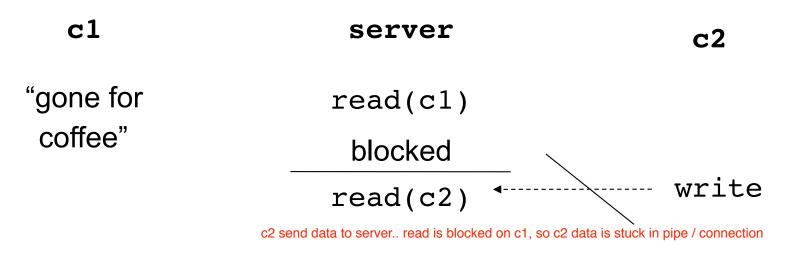
### Status

- A4 will be out later today
  - No extensions are possible
  - You may use late penalties
- Exercises 10 and 11 are posted
  - We will take the best 9 out of 10 marks for labs.
  - Exercise 7 was not for credit.
  - This sort of makes exercise 11 optional

# I/O Multiplexing

Kerrisk Ch 63

## The problem



When reading from multiple sources, blocking on one of the sources could be bad.

An example of denial of service.

One solution: one process for every client. What are the pros and cons of this solution?

<sup>1.</sup>accept call generate socket.

<sup>2.</sup>incoming connection is being forked and accept data in a separate process

<sup>3.</sup> may be problematic

a. too many child process

b. shared data between child processes

## Another way to look at the problem

#### Server

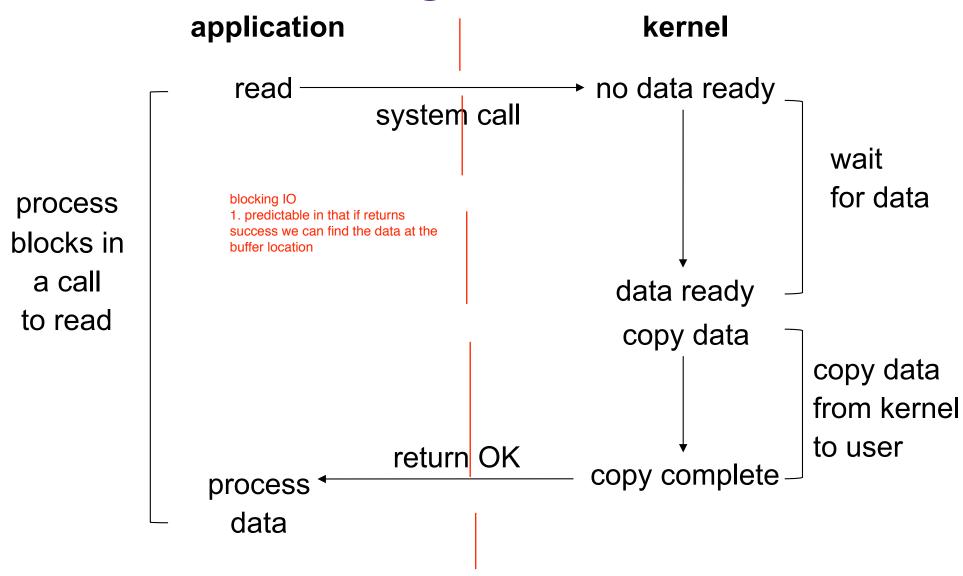
```
while(1)
  accept a new connection
  for each existing connection
    read
    write
```

Which of the system calls might block indefinitely? read and accept

So what happens if there is only one connection?

blocked in accept call ...

