

# Agenda





Version	Reviewed by	Approved by	Remarks
1.0			



### **Learning Outcome**

- Linux Driver Model
- Registering Char Device Drivers
- Device Id, Major, Minor numbers
- Implement Driver Operations
- Data Structures in Kernel Space
- Handling multiple devices
- Concurrency & IPC in Kernel
  - Mutual Exclusion, Synchronization
- IOCTL Operations
- Misc Devices, Implementing Proc Files

#### Pre-Requites

- Operating System Basics Kernel, System Calls, Process Management etc
- Concurrency multithreading etc
- IPC Concepts and User space Programming (System Calls)
- File handling in user space using system calls open, read, write, close
- Virtual Memory Concepts Overview
- Understanding on general data structures, especially Linked List
- Clarity on Blocking vs Non-Blocking calls



# **Major and Minor Numbers**

- Major number Associates with set/family of devices, managed by common driver typically
- Minor number Distinguish each device
- Static or Dynamic allocation of major, minor numbers
- Device Id Major + Minor number
- "In 32 bit device id ==> 12 bit major number (MSBits) + 20 bit minor numbes (LSBits)
- Macros MAJOR, MINOR, MAKEDEV

# Userspace Interface

- "Unix/Linux Philosophy Everything is a file
- Device special files in /dev directory, managed by devtmpfs mount
- "Creation of device file Manual / APIs
  - // mknod command (manual) , device\_create API
- ✓ Checking attributes of device files Is, stat commands (or) Istat system call
- # File operations in userspace open, read, write, close
- # Kernel configuration CONFIG\_DEVTMPS\_MOUNT
- **"VFS Role** 
  - In memory i-node (in-core)
  - open file descriptor







# Step-1: Register Char Driver

```
#include <linux/fs.h>
dev t pdevid;
int ndevices=1;
static int init psuedo init(void)
  int ret;
  ret=alloc chrdev region(&pdevid, 0, ndevices, "pseudo sample");
  if(ret) {
    printk("Pseudo: Failed to register driver\n");
    return -EINVAL:
  printk("Successfully registered,major=%d,minor=%d\n",
                                MAJOR(pdevid), MINOR(pdevid));
  printk("Pseudo Driver Sample..welcome\n");
  return 0;
static void exit psuedo exit(void) {
  unregister chrdev region(pdevid, ndevices);
  printk("Pseudo Driver Sample..Bye\n");
```

Observer assigned major no.

cat /proc/devices
dmesg

# Step-2 : Register File Operations

```
#include <linux/cdev.h>
struct cdev cdev; //global
int ndevices=1;
struct file operations fops = {
        = pseudo open,
  .open
  .release = pseudo close,
  .write = pseudo_write,
  .read = pseudo read
};
//In init
cdev init(&cdev, &fops);
kobject set name(&cdev.kobj,"pdevice%d",i);
ret = cdev add(&cdev, pdevid, 1);
//In exit
cdev del(&cdev);
```

```
Testing driver:-
insmod pseudo.ko
mknod /dev/psample c xxx 0
cat /dev/psample
echo "abc" > /dev/psample
#check dmesg on each operation
# or use open, read, write, close
# on /dev/psample
rmmod pseudo
rm /dev/psample
```

### Step-2 : Dummy File Operations

```
int pseudo open(struct inode* inode , struct file* file)
                                                                           dumpstack usage
  printk("Pseudo--open method\n");
  return 0;
int pseudo close(struct inode* inode , struct file* file)
  printk("Pseudo--release method\n");
  return 0;
ssize t pseudo read(struct file * file, char user * buf , size t size, loff t * off)
  printk("Pseudo--read method\n");
  return 0;
ssize t pseudo write(struct file * file, const char user * buf , size t size, loff t * off)
  printk("Pseudo--write method\n");
 return -ENOSPC;
```

## Step-3: Device file creation

```
//Additional Headers
#include <linux/device.h>
struct device *pdev; //global
struct class *pclass; //global
//In init:-
int i=0;
pclass = class create(THIS MODULE, "pseudo class");
//alloc chrdev region, cdev init, cdev add
pdev = device_create(pclass, NULL, pdevid, NULL, "psample%d",i);
//In exit:-
device destroy(pclass, pdevid);
class destroy(pclass);
```



