|$  targets) is

$$c_{kj}(\mathbf{p}_i) = c_j(\tau_{kj}^*(\mathbf{p}_k \oplus_{n_j^*} j))$$

# Compare bids

For case, when bundle of manager UAV3 was not empty  $|b_3| \neq \emptyset$ 

	Proposal1	Proposal2	Proposal3
UAV1	$c_{11}$	-	-
UAV2	-	$c_{22}$	-
UAV3	-	-	$c_{33}$

### Compare bids

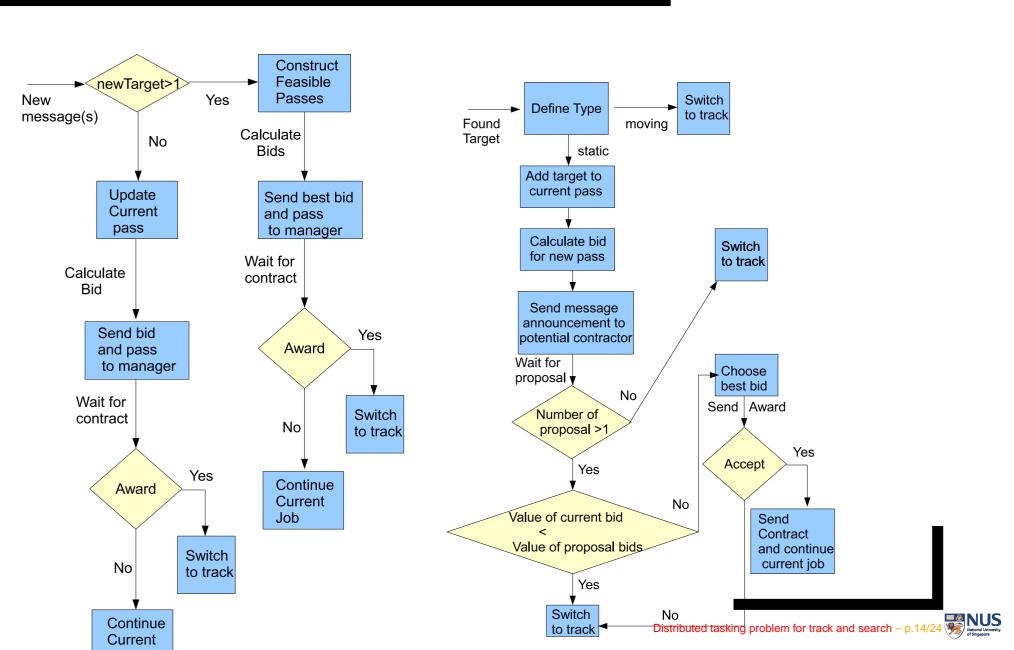
For case, when bundle of manager UAV3 was not empty  $|b_3| \neq \emptyset$ 

	Proposal1	Proposal2	Proposal3
UAV1	$c_{11}$	-	
UAV2	-	$c_{22}$	-
UAV3	-	-	$c_{33}$

For case, when bundle of manager UAV3 was empty  $|b_3|=\emptyset$ 

	Proposal1	Proposal2	Proposal3
UAV1	$c_{11}$	-	-
UAV2	-	$c_{22}$	-
UAV3	$c_{31}$	$c_{32}$	$c_{33}$

### Potential contractors and Manager



#### **Problem statement**

$$\sum_{i=1}^{N_a} \left( \sum_{j=1}^{N_t} c_{ij}(\tau_{ij}(\mathbf{p}_i(\mathbf{x_i}))) x_{ij} \right) \to \max \qquad \text{where } x_{ij} = 1 \text{ if agent } i \text{ is assigned}$$

$$\text{to task } j, \text{ and } \mathbf{x}_i \doteq \{x_{i1}, ..., x_{iN_t}\}$$

$$\text{is a vector of assignments for agent}$$

$$subject \ to: \qquad i, \text{ whose } j\text{-th element is } x_{ij}.$$

$$\begin{array}{ll} \text{where } x_{ij} = 1 \text{ if agent } i \text{ is assigned} \\ \text{to task } j, \text{ and } \mathbf{x}_i \doteq \{x_{i1},...,x_{iN_t}\} \\ \text{is a vector of assignments for agent} \\ subject \ to: \qquad i, \text{ whose } j\text{-th element is } x_{ij}. \end{array}$$

$$\sum_{j=1}^{N_t} x_{ij} \le L_t, \ \forall i \in \mathcal{I}$$

$$\sum_{i=1}^{N_a} x_{ij} \le 1, \ \forall j \in \mathcal{J}$$

$$\sum_{i=1}^{N_a} x_{ij} \leq 1, \ \, \forall j \in \mathcal{J} \qquad \text{in the objective function represents} \\ \text{the local reward for agent } i.$$

 $x_{ij} \in \{0,1\}, \forall (i,j) \in \mathcal{I} \times \mathcal{J}$ 

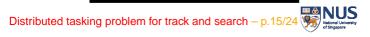
 $N_a$  Number of agents

 $N_t$ - Number of tasks

 $L_t$ - Maximum length of the bundle, i.e. each agent can be assigned a maximum  $L_t$  tasks

 $\mathcal{I}$ - Index set of agents where  $\mathcal{I} \doteq \{1,...,N_a\}$ 

 $\mathcal{J}$ - Index set of tasks where  $\mathcal{J} \doteq \{1,...,N_t\}$ 



#### Problem statement

$$\sum_{i=1}^{N_a} \left( \sum_{j=1}^{N_t} c_{ij}(\tau_{ij}(\mathbf{p}_i(\mathbf{x_i}))) x_{ij} \right) \to \max \qquad \text{where } x_{ij} = 1 \text{ if agent } i \text{ is assigned}$$

$$\text{to task } j, \text{ and } \mathbf{x}_i \doteq \{x_{i1}, ..., x_{iN_t}\}$$
is a vector of assignments for agent 
$$subject \ to: \qquad i, \text{ whose } j\text{-th element is } x_{ij}.$$

where  $x_{ij} = 1$  if agent i is assigned  $subject \ to: i$ , whose j-th element is  $x_{ij}$ .

$$\sum_{j=1}^{N_t} x_{ij} \le L_t, \ \forall i \in \mathcal{I}$$

$$\sum_{i=1}^{N_a} x_{ij} \le 1, \ \forall j \in \mathcal{J}$$

$$x_{ij} \in \{0,1\}, \forall (i,j) \in \mathcal{I} \times \mathcal{J}$$

The summation term in brackets  $\sum_{i=1}^{N_a} x_{ij} \leq 1, \ \, \forall j \in \mathcal{J} \qquad \text{in the objective function represents}$ the local reward for agent i.