## Blocking I/O Model

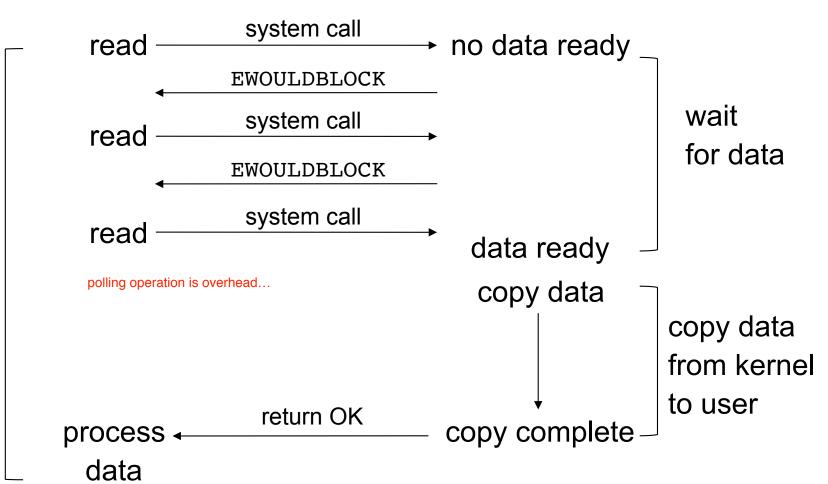


# Nonblocking I/O Model

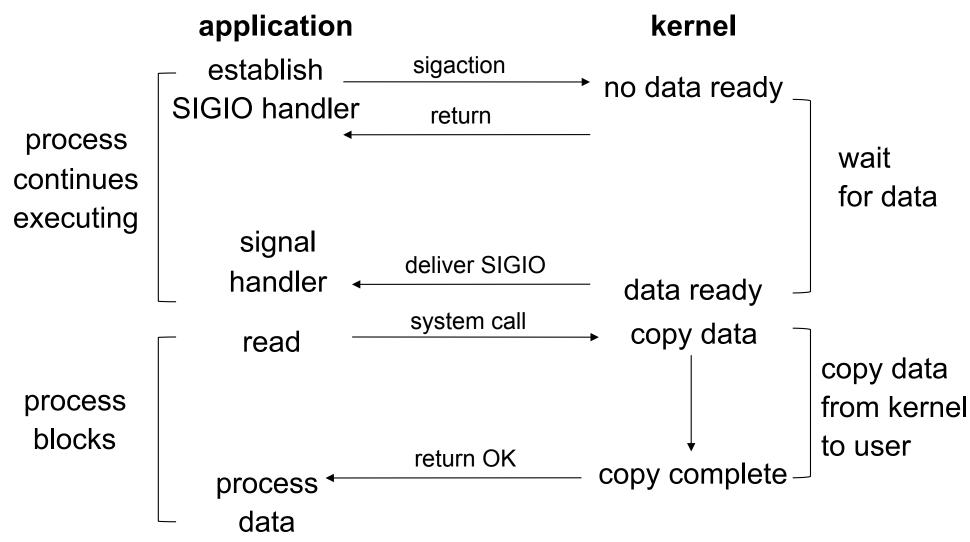
#### application

#### kernel

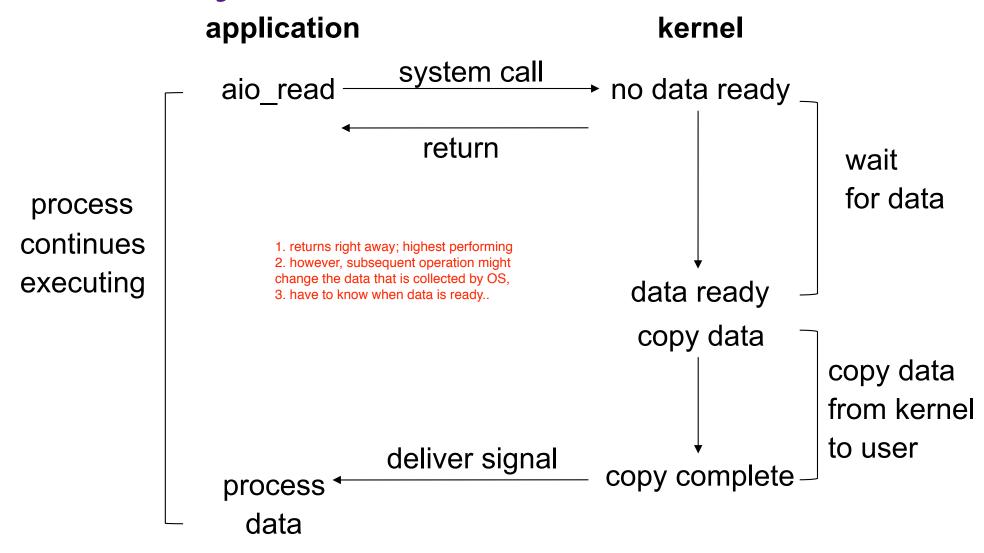
process
repeatedly
calls read
waiting for
an OK
(polling)



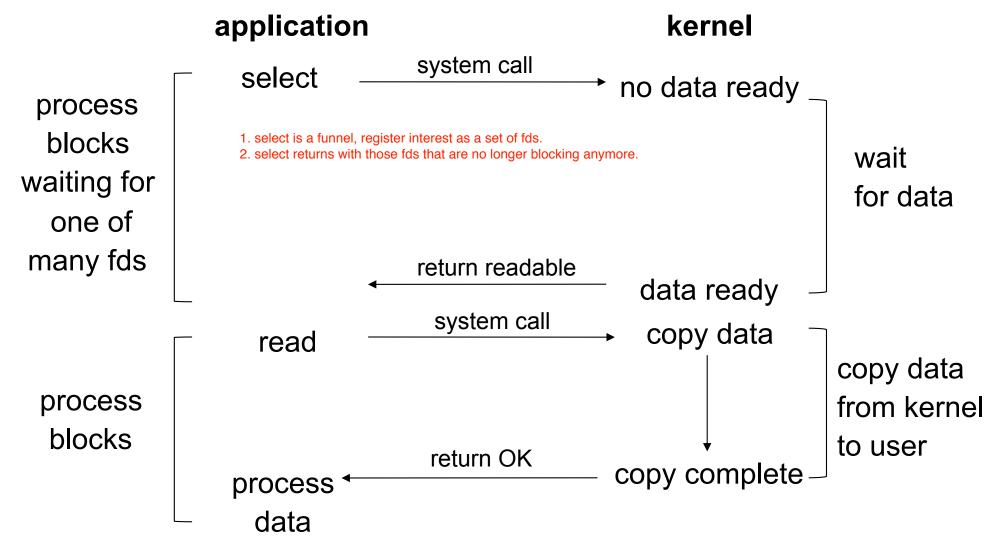
# Signal Driven I/O Model



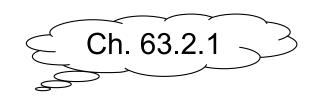
## Asynchronous I/O Model



# I/O Multiplexing Model



## select()



how far into set we have to search

A call to select returns when one of the file descriptors in one of the sets is ready for I/O.

