

# Accelerating Similarity Search for Elastic Measures: A Study and New Generalization of Lower Bounding Distances

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## **ABSTRACT**

Similarity search is a core analytical task, and its performance critically depends on the choice of distance measure. For time-series querying, elastic measures achieve state-of-the-art accuracy but are computationally expensive. Thus, fast lower bounding (LB) measures prune unnecessary comparisons with elastic distances to accelerate similarity search. Despite decades of attention, there has never been a study to assess the progress in this area. In addition, the research has disproportionately focused on one popular elastic measure, while other accurate measures have received little or no attention. Therefore, there is merit in developing a framework to accumulate knowledge from previously developed LBs and eliminate the notoriously challenging task of designing separate LBs for each elastic measure. In this paper, we perform the first comprehensive study of 11 LBs spanning 5 elastic measures using 128 datasets. We identify four properties that constitute the effectiveness of LBs and propose the Generalized Lower Bounding (GLB) framework to satisfy all desirable properties. GLB creates cache-friendly summaries, adaptively exploits summaries of both query and target time series, and captures boundary distances in an unsupervised manner. GLB outperforms all LBs in speedup (e.g., up to 13.5× faster against the strongest LB in terms of pruning power), establishes new state-ofthe-art results for the 5 elastic measures, and provides the first LBs for 2 elastic measures with no known LBs. Overall, GLB enables the effective development of LBs to facilitate fast similarity search.

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## **PVLDB Artifact Availability:**

The source code, data, and/or other artifacts have been made available at https://www.timeseries.org/glb.

# 1 INTRODUCTION

The ubiquity and unprecedented growth of time-varying measurements across scientific and industrial settings are responsible for

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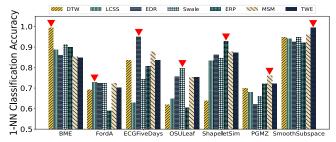


Figure 1: Classification accuracy of 7 elastic measures on sampled UCR datasets [26]. The triangle indicates the measure with the best accuracy: different winner in each dataset.

the increasing popularity of time-series analysis [35, 42, 43, 52, 59, 61, 64, 77, 79, 80]. The backbone of most time-series analytical tasks requires the detection of similar pairs of time series [27, 78]. Specifically, identifying similarities plays a central role in querying [24, 45, 50, 51, 55, 65, 68, 69, 87], indexing [17, 18, 21, 32, 48, 72, 91, 96, 103], motif discovery [7, 23, 57, 67, 99, 100], anomaly detection [9–15, 60, 71, 76, 84, 98, 99], classification [4, 29, 37, 40, 58, 73, 81], and clustering [1, 5, 6, 28, 44, 47, 56, 74, 75]. Consequently, *similarity search*, the process of retrieving the nearest neighbors to a query from a database under a certain distance measure, becomes one of the most fundamental building blocks in time-series analysis.

Unfortunately, the rising volumes of time series and their high dimensionality introduce severe challenges for similarity search [2, 25]. Specifically, the computational and storage costs of retrieving the nearest neighbors become prohibitively high even when the Euclidean distance (ED) is used [30, 31, 93, 94]. However, there is strong evidence that ED might not always be suitable for comparing time series [4, 27, 78, 95]. We often need to handle several distortions for effective time-series comparison, such as misalignments, stretching of observations, and fluctuations. To handle these distortions, dozens of distance measures have been proposed in the time-series literature [3, 8, 19, 20, 22, 27, 32, 33, 66, 73, 74, 83, 88, 92].

A recent study has evaluated over 70 time-series distance measures [78] and reaffirmed that *elastic* measures, which create a non-linear mapping between time series to align or stretch their points, achieve state-of-the-art performance in terms of accuracy. Importantly, the same study has shown that the nearest neighbor (classification) accuracy of ED may not always converge to the high accuracy of elastic measures with increasing dataset sizes, which contradicts the previous belief [86]. Therefore, supporting

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similarity search under different elastic measures is necessary for enabling effective large-scale time-series analytics.

Unfortunately, elastic distance measures scale quadratically (i.e.,  $O(L^2)$ ) to the length L of the time series. Compared to ED, which has linear complexity (i.e., O(L)), elastic distance measures incur an additional runtime overhead, often between one to three orders of magnitude [4, 78, 90]. This cost would prevent applications from using elastic measures in large-scale settings and favor sacrificing the high accuracy by relying on faster but less accurate measures. To alleviate this issue, the idea of *lower bounding* was developed to filter out unpromising candidates before carrying out the expensive elastic distance measure computation [32, 48, 49]. In simple terms, a lower bound (LB) is a fast distance measure that approximates an expensive elastic distance measure and is computed over some summaries of the time series instead of the actual time series.

Numerous LBs have been developed for elastic distance measures [19, 48, 49, 54, 85, 89, 90], with the goal to improve their pruning power (i.e., tightness of LB). However, a tighter LB does not always translate into higher speedup due to the potentially high computational cost necessary to calculate the LB. Despite over two decades of attention, there has never been, to the best of our knowledge, a comprehensive study to assess the progress in this area. In addition, the research effort on LBs has been disproportionally concentrated on Dynamic Time Warping (DTW) [82, 83], which is the oldest elastic measure with at least eight established LBs [48, 49, 54, 85, 89, 97, 101]. In contrast, other useful elastic distance measures, such as EDR [20] or SWALE [66], do not have any LBs reported in the literature, to the best of our knowledge. Importantly, as shown in Figure 1, no single elastic measure always wins on every dataset (as confirmed by a recent study [78]). Considering that all elastic measures are useful in practical time-series tasks, there is merit in developing LBs for elastic distance measures that have received relatively little attention.

Unfortunately, developing LBs is a challenging task. It took about two decades to improve the tightness of DTW LBs to over 80% [54, 85, 89, 90]. We believe it is unsustainable to expect a similar research effort for each elastic measure. Instead, a generalized framework that accumulates the knowledge from previously developed LBs would eliminate the need for designing separate LBs for each elastic measure. Based on our review and evaluation of existing LBs for DTW, we identified four critical properties that constitute the effectiveness of LBs: (\$P1\$) Query Dependence, which indicates the LB extracts summaries only from the query time series; (\$P2\$) Data Dependence, which indicates the LB extracts summaries from both the query and the target data time series; (\$P3\$) Boundary Dependence, which indicates the LB explicitly captures distances of the first and last elements of time series; and (\$P4\$) Reusability, which indicates the ability to cache and reuse results in future calculations.

This paper performs the first comprehensive study of LBs for elastic measures. Specifically, we evaluate 11 state-of-the-art LBs spanning 5 elastic measures using 128 datasets. In addition, we present the Generalized Lower Bounding (GLB) framework, the first framework to create LBs to satisfy all four previously mentioned desirable properties and adapt to all elastic measures (see Table 1 for comparison with baselines). LBs created within the GLB framework are query- and data-dependent and extract summaries (i.e., envelopes, see Section 2) from both the query and target data

Table 1: Analysis of LBs on four properties. GLB matches all properties, while competitors miss one or more.