 $\mathbf{p}_i \doteq \{p_{i1},...,p_{i|\mathbf{p}_i|}\}$  - The variable length vector represent the path for agent i,an ordered sequence of tasks where the elements are the task indices,  $p_{in} \in \mathcal{J}$  for  $n = 1, ..., |\mathbf{p}_i|$ , i.e. its n-th element is  $j \in \mathcal{J}$ if agent i conducts task j at the n-th point along the path. The current length of the path is denoted by  $|\mathbf{p}_i| \leq L_t$ .

#### Problem statement

$$\sum_{i=1}^{N_a} \left( \sum_{j=1}^{N_t} c_{ij}(\tau_{ij}(\mathbf{p}_i(\mathbf{x_i}))) x_{ij} \right) \to \max \qquad \text{where } x_{ij} = 1 \text{ if agent } i \text{ is assigned} \\ \text{to task } j, \text{ and } \mathbf{x}_i \doteq \{x_{i1}, ..., x_{iN_t}\} \\ \text{is a vector of assignments for agent} \\ subject \ to: \qquad i, \text{ whose } j\text{-th element is } x_{ij}.$$

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$$\sum_{j=1}^{N_t} x_{ij} \le L_t, \ \forall i \in \mathcal{I}$$

$$\sum_{j=1}^{N_a} x_{ij} \le 1, \ \forall j \in \mathcal{J}$$

$$\sum_{i=1}^{N_a} x_{ij} \leq 1, \ \, \forall j \in \mathcal{J} \qquad \text{in the objective function represents} \\ \text{the local reward for agent } i.$$

 $x_{ij} \in \{0,1\}, \forall (i,j) \in \mathcal{I} \times \mathcal{J}$ 

An assignment is said to be free of conflicts if each task is assigned to no more than one agent.

### Key assumptions

$$\sum_{i=1}^{N_a} \left( \sum_{j=1}^{N_t} c_{ij}(\tau_{ij}(\mathbf{p}_i(\mathbf{x_i}))) x_{ij} \right) \to \max$$

$$subject \ to:$$

$$\sum_{j=1}^{N_t} x_{ij} \le L_t, \ \forall i \in \mathcal{I}$$

$$\sum_{i=1}^{N_a} x_{ij} \le 1, \ \forall j \in \mathcal{J}$$

$$x_{ij} \in \{0,1\}, \forall (i,j) \in \mathcal{I} \times \mathcal{J}$$

- The score  $c_{ij}$  that agent i obtains by performing task j is defined as a function of the arrival time  $\tau_{ij}$  at which the agent executes the task (or possibly the expected arrival time in a probabilistic setting).
- The arrival time  $\tau_{ij}$  is uniquely defined as a function of the path  $\mathbf{p}_i$  that agent i takes.
- The path  $\mathbf{p}_i$  is uniquely defined by the assignment vector of agent  $i, \mathbf{x}_i$ .

# **Key assumptions**

$$\sum_{i=1}^{N_a} \left( \sum_{j=1}^{N_t} c_{ij}(\tau_{ij}(\mathbf{p}_i(\mathbf{x_i}))) x_{ij} \right) \to \max$$

$$subject \ to:$$

$$\sum_{j=1}^{N_t} x_{ij} \le L_t, \ \forall i \in \mathcal{I}$$

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$$x_{ij} \in \{0,1\}, \forall (i,j) \in \mathcal{I} \times \mathcal{J}$$

- Probabilistic setting).

  The score  $c_{ij}$  that agent i obtains by performing task j is defined as a function of the arrival time  $\tau_{ij}$  at which the agent executes the task (or possibly the expected arrival time in a probabilistic setting).
- The arrival time  $\tau_{ij}$  is uniquely defined as a function of the path  $\mathbf{p}_i$  that agent i takes.
- The path  $p_i$  is uniquely defined by the assignment vector of agent  $i, x_i$ .

An example is the problem involving time-discounted values of targets, in which the sooner an agent arrives at the target, the higher the reward it obtains. Or for scenario involves re-visit tasks, where previously observed targets must be revisited at some scheduled time. In this case the score function would have its maximum at the desired re-visiting time and lower values at other re-visit times.

- lacksquare A bundle,  $\mathbf{b}_i \doteq \{b_{i1},...,b_{i|\mathbf{b}_i|}\}$ 
  - of variable length whose elements are defined by  $b_{in} \in \mathcal{J}$  for  $n=1,...,|\mathbf{b}_i|$ . The current length of the bundle is denoted by  $b_i$ , which cannot exceed the maximum length  $L_t$ , and an empty bundle is represented by  $b_i = \emptyset$  and  $|\mathbf{b}_i| = 0$ . The bundle represents the tasks that agent i has selected to do, and is ordered chronologically with respect to when the tasks were added (i.e. task  $b_{in}$  was added before task  $b_{i(n+1)}$ ).
- lacksquare A corresponding path,  $\mathbf{p}_i \doteq \{p_{i1},...,p_{i|\mathbf{p}_i|}\}$
- lacksquare A vector of times  $au_i \doteq \{ au_{i1}, ..., au_{i| au_i|}\}$
- lacksquare A winning agent list  $\mathbf{z}_i \doteq \{z_{i1},...,z_{iN_t}\}$  of size  $N_t$
- ullet A winning bid list  $\mathbf{y}_i \doteq \{y_{i1},...,y_{iN_t}\}$  of size  $N_t$
- ullet Vector of timestamps  $\mathbf{s}_i \doteq \{s_{i1},...,s_{iN_a}\}$ , of size  $N_a$

- lacksquare A bundle,  $\mathbf{b}_i \doteq \{b_{i1},...,b_{i|\mathbf{b}_i|}\}$
- lacksquare A corresponding path,  $\mathbf{p}_i \doteq \{p_{i1},...,p_{i|\mathbf{p}_i|}\}$

whose elements are defined by  $p_i \doteq \{p_{i1},...,p_{i|\mathbf{p}_i|}\}$  for  $n=1,...,|\mathbf{b}_i|$ . The path contains the same tasks as the bundle, and is used to represent the order in which agent i will execute the tasks in its bundle. The path is therefore the same length as the bundle, and is not permitted to be longer than  $L_t$ ;  $|\mathbf{p}_i| = |\mathbf{b}_i| \leq L_t$ .

