PG4200: Algorithms And Data Structures

Lesson 12:
Data Compression

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Data

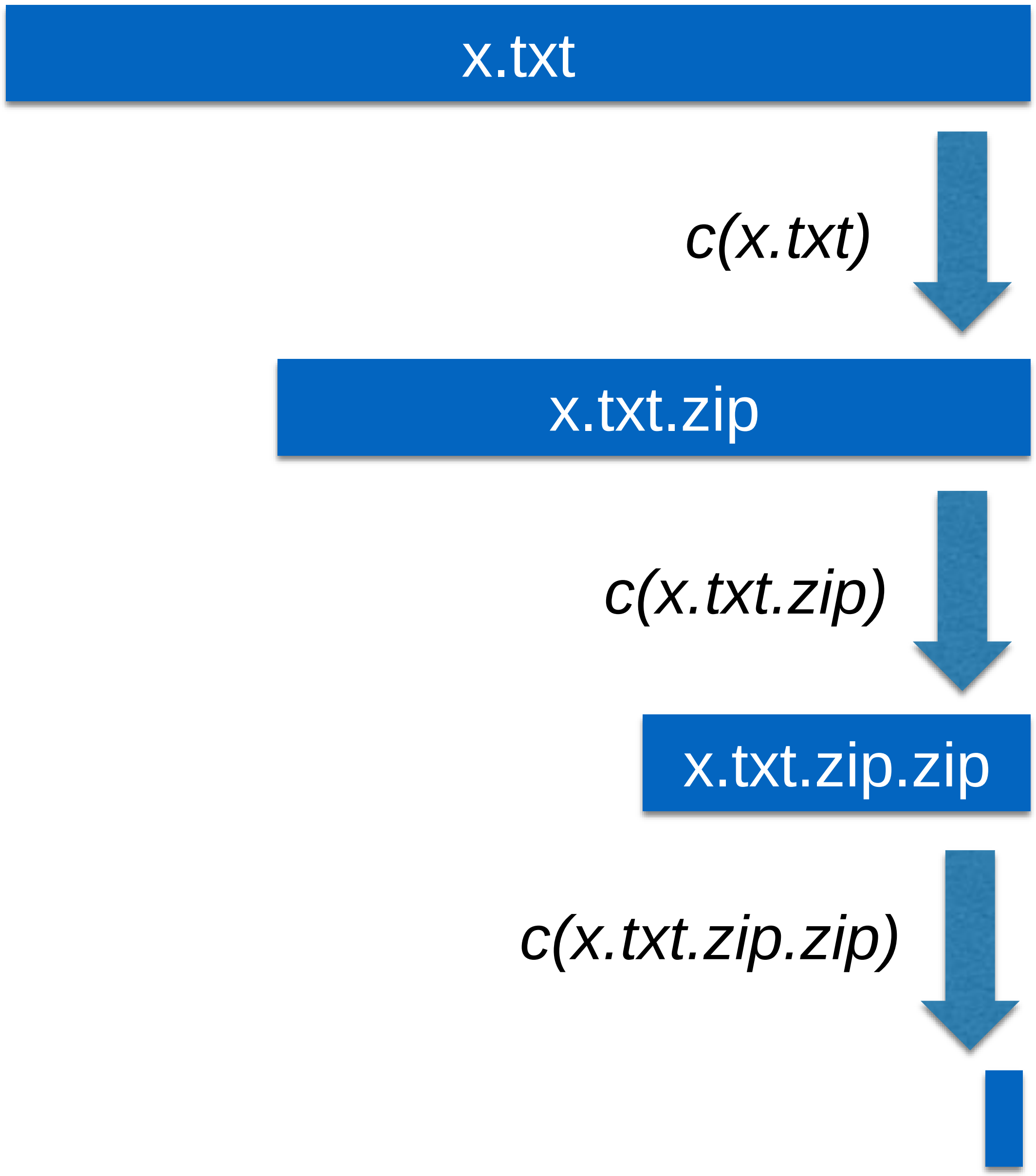
- In current world, quintillions of data is generated every day
- Issue when storing files/data on disk
- Issue when sending data on network
  - Eg, videos on YouTube, voice on Skype, etc

Compression

- Consider data as bitstring of 0/1
- Compression algorithm *c* from data x with length |x| to data y with length |y| < |x|
- Need function *d* to decompress y into x, ie $y = c(x)$, $d(c(x)) = d(y) = x$
- Ie, think about Zip and Unzip commands

How Much Can I Compress???

