

Definition.

1. **directed graph** G is a pair (V, E) , where V is a finite set and E is a binary relation on V . V is a **vertex set** of G , and its elements are called **vertices**. The set E is called the edge set of G and its elements are called edges
 - (a) If (u, v) is an edge in a directed graph, (u, v) is **incident from** or **leaves** vertex u and is **incident to** or **enters** vertex v
 - (b) In the case of directed graph, if v is **adjacent** to u , then $u \rightarrow v$
 - (c) the **out-degree** of a vertex is the number of edges leaving it, and the **in-degree** of a vertex is the number of edges entering it. The **degree** of a vertex in a directed graph is its in-degree plus its out-degree
 - (d) In directed graph, a path $\langle v_0, \dots, v_k \rangle$ forms a **cycle** if $v_0 = v_k$ and the path contains at least one edges;
 - i. The **cycle is simple** if, in addition, v_1, \dots, v_k are distinct
 - ii. A self-loop is a cycle of length 1, i.e. (v, v)
 - iii. A directed graph with no self-loop is **simple**
 - iv. A graph with no cycles is **acyclic**
 - (e) A directed graph is **strongly connected** if
 - i. every two vertices are reachable from each other
 - ii. has exactly one strongly connected component
 - (f) the **strongly connected component** of a directed graph are equivalence classes of vertices under the 'are mutually reachable' relation
 - (g) Given a directed graph $G = (V, E)$, the **undirected version** of G is the undirected graph $G' = (V, E')$, where $(u, v) \in E'$ if and only if $u \neq v$ and $(u, v) \in E$ (that is, undirected version contains edges of G with directions removed and self-loop eliminated)
 - (h) the **neighbor** of a vertex u is any vertex that is adjacent to u in the undirected version of G (that is, v is a neighbor of u if $u \neq v$ and either $(u, v) \in E$ or $(v, u) \in E$).
2. **undirected graph** $G = (V, E)$, the set E consists of unordered pairs of vertices, rather than ordered pairs, i.e. exists set $\{u, v\}$, where $u, v \in V$ and $u \neq v$. (by convention, we use $(u, v) = (v, u)$ to denote a set)
 - (a) If (u, v) is an edge in an undirected graph, then (u, v) is **incident on** vertices u and v
 - (b) If (u, v) is an edge in $G = (V, E)$, then vertex v is **adjacent** to vertex u
 - (c) The **degree** of a vertex in an undirected graph is the number of edges incident on it; A vertex with degree 0 is **isolated**

- (d) In undirected graph, a path $\langle v_0, \dots, v_k \rangle$ forms a **cycle** if $k \geq 3$ and $v_0 = v_k$
 - i. The **cycle is simple** if v_1, \dots, v_k are distinct
 - ii. A graph with no cycles is **acyclic**
 - (e) An undirected graph is **connected** if
 - i. every vertex is reachable from all other vertices
 - ii. has exactly one connected component
 - (f) Given undirected $G = (V, E)$, the **directed version of G** is the directed graph $G' = (V, E')$, where $(u, v) \in E'$ if and only if $(u, v) \in E$. (replacing each undirected edge (u, v) by two directed edges (u, v) and (v, u))
 - (g) u and v are **neighbors** if they are adjacent
3. The **connected component** of a graph are equivalence classes of vertices under the 'is reachable from' relation
- (a) The **edges of a connected component** are those that are incident on only the vertices of the component, i.e. (u, v) is an edge of a connected component if both u and v are vertex of the component

Definition. A **path** of **length** k from a vertex u to a vertex u' in a graph $G = (V, E)$ is a sequence $\langle v_0, v_1, \dots, v_k \rangle$ of vertices such that $u = v_0$ and $u' = v_k$, and $(v_{i-1}, v_i) \in E$ for $i = 1, 2, \dots, k$

- 1. the **length** of the path is the number of edges in the path
- 2. If there is a path from u to u' , we say that u' is **reachable** from u via p , which we sometimes write as $u \xrightarrow{p} u'$ if G is directed
- 3. A path is **simple** if all vertices in the path are distinct
- 4. A **subpath** of path $p = \langle v_0, v_1, \dots, v_k \rangle$ is a contiguous subsequence of its vertices. That is, for any $0 \leq i \leq j \leq k$, subsequence of vertices $\langle v_i, \dots, v_j \rangle$ is a subpath of p

Definition. Isomorphism and Subgraphs

- 1. Two graphs $G = (V, E)$ and $G' = (V', E')$ are **isomorphic** if there exists a bijection $f : V \rightarrow V'$ such that $(u, v) \in E$ if and only if $(f(u), f(v)) \in E'$. In other words, we can relabel the vertices of G and be the vertices of G' , maintaining the corresponding edges in G and G'
- 2. A graph $G' = (V', E')$ is a **subgraph** of $G = (V, E)$ if $V' \subseteq V$ and $E' \subseteq E$. Given a set $V' \subseteq V$, the subgraph of G **induced** by V' is the graph $G' = (V', E')$ where

$$E' = \{(u, v) \in E : u, v \subseteq V'\}$$

Definition. Types of graphs

1. A **Complete graph** is an undirected graph in which every pair of vertices is adjacent
2. A **Bipartite graph** is an undirected graph $G = (V, E)$ in which V can be partitioned into two sets V_1 and V_2 such that $(u, v) \in E$ implies either $u \in V_1$ and $v \in V_2$ or $u \in V_2$ and $v \in V_1$
3. A **Weighted graph** is a graph for which each edge has an associated **weight**, given by a weight function $w : E \rightarrow \mathbb{R}$.
4. A **forest** is an acyclic, undirected graph
5. A **tree** is a connected, acyclic, undirected graph
6. A **directed acyclic graph (DAG)** is as its name suggests