- lacksquare A vector of times  $au_i \doteq \{ au_{i1}, ..., au_{i| au_i|}\}$
- ullet A winning agent list  $\mathbf{z}_i \doteq \{z_{i1},...,z_{iN_t}\}$  of size  $N_t$
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- lacksquare A corresponding path,  $\mathbf{p}_i \doteq \{p_{i1},...,p_{i|\mathbf{p}_i|}\}$
- A vector of times  $\tau_i \doteq \{\tau_{i1},...,\tau_{i|\tau_i|}\}$  whose elements are defined by  $\tau_{in}$  for  $n=1,...,|\tau_i|$ . The times vector represents the corresponding times at which agent i will execute the tasks in its path, and is necessarily the same length as the path.
- lacksquare A winning agent list  $\mathbf{z}_i \doteq \{z_{i1},...,z_{iN_t}\}$  of size  $N_t$
- ullet A winning bid list  $\mathbf{y}_i \doteq \{y_{i1},...,y_{iN_t}\}$  of size  $N_t$
- lacksquare Vector of timestamps  $\mathbf{s}_i \doteq \{s_{i1},...,s_{iN_a}\}$ , of size  $N_a$

- lacksquare A bundle,  $\mathbf{b}_i \doteq \{b_{i1},...,b_{i|\mathbf{b}_i|}\}$
- lacksquare A corresponding path,  $\mathbf{p}_i \doteq \{p_{i1},...,p_{i|\mathbf{p}_i|}\}$
- lacksquare A vector of times  $au_i \doteq \{ au_{i1}, ..., au_{i| au_{i}|}\}$
- A winning agent list  $\mathbf{z}_i \doteq \{z_{i1},...,z_{iN_t}\}$  of size  $N_t$  where each element  $z_{ij} \in \{\mathcal{I} \cup \emptyset\}$  for  $j=1,...,N_t$  indicates who agent i believes is the current winner for task j. Specifically, the value in element  $z_{ij}$  is the index of the agent who is currently winning task j according to agent i, and is  $z_{ij} = \emptyset$ ; if agent i believes that there is no current winner.
- ullet A winning bid list  $\mathbf{y}_i \doteq \{y_{i1},...,y_{iN_t}\}$  of size  $N_t$
- Vector of timestamps  $\mathbf{s}_i \doteq \{s_{i1}, ..., s_{iN_a}\}$ , of size  $N_a$

- lacksquare A bundle,  $\mathbf{b}_i \doteq \{b_{i1},...,b_{i|\mathbf{b}_i|}\}$
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- lacksquare A vector of times  $au_i \doteq \{ au_{i1}, ..., au_{i|\tau_i|}\}$
- ullet A winning agent list  $\mathbf{z}_i \doteq \{z_{i1},...,z_{iN_t}\}$  of size  $N_t$
- A winning bid list  $\mathbf{y}_i \doteq \{y_{i1},...,y_{iN_t}\}$  of size  $N_t$  where the elements  $y_{ij} \in [0,\infty)$  represent the corresponding winners bids and take the value of 0 if there is no winner for the task.
- Vector of timestamps  $\mathbf{s}_i \doteq \{s_{i1}, ..., s_{iN_a}\}$ , of size  $N_a$

- lacksquare A bundle,  $\mathbf{b}_i \doteq \{b_{i1},...,b_{i|\mathbf{b}_i|}\}$
- lacksquare A corresponding path,  $\mathbf{p}_i \doteq \{p_{i1},...,p_{i|\mathbf{p}_i|}\}$
- lacksquare A vector of times  $au_i \doteq \{ au_{i1}, ..., au_{i| au_{i}|}\}$

through a neighboring agent.

- ullet A winning agent list  $\mathbf{z}_i \doteq \{z_{i1},...,z_{iN_t}\}$  of size  $N_t$
- ullet A winning bid list  $\mathbf{y}_i \doteq \{y_{i1},...,y_{iN_t}\}$  of size  $N_t$
- Vector of timestamps  $\mathbf{s}_i \doteq \{s_{i1},...,s_{iN_a}\}$ , of size  $N_a$  where each element  $s_{ik} \in [0,\infty)$  for  $k=1,...,N_a$  represents the timestamp of the last information update agent i received about agent k, either directly or

- lacksquare A bundle,  $\mathbf{b}_i \doteq \{b_{i1},...,b_{i|\mathbf{b}_i|}\}$
- lacksquare A corresponding path,  $\mathbf{p}_i \doteq \{p_{i1},...,p_{i|\mathbf{p}_i|}\}$
- lacksquare A vector of times  $au_i \doteq \{ au_{i1}, ..., au_{i|\tau_i|}\}$
- ullet A winning agent list  $\mathbf{z}_i \doteq \{z_{i1},...,z_{iN_t}\}$  of size  $N_t$
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Each agent must carry these vectors of information in order to be able to perform decentralized algorithm which consists of iterations between two phases:

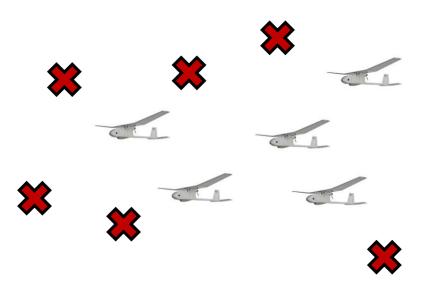
a bundle building phase where each vehicle greedily generates an ordered bundle of tasks, and a

consensus phase where conflicting assignments are identified and resolved through local communication between neighboring agents

- lacksquare A bundle,  $\mathbf{b}_i \doteq \{b_{i1},...,b_{i|\mathbf{b}_i|}\}$
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Algorithm will iterates between these two phases until no changes to the information vectors occur anymore.