- device file created in /dev
- sysfs entries for class and device file

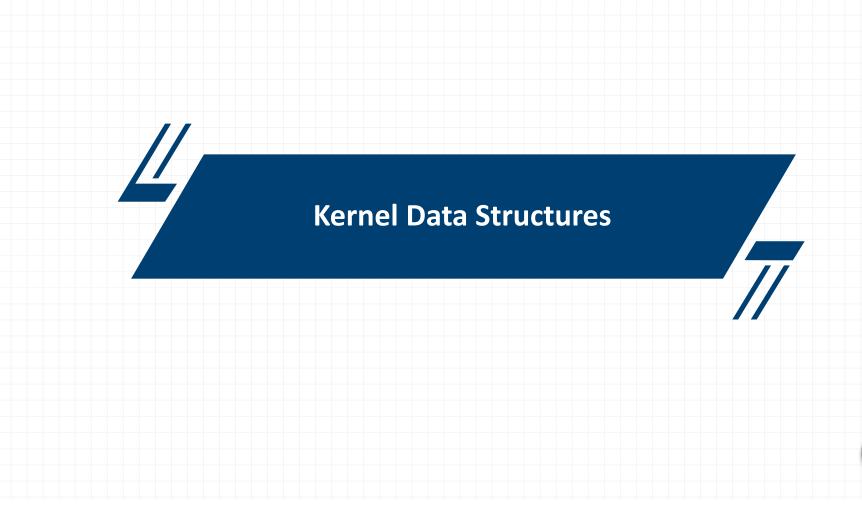
# Step-4: Buffer as pseudo device

```
//Additional Headers
#include <linux/slab.h>
#include <linux/uaccess.h>
//Global
unsigned char *pbuffer;
int rd_offset=0;
int wr offset=0
int buflen=0;
//In init:-
pbuffer = kmalloc(MAX_SIZE, GFP_KERNEL);
//In exit:-
kfree(pbuffer);
```

# Step-4: Implement read, write operations

```
//Write method:-
if(wr offset >= MAX SIZE)
 printk("buffer is full\n");
 return -ENOSPC;
wcount = size;
if(wcount > MAX SIZE - wr offset)
 wcount = MAX SIZE - wr offset;
                                      //min
ret=copy from user(ubuf, pbuffer + wr offset,
                                 wcount);
if(ret)
 printk("copy from user failed\n");
 return -EFAULT;
wr offset+=wcount;
buflen += wcount;
```

```
//Read method:-
if(buflen==0)
                    //wr offset-rd offset==0
  printk("buffer is empty\n");
 return 0;
rcount = size;
if(rcount > buflen)
  rcount = buflen;
                       //min of buflen, size
ret=copy to user(ubuf, pbuffer + rd offset,
                                rcount);
if(ret)
  printk("copy to user failed\n");
  return -EFAULT;
rd offset+=rcount;
buflen -= rcount;
```



#### Kfifo API



```
//Prototypes in kfifo.h
struct __kfifo {
  unsigned int in;
  unsigned int out;
  unsigned int mask;
  unsigned int esize;
  void *data;
};
//Definitions in lib/kfifo.c
```

# Step 5 : kfifo usage in Pseudo Driver

```
#include<linux/kfifo.h>
unsigned char *pbuffer;
#define MAX SIZE 1024
struct kfifo kfifo;
//pseudo_init
pbuffer=kmalloc(MAX_SIZE, GFP_KERNEL);
kfifo init(&kfifo, phuffer);
//kfifo alloc(&kfifo, MAX SIZE, GFP_KERNEL);
//pseudo exit
Kfifo free(kfifo);
```

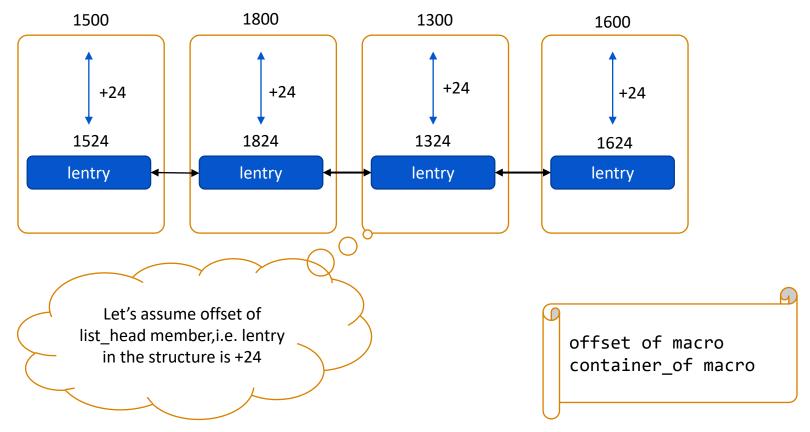
No need to call kfree(pbuffer) in case of kfifo\_init kfifo\_free will internally release slab memory assigned to underlying buffer

