Software Architectures

Lecture 8. P2P architecture

Professor: Yushan (Michael) Sun Fall 2020

Contents:

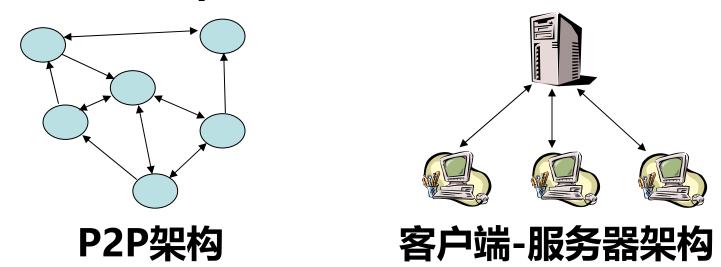
- 1. P2P架构引言
- 2. <u>带有检索服务器的非结构化P2P架构-第1</u> 代P2P架构
- 3. 非结构化纯P2P架构-第2代P2P架构
- 4. 非结构化层次P2P架构-第3代P2P架构
- 5. 结构化P2P架构简介*
- 6. JXTA-P2P 协议

Introduction to P2P Architecture

P2P架构引言

1. Introduction to P2P

P2P (Peer to peer) 思想:



- 1.P2P使得任何的网络感知设备可以为另外一个网络感知设备提供服务。P2P enables any network-aware device to provide services to another network-aware device
- 2. P2P网络的一个节点既是客户端,又是服务器。A peer in P2P network acts as both a client and a server

1. Introduction to P2P

P2P的概念:

Peer-to-Peer computing is described as the sharing of computer resources and services by direct exchange between systems.

- These resources include the exchange of
 - a) information,
 - b) processing cycles,
 - c) cache storage, and
 - d) disk storage for files.

1. Introduction to P2P

Advantage of P2P Architecture

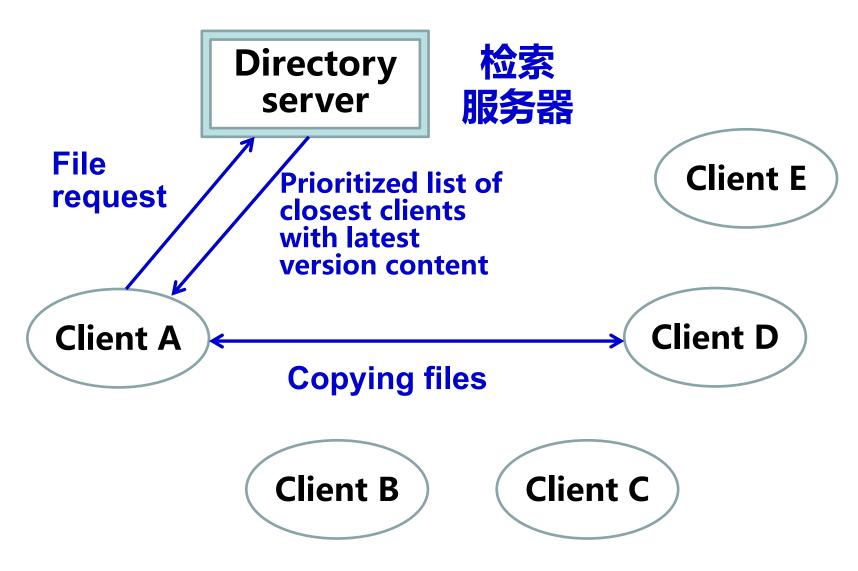
- a) 便宜的设备协作、利用集体力量. It allows economical clients (desktop computers, etc) to take advantage of their collective power to benefit the entire enterprise
- b) 节点既是客户端又是服务器. Clients in a P2P network are also servers
- c) 无传统意义下的服务器超载现象. The load on servers in the traditional sense has reduced



