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# GSM – SRSLY?

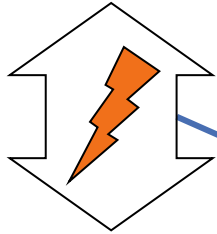
H4RDW4RE

# Summary:

## GSM Encryption needs to be shown insecure

### **GSM is constantly under attack:**

- A5/1 cipher shown insecure repeatedly
- Lack of network authentication allow MITM intercept (IMSI Catcher)



Security expectations divert from reality

### **However, GSM is used in a growing number of sensitive applications:**

- Voice calls, obviously
- SMS for banking
- Seeding RFID/NFC secure elements for access control, payment and authentication

- To rectify the perception of GSM's security, we demonstrate its weaknesses
- The community has computed the cryptographic base for a public demonstration of cracking GSM
- This presentation details motives, approach and next steps of the "A5/1 Security Project"

# GSM is global, omnipresent and insecure

**80% of  
mobile  
phone  
market**

**200+  
countries**

**4 billion  
users!**



**GSM  
encryption  
introduced  
in 1987 ...**

**... then  
disclosed  
and  
shown  
insecure  
in 1994**

# We need to publicly demonstrate that GSM uses insufficient encryption

## Public break attempts

A5/1 shown *academically* broken

A5/1 shown more ...

... and more ...

... and more broken.

Broken with massive computation

***Rainbow table computation***

'97

'00

'03

'05

'06

'03/'08

***Tables never released***

Too expensive

Not enough known data in GSM packets

**... that didn't work.**

# GSM encryption is constantly being broken, just not publicly

## All public break attempts of A5/1 have failed so far

- Academic breaks of A5/1 cipher are not practical [EC1997, FSE2000, Crypto2003, SAC2005]
- Cracking tables computed in 2008 were never released

- 15 years of A5/1 research have not produced a proof of concept
- (until today)

## Meanwhile ...



... A5/1 is constantly being circumvented by intelligence, law enforcement, and criminals

# Active and passive intercept is common as attack devices are readily available

## Two flavors of attack devices

A

Active intercept:

- Phones connect through fake base station
- Easily spottable (but nobody is looking)



B

Passive key cracking:

- Technically challenging
  - Non-trivial RF setup
  - Heavy pre-computation
- Allows hidden operation



This talk demonstrates that GSM intercept is practical to raise awareness

# IMSI catching routes calls through a fake base station



← Advertise base station on beacon channel

→ IMSI: Subscriber Identity (~= username)

- Sort-of secret (replaced by TMSI asap)

← MCC\*: Mobile Country Code

- 262 for .de, 310-316 for USA

MNC\*: Mobile Network Code

- Country-specific, usually a tuple with MCC
- 262-01 for T-Mobile Germany



Phones will connect to any base station with spoofed MNC/MCC

- If you claim it, they will come
- Strongest signal wins
- IMSI catching is detectable from phone, but no detect apps exists !
- Crypto is completely optional and set by the base station !!

\* Full list of MNC/MCCs available on Wikipedia



# IMSI catcher could even be built from open source components


- A Setup**
  - OpenBTS + USRP + 52MHz clock
    - Easy to set up, Asterisk is hardest part
    - On-board 64MHz clock is too unstable
  - Software side is easy
    - ./configure && make
    - Libraries are the only difficulty
- B Configure**
  - Set MCC/MNC to target network
  - Find and use an open channel (ARFCN in GSM-ese)
- C Collect, Decode**
  - Wireshark has a Built-in SIP analyzer
  - Or: capture data on air with Airprobe and decode GSM packets



# The iPhone that wouldn't quit

What if we want to test and *not* “catch” IMSIs?

- Set MCC/MNC to 001-01 (Test/Test)
- Phones camp to strongest signal
  - Remove transmit antenna
  - Minimize transmit power
- GSM-900 in .eu overlaps ISM in USA
  - 902-928MHz is not a GSM band in the USA



Despite all of this we could not shake an iPhone 3G\* ...

\* Other iPhones would not connect at all.

