

**Q1:** Given six memory partitions of M1: 300 KB, M2: 600 KB, M3: 350 KB, M4: 200 KB, M5: 750 KB, and M6: 125 KB (in order) and five processes P1: 115 KB, P2: 500 KB, P3: 358 KB, P4: 200 KB, and P5: 375 KB (in order).

Find the allocation mapping for **First-fit**, **Best-fit** and **Worst-fit** algorithms. Also indicate if any requests cannot be satisfied.

Allocation Technique	P1	P2	P3	P4	P5
First fit	M1	M2	M5	M3	M5/M9 (newly created)

Allocation Technique	P1	P2	P3	P4	P5
Best fit	M6	M2	M5	M4	M5/M9 (newly created)

Allocation Technique	P1	P2	P3	P4	P5
Worst fit	M5	M5/ M7 (newly created)	M2	M3	NA

**Q2:** Assuming a 1-KB page size, what are the page numbers and offsets for the following address references (provided as decimal numbers)?

- (a) 21205                      (b) 164250                      (c) 121357

Page number = Floor\_of(address/page\_size)

Offset = address mod page\_size

- a. 21205: page number: 20, offset: 725
- b. 164250: page number: 160, offset = 410
- c. 121357: page number: 118, offset: 525

**Q3:** A processor has a cycle time of 1 $\mu$ s. It costs an additional 1 $\mu$ s to access a page other than the current one. One percent of all instructions executed access a page other than the current page. Of these 80% access a page already in memory, while the remaining 20% causes a page fault. The page table is memory resident, and the memory access time is 1 ms.

Each page is 4KB in size. If a page is not in memory, it must be fetched from a secondary drum memory. The drum rotates at 3000 RPM and transfers data at 1 million words/second. Each word is 4 bytes, and assume the average rotational latency.

The system also has a translation lookaside buffer (TLB) with a 99% hit ratio and an access time of 100μs.

**Calculate and clearly mention your answer with in a box**

- a. The total time required to fetch a page from the drum when a page fault occurs (seek time and controller overhead are negligible).
- b. Calculate the effective access time (EAT) for an instruction without TLB
- c. Calculate the effective access time (EAT) with TLB

This solution assumes that the only component contributing to page fault service time is the drum access, with all other factors considered negligible.

**Step 1:** Time to fetch a page from the Drum = Average Rotational latency + Transfer time

Rotational latency =  $60/3000 = 0.02$  seconds = 20 ms

Average rotational latency =  $20 \text{ ms} / 2 = 10 \text{ ms}$

Page size is 4 KB = 1024 words ( 1 word = 4 bytes)

Transfer time =  $1024 \text{ words} / (1,000,000 \text{ words per second}) = 0.001024 \text{ sec} = 1.024 \text{ ms}$

Total time to fetch a page from the drum is  $10 \text{ ms} + 1.024 \text{ ms} = \mathbf{11.024 \text{ ms}}$

**Step 2:** Effective access time without TLB

1. Instruction accesses the current page:

This happens with a probability of 99% i.e.,  $0.99 * 1 \text{ microsecond}$

2. Instruction access to a page other than the current page

This occurs with a probability of 1 % = 0.01

- a. Page already in memory (80% = 0.8) i.e.,  $0.8 * (1 \text{ ms} + 1 \text{ ms})$ . Page table access time and memory access time
- b. Page not in memory (20% = 0.2) i.e.,  $0.2 * (1 \text{ ms} + 1 \text{ ms} + \text{page fault service time})$ . Page table access time, memory access time and page fault service time.

$$\text{EAT} = 0.01 * (1 \text{ microsecond} + 0.8 * (1 \text{ ms} + 1 \text{ ms}) + 0.2 * (1 \text{ ms} + 1 \text{ ms} + 11.024 \text{ ms}))$$

Final EAT without TLB is  $0.99 * 1 \text{ microsecond} + 0.01 * (1 \text{ microsecond} + 0.8 * (1 \text{ ms} + 1 \text{ ms}) + 0.2 * (1 \text{ ms} + 1 \text{ ms} + 11.024 \text{ ms})) = \mathbf{43.048 \text{ microseconds}}$