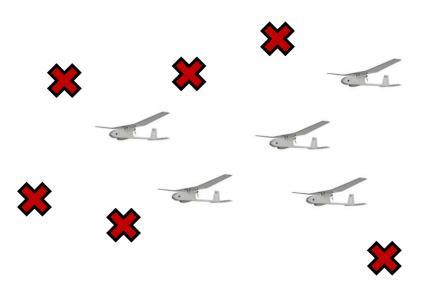
$$\sum_{i=1}^{N_a} \left( \sum_{j=1}^{N_t} c_{ij}(\tau_{ij}(\mathbf{p}_i(\mathbf{x_i}))) x_{ij} \right) \to \max$$

$$\sum_{j=1}^{N_t} x_{ij} \le L_t, \ \forall i \in \mathcal{I}$$

$$\sum_{j=1}^{N_a} x_{ij} \le 1, \ \forall j \in \mathcal{J}$$

$$x_{ij} \in \{0, 1\}, \forall (i, j) \in \mathcal{I} \times \mathcal{J}$$

i	$Task_1$	$Task_2$	$Task_k$	$Task_{N_t}$	Values
Bandle	✓		✓		
Path					
Time					



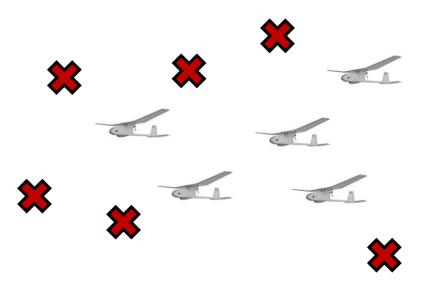
$$\sum_{i=1}^{N_a} \left( \sum_{j=1}^{N_t} c_{ij}(\tau_{ij}(\mathbf{p}_i(\mathbf{x_i}))) x_{ij} \right) \to \max$$

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$$x_{ij} \in \{0, 1\}, \forall (i, j) \in \mathcal{I} \times \mathcal{J}$$

i	$Task_1$	$Task_2$	$Task_k$	$Task_{N_t}$	Values
Bandle	1		2		$b_i = [b_{i1}, b_{i2}]$
Path					
Time					



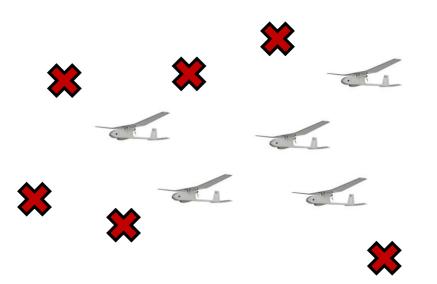
$$\sum_{i=1}^{N_a} \left( \sum_{j=1}^{N_t} c_{ij}(\tau_{ij}(\mathbf{p}_i(\mathbf{x_i}))) x_{ij} \right) \to \max$$

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i	$Task_1$	$Task_2$	$Task_k$	$Task_{N_t}$	Values		
Bandle	1		2		$b_i = [b_{i1}, b_{i2}]$		
Path	2		1		$p_i = [p_{i1}, p_{i2}]$		
Time							



$$\sum_{i=1}^{N_a} \left( \sum_{j=1}^{N_t} c_{ij}(\tau_{ij}(\mathbf{p}_i(\mathbf{x_i}))) x_{ij} \right) \to \max$$

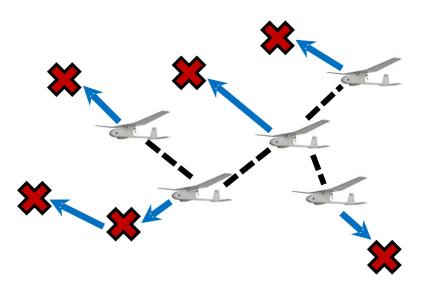
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	$v_{J} = ( \ ) \ ) \ ( \ ) \ v_{J} $						
i	$Task_1$	$Task_2$	$Task_k$	$Task_{N_t}$	Values		
Bandle	1		2		$b_i = [b_{i1}, b_{i2}]$		
Path	2		1		$p_i = [p_{i1}, p_{i2}]$		
Time	20		10		$\tau_i = [\tau_{i1}, \tau_{i2}]$		

# **Agent Information**



$$\sum_{i=1}^{N_a} \left( \sum_{j=1}^{N_t} c_{ij}(\tau_{ij}(\mathbf{p}_i(\mathbf{x_i}))) x_{ij} \right) \to \max$$

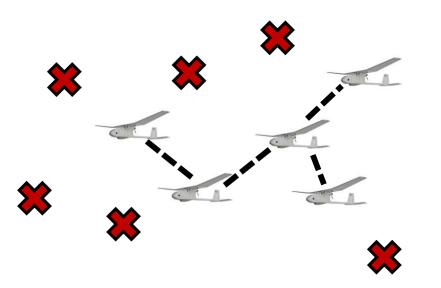
$$\sum_{j=1}^{N_t} x_{ij} \le L_t, \ \forall i \in \mathcal{I}$$

$$\sum_{j=1}^{N_a} x_{ij} \le 1, \ \forall j \in \mathcal{J}$$

$$x_{ij} \in \{0, 1\}, \forall (i, j) \in \mathcal{I} \times \mathcal{J}$$

i	$Task_1$	$Task_2$	$Task_k$	$Task_{N_t}$	Values
Bandle	1		2		$b_i = [b_{i1}, b_{i2}]$
Path	2		1		$p_i = [p_{i1}, p_{i2}]$
Time	20		10		$\tau_i = [\tau_{i1}, \tau_{i2}]$

# **Agent Information**



$$\sum_{i=1}^{N_a} \left( \sum_{j=1}^{N_t} c_{ij}(\tau_{ij}(\mathbf{p}_i(\mathbf{x_i}))) x_{ij} \right) \to \max$$

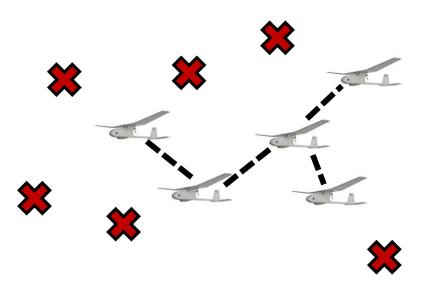
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$$\sum_{i=1}^{N_a} x_{ij} \le 1, \ \forall j \in \mathcal{J}$$

$$x_{ij} \in \{0, 1\}, \forall (i, j) \in \mathcal{I} \times \mathcal{J}$$

i  $Task_1$   $Task_2$   $Task_k$   $Task_{N_t}$  Values  $Winning\ Agent$  2 4 i k  $z_i = [z_{21}, z_{42}, z_{ik}, z_{kN_t}]$   $Winning\ Bids$ 

