COMP3121/9101: Greedy Algorithms

Module 3 Lecture Notes

Term 2, 2025

1 Overview

- Motivation: Introduction to greedy algorithms and their applications.
- Optimal Selection: Problems like activity selection and cell tower placement.
- Optimal Ordering: Problems like tape storage and minimizing job lateness.
- Optimal Merging: Huffman coding for efficient data compression.
- Applications to Graphs: Shortest paths, minimum spanning trees, and directed graph structures.
- Puzzle: A problem involving secure transmission using padlocks and boxes.

2 Greedy Algorithms

2.1 Definition

A greedy algorithm solves a problem by making locally optimal choices at each stage, hoping to find a global optimum. However, these choices may not always lead to the best overall solution.

2.2 When to Use

- Problems where local optima lead to global optima (e.g., fractional knapsack).
- Problems with optimal substructure (e.g., shortest paths).

2.3 Proof Techniques

- **Greedy Stays Ahead**: Show that the greedy choice is never worse than any other choice at each step.
- Exchange Argument: Transform any alternative solution into the greedy solution without worsening the objective.

3 Key Problems and Solutions

3.1 Fractional Knapsack

- **Problem**: Select items with weights w_i and values v_i to maximize total value without exceeding weight limit W. Items can be split.
- Solution: Sort items by $\frac{v_i}{w_i}$ in descending order and take as much as possible of the highest ratio item first.
- Time Complexity: $O(n \log n)$ due to sorting.

3.2 0-1 Knapsack

- Problem: Similar to fractional knapsack, but items cannot be split.
- Challenge: Greedy approach fails; dynamic programming is required.
- Complexity: NP-hard.

3.3 Activity Selection

- **Problem**: Select the maximum number of non-overlapping activities with given start and finish times.
- **Solution**: Sort activities by finish time and greedily select the earliest finishing activity that doesn't conflict with the last selected activity.
- Proof: Exchange argument shows optimality.
- Time Complexity: $O(n \log n)$ for sorting.

3.4 Cell Tower Placement

- **Problem**: Place the minimum number of cell towers to cover all houses along a straight road, each with a 5km range.
- **Solution**: Place towers 5km east of the leftmost uncovered house iteratively.
- **Proof**: Greedy stays ahead shows optimality.
- Time Complexity: O(n) if houses are sorted; $O(n \log n)$ otherwise.

3.5 Tape Storage

- Problem: Order files on a tape to minimize average retrieval time.
- Solution: Sort files by increasing length (or by $\frac{p_i}{L_i}$ if probabilities differ).
- **Proof**: Adjacent inversion argument.
- Time Complexity: $O(n \log n)$ for sorting.

3.6 Minimizing Job Lateness

- **Problem**: Schedule jobs to minimize maximum lateness, where lateness is the difference between finish time and deadline.
- Solution: Sort jobs by deadlines (earliest deadline first).
- Proof: Exchange argument resolves adjacent inversions.
- Time Complexity: $O(n \log n)$ for sorting.

3.7 Huffman Coding

- **Problem**: Construct a prefix code to minimize expected message length given symbol frequencies.
- **Solution**: Build a binary tree by repeatedly combining the two least frequent symbols.
- **Proof**: Optimality follows from the properties of prefix codes and greedy choices.
- Time Complexity: $O(n \log n)$ using a priority queue.

4 Applications to Graphs

4.1 Strongly Connected Components (SCCs)

- **Problem**: Find SCCs in a directed graph where every vertex is reachable from every other vertex in the same component.
- Solution: Use Kosaraju's or Tarjan's algorithm (DFS-based).
- Time Complexity: O(|V| + |E|).

4.2 Dijkstra's Algorithm

- **Problem**: Find shortest paths from a single source in a graph with nonnegative edge weights.
- Solution: Greedily select the closest vertex and relax its outgoing edges.
- Proof: Correctness relies on non-negative weights and greedy choices.
- Time Complexity: $O((n+m)\log n)$ with a priority queue.

4.3 Kruskal's Algorithm

- **Problem**: Find a minimum spanning tree (MST) in a weighted graph.
- **Solution**: Sort edges by weight and add them if they don't form a cycle (using Union-Find).
- Proof: Correctness follows from the cut property.
- Time Complexity: $O(m \log n)$ for sorting and Union-Find operations.

5 Puzzle: Secure Transmission

5.1 Problem

Bob wants to send a teddy bear to Alice using locked boxes and padlocks. Constraints:

- Non-empty boxes must be locked during transit.
- Keys cannot be sent; they must remain with the sender.
- Padlocks can only be sent if locked.

5.2 Solutions

- AND Solution (3 mailings):
 - 1. Bob locks the box with his padlock and sends it to Alice.
 - 2. Alice adds her padlock and sends it back to Bob.
 - 3. Bob removes his padlock and sends the box back to Alice.
- OR Solution (2 mailings):
 - 1. Bob sends an empty box with his padlock to Alice.
 - 2. Alice places the teddy bear in the box, locks it with her padlock, and sends it back to Bob.