- If *c* existed that always compress any input x, could just recursively apply *c* on its output *c(x)* until get |y|=1



Compression vs. Hash

- In hash, usually we map from bigger domain X into smaller domain Y, but:
  - Should not be able to get back x from y=h(x)
  - There can be collisions, ie h(x)=h(x’) for different values in X
- In compression, still dealing with *mapping* to smaller domain Y (ie all possible shorter bitstrings), but:
  - We need function *d* to get back x from y=c(x), ie x=d(y)
  - There can be no collisions, otherwise d(y) would not be deterministic

Domain Size

- If X is the set of all possible bitstrings of size n, then such set has $|\mathrm{X}| = 2^n$ elements x

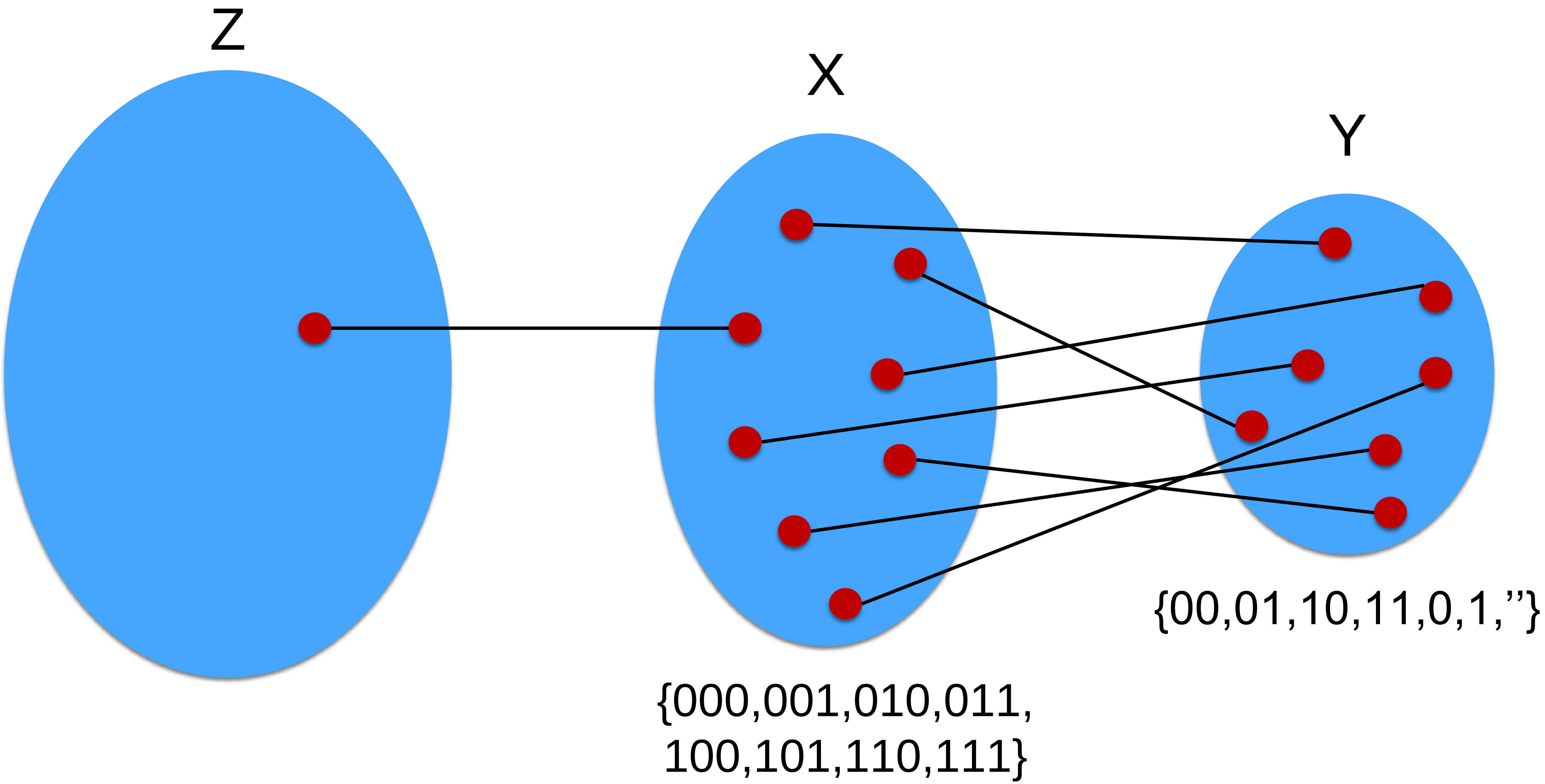
- The set |Y| of all possible shorter bitstrings contains the set of bitstrings of size n-1, n-2, n-3, … 1

