

## 1 Polynomials Intro

**Note 8** **Polynomial:**  $f(x) = a_0 + a_1x + a_2x^2 + \cdots + a_nx^n$ ; in terms of roots,  $f(x) = a(x - r_1)(x - r_2) \cdots (x - r_k)$

**Degree of a polynomial:** the highest exponent in the polynomial

**Galois Field:** denoted as  $\text{GF}(p)$ , it's basically just a fancy way of saying that we're working modulo  $p$ , for a prime  $p$

**Properties** (true over  $\mathbb{R}$  and also over  $\text{GF}(p)$ ):

- Polynomial of degree  $d$  has at most  $d$  roots.
- Exactly one polynomial of degree at most  $d$  passes through  $d + 1$  points.

**Lagrange Interpolation:** Given  $d + 1$  points  $(x_1, y_1), (x_2, y_2), \dots, (x_{d+1}, y_{d+1})$ , we define

$$\Delta_i(x) = \frac{\prod_{j \neq i} (x - x_j)}{\prod_{j \neq i} (x_i - x_j)}.$$

The unique polynomial through all points is  $f(x) = \sum_{i=1}^{d+1} y_i \cdot \Delta_i(x)$

**Secret Sharing:** We make use of the fact that there is a unique polynomial of degree  $d$  passing through a given set of  $d + 1$  points. This means that if we require  $k$  people to come together in order to find a secret, we should use a polynomial of degree  $k - 1$ , and give each person one point. There are more complicated schemes if there are more conditions, but they all use the same concept.

- Consider the  $\Delta_i(x)$  polynomials in Lagrange interpolation. What is the value of  $\Delta_i(x)$  for  $x = x_i$ , and what is its value for  $x = x_j$ , where  $j \neq i$ ? How is this similar to the process of computing a solution with CRT?
- If we perform Lagrange interpolation over  $\text{GF}(p)$  instead of over  $\mathbb{R}$ , what is different?

## 2 Polynomial Practice

Note 8

- (a) If  $f$  and  $g$  are non-zero real polynomials, how many real roots do the following polynomials have at least? How many can they have at most? (Your answer may depend on the degrees of  $f$  and  $g$ .)
- (i)  $f + g$
  - (ii)  $f \cdot g$
  - (iii)  $f/g$ , assuming that  $f/g$  is a polynomial
- (b) Now let  $f$  and  $g$  be polynomials over  $\text{GF}(p)$ .
- (i) We say a polynomial  $f = 0$  if  $\forall x, f(x) = 0$ . Show that if  $f \cdot g = 0$ , it is not always true that either  $f = 0$  or  $g = 0$ .
  - (ii) How many  $f$  of degree *exactly*  $d < p$  are there such that  $f(0) = a$  for some fixed  $a \in \{0, 1, \dots, p - 1\}$ ?
- (c) Find a polynomial  $f$  over  $\text{GF}(5)$  that satisfies  $f(0) = 1, f(2) = 2, f(4) = 0$ . How many such polynomials of degree at most 4 are there?

### 3 Lagrange Interpolation in Finite Fields

Note 8

In this problem, we will break down the terms of Lagrange interpolation by working through an example, where we want to find a unique polynomial  $p(x)$  of degree at most 2 that passes through points  $(-1, 3)$ ,  $(0, 1)$ , and  $(1, 2)$  in modulo 5 arithmetic.

- (a) First, assume we have polynomials  $p_{-1}(x)$ ,  $p_0(x)$ , and  $p_1(x)$  satisfying:

$$\begin{aligned} p_{-1}(0) &\equiv p_{-1}(1) \equiv 0 \pmod{5}; & p_{-1}(-1) &\equiv 1 \pmod{5} \\ p_0(-1) &\equiv p_0(1) \equiv 0 \pmod{5}; & p_0(0) &\equiv 1 \pmod{5} \\ p_1(-1) &\equiv p_1(0) \equiv 0 \pmod{5}; & p_1(1) &\equiv 1 \pmod{5} \end{aligned}$$

Construct  $p(x)$  using a linear combination of  $p_{-1}(x)$ ,  $p_0(x)$ , and  $p_1(x)$ .

- (b) Find  $p_{-1}(x)$ . In other words, find a degree 2 polynomial that has roots at  $x = 0$  and  $x = 1$  and evaluates to 1 at  $x = -1$  (all in modulo 5).

- (c) Find  $p_0(x)$ .

- (d) Find  $p_1(x)$ .

- (e) Now, lets put it all together! Create a suitable polynomial  $p(x)$  by using the linear combination and polynomials constructed above.