Unstructured P2P Architecture with Directory Server- 1st Generation P2P

第一代非结构化P2P架构-带有检索服务器

2. Unstructured P2P Architecture with Directory Server



P2P Architecture with directory server

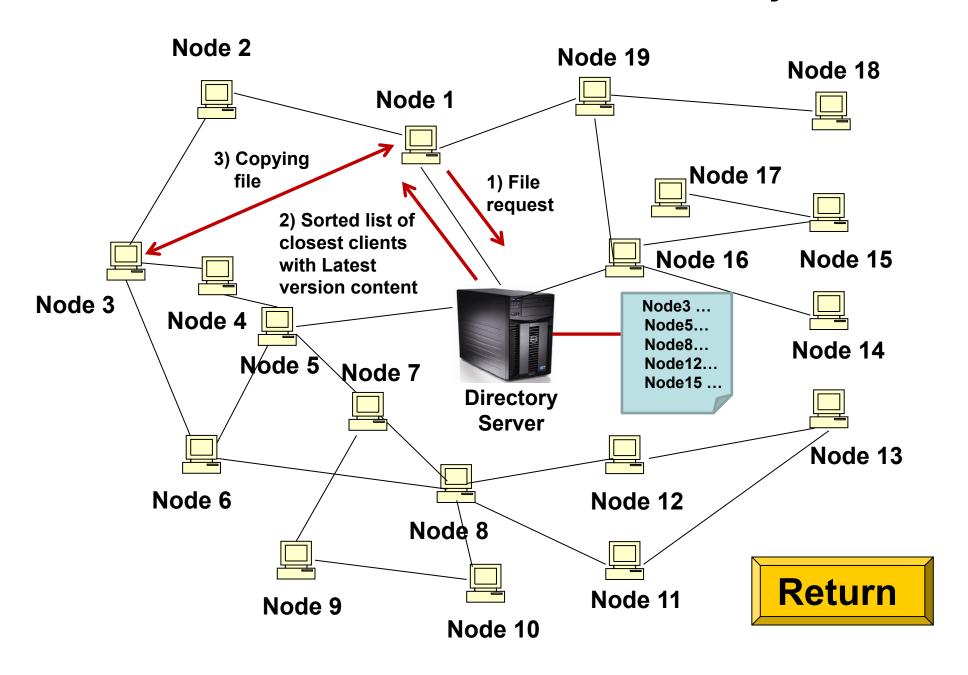
2. Unstructured P2P Architecture with Directory Server

Client A怎样查找其它有用的节点?

The interaction between clients by referencing a Directory Server

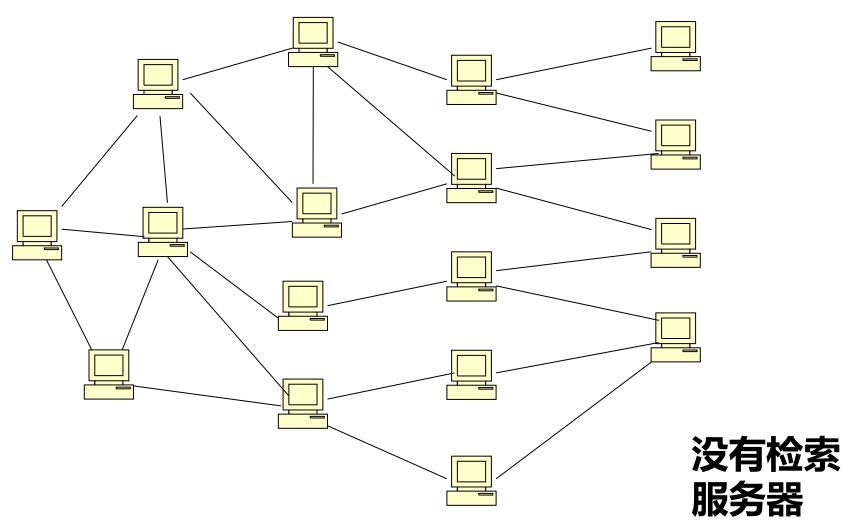
- a) 发送请求. Send request: Client A issues a request to the Directory server
- b) 目录服务器查找对等节点. Find peers. The Directory server then uses the lists it keep to find the peer that contains the content that Client A interests in and tells Client A
- c) 开始互动. Interactions. Client A can then directly interact with that client, in this case Client D which services his request.

2. Unstructured P2P Architecture with Directory Server



Pure P2P Architecture – 2nd Generation P2P

第二代非结构化P2P架构-纯P2P架构



Pure P2P architecture

纯P2P架构的性质(characteristics of pure P2P Architecture)

- a) 节点既是客户端又是服务器. Peers can act as clients and servers and have the same capability as its neighbors.
- b) 没有中心服务器. It has no central servers. It has every node as a Peer and has no central router.
- c) 两种路由结构. There are two routing structures,
 - 1. 分布式目录. One is a distributed catalogue;
 - 2. 直接发送消息. The other direct messaging.