# **Agent Information**



$$\sum_{i=1}^{N_a} \left( \sum_{j=1}^{N_t} c_{ij}(\tau_{ij}(\mathbf{p}_i(\mathbf{x_i}))) x_{ij} \right) \to \max$$

$$\sum_{j=1}^{N_t} x_{ij} \le L_t, \ \forall i \in \mathcal{I}$$

$$\sum_{i=1}^{N_a} x_{ij} \le 1, \ \forall j \in \mathcal{J}$$

$$x_{ij} \in \{0, 1\}, \forall (i, j) \in \mathcal{I} \times \mathcal{J}$$

i	$Task_1$	$Task_2$	$Task_k$	$Task_{N_t}$	Values
Winning Agent	i	4	i	k	$z_i = [z_{i1}, z_{42}, z_{ik}, z_{kN_t}]$
WinningBids	9	5	8	7	$y_i = [y_{i1}, y_{42}, y_{ik}, y_{kN_t}]$

i	$Task_1$	$Task_2$	$Task_k$	$Task_{N_t}$	Values
Bandle	1		2		$b_i = [b_{i1}, b_{i2}]$
Path	2		1		$p_i = [p_{i1}, p_{i2}]$
Time	20		10		$\tau_i = [\tau_{i1}, \tau_{i2}]$

i	$Task_1$	$Task_2$	$Task_k$	$Task_{N_t}$	Values
Winning Agent	i	4	i	k	$z_i = [z_{i1}, z_{42}, z_{ik}, z_{kN_t}]$
WinningBids	9	5	8	7	$y_i = [y_{i1}, y_{42}, y_{ik}, y_{kN_t}]$

i	$Task_1$	$Task_2$	$Task_k$	$Task_{N_t}$	Values
Bandle	1	¥	2		$b_i = [b_{i1}, b_{i2}]$
Path	2		1		$p_i = [p_{i1}, p_{i2}]$
Time	20		10		$\tau_i = [\tau_{i1}, \tau_{i2}]$

i	$Task_1$	$Task_2$	$Task_k$	$Task_{N_t}$	Values
Winning Agent	i	4	i	k	$z_i = [z_{i1}, z_{42}, z_{ik}, z_{kN_t}]$
WinningBids	9	5	8	7	$y_i = [y_{i1}, y_{42}, y_{ik}, y_{kN_t}]$

i	$Task_1$	$Task_2$	$Task_k$	$Task_{N_t}$	Values
Bandle	1	¥	2		$b_i = [b_{i1}, b_{i2}]$
Path	2		1		$p_i = [p_{i1}, p_{i2}]$
Time	20		10		$\tau_i = [\tau_{i1}, \tau_{i2}]$

Calculate a score  $c_{ij} = c_{i2}$  and compare with current

i	$Task_1$	$Task_2$	$Task_k$	$Task_{N_t}$	Values
Winning Agent	i	4	i	k	$z_i = [z_{i1}, z_{42}, z_{ik}, z_{kN_t}]$
WinningBids	9	5	8	7	$y_i = [y_{i1}, y_{42}, y_{ik}, y_{kN_t}]$

i	$Task_1$	$Task_2$	$Task_k$	$Task_{N_t}$	Values
Bandle	1	¥	2		$b_i = [b_{i1}, b_{i2}]$
Path	2		1		$p_i = [p_{i1}, p_{i2}]$
Time	20		10		$ au_i = [ au_{i1},  au_{i2}]$

Calculate a score  $c'_{ij} = c_{i2}$  and compare with current

i	$Task_1$	$Task_2$	$Task_k$	$Task_{N_t}$	Values
Winning Agent	i	4	i	k	$z_i = [z_{i1}, z_{42}, z_{ik}, z_{kN_t}]$
WinningBids	9	5	8	7	$y_i = [y_{i1}, y_{42}, y_{ik}, y_{kN_t}]$

i	$Task_1$	$Task_2$	$Task_k$	$Task_{N_t}$	Values
Bandle	1	¥	2		$b_i = [b_{i1}, b_{i2}]$
Path	2		1		$p_i = [p_{i1}, p_{i2}]$
Time	20		10		$\tau_i = [\tau_{i1}, \tau_{i2}]$

Calculate a score  $c_{ij} = c_{i2}$  and compare with current

i	$Task_1$	$Task_2$	$Task_k$	$Task_{N_t}$	Values
Winning Agent	i	4	i	k	$z_i = [z_{i1}, z_{42}, z_{ik}, z_{kN_t}]$
WinningBids	9	5	8	7	$y_i = [y_{i1}, y_{42}, y_{ik}, y_{kN_t}]$

i	$Task_1$	$Task_2$	$Task_k$	$Task_{N_t}$	Values
Bandle	1	¥	2		$b_i = [b_{i1}, b_{i2}]$
Path	2		1		$p_i = [p_{i1}, p_{i2}]$
Time	20		10		$\tau_i = [\tau_{i1}, \tau_{i2}]$

To calculate best score for task j, we "insert" the task in some location  $n_j$ 

i	$Task_1$	$Task_2$	$Task_k$	$Task_{N_t}$	Values
Bandle	1	¥	2		$b_i = [b_{i1}, b_{i2}]$
Path	2		1		$p_i = [p_{i1}, p_{i2}]$
Time	20		10		$\tau_i = [\tau_{i1}, \tau_{i2}]$

To calculate best score for task j, we first "insert" the task in some location  $n_j$ 

And new path becomes  $(\mathbf{p}_i \oplus_{n_j} j)$ 

i	$Task_1$	$Task_2$	$Task_k$	$Task_{N_t}$	Values
Bandle	1	¥	2	2	
Path	2		1		$p_i = [p_{i1}, p_{i2}]$
Time	20		10		$\tau_i = [\tau_{i1}, \tau_{i2}]$

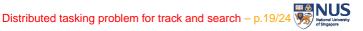
To calculate best score for task j, we first "insert" the task in some location  $n_j$ 

And new path becomes  $(\mathbf{p}_i \oplus_{n_j} j)$  and second calculate the optimal execution time for this new path:

$$\tau_{ij}^*(\mathbf{p}_i \oplus_{n_j} j) = \max_{\tau_{ij} \in [0,\infty)} c_j(\tau_{ij})$$

$$subject \ to:$$

$$\tau_{ik}^*(\mathbf{p}_i \oplus_{n_j} j) = \tau_{ik}^*, \forall k \in \mathbf{p}_i$$



i	$Task_1$	$Task_2$	$Task_k$	$Task_{N_t}$	Values
Bandle	1	¥	2		$b_i = [b_{i1}, b_{i2}]$
Path	2		1		$p_i = [p_{i1}, p_{i2}]$
Time	20		10		$\tau_i = [\tau_{i1}, \tau_{i2}]$