Elastic	Lower	Query	Data	Boundary	n 11	Year &
Measure	Bounds	Dependent	Dependent	Dependent	Reusable	Reference
	LB_Yi	~	-	-	~	1998 [101]
	LB_Kim	· ·	-	· ·	·	2001 [49]
DTW	LB_Keogh	· ·	-	-	·	2002 [46, 48]
	LB_Improved	· ·	~	-	-	2009 [54]
	LB_New	· ·	-	· ·	-	2018 [85]
	LB_Enhanced	· ·	-	· ·	-	2019 [89]
	LB_Petitjean	· ·	~	· ·	-	2021 [97]
	LB_Webb	· ·	~	· ·	-	2021 [97]
LCSS	LB_LCSS	~	-	-	~	2002 [62, 91, 92]
	LB_Keogh-ERP	~	-	-	~	2004 [19]
ERP	LB_Kim-ERP	· ·	-	-	V	2004 [19]
	LB_ERP	· ·	-	· ·	\ \	2004 [19]
MSM	LB_MSM	~	-	-	~	2020 [90]
TWED	LB_TWED	~	-	-	~	2020 [90]
EDR	N/A	-	-	-	-	-
SWALE	N/A	-	-	-	-	-
		Proposed G	eneralized Lo	wer Bound		
All Elastic	GLB	V	./	-	-/	2023
Measures	GLD	_				(this work)

time series. In the GLB framework, LBs compute distances between all envelopes and the target time series, and adaptively select the envelopes that maximize the LB tightness for each pairwise comparison. To satisfy the boundary dependence, LBs under the GLB framework avoid the need to tune parameters and focus only on the leading and trailing time-series points, whereas existing LBs rely on supervised solutions. GLB's ability to reuse the extracted envelopes also avoids the significant overhead introduced by existing solutions focusing on complex and expensive transformations to capture characteristics from the target time series. Finally, GLB is adaptable to all elastic measures due to its abstraction of the different cost functions used internally in elastic distances.

We compare GLB against the 11 state-of-the-art LBs, and we make all source codes available to ensure reproducibility. Compared to the strongest LBs for DTW, GLB\_DTW achieves state-of-the-art pruning power and outperforms *all* LBs in terms of speedup. Specifically, GLB\_DTW is up to 13.5× faster (6.8× faster on average) when compared against the strongest LB in terms of pruning power and wins in at least 115 out of 128 (90%) datasets. For elastic measures other than DTW, GLB establishes new state-of-the-art results in both pruning power and speedup, outperforming all baselines significantly. Importantly, for elastic measures without known LBs, GLB achieves performance comparable to the results for the other elastic measures, which demonstrates the generalizability of GLB.

We present our contributions as follows:

- We provide the first thorough study of elastic measures and LBs, summarizing two decades of progress (Section 2).
- We propose the GLB framework for developing LBs for elastic measures while satisfying four properties: query, data, and boundary dependence, and reusability (Section 3.1).
- We formally define GLB by abstracting costs and combining elements that constitute the effective LBs (Sections 3.2-3.3).
- We prove the correctness of GLB (Section 3.4).
- We present new LBs for 7 elastic distance measures, including two elastic measures without known LBs (Section 4).
- We conduct the most extensive evaluation of LBs until now to demonstrate the robustness of GLB (Sections 5 and 6).

Finally, we summarize the implications of our work (Section 7).

 $<sup>^{1}</sup>www.timeseries.org/glb \\$ 

#### 2 PRELIMINARIES AND RELATED WORK

In this section, we first review the development of elastic distance measures starting with DTW, the earliest and most popular elastic measure. We provide a generalized formula to showcase the recursive (dynamic programming) computation in all elastic measures, which highlights the different cost functions adopted by each elastic measure (Section 2.1). Then, we present LBs to accelerate similarity search for elastic measures and our problem of focus (Section 2.2).

# 2.1 Elastic Time Series Distance Measures

We now review the development of elastic distance measures.

**Dynamic Time Warping (DTW) [82, 83]**: We consider a timeseries  $\mathbf{x} = [x_1, x_2, ..., x_L]$ , an ordered sequence of L datapoints. To measure the similarity between two time series, the most com-

measure the similarity between two time series, the most common distance measure is ED, which offers a one-to-one (linear) alignment between two time series, as shown in Figure 2(a). ED defines the distance between time series  $\mathbf{x} = [x_1, x_2, ..., x_{L_{\mathbf{x}}}]$  and  $\mathbf{y} = [x_1, x_2, ..., x_{L_{\mathbf{x}}}]$ 

Innes the distance between time series  $\mathbf{x} = [x_1, x_2, ..., x_{L_\mathbf{x}}]$  and  $\mathbf{y} = [y_1, y_2, ..., y_{L_\mathbf{y}}]$ , with  $L_\mathbf{x} = L_\mathbf{y} = L$ , as  $ED(\mathbf{x}, \mathbf{y}) = \sqrt{\sum_{i=1}^L (x_i - y_i)^2}$ . The disadvantage of ED is its inability to recognize the similarity between time series that have comparable shapes but are stretched or differ in phase or length. The need to capture shape similarity despite these distortions inspired the development of *elastic measures*, which create a non-linear mapping between time series by comparing one-to-many points in order to align or stretch points. For example, DTW enables local alignments by permitting one-to-many points matching, as shown in Figure 2(b), where the peaks and troughs of two time series are aligned correspondingly. To find the local alignments, DTW first computes a distance matrix, D, using the following recursive computation:

$$D(i,j) = \begin{cases} (x_i - y_j)^2 & \text{if } i, j = 1\\ D(i-1,j) + (x_i - y_j)^2 & \text{if } i \neq 1 \text{ and } j = 1\\ D(i,j-1) + (x_i - y_j)^2 & \text{if } i = 1 \text{ and } j \neq 1\\ D(i-1,j-1) + (x_i - y_j)^2 & \text{if } i = 1 \text{ and } j \neq 1\\ D(i-1,j) + (x_i - y_j)^2 & \text{if } i, j \neq 1\\ D(i,j-1) + (x_i - y_j)^2 & \text{if } i, j \neq 1 \end{cases}$$

DTW determines the alignment path,  $W = \{w_1, w_2, ..., w_p\}$ , which starts from the bottom-left corner and ends at the top-right corner in the matrix (gray path in Figure 2) where the distance of alignments add up to the cell in the top-right corner, or  $D(L_{\mathbf{x}}, L_{\mathbf{y}}) = \sum_{i=1}^{p} (\mathbf{x}_{W_k^{-1}} - \mathbf{y}_{W_k^{-2}})^2$ . The warping path follows two properties [89]:

- Boundary Constraints:  $w_1 = D(1, 1)$  and  $w_p = D(L_x, L_y)$ , meaning the optimal warping path has to start on the bottom-left corner of D and to end on the upper-right corner of D.
- Continuity and Monotonicity: if  $w_k = D(i, j)$  for  $k \in [2, p-1]$ ,  $w_{k+1} \in \{D(i+1, j), D(i, j+1), D(i+1, j+1)\}$ , meaning that the warping path, starting from bottom-left, only moves vertically upwards, horizontally towards the right, or diagonally towards top-right continuously until arriving at the top-right corner, as shown on the right side of Figure 2(b). DTW uses the same distance function in each matrix cell regardless of whether the optimal path arrives at that cell horizontally, vertically, or diagonally (Formula 1).

In contrast to DTW, alternative elastic measures that we review next have three different distance functions, each corresponding to *diagonal* movements, *horizontal* movements, and *vertical* movements, respectively (Figure 2(b) depicts these three costs). As we discuss in Section 3, having different distance functions for diagonal

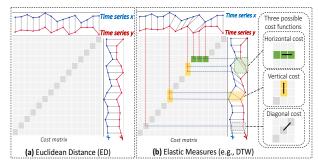


Figure 2: Cost matrix for (a) ED and (b) elastic measures.

and horizontal/vertical movements presents a unique challenge for proposing effective lower bounds for particular elastic measures. GLB provides a principled solution to tackle this critical problem. After computing distance matrix D and optimal warping path W, the distance between the time series  $\mathbf{x}$  and  $\mathbf{y}$  of lengths  $L_{\mathbf{x}}$  and  $L_{\mathbf{y}}$ 

is defined as 
$$DTW(\mathbf{x}, \mathbf{y}) = \sqrt{D(L_{\mathbf{x}}, L_{\mathbf{y}})} = \sqrt{\sum_{i=1}^{p} w_i}$$
.