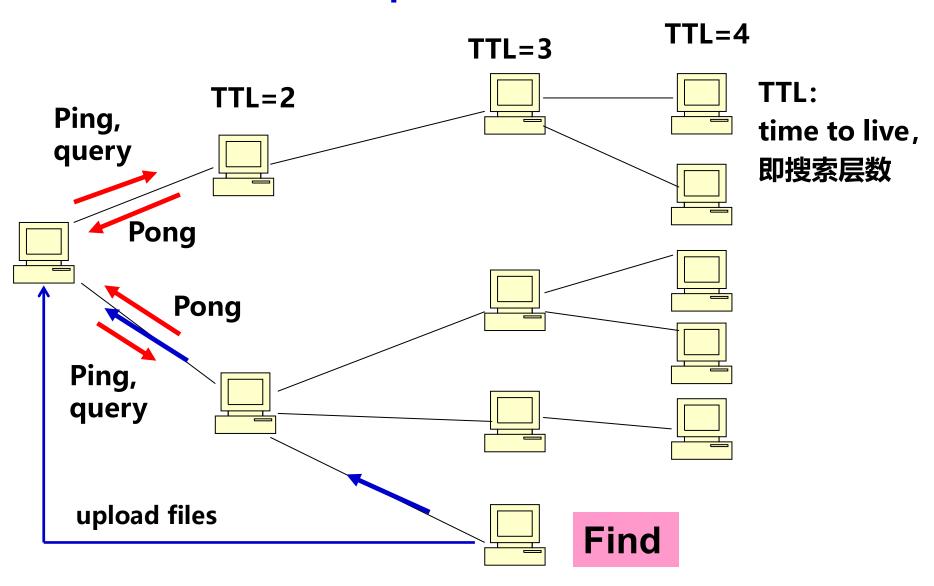
Gnutella: Example of Pure P2P Architecture

- Gnutella简史
- Gnutella, a Pure P2P Model, is a file sharing application and protocol:
 - the end hosts join Gnutella by connecting to existing end hosts already on the Gnutella.
 - by Justin Frankel, published in 2000/03
 - Gnutella has become a Protocol now

Gnutella资源查找方式

To facilitate file sharing, messages are sent between end hosts.

- Queries for files are broadcasted on the overlay network (覆盖网络(应用层网络)), and
- Replies are routed back to the host that originally generated the query through the overlay network.

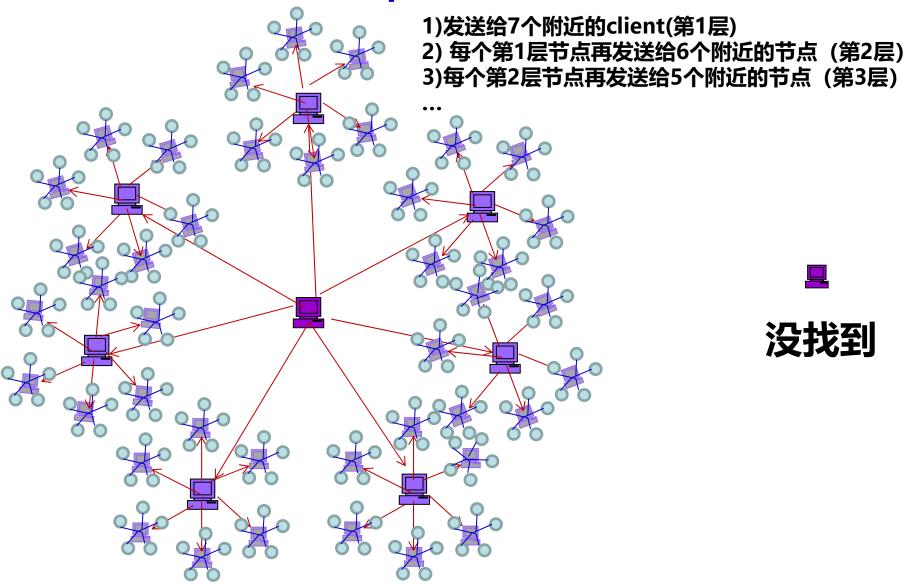


Gnutella网络资源查找原理

查询步骤 (Query steps):

- 1. 连接到P2P网络. Node A is first connected to the network.
- 2. 查找新节点. Ping Node A Pings to the other Nodes (B and C) to discover new nodes on the network.
- 3. 新节点回答. Pong A pong message is sent as a reply to a ping and provides information on a network node, including
 - > IP address,
 - ➤ port number (端口号), and
 - > number of files shared.