# Step-5: kfifo usage in pseudo driver

```
//Write method:-
if(kfifo is full(&myfifo))
 printk("buffer is full\n");
 return -ENOSPC;
wcount = size;
if(wcount > kfifo avail(&myfifo);)
 wcount = kfifo avail(&myfifo);
                                      //min
char *tbuf=kmalloc(wcount, GFP KERNEL);
ret=copy from user(tbuf, ubuf, wcount);
//error handling if copy form user
kfifo in(&myfifo, tbuf, wcount);
kfree(tbuf);
```

```
//Read method:-
if(kfifo is empty(&myfifo)) {
  printk("buffer is empty\n");
  return 0;
rcount = size;
if(rcount > kfifo len(&myfifo))
  rcount = kfifo len(&myfifo);
                                     //min
tbuf = kmalloc(rcount, GFP KERNEL);
kfifo out(&myfifo, tbuf, rcount);
ret=copy to user(ubuf, tbuf, rcount);
//error handling
kfree(tbuf);
```

# List implementation in Kernel



#### List APIs

```
LIST_HEAD
list_add_tail
list_del
list_for_each
list_for_each_safe
list_for_each_entry
list_for_each_entry_safe
list_entry
```

```
LIST_HEAD(mylist);
```

```
sample_t *pnew = kmalloc(sizeof(sample_t), GFP_KERNEL);
//fill other members of pnew
mylist_addd(&new->lentry, &mylist);
```

#### List APIs

```
struct list_head *pcur;
list_for_each(pcur, &mylist) {
  ptr = list_entry(pcur, descriptor_t, lentry);
  //access other members of ptr
}
```

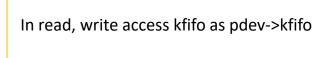
```
struct list_head *pcur, *prev;
list_for_each_safe(pcur, prev, &mylist) {
  ptr = list_entry(pcur, descriptor_t, lentry);
  kfree(ptr);
}
```

```
sample_t *ptr;
list_for_each_entry(ptr, &mylist, lentry) {
  //access other members of ptr
}
```

```
sample_t *ptr, *qtr;
list_for_each_entry_safe(ptr, qtr,&mylist,lentry)
{
   kfree(ptr);
}
```

### Step-6: Private Object in Pseudo Driver

```
typedef struct priv obj {
  struct cdev cdev;
  struct kfifo kfifo;
  strcut device* pdev;
}PRIV OBJ;
PRIV OBJ* pobj;
//init:-
pobj=kmalloc(sizeof(PRIV OBJ), GFP KERNEL);
//pobj->cdev,pobj->kfifo, pobj->pdev
//exit:-
//pobj->cdev,pobj->kfifo, pobj->pdev
kfree(pobj);
```



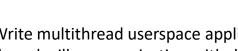
### Step-7: list API usage in Pseudo Driver

```
typedef struct priv obj {
  struct cdev cdev;
  struct kfifo kfifo;
  strcut device* pdev;
  struct list head *lentry;
}PRIV OBJ;
LIST_HEAD(mydevlist);
//init:-
PRIV OBJ *pobj; //local, not global this time
pobj=kmalloc(sizeof(PRIV OBJ), GFP KERNEL);
//pobj->cdev,pobj->kfifo, pobj->pdev
list add tail(&pobj->lentry, &mydevlist);
//exit:-
PRIV OBJ *ptr, *qtr;
list for each entry safe(ptr, qtr,&mydevlist,lentry) {
 kfree(ptr);
```