- $|Y| = 2^{n-1} + 2^{n-2} + \cdots + 2^1 + 2^0 = 2^n - 1 < 2^n$

- In other words, |Y| < |X|

Mapping X -> Y and Back

- length(z) > length(x) > length(y),   length(x)=2^3, n=3 , |Y|=7
- When using c(x), there exists at least 1 case in which size does not decrease



What About Compression of At Least 50%?

- $|X| = 2^n,\ |Y| = 2^{n/2} + 2^{\frac{n}{2}-1} + \cdots 2^0 = 2^{\frac{n}{2}+1} - 1$

- $\frac{|X|}{|Y|} = \frac{2^n}{2^{\frac{n}{2}+1}-1} > \frac{2^n}{2^{\frac{n}{2}+1}} = 2^{n-(\frac{n}{2}+1)} = 2^{\frac{n}{2}-1}$

- So, for each element x that can be compressed by at least 50%, there are $\Omega(2^n)$ elements in X that cannot be compressed by 50%
  - Eg, n=20, |X|=1_048_576, |Y|=2047, |X|/|Y|~= 512

- In other words, it is *super extremely* unlikely to compress a random element from X, for increasing values of *n*

But, but…

- …when I compress my files with Zip I always reduce size by quite a lot!

- Point is, you are not compressing random files, but usually files with specific structures and properties:
  - Text files (English and Norwegian language)
  - Videos
  - Etc.

Exploit Redundancy

- Consider a bitstring with one trillion ‘1’s and no ‘0’
- 11111111111111111111111111111111111111111
11111111111111111111111111111111111111111
11111111111111111111111111111111111111111
11111111111111111111111111111111111111111
11111111111111111111111111111111111111111
11111111111111111111111111111111111111111
11111111111111111111111111111111111111111
11111111111111111… 1
- How to compress it?

We have already compressed it…

- I did not write a trillion 1s in the previous slide

- The sentence “Consider a bitstring with one trillion ‘1’s and no ‘0’” was an instruction to (de)compress it, and it only consisted of 54 characters

- If 8 bits per character, we can compress those 1 trillion bits with just 8x54=432 bits

- Who is *d(x)* here? It is the program that reads that instruction and create the trillion 1s

Another Example

- “Consider a bitstring with ‘01’ repeated half a trillion times”
- 01010101010101010101010101010101010101010101010101
01010101010101010101010101010101010101010101010101
01010101010101010101010101010101010101010101010101
010101010101010101010101… 01

No Structure

- Which instruction can express the following?