 $\Longrightarrow$  optimal score for the task at location  $n_j$  is  $c_j(\tau_{ij}^*(\mathbf{p}_i \oplus_{n_j} j))$ .

$$\tau_{ij}^*(\mathbf{p}_i \oplus_{n_j} j) = \max_{\tau_{ij} \in [0,\infty)} c_j(\tau_{ij})$$

$$subject \ to:$$

$$\tau_{ik}^*(\mathbf{p}_i \oplus_{n_j} j) = \tau_{ik}^*, \forall k \in \mathbf{p}_i$$

_						
	i	$Task_1$	$Task_2$	$Task_k$	$Task_{N_t}$	Values
	Bandle	1	*	2		$b_i = [b_{i1}, b_{i2}]$
	Path	2		1		$p_i = [p_{i1}, p_{i2}]$
	Time	20		10		$\tau_i = [\tau_{i1}, \tau_{i2}]$

 $<sup>\</sup>Longrightarrow$  optimal score for the task at location  $n_j$  is  $c_j(\tau_{ij}^*(\mathbf{p}_i \oplus_{n_j} j))$ .

 $<sup>\</sup>Longrightarrow$  And optimal location  $n_j^*$  is then given by  $n_j^* = \max_{n_j} c_j (\tau_{ik}^*(\mathbf{p}_i \oplus_{n_j} j))$ 

i	$Task_1$	$Task_2$	$Task_k$	$Task_{N_t}$	Values
Bandle	1	¥	2		$b_i = [b_{i1}, b_{i2}]$
Path	2		1		$p_i = [p_{i1}, p_{i2}]$
Time	20		10		$\tau_i = [\tau_{i1}, \tau_{i2}]$

- $\implies$  optimal score for the task at location  $n_j$  is  $c_j(\tau_{ij}^*(\mathbf{p}_i \oplus_{n_j} j))$ .  $\implies$  And optimal location  $n_j^*$  is then given by  $n_j^* = \max_{n_j} c_j(\tau_{ik}^*(\mathbf{p}_i \oplus_{n_j} j))$   $\implies$  Final score for task j is  $c_{ij}(\mathbf{p}_i) = c_j(\tau_{ij}^*(\mathbf{p}_i \oplus_{n_j^*} j))$

i	$Task_1$	$Task_2$	$Task_k$	$Task_{N_t}$	Values
Bandle	1	<b>→</b> ⅓	2	¥.	$b_i = [b_{i1}, b_{i2}]$
Path	2		1		$p_i = [p_{i1}, p_{i2}]$
Time	20		10		$\tau_i = [\tau_{i1}, \tau_{i2}]$

- $\Longrightarrow$  optimal score for the task at location  $n_j$  is  $c_j(\tau_{ij}^*(\mathbf{p}_i \oplus_{n_j} j))$ .
- $\Longrightarrow$  And optimal location  $n_j^*$  is then given by  $n_j^* = \max_{n_j} c_j (\tau_{ik}^*(\mathbf{p}_i \oplus_{n_j} j))$
- $\Longrightarrow$  Final score for task j is  $c_{ij}(\mathbf{p}_i) = c_j(\tau_{ij}^*(\mathbf{p}_i \oplus_{n_j^*} j))$
- Final step is to select the highest scoring task to add to the bundle
- $j^* = \max_{j \notin \mathbf{p}_i} c_{ij}(\mathbf{p}_i) h_{ij}$ , where  $h_{ij} = \mathbf{I}(c_{ij}(\mathbf{p}_i) > y_{ij})$  the indicator function

i	$Task_1$	$Task_2$	$Task_k$	$Task_{N_t}$	Values
Bandle	1	¥	2		$b_i = [b_{i1}, b_{i2}]$
Path	2		1		$p_i = [p_{i1}, p_{i2}]$
Time	20		10		$\tau_i = [\tau_{i1}, \tau_{i2}]$

i	$Task_1$	$Task_2$	$Task_k$	$Task_{N_t}$	Values
Winning Agent	i	4	i	k	$z_i = [z_{i1}, z_{42}, z_{ik}, z_{kN_t}]$
WinningBids	9	5	8	7	$y_i = [y_{i1}, y_{42}, y_{ik}, y_{kN_t}]$

i	$Task_1$	$Task_2$	$Task_k$	$Task_{N_t}$	Values
Bandle	1	<b>√</b>	2		$\mathbf{b}_i \leftarrow (\mathbf{b}_i \oplus_{end} j^*)$
Path	2		1		$p_i = [p_{i1}, p_{i2}]$
Time	20		10		$\tau_i = [\tau_{i1}, \tau_{i2}]$

i	$Task_1$	$Task_2$	$Task_k$	$Task_{N_t}$	Values
Winning Agent	i	4	i	k	$z_i = [z_{i1}, z_{42}, z_{ik}, z_{kN_t}]$
WinningBids	9	5	8	7	$y_i = [y_{i1}, y_{42}, y_{ik}, y_{kN_t}]$

i	$Task_1$	$Task_2$	$Task_k$	$Task_{N_t}$	Values
Bandle	1	<b>√</b>	2		$\mathbf{b}_i \leftarrow (\mathbf{b}_i \oplus_{end} j^*)$
Path	2		1		$\mathbf{p}_i \leftarrow (\mathbf{p}_i \oplus_{n_{j^*}} j^*)$
Time	20		10		$\tau_i = [\tau_{i1}, \tau_{i2}]$

i	$Task_1$	$Task_2$	$Task_k$	$Task_{N_t}$	Values
Winning Agent	i	4	i	k	$z_i = [z_{i1}, z_{42}, z_{ik}, z_{kN_t}]$
WinningBids	9	5	8	7	$y_i = [y_{i1}, y_{42}, y_{ik}, y_{kN_t}]$