# Fun bugs exposed by OpenBTS

During testing, we saw bugs in OpenBTS and phones:

- Persistent MNO shortnames
  - Chinese student spoofed local MNO
  - Classmates connected
  - Network name of “OpenBTS”, even after BTS was removed & phones hard rebooted!
- Open / Closed registration
  - Separate from SIP-level HLR auth
  - Supposed to send “not authorized” message
  - Instead sent “You’ve been stolen” message
  - Hard reboot required, maybe more

- Still many bugs in GSM stacks
- They are being found thanks to open source

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# A5/1 is vulnerable to generic pre-computation attacks

## Code book attacks

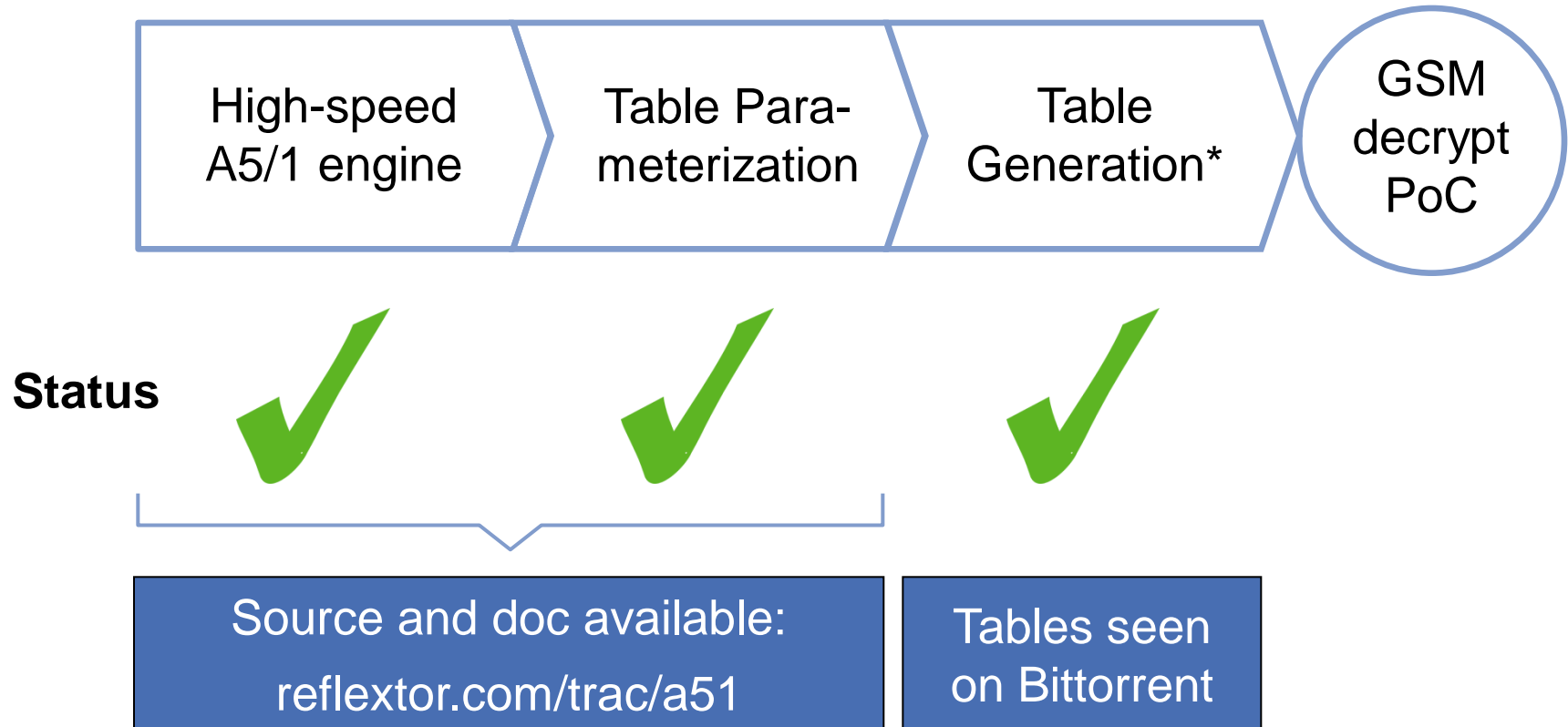
- For ciphers with small keys, code books allow decryption

Secret state	Output
A52F8C02	52E91001
62B9320A	52E91002
C309ED0A	52E91003

- Code book provides a mapping from known output to secret state
- An A5/1 code book is 128 Petabyte and takes 100,000+ years to be computed on a PC

This talk revisits techniques for computing and storing a A5/1 code book efficiently

# Groundwork for table generation is complete and released as open source



\* Community provided: fast graphics cards (NVidia or ATI) and Cell processors (Playstation)

# Key requirement of code book generation is a fast A5/1 engine

Time on single threaded CPU:  
**100,000+ years**



- A** Parallelization
- GPUs\*: hundreds of threads
  - FPGA: thousands of engines