- 11011001100000000011110110111000010011000
  11110110010001110100110101111000 10110100
  00111000110111111111111110 001101110001111
  0001011010010100010001001101000110110100
  00101000011101000000011011010001100010111
  000100110100001101001100111000001100000100
  10110111101000

- A description / set of instructions for a bitstring with no structure can be longer than the bitstring itself

Structure in Data

- Most data you deal with has “structure”

- Eg, in English text, there is a specific set of words from a dictionary that are used, with grammatical rules, and some letters that are more common

- Usually you do not deal with compressing text like: “Ng1LCxc7Q1a9EdkOzfV9gK3GD4DujgIcMTrpPdy KVVtSy0Oo6U4eZG3Z3QbbTC7PRAhGUw78Yo09 3ixlf4McI9SBD483k8L1awMm”

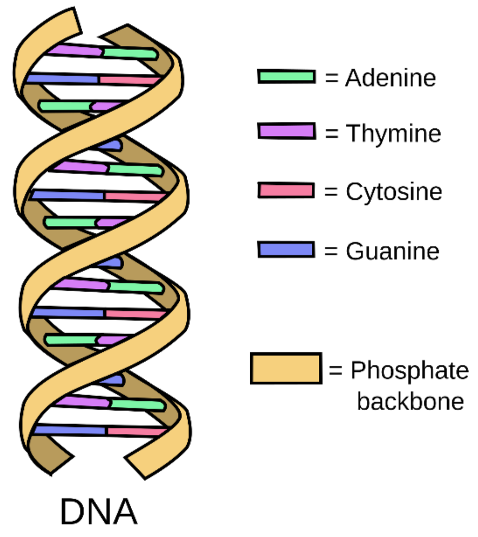
Compression Algorithm

- Mapping function from X to Y (with and existing decompression from Y to X)

- Exploit knowledge of the domain X of the data to compress, ie make use of its structure (if any)

- General compression algorithm: try to automatically find structure in the data and special properties, and exploit them

DNA Data



- DNA represented as sequence of 4 types of nucleobases
  - **A**denine, **T**hymine, **C**ytosine and **G**uanine
- How to store such data?
- Can use a *text file*, in which we use letter for each nucleobase: **A**, **T**, **C** and **G,** eg, AATCGTCGGATCGGCCCATCG…
- Lot of data: human genome is about 3.2 billion bases
- Using 1 byte / 8 bits per character, means 3G per human genome
- Can we do better?

Finite Set

- The data in DNA sequence is from a finite set {A,T,C,G}
- Instead of 8 bits (ie 1 byte) per letter, we can use custom mapping of using just 2 bits
  - A (00), T(01), C(10) and G(11)
- So, no longer using text file, but our custom compressed format
- Going from 8 to 2 bits per nucleobase gives us a 75% compression saving

ASCII Code: Character to Binary

0	0011 0000	O	0100 1111	m	0110 1101
1	0011 0001	P	0101 0000	n	0110 1110
2	0011 0010	Q	0101 0001	o	0110 1111
3	0011 0011	R	0101 0010	p	0111 0000
4	0011 0100	S	0101 0011	q	0111 0001
5	0011 0101	T	0101 0100	r	0111 0010
6	0011 0110	U	0101 0101	s	0111 0011
7	0011 0111	V	0101 0110	t	0111 0100
8	0011 1000	W	0101 0111	u	0111 0101
9	0011 1001	X	0101 1000	v	0111 0110
A	0100 0001	Y	0101 1001	w	0111 0111
B	0100 0010	Z	0101 1010	x	0111 1000
C	0100 0011	a	0110 0001	y	0111 1001
D	0100 0100	b	0110 0010	z	0111 1010
E	0100 0101	c	0110 0011	.	0010 1110
F	0100 0110	d	0110 0100	,	0010 0111
G	0100 0111	e	0110 0101	:	0011 1010
H	0100 1000	f	0110 0110	;	0011 1011
I	0100 1001	g	0110 0111	?	0011 1111
J	0100 1010	h	0110 1000	!	0010 0001
K	0100 1011	I	0110 1001	'	0010 1100
L	0100 1100	j	0110 1010	"	0010 0010
M	0100 1101	k	0110 1011	(	0010 1000
N	0100 1110	l	0110 1100	)	0010 1001
				space	0010 0000

Char Representation

- Considering char encoding for just ASCII characters
- A -> 01000001
- T -> 01010100
- C -> 01000011
- G -> 01000111