DTW is commonly used together with a locality constraint to limit the range of warping allowed and thereby avoiding unreasonably far-reaching alignments; this approach is referred to as Constrained DTW (cDTW) [36]. cDTW is faster than DTW and often results in higher classification accuracy [78]. The most widely adopted locality constraint is the Sakoe-Chiba band [36], which we refer to as the warping window. The warping window, w, is the maximum possible deviation of the alignment path from the diagonal of D, and cells further away are not computed.

**Generalization of Elastic Measures**: Since the development of DTW, numerous alternative elastic measures have been proposed to address certain limitations. These measures have employed different distance functions for diagonal and vertical/horizontal movements and introduced additional parameters to tackle weaknesses of DTW.

Despite the differences in parameters and distance functions, subsequent elastic measures and DTW share a common goal of finding the warping path through dynamic programming to compute D, so elastic measures can be generalized as follows:

$$D(i,j) = \begin{cases} initial\_distance(x_i, y_j) & \text{if } i, j = 1 \\ D(i-1, j) + dist^V(x_i, y_j) & \text{if } i \neq 1 \text{ and } j = 1 \\ D(i, j-1) + dist^H(x_i, y_j) & \text{if } i = 1 \text{ and } j \neq 1 \\ 0(i-1, j-1) + dist^D(x_i, y_j) & \text{if } i, j \neq 1 \\ D(i-1, j) + dist^V(x_i, y_j) & \text{if } i, j \neq 1 \\ D(i, j-1) + dist^H(x_i, y_j) & \text{otherwise} \end{cases}$$

$$(2)$$

where  $dist^D(x_i, y_j)$ ,  $dist^V(x_i, y_j)$ , and  $dist^H(x_i, y_j)$  are the distance functions for diagonal, vertical, and horizontal movements, respectively, and  $initial\_distance(x_i, y_j)$  is the distance function for initial alignment in D(1, 1). After computing  $D(L_{\mathbf{X}}, L_{\mathbf{Y}})$ , each elastic measure adopts a function,  $transD(L_{\mathbf{X}}, L_{\mathbf{Y}})$ , that transforms distance in  $D(L_{\mathbf{X}}, L_{\mathbf{Y}})$  to calculate distance between  $\mathbf{x}$  and  $\mathbf{y}$  based on  $D(L_{\mathbf{X}}, L_{\mathbf{Y}})$  (e.g., for DTW,  $trans(D(L_{\mathbf{X}}, L_{\mathbf{Y}})) = \sqrt{D(L_{\mathbf{X}}, L_{\mathbf{Y}})}$ ; for MSM,  $trans(D(L_{\mathbf{X}}, L_{\mathbf{Y}})) = D(L_{\mathbf{X}}, L_{\mathbf{Y}})$ ).

**Threshold-based Elastic Measures**: To mitigate the disproportionate impact the outlier data points have in DTW distance, a group of elastic measures use a threshold parameter  $\epsilon$  to decide whether two elements should match or not. Adopting a threshold restricts the relationship between two elements to be either a match or a mismatch, regardless of how close or distant they are.

Table 2: Summary of distances ( $dist^D(x_i, y_j)$ ,  $dist^H(x_i, y_j)$ , and  $dist^V(x_i, y_j)$ ), transformation functions ( $trans(D(L_x, L_y))$ ), and LBs for threshold-based elastic distances.

	$dist^D(x_i, y_j)$	$\begin{cases} 1 \text{ if }  x_i - y_j  \le \epsilon \\ 0 \text{ otherwise} \end{cases}$
LCSS	$dist^V(x_i, y_j)$	0
	$dist^H(x_i, y_j)$	0
	$trans(D(L_x, L_y))$	$1 - \frac{D(L_{\mathbf{x}}, L_{\mathbf{y}})}{min(L_{\mathbf{x}}, L_{\mathbf{y}})}$
	LB	$UE(\mathbf{y})_i = max(y_{i-w} : y_{i+w})$ $LE(\mathbf{y})_i = min(y_{i-w} : y_{i+w})$ $LB\_LCSS(\mathbf{x}, \mathbf{y}) = 1 - \frac{1}{L_x} \sum_{i=1}^{L_x} \begin{cases} 1 \text{ if } LE(\mathbf{y})_i \le x_i \le UE(\mathbf{y})_i \\ 0 \text{ otherwise} \end{cases}$
	$dist^D(x_i, y_j)$	$\begin{cases} 0 \text{ if }  x_i - y_j  \le \epsilon \\ 1 \text{ otherwise} \end{cases}$
EDR	$dist^V(x_i, y_j)$	1
	$dist^H(x_i, y_j)$	1
	$trans(D(L_x, L_y))$	$D(L_{\mathbf{x}}, L_{\mathbf{y}})$
	LB	No Existing Lower Bound
	$dist^D(x_i, y_j)$	$\begin{cases} r \text{ if }  x_i - y_j  \le \epsilon \\ p \text{ otherwise} \end{cases}$
SWALE	$dist^V(x_i, y_j)$	p
	$dist^H(x_i, y_j)$	p
	$trans(D(L_x, L_y))$	$D(L_{\mathbf{x}}, L_{\mathbf{y}})$
	LB	No Existing Lower Bound

Such quantization makes threshold-based approaches more robust against outliers and noisy time series than other elastic measures. A match between a new pair of elements indicates a diagonal movement in the accumulated distance matrix D, whereas a mismatch is associated with either a horizontal or vertical movement. Table 2 summarizes distance functions, transformation functions, and LBs for three popular threshold-based elastic distance measures:

- Longest Common Subsequence (LCSS) [92] was originally used for pattern matching in text strings and was adapted to measure similarity between time series. LCSS increases similarity by 1 when two elements match and 0 in cases of mismatch. The LCSS distance is the LCSS similarity normalized by the length of the shorter time series.
- Edit Distance on Real Sequences (EDR) [20] is an adaption of the edit distance for strings. EDR computes the distance by penalizing mismatches instead of rewarding matches.
- **Sequence Weighted Alignment (SWALE)** [27] parameterizes EDR using a parameter *r* for a match and a punishment parameter *p* for mismatch, instead of fixed 1 and 0.

Metric Elastic Measures: The aforementioned elastic measures, including DTW and threshold-based measures, are all non-metric, and thus do not satisfy the triangle inequality [88]. Being a metric enables elastic measures to use generic indexing methods [38, 39, 102] and clustering methods [16, 34, 41] designed for metrics. In addition, in nearest neighbor search, triangle inequality can be applied to efficiently prune comparisons [19]. As summarized in Table 3, there are three popular metric elastic measures:

• Edit Distance with Real Penalty (ERP) [19] is similar to DTW as ERP uses the squared difference  $((x_i - y_j)^2)$  between two elements as the distance function for diagonal movement. Different from DTW, ERP introduces an additional gap

Table 3: Summary of distances  $(dist^D(x_i, y_j), dist^H(x_i, y_j))$ , and  $dist^V(x_i, y_j)$ , transformation functions  $(trans(D(L_x, L_y)))$ , and LBs for metric elastic distances.