- B** Algorithmic tweaks
- GPUs\*: compute 4 bits at once
  - FPGA: minimize critical path

- C** Implementation tweaks
- GPUs\*: bitslicing

**3 months on 40 CUDA nodes**

Latest speeds [chains/sec]

- Nvidia GTX280 – 500
- ATI HD5870 – 500
- PS3 Cell – 120

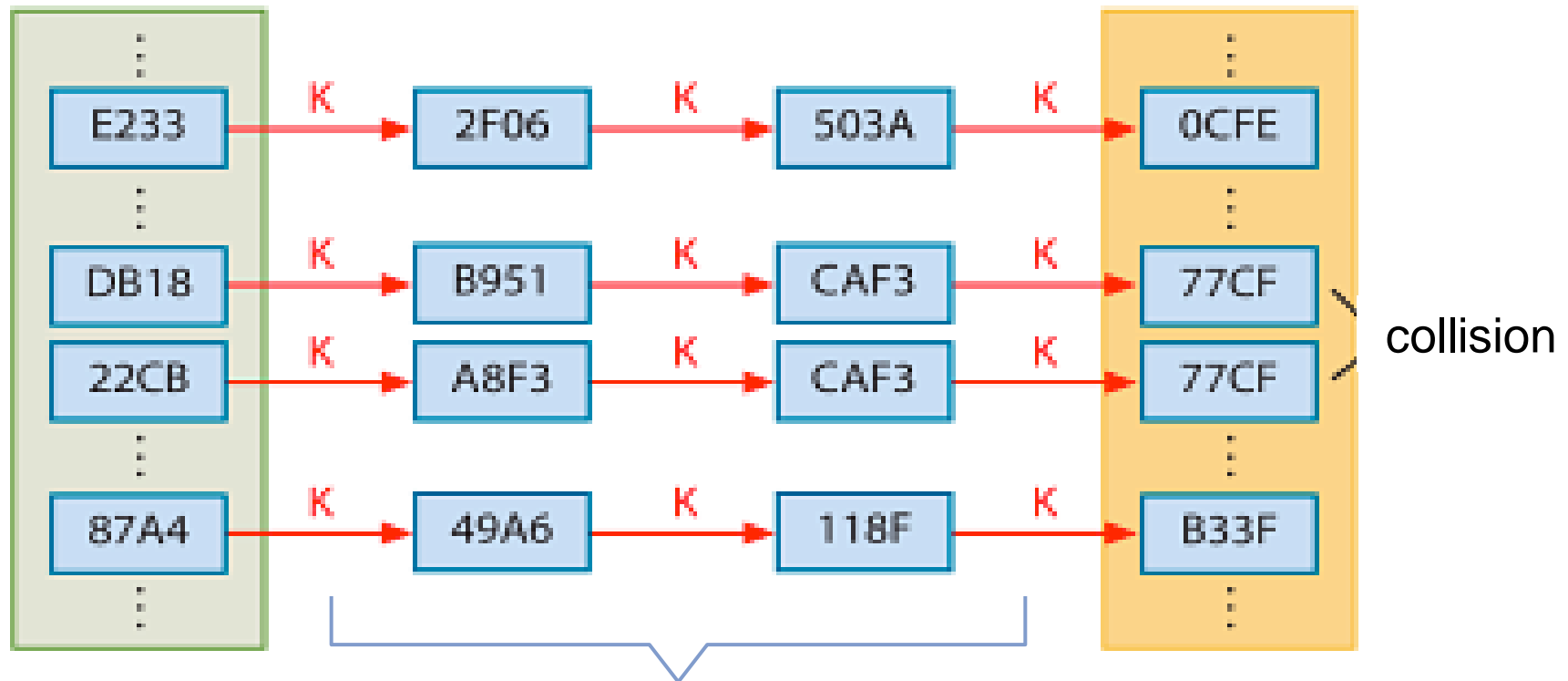
\* NVidia CUDA and ATI Brook GPUs are supported

# Cracking to be demonstrated on Wednesday

- The first tables started showing up on the congress FTPs and Bittorrents;
    - check *reflextor.com/trac/a51* for up-to-date details
  - We want more!
    - Please sort your tables before uploading (tutorial on *reflextor.com/trac/a51*)
    - After the congress, keep sharing through Bittorrent
- 
- We continue to collect tables until Tuesday evening
  - Current state to be demonstrated in **workshop**
    - Wednesday Dec 30, 13:00**, Large workshop room (A03)
    - Bring encrypted GSM sniffs you want to decrypt