Text File to Our Format

- Text: “TCGA”
- Binary of text: 01010100010000110100011101000001
- Compressed: *c*(“TCGA”) = 01101100
- Note: if I read our custom format as a text file, we would get 01101100 -> “l”

Charsets

- Before we go into the details of how to compress text files, we need to go into *more details* on how characters are represented on computers

- Each character is mapped to a bitstring representation, which can be seen as a number

- But there are many types of mappings, called *Charsets*

ASCII Codes

- **A**merican **S**tandard **C**ode for **I**nformation **I**nterchange (ASCII)

- Mapping for 128 characters commonly used in English
  - Eg, a-z, A-Z, 0-9, ?, !, #, %, …

- As $128 = 2^7$, we just need 7 bits, which can be store in 1 byte

- Problem: how to represent special characters like the Norwegian øæåØÆÅ, or Japanese 私はアンドレアです???

Unicode

- Standard for encoding of characters
- Representing up to 1,114,112 possible characters
- Currently mapping 136,755 characters used in most languages around the world
- It implies we might need at least log2(1114112) = 20.08746 bits for the mapping, ie 3 bytes
  - Ie, using a single byte is not enough

Common Charsets

- **ISO/IEC 8859-1**: using 1 byte, representing up to 256 characters, including Norwegian and Swedish ones, but not full Unicode (eg, no Japanese)

- **UTF-8**: *most used encoding*. Multi-byte representation, up to 4 bytes. Can represent whole Unicode. Ascii (most common) codes need 1 byte, but Norwegian need 2

- **UTF-16**: used internally by Java (eg, “char” variables). Each character takes *at least* 2 bytes. Covers whole Unicode.

Parsing ISO-8859-1

- As each character is 1 byte, I just read 1 byte (ie 8 bits) at a time
- Direct mapping from 8 bits to a specific character

Parsing UTF-8

- Read 1 byte at a time, but need to find out if single byte character (eg, “A”), or beginning of multi-byte one (eg, “Ø” or “す”)
- If multi-byte character, need to read all of them before being able to map them to a single char
- Look at first 2 bits in each byte:
  - 0xxxxxxx -> single byte character (using remaining 7 bits)
  - 11xxxxxx -> beginning of a multi-byte character
  - 10xxxxxx -> continuation of a multi-byte character

Generic Text Compression

- Analyze the alphabet of the text to compress
  - Ie, the {A,T,C,G} in the DNA example
- Automatically create a custom encoding of char to bits
- Idea: often used characters should have smaller bit representation than seldom used characters

Letters in *The Odyssey*

- *“T**e**ll m**e**, O mus**e**, of that ing**e**nious h**e**ro who trav**e**ll**e**d far and wid**e** aft**e**r h**e** had sack**e**d th**e** famous town of Troy…”*
  - ' ' (108792): 111
  - 'e' (59526): 001
  - 't' (39156): 1010
  - ...
  - '1' (1): 0101001100111110010
  - '2' (1): 0101001100111110011
- Spaces and letter ‘e’ are the most common (108k and 59k occurrencies), can use just few bits (ie 3) for them
  - Eg, already in opening sentence, ‘e’ used 12 times out of 110 chars
- Numbers are seldom used in that book, so use more bits (ie 17)

Prefix-Free Codes

- Whatever encoding (ie mapping from char to bitstring) we choose, it **must** be *prefix-free*