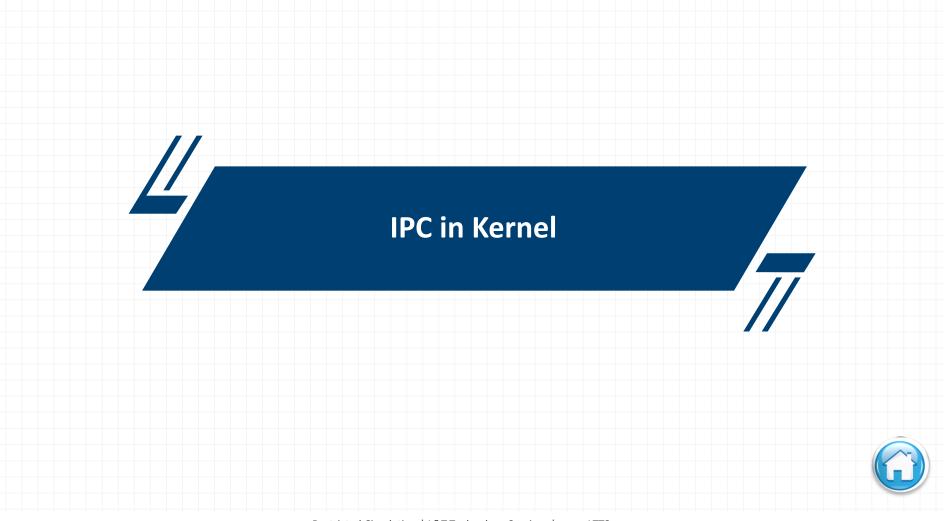
# Step-7: Handling multiple Devices

```
//pseudo open
PRIV OBJ *pobj = container of(inode->i cdev, PRIV OBJ, cdev);
file->private data=pobj;
```

```
//pseudo_read, pseudo_write
PRIV_OBJ *pobj = file->private_data;
//access kfifo as pobj->kfifo
```



Write multithread userspace application, where each thread will communication with different devices



# Concurrency & IPC Support in Linux Kernel

# Concurrency

Kernel Threads

# Locking & Synchronization

- Semaphores
- Mutex
- Spinlocks
- Atomic Operations
- Wait queues

#### Kernel Threads

```
//Global
static struct task struct *task1;
static struct task struct *task2;
//tdemo init:-
task1=kthread run(thread one, NULL, "thread A");
//kthread create + wake up process
task2=kthread_run(thread_two, NULL, "thread B");
//tdemo exit:-
if(task1)
  kthread stop(task1);
if(task2)
  kthread stop(task2);
```

```
static int thread_one(void *pargs){
  int i;
  while(!kthread_should_stop())
  {
    printk("Thread A--%d\n",k++);
    msleep(1000); //ssleep, usleep
  }
  do_exit(0);
  return 0;
}
```

```
//Similarly write code for thread_two
```

# Signal Handling

```
static int thread_one(void *pargs){
  int i;
  allow_signal(SIGKILL);
  while(!kthread_should_stop())
  {
    printk("Thread A--%d\n",k++);
    if (signal_pending(task1))
       break;
    msleep(1000); //ssleep, usleep
  }
  do_exit(0);
  return 0;
}
```

- Do kernel threads will have task\_struct and thread\_info structures
- What about address space of kernel threads, do mm field in task\_struct applicable for them?
- ☐ Can you observe kernel threads from userspace using utilities like ps command
- ☐ Send SIGKILL to running kernel thread and observer the behavior

#### Race Conditions & Concurrent Scenarios

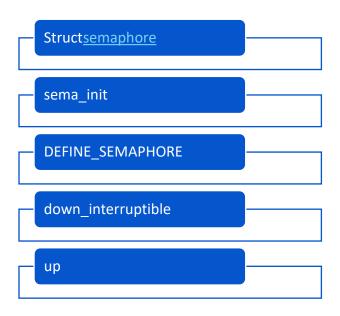
```
int val=100;
const int max=10000;
static int thread_one(void *pargs) {
  int i;
  for(i=1;i<=max;i++)</pre>
    val++;
    if(kthread should stop())
      break;
  do exit(0);
 return 0;
//Similarly write thread two with val--
```

```
int val=100;
const int max=20;
static int thread one(void *pargs) {
  int i;
  for(i=1;i<=max;i++)</pre>
     printk("Thread A--%d\n",k++);
     if(kthread should stop())
       break;
     msleep(100);
  do exit(0);
 return 0;
//Similarly write thread two which
//prints max times with some delay.
```

