

Error Correcting Codes and Secret Sharing Intro

Secret Sharing: We make use of the fact that there is a unique polynomial of degree d passing through a given set of $d + 1$ points. This means that if we require k people to come together in order to find a secret, we should use a polynomial of degree $k - 1$, and give each person one point. There are more complicated schemes if there are more conditions, but they all use the same concept.

Erasure errors: A packet/point $(x, P(x))$ is *lost* in the communication channel.

$$\boxed{1 \mid 5 \mid 3 \mid 4 \mid 3} \quad \longrightarrow \quad \boxed{1 \mid 5 \quad \quad 4 \mid 3}$$

Protection: (n packets, k errors) interpolate polynomial through the message points, send $n + k$ packets on the polynomial

General errors: A packet is *modified* in the communication channel.

$$\boxed{1 \mid 5 \mid 3 \mid 4} \quad \longrightarrow \quad \boxed{1 \mid 5 \mid 6 \mid 4}$$

Protection: (n packets, k errors) interpolate polynomial through the message points, send $n + 2k$ packets on the polynomial

Berlekamp–Welch Algorithm:

Variables: sent message packets m_i , received packets r_i , error locations e_i

Polynomials:

- $P(x)$: original polynomial through message (this is what we want)
- $E(x) = (x - e_1)(x - e_2) \cdots (x - e_k)$: error locator polynomial
- $Q(x) = P(x)E(x)$, or $Q(i) = P(i)E(i) = r_iE(i)$ for all i

$Q(x)$ and $E(x)$ are unknown, but we can solve for them using a system of equations.

1 Berlekamp–Welch Warm Up

Note 8
Note 9

Let $P(i)$, a polynomial applied to the input i , be the original encoded polynomial before sent, and let r_i be the received info for the input i which may or may not be corrupted.

(a) If you want to send a length- n message, what should the degree of $P(x)$ be? Why?

(b) When does $r_i = P(i)$? When does r_i not equal $P(i)$?

- (c) If there are at most k erasure errors, how many packets should you send? If there are at most k general errors, how many packets should you send? (We will see the reason for this later.) Now we will only consider general errors.
- (d) What do the roots of the error polynomial $E(x)$ represent? Does the receiver know the roots of $E(x)$? If there are at most k errors, what is the maximum degree of $E(x)$? Using the information about the degree of $P(x)$ and $E(x)$, what is the degree of $Q(x) = P(x)E(x)$?
- (e) Why is the equation $Q(i) = P(i)E(i) = r_i E(i)$ always true? (Consider what happens when $P(i) = r_i$, and what happens when $P(i)$ does not equal r_i .)
- (f) In the polynomials $Q(x)$ and $E(x)$, how many total unknown coefficients are there? (These are the variables you must solve for. Think about the degree of the polynomials.) When you receive packets, how many equations do you have? Do you have enough equations to solve for all of the unknowns? (Think about the answer to the earlier question - does it make sense now why we send as many packets as we do?)
- (g) If you have $Q(x)$ and $E(x)$, how does one recover $P(x)$? If you know $P(x)$, how can you recover the original message?

2 Berlekamp-Welch Algorithm

Note 8

Note 9

In this question we will send the message $(m_0, m_1, m_2) = (1, 1, 4)$ of length $n = 3$. We will use an error-correcting code for $k = 1$ general error, doing arithmetic over GF(5).

- (a) Construct a polynomial $P(x) \pmod{5}$ of degree at most 2, so that

$$P(0) = 1, \quad P(1) = 1, \quad P(2) = 4.$$

What is the message $(c_0, c_1, c_2, c_3, c_4)$ that is sent?

- (b) Suppose you receive the message $(0, 1, 4, 0, 4)$ and know that one packet was corrupted. Set up the system of linear equations in the Berlekamp-Welch algorithm to find $Q(x)$ and $E(x)$.

- (c) Assume that after solving the equations in part (b) we get $Q(x) = 4x^3 + x^2 + x$ and $E(x) = x$. Show how to recover the original message from Q and E .

3 Secrets in the United Nations

Note 8

A vault in the United Nations can be opened with a secret combination $s \in \mathbb{Z}$. In only two situations should this vault be opened: (i) all 193 member countries must agree, or (ii) at least 55 countries, plus the U.N. Secretary-General, must agree.

- (a) Propose a scheme that gives private information to the Secretary-General and all 193 member countries so that the secret combination s can only be recovered under either one of the two specified conditions.

 - (b) The General Assembly of the UN decides to add an extra level of security: each of the 193 member countries has a delegation of 12 representatives, all of whom must agree in order for that country to help open the vault. Propose a scheme that adds this new feature. The scheme should give private information to the Secretary-General and to each representative of each country.