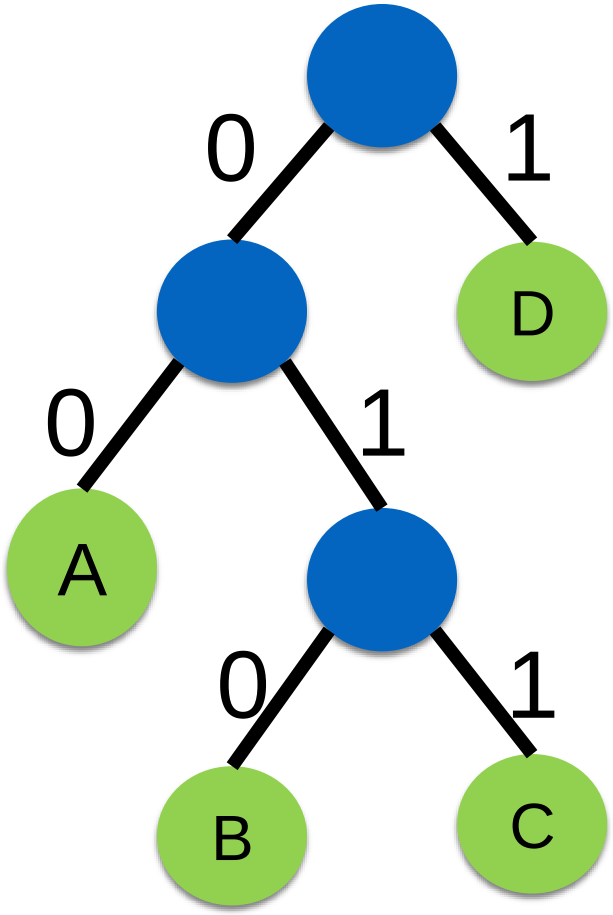
- No code should be the starting (ie prefix) of another code

- Eg, consider A=0, B=1, C=01
  - This is wrong, as A(0) is a prefix for C(01)

- How to decode 01? AB or C?
  - The decoding has to be non-ambiguous
  - If encoding is prefix-free, we know exactly each bit token to decode

Creating Prefix-Free Codes

- Using data structure called “Trie”
- Binary tree
- Labelled edges: 0 (left) and 1 (right)
- Letters in the leaves, not internal nodes
- Key for a leaf is codes on path to it
  - A = 00
  - B = 010
  - C = 011
  - D = 1
- This guarantees prefix-free mapping, as letters are ONLY on the leaves



Optimal Trie

- We want most common letters having shortest bit encoding

- Given an alphabet (eg {A,B,C,D}), there can be many different tries for it

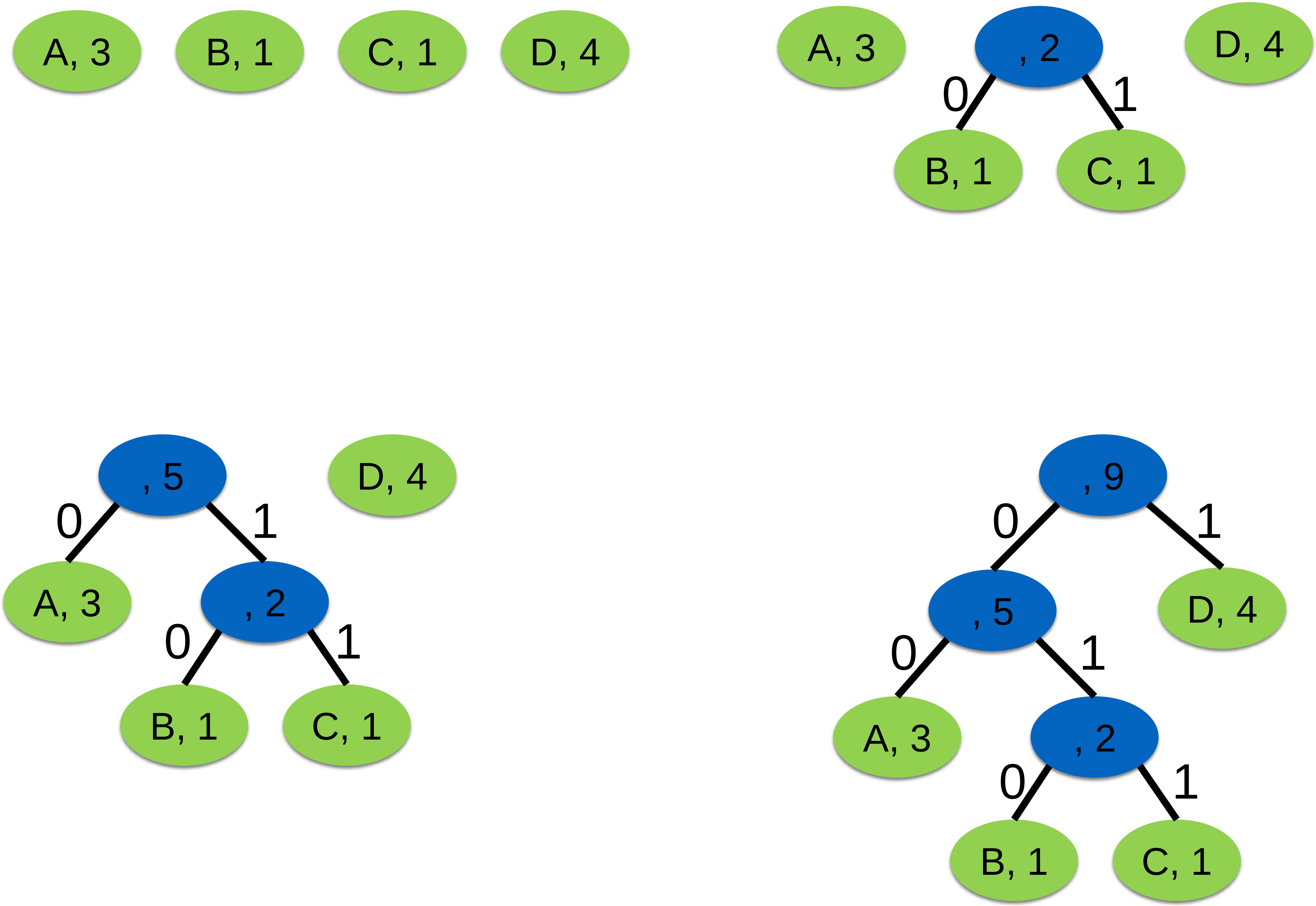
- How to build the optimal trie?