	$dist^D(x_i,y_j)$	$(\mathbf{x}_i - \mathbf{y}_j)^2$					
ERP	$dist^V(x_i, y_j)$	$(\mathbf{x}_i - g)^2$					
	$dist^H(x_i, y_j)$	$(y_j - g)^2$					
	$trans(D(L_{\mathbf{x}}, L_{\mathbf{y}}))$	$D(L_{\mathbf{x}}, L_{\mathbf{y}})$					
	LB	$LB\_Kim - ERP(\mathbf{x}, \mathbf{y}) = \max \begin{cases}  x_1 - y_1  \\  x_2 - y_{1-y}  \\  max(\mathbf{x}) - max(\mathbf{y})  \\  min(\mathbf{x}) - min(\mathbf{y})  \end{cases}$					
		$LB\_ERP(\mathbf{x}, \mathbf{y}) =  \sum \mathbf{y} - \sum \mathbf{x} $					
		$UE_i = max(g, max(c_{i-w} : c_{i+w}))$					
		$LE_i = min(g, min(c_{i-w} : c_{i+w}))$ $(u_i - UE_i)^2 \text{ if } u_i > UE_i$					
		$LB\_Keogh - ERP(\mathbf{x}, \mathbf{y}) = \begin{cases} (y_i - UE_i)^{-1} & \text{if } y_i > UE_i \\ \sum_{i=1}^{n} \left\{ (y_i - LE_i)^2 & \text{if } y_i < LE_i \end{cases}$					
		$ \sqrt{\frac{2i}{0}} \text{ otherwise} $					
	$dist^D(x_i, y_j)$	$ \mathbf{x}_i - y_i $					
MSM	$dist^V(x_i, y_j)$	$\begin{cases} c \text{ if } x_{i-1} \leq x_i \leq y_j \text{ or } x_{i-1} \geq x_i \geq y_j \\ c + min \begin{cases}  x_i - x_{i-1}  & \text{otherwise} \\  x_i - y_j  & \end{cases} \end{cases}$					
	$dist^H(x_i, y_j)$	$\begin{cases} c \text{ if } y_{j-1} \leq y_j \leq x_i \text{ or } y_{j-1} \geq y_j \geq x_i \\ c + min \begin{cases}  y_j - y_{j-1}  & \text{otherwise} \\  y_i - x_i  & \text{otherwise} \end{cases}$					
	$trans(D(L_{\boldsymbol{x}},L_{\boldsymbol{y}}))$	$D(L_{\mathbf{x}}, L_{\mathbf{y}}))$					
	LB	$LB\_MSM(\mathbf{x}, \mathbf{y}) =  \mathbf{x}_1 - \mathbf{y}_1  + \sum_{i=2}^{L_Y} \begin{cases} min( y_i - max(\mathbf{x}) , c) \text{ if } y_{i-1} \ge y_i > max(\mathbf{x}) \\ min( y_i - min(\mathbf{x}) , c) \text{ if } y_{i-1} \le y_i < min(\mathbf{x}) \\ 0 \text{ otherwise} \end{cases}$					
	$dist^D(x_i, y_j)$	$(x_i - y_j)^2 + (x_{i-1} - y_{j-1})^2 + \nu( t_{x_i} - t_{x_{i-1}}  +  t_{y_j} - t_{y_{j-1}} )$					
TWED	$dist^V(x_i, y_j)$	$(x_i - x_{i-1})^2 + \nu( t_{x_i} - t_{x_{i-1}} ) + \lambda$					
	$dist^H(x_i, y_j)$	$(y_j - y_{j-1})^2 + v( t_{y_j} - t_{y_{j-1}} ) + \lambda$					
	$trans(D(L_x, L_y))$	$D(L_{\mathbf{x}}, L_{\mathbf{y}})$					
	LB	$LB\_TWED(\mathbf{x}, \mathbf{y}) = \min \left\{ \begin{aligned} &(\mathbf{x}_1 - y_1)^2 \\ &\mathbf{x}_1^2 + v + \lambda \\ &y_1^2 + v + \lambda \end{aligned} \right. & + \sum_{i=2}^{Ly} \left\{ \begin{aligned} &\min(v_i (y_i - max(max(\mathbf{x}), y_{i-1}))^2 \\ &\text{if } y_i > max(max(\mathbf{x}), y_{i-1}) \\ &\min(v_i (y_i - min(min(\mathbf{x}), y_{i-1}))^2 \\ &\text{if } y_i < min(min(\mathbf{x}), y_{i-1}) \end{aligned} \right\}$					

value parameter, g, to compute the distance for horizontal and vertical movements. A drawback of ERP is the inability to handle vertically shifted time series.

- Move-Split-Merge (MSM) [88] combines advantages of several elastic measures. MSM is a translation-invariant metric. A constant distance c is associated with split (replicating the previous element) and merge (merging two identical elements into a single one) operations, which correspond to the horizontal and vertical movements in the distance matrix. The move operation is a diagonal movement in the distance matrix, where the distance is  $|x_i y_j|$ .
- Time Warp Edit Distance (TWED) [63] penalizes the difference in timestamps in addition to the difference in numerical values. TWED penalizes timestamp difference with parameter *v* in all three types of movements. TWED also employs an additional stiffness parameter *λ* in horizontal and vertical movements to control warping.

## 2.2 Lower Bounds for Elastic Measures

As described in the previous section, all elastic measures compute a distance matrix containing  $L^2$  cells where L is the length of two-time series, resulting in  $O(L^2)$  complexity (we ignore the constrained variants as in the worst case they also have similarly high complexity). This runtime complexity is rather time-consuming for similarity search, considering that even for moderate-size databases, millions of such comparisons have to be performed. As a result,

## Algorithm 1: LB-Accelerated Nearest Neighbor Search

```
Require: y, a query series
Require: X. a set of data time series
Require: L, labels of data time series L
Ensure: label_y, label of query time series
 1: for x in X do
       lb\_list[i] = LowerBound(\mathbf{x}, \mathbf{y})
3: end for
 4: best\_so\_far = +\infty
5: ordering = sort(lb_list)
6: X = X[ordering]
 7: for x in X do
       if lb_list[x] < best_so_far then</pre>
             ctual_dist = ElasticMeasure(x, y)
10:
            if actual_dist < best_so_far then
               best\_so\_far = actual\_dist
11:
12:
               label_{y} = L[x]
            end if
14:
       end if
15: end for
```

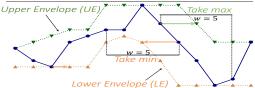


Figure 3: Construction of upper and lower envelopes: the solid line represents the original time series, and two dashed lines represent the upper and lower envelopes, respectively.

similarity search under elastic measures becomes orders of magnitude slower than when ED is used [78]. The most popular approach to accelerate elastic measures is through lower bounding. LB approximates the elastic measure distance without computing the full distance matrix. Specifically, LB uses the approximated value to filter out unpromising candidates in tasks such as nearest neighbor search. Algorithm 1 illustrates how LB is applied to accelerate the 1-NN classification of time series y. First, the list of time series is reordered based on their LB distance to the query time series y (lines 1-6). Then, going through the list, the actual elastic measure distance is only computed for promising time series whose LB distance with query time series y is less than the actual distance between y and the current nearest neighbor (lines 7-9). Finally, if the actual distance is less than the  $best\_so\_far$  distance, the  $best\_so\_far$  distance and the nearest neighbor are updated (lines 10-12).

Research efforts in lower bounds for elastic measures have concentrated on developing tighter lower bounds of DTW due to its popularity. LB\_Kim [49] is a cheap DTW LB with O(1) complexity, which takes the maximum among  $\{|X_1-Y_1|,|X_L-Y_L|,|X_{max}-Y_{max}|,|X_{min}-Y_{min}|\}$ . Although LB\_Kim is fast to compute, its looseness makes it ineffective in filtering out less obvious candidates. LB\_Yi exploits the fact all points in  ${\bf x}$  that are either larger than  $max({\bf y})$  or smaller than  $min({\bf y})$  necessarily contribute to the final DTW distance. Subsequent to LB\_Yi, LB\_Keogh [48] obtains a much higher tightness than LB\_Yi and LB\_Kim by utilizing envelopes and a warping window  ${\bf w}$  (Table 4, Equation 4). As shown in Figure 3, LB\_Keogh first constructs the upper and lower envelopes of the query time series and computes the distance between these query envelopes and the target (candidate) time series.