- 4. 发送请求消息. A query message is used to search for files shared by other nodes on the network. It contains a query string and a minimum requested link speed.
- 5. 获得回答. A query reply message contains
 - > a list of files which match a given query,
 - > the size of each file, and
 - > the link speed of the responding node.
- 6. 上载文件. A push message is used to upload file to clients behind a firewall who can not download files themselves.



Search is not guaranteed if object is outside the search range!

早期版本的Gnutella的缺点:

- 1) 不稳定的连接. Unstable connectivity of the servants
 - performance management difficult
- 2) 伸缩性问题. Scalability issues: e.g. when TTL=10, will produce huge networks
- 3) <mark>搜索范围不够大</mark>,本应该存在的资源没有被 覆盖

关于Gnutella协议

- Gnutella 采用分布式的协议,每个节点既是服务器又是客户端. Gnutella is a distributed protocol such that the clients become servers and they become clients at the same time. A node
 - is a Client when it is looking for some data
 - is a Sever if it is servicing a request of another node.
- The communication between the nodes is done by the TCP/IP protocol.

Gnutella的性质

- a) 可靠性: If one computer goes down, the network is unaffected, hence more reliable.
- b) 匿名性: Gnutella is anonymous-there is no need to provide a name or e-mail address to use Gnutella
- c) 共享任何种类资源: Gnutella provides a mechanism to share any type of resource: photos, computer program, and movies, etc.
- d) 纯P2P: Gnutella is Pure P2P and is completely decentralized.

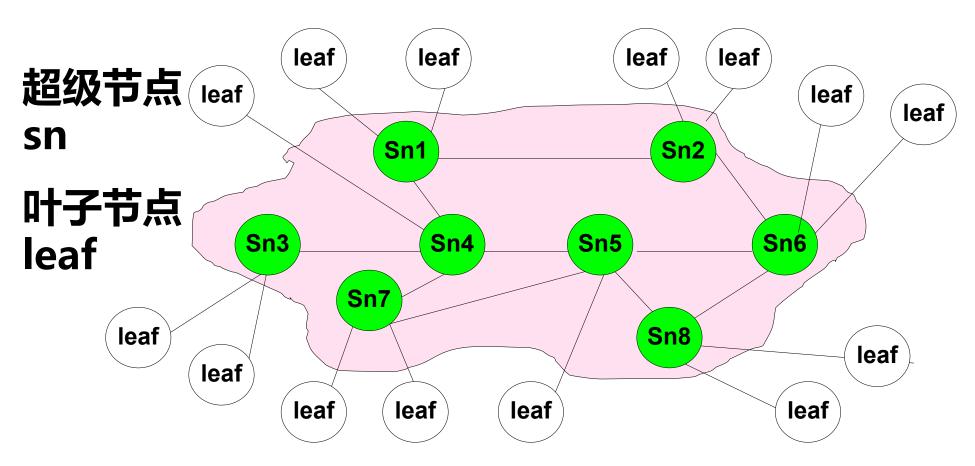
Return

Unstructured Hierarchical P2P Architecture— 3nd Generation P2P

第三代P2P架构-非结构化层次化P2P架构

- 纯P2P Gnutella V0.4存在的问题
- 网络大了、搜索量大增。Large scale P2P networks around the world caused the ever-increasing search traffic.
- 对于大网络,搜索应答率比的越来越低。The search recall ratio (the number of search results: the total number of available copies of the searched object) was significantly decreased as the network scale became larger

- 需要新设计,以便改善性能
- To tackle this problem, Gnutella v0.6 introduced two-layer hierarchy in its architecture as below