i	$Task_1$	$Task_2$	$Task_k$	$Task_{N_t}$	Values
Bandle	1		2		$\mathbf{b}_i \leftarrow (\mathbf{b}_i \oplus_{end} j^*)$
	✓	<b>√</b>	✓		
Path					$\mathbf{p} \cdot \leftarrow (\mathbf{p} \cdot \mathbf{p} \cdot \mathbf{p} \cdot \mathbf{p}^*)$
	2		1		$\mathbf{p}_i \leftarrow (\mathbf{p}_i \oplus_{n_{j^*}} j^*)$
Time					- / (- · · · · · · · · · · · · · · · · · ·
$ $ $ $ $ $ $ $ $ $ $ $ $ $ $ $ $ $	20		10		$\tau_i \leftarrow (\tau_i \oplus_{n_{j^*}} \tau_{ij^*}^* (\mathbf{p}_i \oplus_{n_{j^*}} j^*)$

i	$Task_1$	$Task_2$	$Task_k$	$Task_{N_t}$	Values
Winning Agent	i	4	i	k	$z_i = [z_{i1}, z_{42}, z_{ik}, z_{kN_t}]$
WinningBids	9	5	8	7	$y_i = [y_{i1}, y_{42}, y_{ik}, y_{kN_t}]$

i	$Task_1$	$Task_2$	$Task_k$	$Task_{N_t}$	Values
Bandle	1		2		$\mathbf{b}_i \leftarrow (\mathbf{b}_i \oplus_{end} j^*)$
	<b>√</b>	<b>√</b>	<b>√</b>		
Path					$\mathbf{p}_i \leftarrow (\mathbf{p}_i \oplus_{n_{j^*}} j^*)$
	2		1		J T
Time					$\tau_i \leftarrow (\tau_i \oplus_{n_{j^*}} \tau_{ij^*}^* (\mathbf{p}_i \oplus_{n_{j^*}} j^*)$
	20		10		

i	$Task_1$	$Task_2$	$Task_k$	$Task_{N_t}$	Values
Winning Agent	i	i	i	k	$z_i = [z_{i1}, z_{i2}, \ldots]$
WinningBids	9	$c_{ij*}(\mathbf{p}_i)$	8	7	$y_i = [y_{i1},  extbf{y_{i2}},]$

i	$Task_1$	$Task_2$	$Task_k$	$Task_{N_t}$	Values
Bandle	1		2		$\mathbf{b}_i \leftarrow (\mathbf{b}_i \oplus_{end} j^*)$
Danac	$\checkmark$	<b>√</b>	<b>√</b>		$\mathbf{D}_{i} \leftarrow (\mathbf{D}_{i} \cup end \mathbf{J}^{-})$
Path					$\mathbf{p}_i \leftarrow (\mathbf{p}_i \oplus_{n_{j^*}} j^*)$
1 acre	2		1		$\mathbf{p}_i \leftarrow (\mathbf{p}_i \cup n_{j^*} J)$
Time					$\tau \leftarrow (\tau \cdot \Phi - \tau^*) (\mathbf{p} \cdot \Phi - i^*)$
1 11116	20		10		$\tau_i \leftarrow (\tau_i \oplus_{n_{j^*}} \tau_{ij^*}^* (\mathbf{p}_i \oplus_{n_{j^*}} j^*)$

Bundle recursion continues until  $|\mathbf{b}_i| = L_t$  or  $h_{ij} = 0$  for all  $j \notin \mathbf{p}_i$ 

i	$Task_1$	$Task_2$	$Task_k$	$Task_{N_t}$	Values
Winning Agent	i	i	i	k	$z_i = [z_{i1},  extbf{z_{i2}},]$
WinningBids	9	$c_{ij*}(\mathbf{p}_i)$	8	7	$y_i = [y_{i1}, \textcolor{red}{y_{i2}},]$

#### Consensus

i, (receiver)	$Task_1$	$Task_2$	$Task_k$	$Task_{N_t}$	Values
Winning Agent					$z_i = [z_{i1}, z_{i2},]$
WinningBids					$y_i = [y_{i1}, y_{i2}, \dots]$

 $Update: z_{ij} = z_{kj}, \quad y_{ij} = y_{kj}$ 

Reset:  $z_{ij} = \emptyset$ ,  $y_{ij} = 0$ 

Leave:  $z_{ij} = z_{ij}, \quad y_{ij} = y_{ij}$ 

k, (sender)	$Task_1$	$Task_2$	$Task_k$	$Task_{N_t}$	Values
Winning Agent					$z_k = [z_{k1}, z_{k2}, \ldots]$
WinningBids					$y_k = [y_{k1}, y_{k2}, \dots]$

#### **Decision Rules**

Agent $k$ thinks $z_{kj}$ is	Agent $i$ thinks $z_{ij}$ is	Receiver Action
k	i	if $y_{kj} > y_{ij} \rightarrow update$
k	k	update
k	$m \not \in \{i,k\}$	$if \ s_{km} > s_{im} \ \text{or} \ y_{kj} > y_{ij} \rightarrow update$
k	none	update

$$s_{ik} = \begin{cases} \tau_r(i.e. \ message \ reception \ time), & if \ g_{ik} = 1; \\ \max\{s_{mk} | m \in \mathcal{I}, g_{im} = 1\}, & otherwise \end{cases}$$

Agent $k$ thinks $z_{kj}$ is	Agent $i$ thinks $z_{ij}$ is	Receiver Action
i	i	leave
i	k	reset
i	$m \not \in \{i,k\}$	$if \ s_{km} > s_{im} \to reset$
i	none	leave

# **Decision Rules**

$m \not \in \{i,k\}$	i	if $s_{km} > s_{im}$ and $y_{kj} > y_{ij} \rightarrow update$
$m \not \in \{i,k\}$	k	$if \ s_{km} > s_{im} \rightarrow update$ $else \rightarrow reset$
$m \not \in \{i,k\}$	m	$s_{km} > s_{im} \rightarrow update$
$m  ot\in \{i,k\}$	$n \not \in \{i,k,m\}$	$if \ s_{km} > s_{im} \ and \ s_{kn} > s_{in} \rightarrow update$ $if \ s_{km} > s_{im} \ and \ y_{kj} > y_{ij} \rightarrow update$ $if \ s_{kn} > s_{in} \ and \ s_{im} > s_{km} \rightarrow reset$
$m \not \in \{i,k\}$	none	$if \ s_{km} > s_{im} \rightarrow update$

Agent $k$ thinks $z_{kj}$ is	Agent $i$ thinks $z_{ij}$ is	Receiver Action	
none	i	leave	
none	k	update	
none	$m \not \in \{i,k\}$	$if s_{km} > s_{im} \rightarrow update$	
none	none	leave	