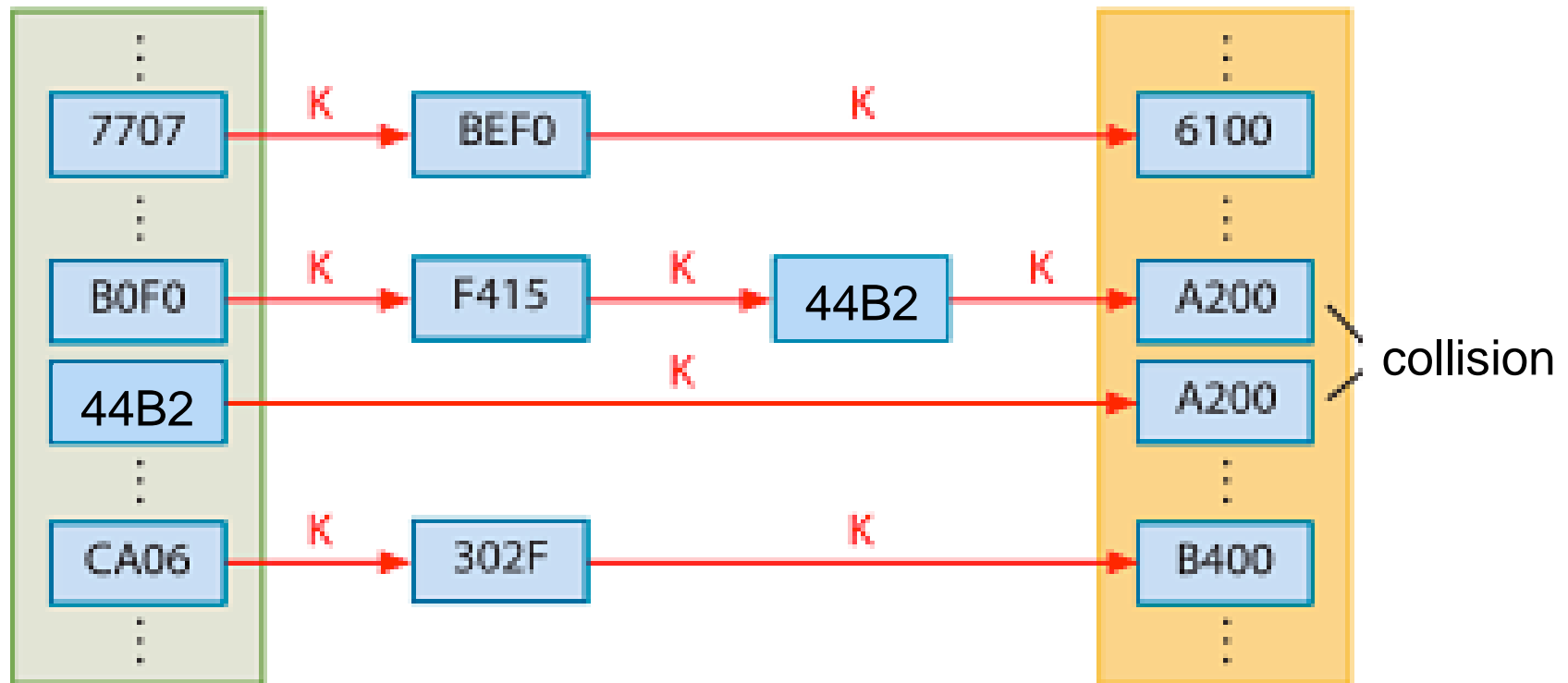


# Pre-computation tables store the code book condensed



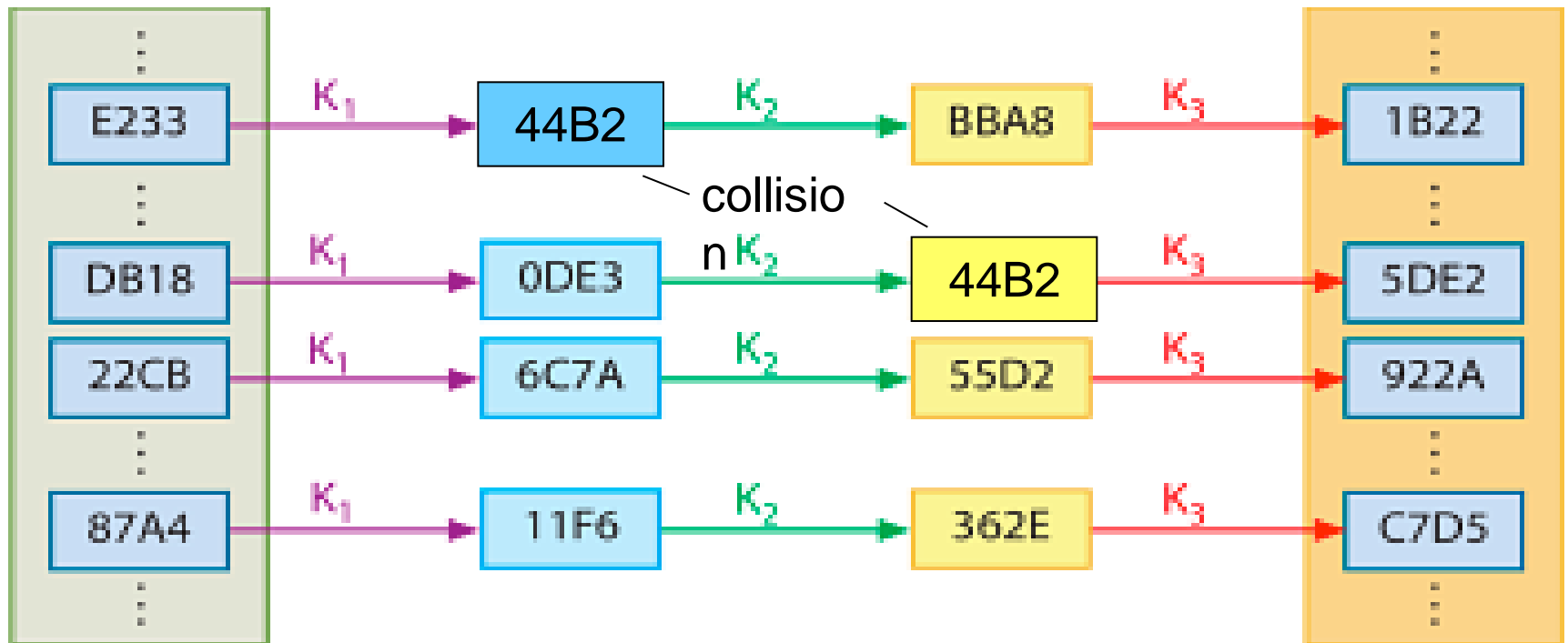
Longer chains := a) less storage, b) longer attack time