LB\_Improved [54] computes the ordinary LB\_Keogh as well as the LB\_Keogh between the query time series and the projection of target time series on the query envelopes (Table 4, Equations 6 and 7). This makes LB\_Improved tighter than LB\_Keogh but also adds

Table 4: Key existing LBs for DTW.

	O(1) Complexity	$LB\_Kim(\mathbf{x}, \mathbf{y}) = \max \begin{cases}  x_1 - y_1  \\  x_{L_x} - y_{L_x}  \\  max(\mathbf{x}) - max(\mathbf{y})  \\  min(\mathbf{x}) - min(\mathbf{y})  \end{cases} $ (3)
		$UE(\mathbf{y})_i = max(y_{i-\mathbf{w}} : y_{i+\mathbf{w}})$
		$LE(\mathbf{y})_i = min(y_{i-w} : y_{i+w})$
	Based on	$(x_i - UE(\mathbf{y})_i)^2$ if $x_i > UE(\mathbf{y})_i$
DTW	envelopes	$LB\_Keogh(\mathbf{x}, \mathbf{y}) = \sqrt{\sum_{i=1}^{L_{\mathbf{x}}} \begin{cases} (x_i - UE(\mathbf{y})_i)^2 \text{ if } x_i > UE(\mathbf{y})_i \\ (x_i - LE(\mathbf{y})_i)^2 \text{ if } x_i < LE(\mathbf{y})_i \end{cases}} $ $0 \text{ otherwise}$
		0 otherwise
		$LB_{-}Yi(\mathbf{x}, \mathbf{y}) = \sum_{i=1}^{L} \begin{cases} (x_i - max(\mathbf{y}))^2 & \text{if } x_i > max(\mathbf{y}) \\ (x_i - min(\mathbf{y}))^2 & \text{if } x_i < min(\mathbf{y}) \end{cases} $ (5)
		$UE(\mathbf{y})_i \text{ if } x_i \ge UE(\mathbf{y})_i$
		$H(\mathbf{x}, \mathbf{y})_i = \left\{ LE(\mathbf{y})_i \text{ if } x_i \leq LE(\mathbf{y})_i \right\}$ (6)
		$H(\mathbf{x}, \mathbf{y})_i = \begin{cases} E(\mathbf{y})_i & \text{if } x_i \le E(\mathbf{y})_i \\ E(\mathbf{y})_i & \text{if } x_i \le E(\mathbf{y})_i \end{cases} $ (6)
		$LB\_Improved(\mathbf{x}, \mathbf{y}) = LB\_Keogh(\mathbf{x}, \mathbf{y}) + LB\_Keogh(\mathbf{y}, H(\mathbf{x}, \mathbf{y})) $ (7)
	Oth	$\delta(x, \mathbf{y}) = \min(x - y)^2 \text{ for } y \in \mathbf{y}(8)$
	Others	$LB_{-}New(\mathbf{x}, \mathbf{y}) = \sqrt{(x_1 - y_1)^2 + (x_{L_{\mathbf{x}}} - y_{\mathbf{y}})^2 + \sum_{i=2}^{L_x - 1} \delta(x_i, \mathbf{y})(9)}$

overhead. In addition, the projection of candidate queries on query envelopes cannot be computed in a pre-processing step. LB New [85] (Table 4, Equations 8 and 9) also obtains higher pruning power than LB\_Keogh. However, instead of constructing envelopes and considering only elements outside of envelopes, LB New pairs each candidate element,  $x_i$ , with the closest query element,  $y_i$ , where  $y_j \in Y_i = (y_{max(1,j-w)} : y_{min(j+w,L_y)})$ . Differently from the previous two approaches, LB\_Enhanced [89] constructs alternating bands around the upper-left corner (i.e., Left Bands) and bottom-right corner (i.e., Right Bands), where each band captures one cell in the warping path. The number of bands is decided by a parameter V, which requires a tuning process. LB Keogh is applied for the remaining portion of the time series not covered by the two bands. LB Petitjean [97] improves the projection strategy in LB Improved and incorporates the bands from LB Enhanced, which translates into increased pruning power. LB\_Webb [97] approximates LB Petitjean without directly computing envelopes or projections, resulting in improved performance over LB\_Petitjean. Due to space limitations, we omit the definitions of LB\_Enhanced, LB Petitjean, and LB Webb in Table 4, which require introducing new notation and concepts not shared with the key existing LBs.

In addition to DTW lower bounds, a few lower bounds have been developed for other elastic measures as well. For example, in efforts to develop ERP LBs, LB\_Keogh and LB\_Kim were adapted to form the LB\_Keogh-ERP and LB\_Kim-ERP, respectively. In particular, LB\_Keogh-ERP adjusted the original LB\_Keogh envelopes by adding another parameter g (Table 3). LB\_Kim-ERP also adds a parameter g to LB\_Kim to incorporate possible gaps in the alignment. In addition to adaptation from DTW, authors of ERP also developed LB\_ERP specifically for ERP. However, LB\_ERP is only applicable for g=0, which limits the usability. In addition to adaptation to ERP, LB\_Keogh was also modified to LB\_LCSS by replacing DTW's Euclidean Distance with reward parameter of 1 and penalty parameter of 0 in cases of match and mismatch. Recently, [90] defined LBs for TWED and MSM based on their distance functions to capture their initial boundary distance and query characteristics (Table 3).

After reviewing the literature, we observe that (i) multiple LBs focus on improving the tightness of LB\_Keogh, which does not always translate into higher speed up; and (ii) different distances for different movements make the development of LBs challenging.

#### 3 THE GLB FRAMEWORK

Until now, we have thoroughly reviewed the state-of-the-art elastic measures and their corresponding LB distances. It becomes evident there is merit in developing a framework to ease the process of deriving LBs for all elastic measures. By studying the development of LBs in the past two decades, we identified four critical properties (i.e.,  $\mathcal{P}1-4$ , Section 1) necessary for the effectiveness of LBs. We aim to develop a generalized framework to satisfy all four properties. Next, we introduce GLB, our innovative framework to facilitate the creation of efficient LBs for elastic measures that can be mapped into the generalization of elastic measures described in Section 2.1 (Equation 2). The GLB framework attains its generalizability by introducing a novel abstraction of the various distance functions defined between elements of time series across elastic measures.

#### 3.1 Main Ideas

GLB aims to satisfy four critical properties (i.e.,  $\mathcal{P}1-4$ ): query dependence, data dependence, boundary dependence, and reusability. **Query and Data Dependence:** To consider characteristics of both the query time series and the target time series while being cache-friendly, GLB constructs envelopes for each query time series  $\mathbf{y}$  and data time series  $\mathbf{x}$ . Since  $\mathbf{y}$  and  $\mathbf{x}$  have different numerical values and, therefore, different envelopes, the distance between the data time series and the query envelope is different from the distance between the query time series and the data envelope. GLB adaptively selects between the two envelopes to maximize the LB tightness for each pairwise comparison.