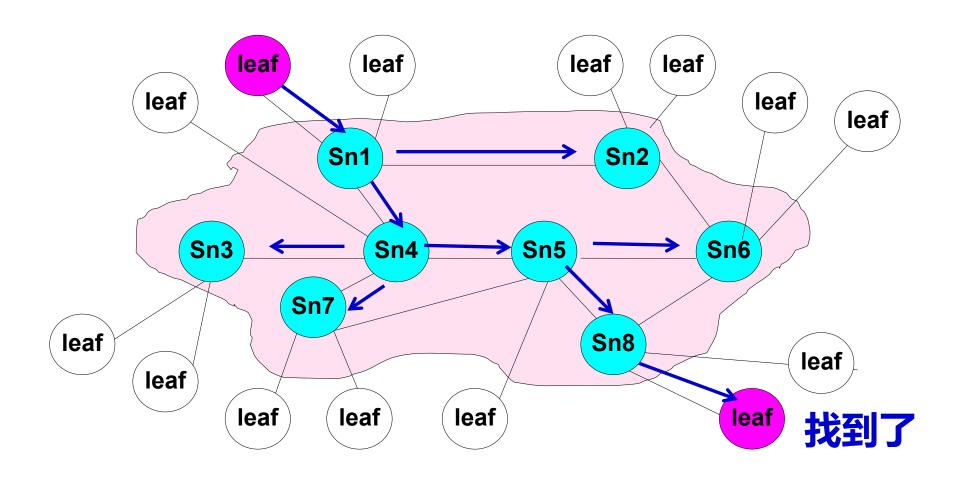


非结构化层次P2P架构-第3代P2P: Gnutella v0.6 network topology

- 架构的解释
- In this architecture, nodes are divided into into leaf nodes and ultrapeer nodes.
- 叶子节点 (A leaf node) : only maintains connection with its own ultrapeer,
- 超级节点 (An ultrapeer):
 - maintains connections to its own leaf nodes and acts as proxy for them
 - connects to the other ultrapeers from the overlay

- 在查找过程中, 叶子节点的责任.
- Leaf nodes are only responsible for
 - initiating lookup requests,
 - receiving correlated lookup response, and
 - responding the requests that they can exactly answer,

- 在查找过程中,超级节点的责任。
- Ultrapeers are responsible for
 - 转发请求到其它超级节点 forwarding lookup requests to other ultrapeers or
 - 转发请求到自己的叶子节点 its own connected leaf nodes, if it knows exactly the leaf node is able to answer the request.
- Ultrapeers can also initiate requests and receive correlated responses.
- At the level of ultrapeer, similar floodingbased mechanism as Gnutella v0.4 is utilized to forward the lookup requests.



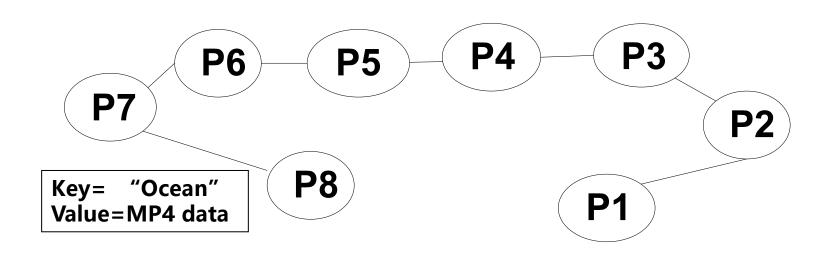
查找过程示意图



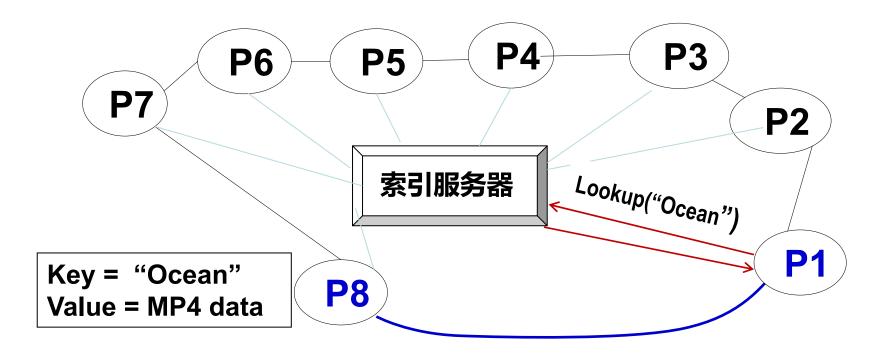
Brief Introduction to Structured P2P Architecture*

结构化P2P软件体系结构简介*

• 问题: P2P网络中的一个节点P1怎样在P2P网络中查找带有所需资源的节点。例如,在如下图所示的P2P网络,怎样查找所需的资源(Key="Ocean", Value = MP4 data) 呢?