# Distinguished point tables save hard disk lookups



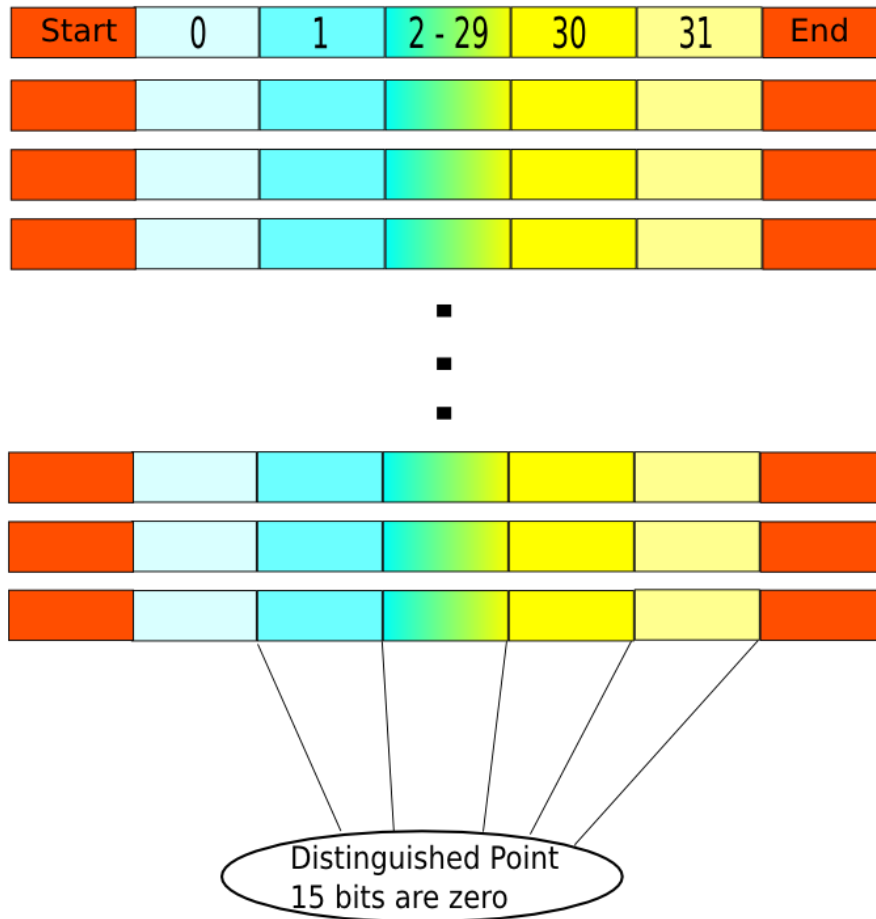
Hard disk access only needed at distinguished points