Scenario-1

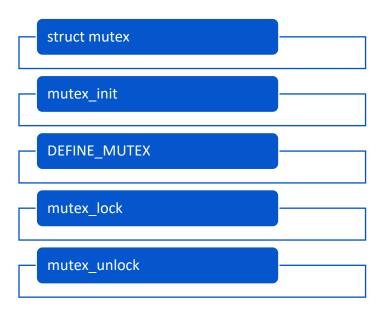
Scenario-2

# Semaphore API



Locate these APIs in kernel source – prototype and definition
 Identify other Semaphore APIs
 Apply semaphore operations for Scenario-1 to avoid race conditions
 Apply semaphore operations for Scenario-2 to ensure only one for loop (any) executes at a time

#### Mutex API



Mutex vs Semaphore, Salient features of Mutex!!

- ☐ Locate these APIs in kernel source prototype and definition
- ☐ Identify other Mutex APIs, especially trylock, recursive locks
- Apply Mutex operations for Scenario-1 to avoid race conditions
- Apply Mutex operations for Scenario-2 to ensure only one for loop (any) executes at a time

# Spinlock API

spin lock init **DEFINE SPINLOCK** spin\_lock spin unlock

■ Locate these APIs in kernel source – prototype and definition

☐ Identify other Spinlock APIs, especially trylock, recursive locks

Apply Spinlock operations for kernel threads

☐ Significant usage of spinlocks in Interrupt Handling, Bottom Halves(TBD Later)

Spinlock vs Semaphore/Mutex, when to use which??

Context switching vs Busy Waiting/Spinning?

### Wait Queue API

struct wait\_queue\_head / wait\_queue\_head\_t

DEFINE\_SPINLOCK

init\_waitqueue\_head

wait\_event\_interruptible

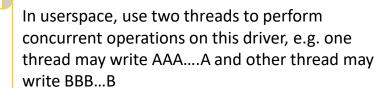
wake\_up\_interruptible

- Locate these APIs in kernel source prototype and definition and refer kernel docs
- ☐ Identify other waitqueue APIs,
- Other IPC techniques built on top of waitqueues
- Apply Waitqueue operations in Scenario-2, to ensure that always for loop of one kernel thread runs before other

### Step-8: Generate Race Conditions in Pseudo Driver

```
// Consider Step-4 code as baseline, using global buffer as Pseudo
device

//pseudo_write:-
for(i=0;i<wcount;i++) {
   get_user(*(pbuffer+wr_offset),ubuf);
   ++wr_offset;
   ubuf++;
   msleep(100);   //intentional delay to generate race conditions
}</pre>
```



# Step-8: Mutual Exclusion in Pseudo Driver

```
//pseudo_write:-
down_interruptible(&s1); //mutex_lock
for(i=0;i<wcount;i++) {
   get_user(*(pbuffer+wr_offset),ubuf);
   ++wr_offset;
   ubuf++;
   msleep(100);
}
up(&s1); //mutex_unlock</pre>
```

☐ When write operation is in progress on behalf of one thread, no other write operation (writing to buffer part) can takes place on behalf of other threads.

## Step-9: Synchronization in Pseudo Driver

```
//Additional header files
                                              Read method blocks when kfifo is empty and unlocks when
#include<linux/wait.h>
                                              some data is written to kfifo by write operation
//inside private object
wait queue head t rd queue;
//init:-
                                               Userspace test code to check synchronization
init waitqueue head(&pobj->rd queue);
//pseudo read:-
if(kfifo is empty(&pobj->kfifo)) // //file->f flags & O NONBLOCK, -EAGAIN
   wait event interruptible(pobj->rd queue,(kfifo len(&pobj->kfifo)>0)); //!kfifo is empty
   //instead of returning 0, if 0 NONBLOCK is not requested from userspace
//pseudo write:-
                                           //on writing some data into kfifo
wake up interruptible(&pobj->rd queue);
```

# Step-9: Synchronization in Pseudo Driver

```
//Additional header files
                                               Write method blocks when kfifo is full and unlocks when
#include<linux/wait.h>
                                               some data is retrieved from kfifo by read operation
//inside private object
wait queue head t wr queue;
//init:-
                                               Userspace test code to check synchronization
init waitqueue head(&pobj->wr queue);
//pseudo write:-
if(kfifo is full(&pobj->kfifo)) // //file->f flags & O NONBLOCK, -EAGAIN
   wait event interruptible(pobj->wr queue,(kfifo avail(&pobj->kfifo)>0)); //!kfifo is full
   //instead of returning -ENOSPC, if O NONBLOCK is not requested from userspace
//pseudo read:-
                                          //on reading some data from kfifo
wake up interruptible(&pobj->wr queue);
```