Boundary Dependence: One deficiency we observed in the construction of LB\_Enhanced is the need for a time-consuming parameter tuning process to decide the appropriate value for its tightness parameter V, which significantly reduces the speedup in the overall process. To build on top of LB\_Enhanced and take advantage of the boundary distances, GLB includes only the distances for aligning leading elements,  $initial\_distance(x_1, y_1)$ , and distances for aligning ending elements, ending\_distance( $x_{L_x}$ ,  $y_{L_y}$ ). In particular,  $initial\_distance(x_1, y_1)$  depends on how each elastic measure computes the distance for the first cell D(1,1) of the accumulated distance matrix. Such distance function varies across each elastic measure; for instance, in DTW and ERP,  $initial\_distance(x_1, y_1) =$  $(x_1 - y_1)^2$ , while initial\_distance $(x_1, y_1) = |x_1 - y_1|$  in MSM. By abstracting the internal costs of elastic measures, GLB adapts the LB computation to different elastic measures, as we will see later. Reusability: As introduced in [54], calculating a data-dependent envelope in addition to the query envelope in LB\_Keogh allows a tighter lower bound. [54] demonstrated this idea by developing LB\_Improved, which computes the query envelope as well as the envelope of a projection time series (see Table 4, Eq. 6). However, we observed that computing the projection time series relies on both the query and data time series, making envelopes of projection time series not reusable. In fact, according to our evaluation across 128 datasets (see Table 7), despite LB\_Improved's impressive pruning power of 84.71% (avoided 84.71% of calls to the DTW distance) over LB\_Keogh's 70.19%, LB\_Improved is only able to reduce runtime by 63%, whereas LB\_Keogh could reduce runtime by 76%. The inconsistency between LB\_Improved's strong pruning power and less impressive speedup suggests that reusability is a key factor. As a result, reusability is a core design principle of GLB.

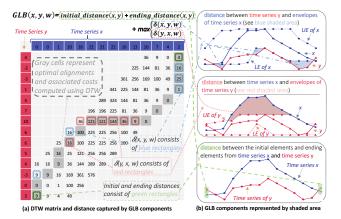


Figure 4: GLB\_DTW for two time series with warping window=3. In the distance matrix: the red column represents query time series y; the blue row represents data time series x; gray cells represent the DTW alignment path; blue rectangles are distances between query time series y and envelope of data time series x; red rectangles are distances between data time series x and envelope of query time series y; green rectangles represent the boundary distances.

# 3.2 Adaptable Distance Functions

In addition to the four aforementioned properties, GLB aims to provide a generalized framework to ease the process of developing new LBs. From our extensive evaluation (see Sections 5 and 6), we observe that it took about two decades to improve the tightness of DTW's LB to over 80% today. Unfortunately, the tightness of the state-of-the-art LBs for other elastic measures, such as for MSM and TWE, are still below 20%. We believe it is unsustainable to expect similar research efforts dedicated to each elastic measure individually. Contrary to the isolated development of different LBs for different elastic measures, we observe that all elastic measures share the same dynamic programming construct, and their only differences lie in the distance functions used to capture movements along the accumulated distance matrix. Thus, for GLB, we aim to leverage this shared structure (see Section 2.1, Equation 2, on how elastic measures can be generalized into the same framework). An important advantage of such a generalized LB is the ability to transfer the research effort devoted to the development of LBs for one elastic measure to the others.

To develop a generalized framework, we need to account for the difference between DTW and other elastic measures. Specifically, when moving horizontally and vertically through the accumulated distance matrix, DTW uses the same distance function as moving diagonally; by contrast, other elastic measures use different distance functions for horizontal, vertical, and diagonal movements. For instance, in ERP, the distance for diagonal movement is  $dist^D(x_i - y_j) = (x_i - y_j)^2$ , but the distance for horizontal movement is  $dist^H(x_i - y_j) = (x_i - g)^2$ . Since we are unsure whether each element is involved in a diagonal or horizontal or vertical movement in the optimal alignment, GLB takes the minimum of the three alternatives. Similarly, when calculating the distance of the last cell,  $D(L_x, L_y)$ , in the accumulated distance matrix, we need to take the minimum among three alternative movements to  $D(L_x, L_y)$ .

# 3.3 Mathematical Formulation

Having introduced the main ideas behind GLB, we now provide a formal definition. GLB includes the pairwise distances between the first and last pairs of elements of two time series and then the maximum of query and data characteristics captured by query and data envelopes. A warping window can also be applied to GLB in cases when the warping windows are used for computing distances in elastic measures. GLB is formally defined as:

$$ending\_distance(\mathbf{x}, \mathbf{y}) = min \begin{cases} dist^D(x_{L_{\mathbf{x}}}, y_{L_{\mathbf{y}}}) \\ dist^V(x_{L_{\mathbf{x}-1}}, y_{L_{\mathbf{y}}}) \\ dist^H(x_{L_{\mathbf{x}}}, y_{L_{\mathbf{y}-1}}) \end{cases}$$
(10)

where  $dist^D(x_i, y_i)$ ,  $dist^H(x_i, y_i)$ ,  $dist^V(x_i, y_i)$  represent distance functions for diagonal, horizontal, and vertical movements. For instance, in the case of ERP with parameter q = 0, there are three possible distances from the alignment between the ending element of time series  ${\bf x}$  and  ${\bf y}$ :  $\{(x_{L_{\bf x}}-y_{L_{\bf y}})^2,(x_{L_{\bf x}}-g)^2,(y_{L_{\bf y}}-g)^2\}$ . Then the ending\_distance in GLB\_ERP is  $min(\{(x_{L_{\bf x}}-y_{L_{\bf y}})^2,(x_{L_{\bf x}}-g)^2\})$ .  $(y_{L_{\mathbf{v}}} - g)^2$ ).  $\delta(\mathbf{x}, \mathbf{y}, w)$  and  $\delta(\mathbf{y}, \mathbf{x}, w)$  illustrate how GLB computes query and data envelopes and capture the distance between the envelopes and the other time series. In the calculation of such distance, the abstracted distance functions  $\delta(\mathbf{x}, \mathbf{y}, w)$  and  $\delta(\mathbf{y}, \mathbf{x}, w)$ were utilized to allow for the use of any elastic measure's custom distance functions. This is a departure from previous LBs in the literature, which required a different LB for each elastic measure. Instead, by specifying the distance functions ( $dist^D(x_i, y_i)$ ,  $dist^{V}(x_{i}, x_{i-1}), dist^{H}(y_{i}, y_{i-1}))$ , each elastic measure can calculate its own LB within the GLB framework. This approach offers greater flexibility and adaptability in computing the distance between time series.  $\delta(\mathbf{x}, \mathbf{y}, w)$  and  $\delta(\mathbf{y}, \mathbf{x}, w)$  are defined as:

$$\delta(\mathbf{x}, \mathbf{y}, w) = \sum_{j=2}^{L_{\mathbf{y}}-1} \begin{cases} \min(dist^{D}(y_{j}, UE(\mathbf{x})_{j}), dist^{H}(y_{j}, y_{j-1})) \text{ for } y_{j} \geq UE(\mathbf{x})_{j} \\ \min(dist^{D}(y_{j}, LE(\mathbf{y})_{j}), dist^{H}(y_{j}, y_{j-1})) \text{ for } y_{j} \leq LE(\mathbf{x})_{j} \\ 0 \text{ otherwise} \end{cases}$$

where  $UE(\mathbf{x})_j = max(x_{max(1,j-w)}: x_{min(L_{\mathbf{x}},j+w)})$  and  $LE(\mathbf{x})_j = min(x_{max(1,j-w)}: x_{min(L_{\mathbf{x}},j+w)})$ , and w is the window.

$$\delta(\mathbf{y}, \mathbf{x}, w) = \sum_{i=2}^{L_{\mathbf{x}}-1} \begin{cases} min(dist^D(x_i, UE(\mathbf{y})_i), dist^V(x_i, x_{i-1})) \text{ for } x_i \geq UE(\mathbf{y})_i \\ min(dist^D(x_i, LE(\mathbf{y})_i), dist^V(x_i, x_{i-1})) \text{ for } x_i \leq LE(\mathbf{y})_i \\ 0 \text{ otherwise} \end{cases}$$

where  $UE(\mathbf{y})_i = max(y_{max(1,i-w)}: y_{min(L_{\mathbf{y}},i+w)})$ 

and  $LE(\mathbf{y})_i = min(y_{max(1,i-w)}: y_{min(L_{\mathbf{y}},i+w)})$ , and window is w.