If timeout is not NULL, then select returns when a descriptor is ready or timeout time has passed.

If timeout is 0, select returns immediately after checking descriptors.

### Readiness

### Ready to read when

there is data in the receive buffer to be read

end-of-file state on file descriptor

the socket is a listening socket and there is a <u>connection</u> pending

a socket error is pending

#### Ready to write when

there is space available in the write buffer a socket error is pending

### Exception condition pending when

TCP out-of-band data

We are typically interested in when bytes are available to be read, but sometimes we use select on write or exception sets

### select timeout

 The timeout specifies how long we're willing to wait for a fd to become ready

- If timeout is <u>NULL</u>, wait forever (or until we catch a signal)
- If timeout is zero, test and return immediately
- Otherwise wait up to specified timeout
- select returns when a fd ready or we timeout

## Descriptor sets

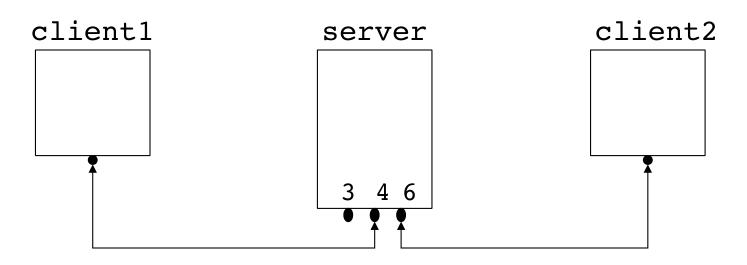
Typically implemented as an array of integers where each bit corresponds to a descriptor (except in Windows).

Implementation is hidden in the fd\_set data type FD\_SETSIZE is the number of descriptors in the data type

maxfdp1 specifies the number of descriptors to test Macros:

```
void FD_ZERO(fd_set *fdset);
void FD_SET(int fd, fd_set *fdset);
void FD_CLR(int fd, fd_set *fdset);
int FD_ISSET(int fd, fd_set *fdset);
```

## Descriptor sets



select search linearly, so efficient to tell select where to stop

fd0 fd1 fd2 fd3 fd4 fd5 fd6

$$maxfd + 1 = 7$$

After select:

rset 0 0 0 1 0 0 0

## select example

```
fd set rfds;
struct timeval tv;
int retval;
FD ZERO(&rfds); /* Watch stdin (fd 0) for input */
FD SET(STDIN_FILENO, &rfds);
tv.tv sec = 5; /* Wait up to five seconds. */
tv.tv usec = 0;
retval = select(1, &rfds, NULL, NULL, &tv);
if (retval == -1)
 perror("select()");
else if (retval > 0)
  printf("Data is available now.\n");
  /* FD ISSET(0, &rfds) will be true, can use read() */
else
 printf("No data within five seconds.\n");
```

```
for(;;) {
  rset = allset;
  nready = Select(maxfd+1, &rset ,NULL,NULL,NULL);
  if(FD ISSET(listenfd, &rset)) {
    connfd = Accept(listenfd, &caddr, &clen);
    for(i = 0; i < FD SETSIZE; i++)
         if(client[i] < 0) {
              client[i] = connfd; break;
    FD SET(connfd, &allset);
    if(connfd > maxfd) maxfd = connfd;
  for(i = 0; i <= maxi; i++) {
    if(sockfd = client[i]) < 0) continue;</pre>
    if(FD ISSET(sockfd, &rset))
         Read(sockfd, line, MAXLINE);
```

```
for(;;) {
  rset = allset;
  nready = Select(maxfd+1, &rset ,NULL,NULL);
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    FD SET(connfd, &allset);
    if(connfd > maxfd) maxfd = connfd;
  for(i = 0; i <= maxi; i++) {
    if(sockfd = client[i]) < 0) continue;</pre>
    if(FD ISSET(sockfd, &rset))
         Read(sockfd, line, MAXLINE);
```

### **End of Line**

- There are two characters that determine end-of-line
  - Carriage return (CR, \r, ^M)
  - Line feed (LF, \n)
- Early operating systems defined their own conventions using one or both of CR and LF.
  - Unix: LF,
  - DOS/Windows: CR LF
  - Mac classic: CR

# **Network Line Ending**

 Transferring data between machines with different operating systems, means deciding on a common line ending.

CR LF is the standard

 (Of course it is possible with regular expression matching to mostly ignore this issue, but still better to conform?)