## Step-10: ioctl

```
struct file_operations fops = {
  .open = pseudo_open,
  .release = pseudo_close,
  .write = pseudo_write,
  .read = pseudo_read,
  .unlocked_ioctl = pseudo_ioctl
};
```

```
static long pseudo ioctl(struct file *file,
unsigned int cmd, unsigned long arg)
  int ret;
  printk("Pseudo--ioctl method\n");
  switch (cmd) {
    case MY IOCTL LEN:
        printk("ioctl--kfifo length is %d\n",
                kfifo len(&myfifo));
        break:
    case MY IOCTL AVAIL:
        printk("ioctl--kfifo avail is %d\n",
        kfifo avail(&myfifo));
        break;
   case MY IOCTL RESET:
        printk("ioctl--kfifo got reset\n");
        kfifo reset(&myfifo);
        break;
  return 0;
```

# Userspace code

```
//similar header file with magic numbers in
userspace

#define IOC_MAGIC 'p'
#define MY_IOCTL_LEN _IO(IOC_MAGIC, 1)
#define MY_IOCTL_AVAIL _IO(IOC_MAGIC, 2)
#define MY_IOCTL_RESET _IO(IOC_MAGIC, 3)
```

```
fd = open("/dev/psample", O_RDWR);
if(fd<0) {
  perror("open");
  exit(1);
ret=ioctl(fd, MY IOCTL LEN);
if(ret<0) {
  perror("ioctl");
 exit(3);
ret=ioctl(fd, MY_IOCTL_AVAIL);
if(ret<0)
  perror("ioctl");
  exit(3);
ret=ioctl(fd, MY_IOCTL_RESET);
if(ret<0) {
  perror("ioctl");
  exit(3);
close(fd);
```

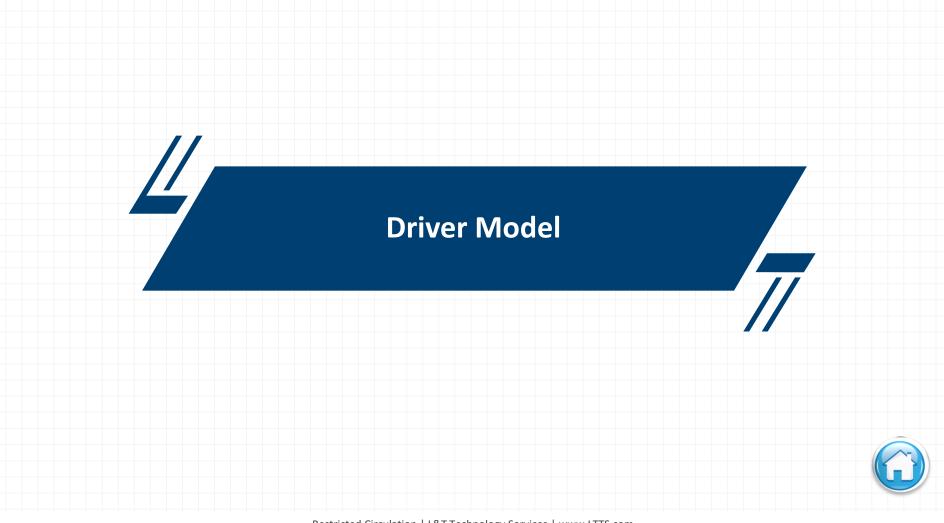
## Step 10 : Another IOCTL Operation

```
struct pseudo_stat {
  int len;
  int avail;
};
#define MY_IOCTL_PSTAT __IOR(IOC_MAGIC, 4, struct pseudo_stat)
```

```
static long pseudo_ioctl(struct file *file, unsigned int cmd, unsigned long arg) {
 switch(cmd) {
   case MY IOCTL PSTAT:
     printk("ioctl--kfifo statistrics\n");
     stat.len=kfifo len(&myfifo);
     stat.avail=kfifo avail(&myfifo);
     ret=copy to user( (char user*)arg, &stat, sizeof(pseudo stat));
     if(ret) {
       printk("error in copy to user\n");
       return -EFAULT;
     break;
 return 0;
```