An implementation of GLB applied to accelerate the nearest neighbor search is provided in Algorithm 2. Firstly, envelopes of query and data time series are computed and stored in a cachefriendly approach (lines 1-11, 21-22) and serve as inputs to compute GLB. The GLB function computes and returns the sum of boundary distances and delta functions (lines 13-19). Finally, the GLB distances are used to accelerate the 1-NN Search (lines 21-41).

# 3.4 Proof of Lower Bounding Property

Having formally introduced GLB, in this section, we provide proof of the LB property of GLB. Due to this property, GLB ensures its correctness and no false positives exist (i.e., ignored comparisons with time series are correctly avoided), which ensures the utility of elastic measures in practice. For example, when GLB is used in classification tasks, the classification accuracy remains unaffected when LBs are applied. The LB property is formally defined as follows:

For any two time series  $\mathbf{x}$  and  $\mathbf{y}$  of length  $L_{\mathbf{x}}$  and  $L_{\mathbf{y}}$  respectively, an optimal alignment path  $W = w_1, w_2, ..., w_p$  where  $w_k = (i, j)$  indicates  $x_i$  is aligned to  $y_j$ , following inequality holds:  $GLB\_EE(\mathbf{x}, \mathbf{y}) \leq EE(\mathbf{x}, \mathbf{y})$  where EE represents an aforementioned elastic measure.

*Proof* Firstly, we observe the actual distance between two time series,  $EE(\mathbf{x}, \mathbf{y})$ , can be generally defined as:

$$\begin{split} \text{EE}(\mathbf{x},\,\mathbf{y}) &= \text{initial\_distance}(\mathbf{x}_1,y_1) + \sum_{k=2}^{P-1} dist(X_{W_k^{-1}},Y_{W_k^{-2}}) + min \begin{cases} dist^D(x_{L_x},y_{L_y}) \\ dist^V(x_{L_x-1},y_{L_y}) \\ dist^H(x_{L_x},y_{L_{x-1}}) \end{cases} \end{split}$$

Since the first and last components of  $EE(\mathbf{x}, \mathbf{y})$  are exactly the same as the first and second term of GLB, proving  $GLB\_EE(\mathbf{x}, \mathbf{y}) \le EE(\mathbf{x}, \mathbf{y})$  is the same as proving:

$$\max \begin{cases} \delta(\mathbf{x}, \mathbf{y}, w) \\ \delta(\mathbf{y}, \mathbf{x}, w) \end{cases} \leq \sum_{k=2}^{P-1} dist(\mathbf{x}_{W_k^{1}}, \mathbf{y}_{W_k^{2}})$$
 (11)

where  $\delta(\mathbf{x}, \mathbf{y}, w)$  and  $\delta(\mathbf{y}, \mathbf{x}, w)$  are defined in Section 3.3. To prove this statement, we first show  $\delta(\mathbf{x}, \mathbf{y}, w) \leq \sum\limits_{k=2}^{P-1} dist(\mathbf{x}_{W_k}{}^{\scriptscriptstyle 1}, \mathbf{y}_{W_k}{}^{\scriptscriptstyle 2})$ 

and then  $\delta(\mathbf{y},\mathbf{x},w) \leq \sum\limits_{k=2}^{P-1} dist(\mathbf{x}_{W_k^1},\mathbf{y}_{W_k^2})$ . Suppose the optimal alignment for  $x_i$  is  $y_j*$ , following must hold: if  $x_i > max(\mathbf{y})$ , then  $dist^D(x_i,y_j*) \geq dist^D(x_i,max(\mathbf{y}))$ ; if  $x_i < min(\mathbf{y})$ , then  $dist^D(x_i,y_j*) \geq dist^D(x_i,min(\mathbf{y}))$ . Thus, regardless of whether the alignment distance for particular x is produced by distance function  $dist^D$  or  $dist^V$  or  $dist^H$  (depending on whether the optimal alignment path arrives at that cell diagonally or horizontally/vertically), we have:

$$\begin{split} \operatorname{dist}(\mathbf{x}_{W_k^{-1}}, \mathbf{y}_{W_k^{-2}}) &\geq \min(\operatorname{dist}^D(y_j, UE(\mathbf{x})_j), \operatorname{dist}^H(y_j, y_{j-1})) \text{ for } y_j \geq UE(\mathbf{x})_j \\ \operatorname{dist}(\mathbf{x}_{W_k^{-1}}, \mathbf{y}_{W_k^{-2}}) &\geq \min(\operatorname{dist}^D(y_j, LE(\mathbf{x})_j), \operatorname{dist}^H(y_j, y_{j-1})) \text{ for } y_j \leq LE(\mathbf{x})_j \end{split}$$

which indicates that  $\delta(\mathbf{x}, \mathbf{y}, w) \leq \sum_{k=2}^{P-1} dist(\mathbf{x}_{W_k}, \mathbf{y}_{W_k})$ . The same reasoning applies to  $\delta(\mathbf{y}, \mathbf{x}, w) \leq \sum_{k=2}^{P-1} dist(\mathbf{x}_{W_k}, \mathbf{y}_{W_k})$ , we know:

$$\max \begin{cases} \delta(\mathbf{x}, \mathbf{y}, w) \\ \delta(\mathbf{y}, \mathbf{x}, w) \end{cases} \leq \sum_{k=2}^{P-1} dist(\mathbf{x}_{W_k^{-1}}, \mathbf{y}_{W_k^{-2}}).$$

## 4 NEW LOWER BOUNDS BASED ON GLB

In this section, we present the new LBs based on GLB. First, we present the design and formal definition of GLB\_DTW, the GLB variant for DTW (Section 4.1); then, we showcase the formal definition of GLB variants for other elastic measures (Section 4.2).

# 4.1 GLB for Dynamic Time Warping

Improvements of GLB\_DTW over LB\_Keogh, are two-fold: (1) GLB\_DTW incorporates the distances for aligning the first and

## Algorithm 2: Nearest Neighbor Search with GLB

```
Require: X is a n \times L_X matrix of n data time series
Require: L is an n \times 1 vector with labels of data time series
Require: Y is an m \times L_{\mathbf{X}} matrix of m query time series
Require: w is the warping window
Ensure: query_label is an m \times 1 vector containing labels of the m time series in Y
 1: function make envelopes(X, w)
 3: LE = []
 4: for i = 1 to Number_of_Rows(X) do
       for j = 1 to Number_of_Columns(X) do

UE[i, j] = \max(X[i, j - w] : X[i, j + w])
            LE[i,j] = \min(X[i]_{j-w} : X[i]_{j+w})
 8:
        end for
 9: end for
10: return UE, LE
11: end function
13: function GLB(x, y, yue, xle, yue, yle)
14: boundary\_distance = initial\_distance(\mathbf{x}, \mathbf{y}) + ending\_distance(\mathbf{x}, \mathbf{y})
15: query distance = delta(x, yue, yle)
16: data\_distance = delta(y, xue, xle)
17: \ GLB\_dist = boundary\_distance + \max(query\_distance, data\_distance)
18: return GLB dist
20:
21: XUE, XLE = make_envelopes(X, window)
22: YUE, YLE = make_envelopes(Y, window)
23: query\_class = []
24: for i = 1to m do
25:
26:
        lb\_list = []
        for j = 1to n do
28:
            lb\_list_j = GLB(X_i, Y_j, XUE_i, XLE_i, YUE_j, YLE_j)
29:
        end for
        ordering = sort(lb\_list)
        X = X[ordering]
        for j = 1to n do

if lb\_dist < best\_so\_far then
32
33:
34:
               actual\_dist = ElasticMeasure(X_i, Y_i)
35:
               if actual\_dist < best\_so\_far then
36:
                   best\_so\_far = actual\_dist
37:
                   query\_class_i = L_i
38
               end if
39:
            end if
40:
        end for
```

last elements from two sequences (the first two terms of Equation 14), whereas LB Keogh will only include these distances if the first and/or last element of the candidate series is outside of envelopes of query series; (2) whereas LB\_Keogh only computes envelopes for query series, GLB\_DTW computes envelopes for both data series (Equation 12) and query series (Equation 13), finds the distances arising from their respective distance to the other time series and then takes the maximum between the two.