# Rainbow tables mitigate collisions



Rainbow tables have no mergers, but an exponentially higher attack time

# The combination of both table optimizations is optimal



## Assumptions

- 2 TB total storage
- 99% success rate when collecting all available keystream\*
- 50% success rate when only obvious keystream is used
- Near real-time decryption with distributed cracking network

## The most resource efficient table for A5/1 is:

- 32 DP segments of length  $2^{15}$
- Merged into one rainbow
- 380 such tables with height  $2^{28.5}$  needed

\* Collecting all available key stream requires data from a registered phone

# GSM discloses more known keystream than assumed in previous crack attempts

		Assignment			Timing known through
		Very early	Early	Late	
<b>Mobile terminated calls</b>	1. Empty Ack after 'Assignment complete'	Known Channel	Unknown Channel	Unknown Channel	<b>"Stealing bits"</b>
	2. Empty Ack after 'Alerting'	Unknown Channel	Unknown Channel	Unknown Channel	
	3. 'Connect Acknowledge'	Unknown Channel	Unknown Channel	Unknown Channel	
	4. Idle filling on SDCCH (multiple frames)			Known Channel	
	5. System Information 5+6 (~1/sec)	Known Channel	Unknown Channel	Known Channel	
<b>Network terminated calls</b>	1. Empty Ack after 'Cipher mode complete'	Known Channel	Known Channel	Known Channel	<b>Counting frames</b>
	2. 'Call proceeding'	Known Channel	Known Channel	Known Channel	
	3. 'Alerting'	Known Channel	Unknown Channel	Known Channel	
	4. Idle filling (multiple frames)			Known Channel	<b>"Stealing bits"</b>
	5. 'Connect'	Known Channel	Unknown Channel	Unknown Channel	
	6. System Information 5+6 (~1/sec)	Known Channel	Unknown Channel	Known Channel	

# Industry responds by creating a new challenge

“... the GSM **call has to be** identified and **recorded** from the radio interface. [...] we strongly suspect **the team** developing the intercept approach **has underestimated its practical complexity**.

**A hacker would need a radio receiver system and the signal processing software necessary to process the raw radio data.”** – GSMA, Aug.‘09

These remaining components of an interceptor could be repurposed from open source projects

# Hypothetically, an interceptor can be built from open source components

## Component

### “Radio receiver system”



– USRP2 –

## Current capabilities

- Capture 25 MHz of GSM spectrum
  - Typically enough for one operator
  - Need separate boards for up- and down-link

## Needed Improvements

- Decrease data rate (from 40MB/s)
  - Ability to detect active channels (energy detector)
  - Execute steps of decode chain in FPGA (Tune and Decimate)

**Bottleneck: USRP programming**

### “Signal processing software”

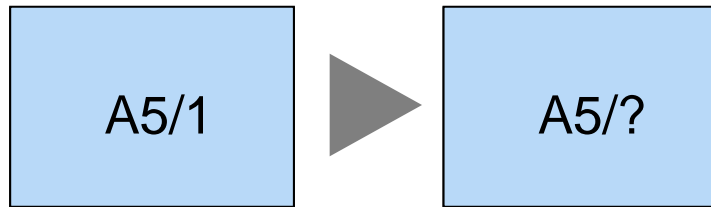
**OpenBTS**

- Decode and record one GSM channel
  - Interpret control channel data
- 
- Decrypt A5/1 packets given the correct key (v.3.0?)
  - Decode several channels (v.3.0)



# GSM's security must be overhauled

Upgrading GSM's encryption function should be a mandatory security patch



However, replacing A5/1 with A5/3 may not be enough:

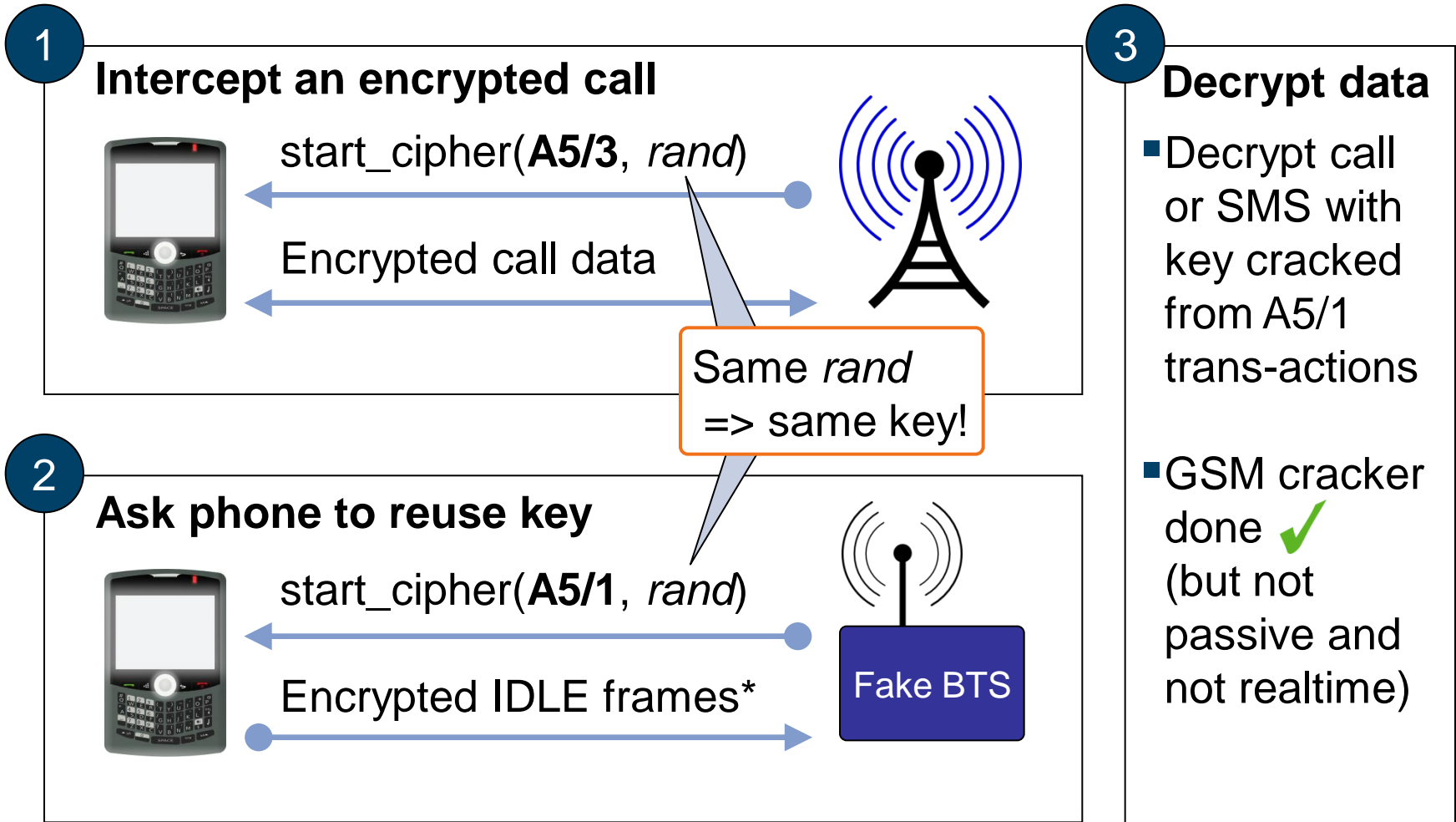
- A** The A5/3 cipher “Kasumi” is academically broken
- B** The same keys are used for A5/1 and A5/3 (weakest link security)