# Userspace Test Code for IOCTL

```
//open
struct pseudo_stat stat;
ret=ioctl(fd, MY_IOCTL_PSTAT, &stat);
if(ret<0) {
   perror("ioctl");
   exit(4);
}
//print stat.len, stat.avail
//close</pre>
```



# Platform Device – Dummy Serial Example

```
//myserial_init:-
platform_device_register(&myuart_device);
//myserial_exit:-
platform_device_unregister(&myuart_device);
```

# Platform Driver - Dummy Serial Example

```
static int myserial_probe(struct platform_device *pfdev)
{
   printk("Dummy serial driver -- probe method\n");
   return 0;
}
static int myserial_remove(struct platform_device *pdev)
{
   printk("Dummy serial driver -- remove method\n");
   return 0;
}
```

```
#include <linux/platform_device.h>
```

```
#define DRIVER_NAME "uart_dummy"
//myserial_init:-
platform_driver_register(&myserial_driver);
//myserial_exit:-
platform_driver_unregister(&myserial_driver);
```

## Platform Driver - Probe

```
int ret;
struct resource *res;
struct myserial dev *dev;
unsigned int baud divisor, uartclk;
pr info("Called feserial_probe\n");
res = platform_get_resource(pdev, IORESOURCE MEM, 0);
if (res == NULL)
    pr_err("Unable to obtain platform memory resource\n");
    return -1;
dev = devm kzalloc(&pdev->dev, sizeof(struct myserial dev), GFP KERNEL);
if (!dev) {
    pr err("Unable to obtain kernel memory\n");
    return - ENOMEM;
dev->irq = platform get irq(pdev, 0);
dev->miscdev.minor = MISC DYNAMIC MINOR;
dev->miscdev.name = kasprintf(GFP KERNEL, "myserial-%x",res->start);
dev->miscdev.fops = &myserial fops;
platform set drvdata(pdev, dev);
```

## Platform Driver - Probe

```
ret = misc_register(&dev->miscdev);
if (ret) {
    pr_err("Failed to register misc console\n");
    return -ENODEV;
}
pr_info("dummy serial: uart base is : %x\n", res->start);
pr_info("dummy serial: irq number is : %d\n", dev->irq);
printk("Dummy serial driver -- probe method\n");
return 0;
```

```
struct myserial_dev {
    struct platform_device *pdev;
    struct miscdevice miscdev;
    int irq;
    void __iomem *regs;
};
```

## **Device & Driver Modules**

Implement Device & Driver Modules Separately,

□ Load device module first and then load driver module
□ Remove driver module first and then remove device module
□ Load driver module first and then load device module
□ When the probe method of driver module got invoked?
□ Unload device module before driver module
□ When the remove method of driver module got executed?

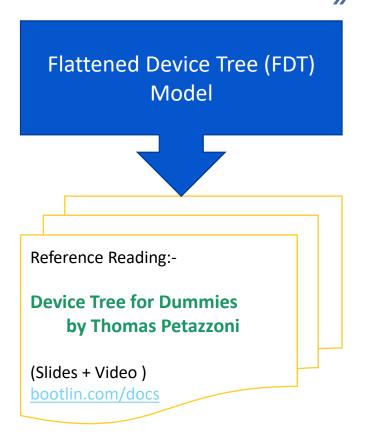
Typically, device module will be configured as static module

### Additional Reading/Ref:-

https://www.kernel.org/doc/html/latest/driver-api/driver-model/platform.html https://lwn.net/Articles/448499/

## **Device Tree Based Driver**

```
static const struct of device id myuart of match[] =
    { .compatible = "arm, myuart", },
    {},
};
MODULE DEVICE TABLE(of, myuart of match);
static struct platform driver myserial driver=
                         = myserial probe,
    .probe
                         = myserial remove,
    .remove
    .driver
          .name = DRIVER NAME,
          .of match table = of match ptr(myuart of match),
          .owner=THIS MODULE,
    },
};
```