**Step 3: Effective access time with TLB**

TLB hit ratio is 99% i.e.,  $0.99 * (100 \text{ microseconds} + 1 \text{ ms})$

TLB miss ratio is 1% i.e.,  $0.01 * (100 \text{ microseconds} + 1 \text{ ms} + 1 \text{ ms})$

$$0.99 * 1 \text{ microsecond} + 0.01 * (1 \text{ microsecond} + 0.8 * (0.99 * (100 \text{ microseconds} + 1 \text{ ms}) + 0.01 * (100 \text{ microseconds} + 1 \text{ ms} + 1 \text{ ms})) + 0.2 * (100 \text{ microseconds} + 1 \text{ ms} + 11.024 \text{ ms})) = \mathbf{35.128 \text{ microseconds}}$$

**Q4:** Suppose that during a seek the disk accelerates the disk arm at a constant rate for the first half of the seek and decelerates the disk arm at the same rate for the second half of the seek. Suppose a disk drive has 5,000 cylinders, numbered 0 to 4999. The drive is currently serving a requisition at cylinder 143, and previous request was at cylinder 125. The queue of pending requests, in FIFO order is:

86, 1470, 913, 1774, 948, 1509, 1022, 1750, 130

Assume that the disk arm can perform a seek to an adjacent cylinder in 1 millisecond and full stroke seek over all 5,000 cylinders in 18 milliseconds.

Calculate the total seek time to serve all the requests for each of the following schedules:

- a. **FCFS**                      b. **SSTF**                      c. **SCAN**

**Hint:**  $T_{\text{seek}} = T_{\text{adjacent}} + (T_{\text{full stroke}} - T_{\text{adjacent}}) * \sqrt{\frac{\text{cylinders to move}}{\text{Total cylinders}}}$

Substituting the values in the formulae gives  $1 + 0.24 * \sqrt{\text{cylinders to move}}$

For FCFS

$$9 + 0.24 (\sqrt{143 - 86} + \sqrt{1470 - 86} + \sqrt{1470 - 913} + \sqrt{1774 - 913} + \sqrt{1774 - 948} + \sqrt{1509 - 948} + \sqrt{1509 - 1022} + \sqrt{1750 - 1022} + \sqrt{1750 - 130}) = \mathbf{66.412}$$

For SJF

$$9 + 0.24 (\sqrt{143 - 130} + \sqrt{130 - 86} + \sqrt{913 - 86} + \sqrt{948 - 913} + \sqrt{1022 - 948} + \sqrt{1470 - 1022} + \sqrt{1509 - 1470} + \sqrt{1750 - 1509} + \sqrt{1774 - 1750}) = \mathbf{33.312}$$

For SCAN

$$9 + 0.24 (\sqrt{913 - 143} + \sqrt{948 - 913} + \sqrt{1022 - 948} + \sqrt{1470 - 1022} + \sqrt{1509 - 1470} + \sqrt{1750 - 1509} + \sqrt{1774 - 1750} + \sqrt{1774 - 130} + \sqrt{130 - 86}) = \mathbf{41.911}$$

**Q5:** In a certain page-replacement algorithm, each page frame is associated with a counter. During a page fault, the page frame with the smallest counter value is replaced, and its counter is **reset**. When a page in memory is referenced, its associated counter is **incremented**. During both page faults and memory references, the counters of all page frames other than the referenced page are **decremented**.

How many page faults occur for the following reference string with **four** page frames and has a current configuration shown in the table on the **next column**.

1, 2, 3, 4, 5, 3, 4, 1, 6, 7, 8, 7, 8, 9, 7, 8, 9, 5, 4, 5, 4, 2

If multiple page frames share the **same** counter value, apply the least recently used (LRU) strategy to select a page frame for replacement. Marks will be awarded only for accurately depicting each page frame along with counter values at every step in a clear pictorial format **and** correctly determining the total number of page faults.

1 2 3 4 5 3 4 1 6 7 8 7 8 9 7 8 9 5 4 5 4 2

Diagram illustrating the page replacement process for the reference string: 1, 2, 3, 4, 5, 3, 4, 1, 6, 7, 8, 7, 8, 9, 7, 8, 9, 5, 4, 5, 4, 2.

The diagram shows the state of the 4 page frames at each step, with counter values for each page. The frames are labeled P (Page Fault) at steps 1, 2, 3, 4, 5, 10, 11, 12, 13, 14, 15, 16, and 17.

Final state of the page frames (Step 21):

4	0
5	-1
8	-6
9	-5

4	-1
5	-2
2	0
9	-6

no. of Page faults = 13