## Summary of the Attack on Kasumi:

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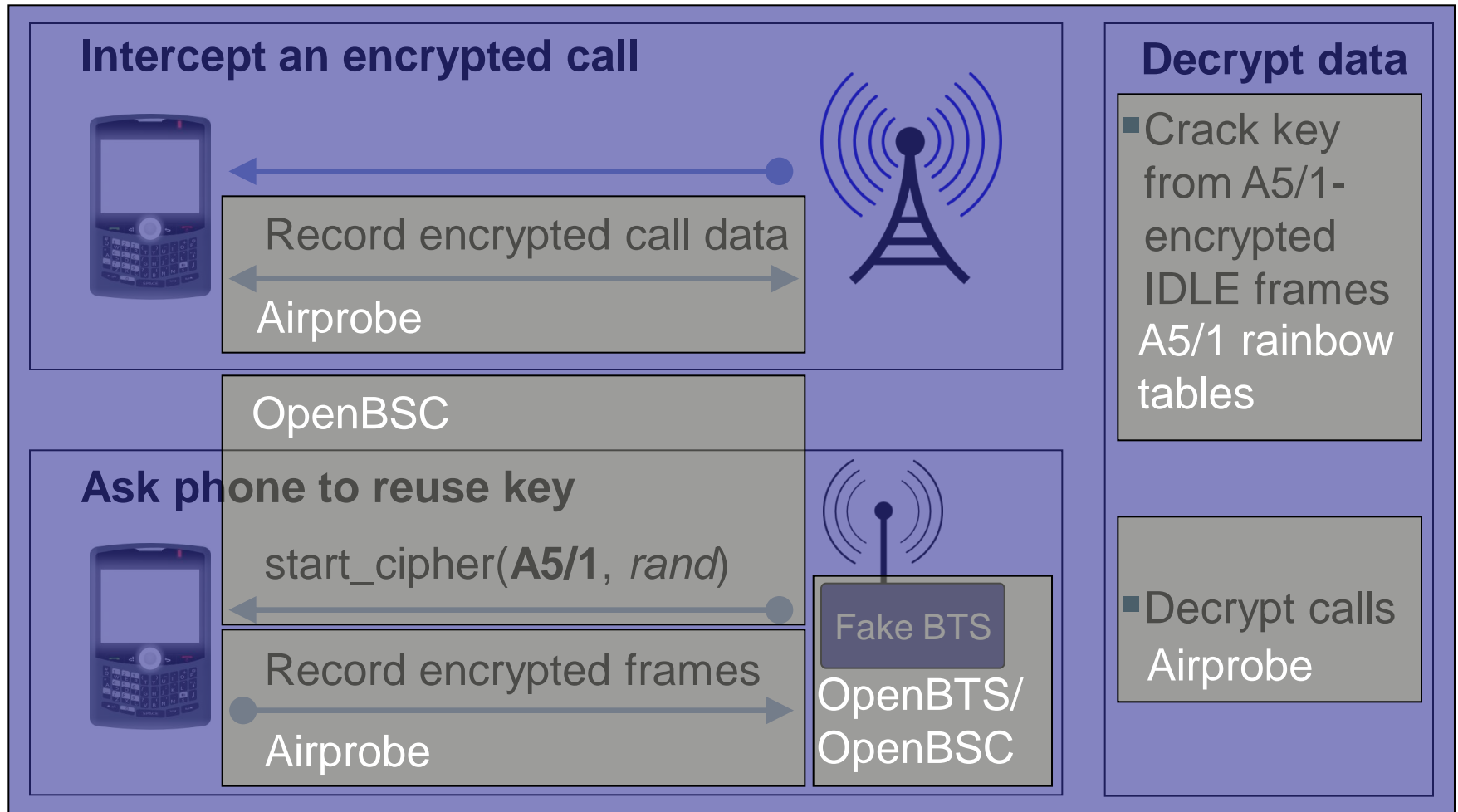
- ◆ Data complexity:  $2^{26}$  plaintexts/ciphertexts
- ◆ Space complexity:  $2^{30}$  bytes (one gigabyte)
- ◆ Time complexity:  $2^{32}$  (hardest part: ex search)
- ◆ Completely practical complexities
- ◆ Attack verified by actual software simulation

## B A5/3 can be cracked in a semi-active attack



\* IDLE frames contain known plaintext

# B All tools needed for the semi-active attack are openly available



\* IDLE frames contain known plaintext

# A5/1 cracking is just the first step ...

- Pre-computation framework build to be generic
    - Any cipher with small key space
    - Flexible table layout
    - Various back ends: CPU, CUDA, ATI, FPGA
  - All tools released open source
- 
- Please get involved
    - Port table generator to cipher in your projects
    - Find data to be decrypted (i.e., through programming the USRP's FPGA)

## Questions?

Documentation, Source	<a href="http://reflextor.com/trac/a51">reflextor.com/trac/a51</a>
Mailing list	<a href="http://tinyurl.com/a51list">tinyurl.com/a51list</a>
c't article	<a href="http://tinyurl.com/ct-rainbows">tinyurl.com/ct-rainbows</a>

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