#### **Decision Rules**

i, (receiver)	$Task_1$	$Task_2$	$Task_k$	$Task_{N_t}$	Values
Winning Agent					$z_i = [z_{i1}, z_{i2}, \ldots]$
WinningBids					$y_i = [y_{i1}, y_{i2}, \dots]$

 $Update: z_{ij} = z_{kj}, \quad y_{ij} = y_{kj}$ 

 $Reset: z_{ij} = \emptyset, \quad y_{ij} = 0$ 

 $\underline{Leave}: \ z_{ij} = z_{ij}, \ y_{ij} = y_{ij}$ 

k, (sender)	$Task_1$	$Task_2$	$Task_k$	$Task_{N_t}$	Values
$Winning\ Agent$					$z_k = [z_{k1}, z_{k2}, \ldots]$
WinningBids					$y_k = [y_{k1}, y_{k2},]$

Calculate marginal score for all tasks

$$c_{ij}(\mathbf{p}_i) = \begin{cases} 0, & if \ j \in \mathbf{p}_i; \\ \max_{n \le l_b} S_{path}(\mathbf{p}_i \oplus_n j) - S_{path}(\mathbf{p}_i), & otherwise \end{cases}$$

- Determine which tasks are winnable
- Select the index of the best eligible task,  $j^*$ , and select best location in the plan to insert the task,  $n_j^*$
- If  $c_{ij^*} \leq 0$ , then return. otherwise, continue
- Update agent information
- Update shared information vectors
- if  $l_b = L_t$ , then return, otherwise, go to 1.

- Calculate marginal score for all tasks
- Determine which tasks are winnable

$$h_{ij} = \mathbf{I}(c_{ij}(\mathbf{p}_i) > y_{ij}), \forall j \in \mathcal{J}$$

- Select the index of the best eligible task,  $j^*$ , and select best location in the plan to insert the task,  $n_j^*$
- If  $c_{ij^*} \leq 0$ , then return. otherwise, continue
- Update agent information
- Update shared information vectors
- if  $l_b = L_t$ , then return, otherwise, go to 1.

- Calculate marginal score for all tasks
- Determine which tasks are winnable
- Select the index of the best eligible task,  $j^*$ , and select best location in the plan to insert the task,  $n_j^*$

$$j^* = \max_{j \in \mathcal{J}} c_{ij} h_{ij}$$
$$n_j^* = \max_{n \in \{0, \dots, l_b\}} S_{path}(\mathbf{p}_i \oplus_n j^*)$$

- If  $c_{ij^*} \leq 0$ , then return. otherwise, continue
- Update agent information
- Update shared information vectors
- if  $l_b = L_t$ , then return, otherwise, go to 1.

- Calculate marginal score for all tasks
- Determine which tasks are winnable
- Select the index of the best eligible task,  $j^*$ , and select best location in the plan to insert the task,  $n_j^*$
- If  $c_{ij^*} \leq 0$ , then return. otherwise, continue
- Update agent information
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- if  $l_b = L_t$ , then return, otherwise, go to 1.

- Calculate marginal score for all tasks
- Determine which tasks are winnable
- Select the index of the best eligible task,  $j^*$ , and select best location in the plan to insert the task,  $n_j^*$
- If  $c_{ij^*} \leq 0$ , then return. otherwise, continue
- Update agent information

$$\mathbf{b}_i \leftarrow (\mathbf{b}_i \oplus_{l_b} j^*)$$

$$\mathbf{p}_i \leftarrow (\mathbf{p}_i \oplus_{n_j^*} j^*)$$

- Update shared information vectors
- if  $l_b = L_t$ , then return, otherwise, go to 1.

- Calculate marginal score for all tasks
- Determine which tasks are winnable
- Select the index of the best eligible task,  $j^*$ , and select best location in the plan to insert the task,  $n_j^*$
- If  $c_{ij^*} \leq 0$ , then return. otherwise, continue
- Update agent information
- Update shared information vectors

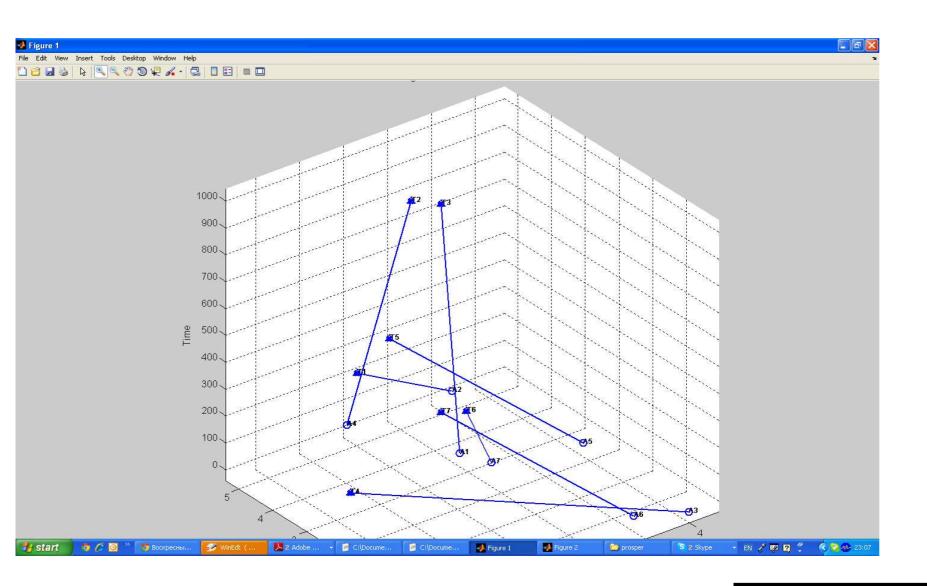
$$y_{i(j^*)} = c_{i(j^*)}$$

$$z_{i(j^*)} = i$$

• if  $l_b = L_t$ , then return, otherwise, go to 1.

- Calculate marginal score for all tasks
- Determine which tasks are winnable
- Select the index of the best eligible task,  $j^*$ , and select best location in the plan to insert the task,  $n_j^*$
- If  $c_{ij^*} \leq 0$ , then return. otherwise, continue
- Update agent information
- Update shared information vectors
- $\blacksquare$  if  $l_b = L_t$ , then return, otherwise, go to 1.

#### Simulation



#### The end

#### Thank you!