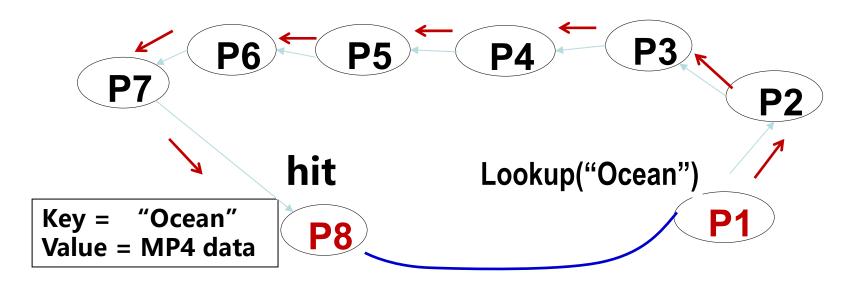


 在第一代无结构化P2P网络中:使用一个中心检索 服务器接收所有的查询,服务器告知去哪下载其所 需要的数据。



非结构化P2P网络Napster资源查找的方式

• 在第2代纯P2P网络中:使用消息洪泛 (message flooding)来定位数据。一个消息被发到系统内每一个节点,直到找到其需要的数据为止。



非结构化P2P网络Gnutella资源查找方式-有可能 查不到资源

- · 结构化P2P拓扑:
- · 全分布式结构化拓扑的P2P网络主要是采用分布式散列表(Distributed Hash Table, DHT)技术来组织网络中的结点,建立拓扑结构。
- · 一个对象的名字或关键词被映射为128位或160 位的散列值。
- · 最经典的案例是Tapestry, Pastry, Chord和 CAN技术。

- · 而DHT的主要思想是: 为网络建立结构
- 用(fileName, machineIP)表示某文件在某计算机上
- 使用SHA-1作为哈希函数
 (fileName, machinelP)←→ (KID, VID)
- 每条文件索引被表示成一个(KID, VID)对, KID称为关键字, VID代表实际存储文件的节点的IP地址(+port)

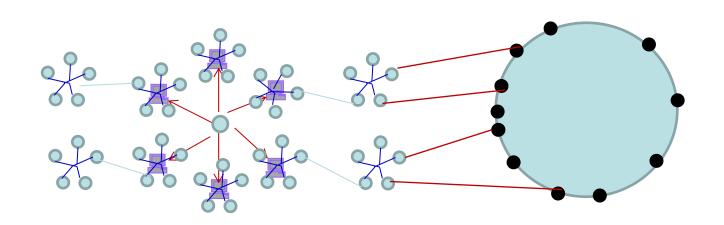
思想:集合{(KID, VID)}组成一张完整的文件索引哈希表。只要输入目标文件的K值,就可以从这张表中查出所有存储该文件的节点地址。

- · Chord算法是全分布式结构化拓扑的一种 实现。
- Chord算法

- Chord算法
- Chord把Node(机器的IP地址+Port)与Key (资源标识)映射到相同的空间。
- · 哈希函数SHA-1会产生一个2¹⁶⁰的整数空间,每项为一个16字节(160bit)的大整数。具体地说:

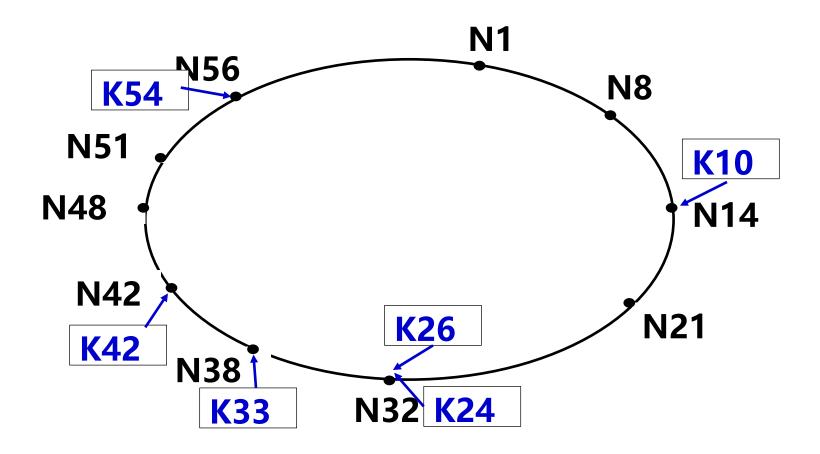
Node←→NID; Key←→KID

这些整数NID和KID首尾相连形成一个按大小顺时针排列的虚拟的环,称之为Chord环。



- Chord的构成:
- 节点ID: NID(node identifier),表示一个物理机器,是m位的一个数字,由节点机器的IP地址通过哈希操作得到;
- 资源ID: KID (key identifiers),表示一个资源,称资源ID (或者key ID),是m位的一个数字,由fileName通过哈希操作得到。