Huffman Algorithm

- Create optimal trie *t* for given input text *x*
- *c(x)=y ->* write binary of trie *t,* plus *x* encoded with *t*
- *d(y) ->* read binary of *t,* recreate it, use *t* to decode *x*

Creating Optimal Trie

- Consider string: “DBCADDADA”
- Compute occurrence of each symbol
  - A = 3, B = 1, C = 1, D = 4
- Create node for each symbol, with storing occurrences
- Choose 2 node roots with least occurrences, merge them in a new subtree with root having sum of their occurrences
- Repeat until only one root remains


A, 3
B, 1
C, 1
D, 4
A, 3
, 2
D, 4
0
1
B, 1
C, 1
, 5
D, 4
0
1
A, 3
, 2
0
1
B, 1
C, 1
, 9
0
1
, 5
D, 4
0
1
A, 3
, 2
0
1
B, 1
C, 1

Encoding The Trie

```
private static void writeTrie(
        Node node, BitWriter buffer) {

    if (node.isLeaf()) {
        buffer.write(true);
        buffer.write(node.ch);
        return;
    }

    buffer.write(false);
    writeTrie(node.left, buffer);
    writeTrie(node.right, buffer);
}
```

- For leaves: write a ‘1’ followed by 16 bits for the char in UTF-16
- For intermediate nodes, write a ‘0’, and then recursively write left and then right nodes
- Doing this gives us a non-ambiguous bitstring that we can decode later on to recreate the exact same trie

Huffman on *The Odyssey*

- Original in UTF-8 is 621737 bytes, ie 621kb
- Compressed with Huffman: 346507 bytes, ie 346kb
- Compression ratio: 0.55, ie we saved 45% of space

LZW Compression

- In Huffman, we have fixed size values (i.e., single chars) that are mapped by multibit key codes

- Another approach is to have fixed size key codes mapping multibit values

- Example: “*have*” is very common in English, and could be mapped by a single bitcode 000, instead of having a code for each of its chars… but need to have bitcode length *n* to represent all the 2^n different words/tokens in a document

End of the Course…

- LZW is a popular compression algorithm based on that approach of fixed-size codes
- Being last class of this (long) course, we will not go into its details
- But for sake of completeness, the Git repository has its source and tests, but those will NOT be part of the exam

Homework

- Study Book Chapter 5.5
- Study code in the *org.pg4200.les12* package
- Do exercises in *exercises/ex12*