## **Device Tree Based Driver**

```
res = platform get resource(pdev, IORESOURCE MEM, 0);
if (res == NULL)
    pr err("Unable to obtain platform memory resource\n");
    return -1;
dev = devm kzalloc(&pdev->dev, sizeof(struct myserial dev), GFP KERNEL);
if (!dev) {
    pr err("Unable to obtain kernel memory\n");
    return - ENOMEM;
dev->irq = platform_get_irq(pdev, 0);
dev->regs = devm ioremap resource(&pdev->dev, res);
if (!dev->regs) {
        dev_err(&pdev->dev, "Cannot remap registers\n");
        return -ENOMEM;
of property read u32(pdev->dev.of node, "clock-frequency", &uartclk);
```

## **Device Tree Based Driver**

```
if (pdev->dev.of_node) {
    struct device_node *np = pdev->dev.of_node;
    ret = of_property_read_u32(np, "current-speed", &baudrate);
    if (ret)
    return ret;
}
else {
    baudrate = 115200;
}
```

```
//vexpress-v2p-ca9.dts:-
dummyserial@9000 {
  compatible = "arm,myuart", "arm,dummyserial";
  reg = <0x9000 0x20>;
  interrupts = <5>;
};
```

- ☐ Regenerate dtb file and test with this code
- Observe dummy serial entry in /proc/device-tree
- Observe new device tree entry in U-Boot shell



## Miscellaneous Devices

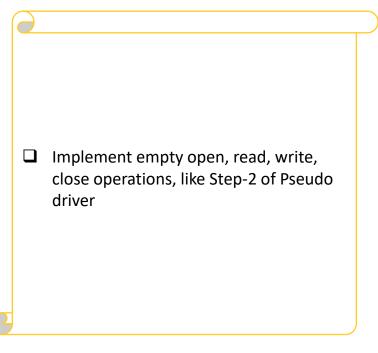
```
#include <linux/miscdevice.h>
static struct file operations fops = {
           = pseudo open,
  .open
  .release = pseudo close,
  .read = pseudo read,
  .write = pseudo write,
};
static struct miscdevice pseudo char misc = {
  .minor = MISC DYNAMIC MINOR,
  .name = "pseudo misc char",
  .fops = &fops,
};
//pseudo misc init
ret = misc register(&pseudo char misc);
//pseudo misc exit
misc deregister(&pseudo char misc);
```

#### Ref:-

https://www.kernel.org/doc/html/v4.14/driverapi/misc\_devices.html https://github.com/opersys/circular-driver

# Simple Proc File

```
static struct proc dir entry *pentry;
static struct proc dir entry *pdir;
struct file_operations fops={
           = pdemo open,
  .open
  .release = pdemo close,
  .read = pdemo read,
  .write = pdemo write,
};
//pdemo init:-
pdir=proc mkdir("ptest", NULL);
//if(pdir==NULL) ...
pentry=proc create("psample",0666,pdir,&fops);
//if(pentry==NULL) ...
proc set user(pentry, KUIDT INIT(0), KGIDT INIT(0));
proc set size(pentry,80);
//pdemo exit:-
remove proc entry("psample",pdir);
remove proc entry("ptest", NULL);
```



# Sequential Proc File – single\_open

```
static struct proc_dir_entry *pentry;
static struct proc dir entry *pdir;
ssize t custom proc show(struct seq file* m, void* p)
 seq printf(m,"psingle--custom show,dummy var=%d\n",100);
 seq printf(m,"current pid=%d\n",current->pid);
 return 0:
int pdemo open (struct inode * inode, struct file * file)
 return single open(file,custom proc show,NULL);
struct file_operations fops={
          = pdemo open,
  .open
  .release = single_release,
 .read = seq read,
 .write = sea write,
};
```

# Assignment/Coding Tasks based on Proc

- Implement simple read, write operations on Proc file
- Implement a proc file which does simple echo operation, i.e. a string written to proc file, same will be read back. You may consider any other string operation like reverse, toggle the case.
- Implement a proc file, on read operation print all possible attributes of invoking process (like system call task given earlier)
- Implement a proc file, which prints pid, ppid of all processes on read operation (like system call task given earlier)
- // Implement a proc file, which prints minimal information about available CPUs (minimal version of /proc/cpuinfo), Hint:- for\_each\_cpu / cputype.h / read\_cupid
- // Implement a proc file as follows
  - ✓ init method append n entries to list
  - " read method traverse the list and print (export data to userspace)
  - write method append one more element to list
  - exit method traverse and delete list nodes