- Chord环(Chord Ring)的构造:将NID和KID分配 到一个大小为2^m的环上,表示资源KID被分配到 NID 上。(注:在这个环上的ID为0-2^{m-1})。
- Chord环上资源的分配。资源被分配到使得NID>= KID的第一个节点(NID)上,如下图所示。
- · 这个节点称为k的后继节点,例如N14是K10的后继节点,是环上从k起顺时针方向的第一个节点,记为successor(k),

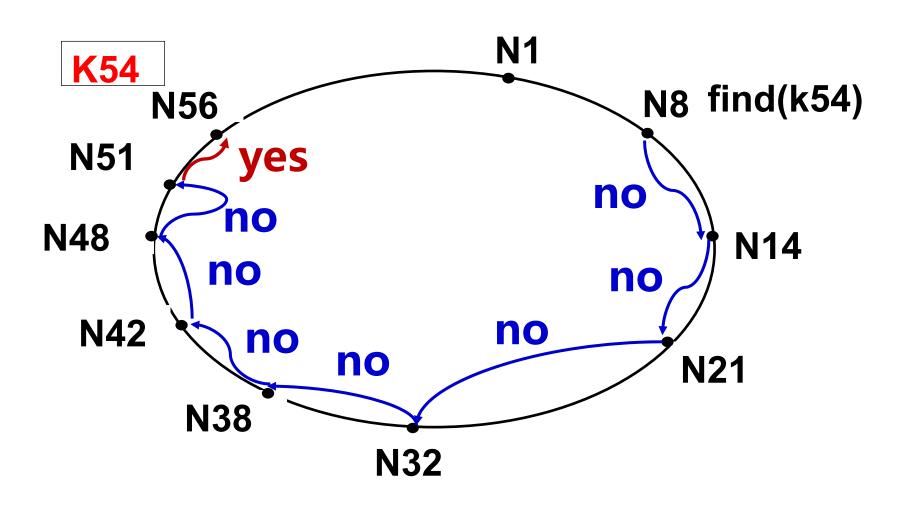


这是一个m=6的环(2⁶=64),其中有10个节点,5个资源,K10的后继节点为N14,也就是说K10被分配给了N14。这样的分配是为了能够有效地查找。

Chord资源查找

- · 资源查找 (定位) 是Chord协议的核心功能。为了便 于理解,我们介绍一个简单的资源定位方法。
- 简单方法:考虑如下场景:节点n寻找KID=id的资源 ,此时节点n首先查询资源是否在下一个节点上(find_successor),即要查看资源k的KID是否在该节 点NID和下一个节点的NID之间。
- 若是,则说明资源k被分配给了下一个节点,若不是则在下一个节点上发起同样的查询,问询下一个点是否有该资源。如此迭代下去。

例如下图



- · 假设节点N8寻找K54这个资源
- N8.find_successor(K54),发现54 ε (8; 14]不成立。
- N14.find successor(K54),发现54 ε (14; 21]不成立。
- N21.find_successor(K54),发现54 ε (21; 32]<mark>不成立</mark>。
- N32.find successor(K54),发现54 ε (32; 38]不成立。
- N38.find_successor(K54),发现54 € (38; 42]<mark>不成立</mark>。
- N42.find successor(K54),发现54 ε (42; 48]不成立。
- N48.find_successor(K54),发现54 ε (48; 51]不成立。
- N51.find_successor(K54),发现54 € (51; 56]成立。
- · 于是得知资源K54在N56这个节点上。

- 缺点: 查找速度为O(N)。
- 有改进的余地。
- Chord实际上使用了可伸缩资源定位的方式来提高效率;算法较为复杂,此课程省略。
- 由于覆盖网络采用了确定性拓扑结构,DHT可以 提供精确的发现。
- 只要目的结点存在于网络中,DHT总能发现它, 发现的准确性得到了保证。

Return

JXTA-P2P Protocol

JXTA Java™ Standard Edition v2.5: Programmers Guide

7. JXTA-P2P Protocol

- What is JXTA Java?
- JXTA™ is a set of open, generalized peer-topeer (P2P) protocols that allow any networked device, such as
 - sensors,
 - Cell phones,
 - computers, etc.,

to communicate and collaborate mutually as peers.

7. JXTA-P2P Protocol