$$\delta_{DTW}(\mathbf{x}, \mathbf{y}, \mathbf{w}) = \sum_{j=2}^{L_{\mathbf{y}}-1} \begin{cases} (y_j - UE(\mathbf{x})_j)^2 \text{ for } y_j \ge UE(\mathbf{x})_j \\ (y_j - LE(\mathbf{x})_j)^2 \text{ for } y_j \le LE(\mathbf{x})_j \end{cases}$$
(12)
$$\delta_{DTW}(\mathbf{y}, \mathbf{x}, \mathbf{w}) = \sum_{i=2}^{L_{\mathbf{x}}-1} \begin{cases} (x_i - UE(\mathbf{y})_i)^2 \text{ for } x_i \ge UE(\mathbf{y})_i \\ (x_i - LE(\mathbf{y})_i)^2 \text{ for } x_i \le LE(\mathbf{y})_i \end{cases}$$
(13)
$$\text{otherwise}$$

$$\text{are 4 shows the LB distance captured by LB Keogh (Figure 1)}$$

$$\delta_{DTW}(\mathbf{y}, \mathbf{x}, w) = \sum_{i=2}^{L_{\mathbf{X}}-1} \begin{cases} (x_i - UE(\mathbf{y})_i)^2 \text{ for } x_i \ge UE(\mathbf{y})_i \\ (x_i - LE(\mathbf{y})_i)^2 \text{ for } x_i \le LE(\mathbf{y})_i \end{cases}$$
(13)

Figure 4 shows the LB distance captured by LB\_Keogh (Figure 4(b-1)) and GLB\_DTW. We observe that by computing envelopes for the data time series and by considering initial and ending boundary distances, GLB\_DTW improves the tightness of the LB (i.e., includes more distances than LB\_Keogh). Formally defined in Equation 14, GLB\_DTW consists of distances captured by a boundary distance function, query envelope, and data envelope:

$$GLB\_DTW = \sqrt{(x_i - y_i)^2 + (x_{L_{\mathbf{X}}} - y_{L_{\mathbf{y}}})^2 + max \begin{cases} \delta_{DTW}(\mathbf{y}, \mathbf{x}, \mathbf{w}) \\ \delta_{DTW}(\mathbf{x}, \mathbf{y}, \mathbf{w}) \end{cases}}$$
(14)

Table 5: Summary of GLB variants for ERP, MSM, TWED, LCSS, EDR, and SWALE.

	Boundaries	$(\mathbf{x}_i - y_j)^2$
ERP	UE	$\overline{\mathrm{UE}(\mathbf{x})_j} = \max(x_{\max(1,j-w)}: x_{\min(L_{\mathbf{x}},j+w)}); UE(\mathbf{y})_i = \max(y_{\max(1,i-w)}: y_{\min(L_{\mathbf{y}},i+w)}) = \max(y_{\max(1,j-w)}: y_{\min(L_{\mathbf{y}},i+w)}) = \min(y_{\max(1,j-w)}: y_{\min(L_{\mathbf{y}},i+w)}) = \min(y_{\min(L_{\mathbf{y}},i+w)}: y_{\min(L_{\mathbf{y}},i+w)}: y_{\min(L_{\mathbf$
ERF	LE	$LE(\mathbf{x})_j = min(x_{max(1,j-w)} : x_{min(L_{\mathbf{x}},j+w)}); LE(\mathbf{y})_i = min(y_{max(1,i-w)} : y_{min(L_{\mathbf{y}},i+w)})$
		$min((x_i - UE(\mathbf{y})_i)^2, (x_i - g)^2)$ for $x_i \ge UE(\mathbf{y})_i$
	$\delta(\mathbf{y}, \mathbf{x}, w)$	$\sum_{i=2}^{L_{\mathbf{x}}-1} \begin{cases} \min((x_i - LE(\mathbf{y})_i)^2, (x_i - g)^2) \text{ for } x_i \leq LE(\mathbf{y})_i \end{cases}$
		0 otherwise
		$min((y_j - UE(\mathbf{x})_j)^2, (y_j - g)^2)$ for $y_j \ge UE(\mathbf{x})_j$
	$\delta(\mathbf{x}, \mathbf{y}, w)$	$\sum_{j=2}^{L_y-1} \begin{cases} \min((y_j - LE(\mathbf{x})_j)^2, (y_j - g)^2) \text{ for } y_j \le LE(\mathbf{x})_j \end{cases}$
		0 otherwise
		$ x_1 - y_1  + C(x_{L_x}, x_{L_x-1}, y_{L_y}, y_{L_y-1})$ where
		$\left(x_{L_{\mathbf{x}}-1} \ge x_{L_{\mathbf{x}}} \ge y_{L_{\mathbf{v}}}\right)$
		$min( x_{L_X} - y_{L_Y} , c)$ if $\begin{cases} x_{L_X} - 1 \le x_{L_X} \le y_{L_Y} \\ y_{U_1-1} \ge y_{U_1} \ge x_{L_1} \end{cases}$
		$\min( x_{L_X} - y_{L_Y} , c)$ if $\begin{cases} y_{L_Y-1} \ge y_{L_Y} \ge x_{L_X} \end{cases}$
	Boundaries	
MSM		$\begin{aligned} &C(\mathbf{x}_{L_{\mathbf{x}}},\mathbf{x}_{L_{\mathbf{x}}-1},y_{L_{\mathbf{y}}},y_{L_{\mathbf{y}}-1}) = min \\ & & \left\{ \begin{aligned} &y_{L_{\mathbf{y}}-1} \leq y_{L_{\mathbf{y}}} \leq x_{L_{\mathbf{x}}} \\ &c +  x_{L_{\mathbf{x}}} - y_{L_{\mathbf{y}}}  \end{aligned} \right. \end{aligned} $ otherwise
MOM		$min \left\{ c +  x_{I_{-}} - x_{I_{1}}  \right\}$ otherwise
		$c +  y_{L_y} - y_{L_{y-1}} $
	UE	$UE(\mathbf{x})_j = max(x_{max(1,j-w)} : x_{min(L_{\mathbf{x}},j+w)}); UE(\mathbf{y})_i = max(y_{max(1,i-w)} : y_{min(L_{\mathbf{y}},i+w)})$
	LE	$LE(\mathbf{x})_{j} = min(x_{max(1,j-w)} : x_{min(L_{\mathbf{x}},j+w)}); LE(\mathbf{y})_{i} = min(y_{max(1,i-w)} : y_{min(L_{\mathbf{y}},i+w)})$
		$min( x_i - UE(\mathbf{v})_i , c)$ for $x_i \ge UE(\mathbf{v})_i$
	$\delta(\mathbf{y}, \mathbf{x}, w)$	$\sum_{i=2}^{L_X-1} \begin{cases} \min( x_i - LE(\mathbf{y})_i , c) \text{ for } x_i \le LE(\mathbf{y})_i \end{cases}$
		0 otherwise
	$\delta(\mathbf{x}, \mathbf{y}, w)$	$\sum_{j=2}^{L_y-1} \min( y_j - UE(\mathbf{x})_j , c) \text{ for } y_j \ge UE(\mathbf{x})_j$ $\min( y_j - LE(\mathbf{x})_j , c) \text{ for } y_j \le LE(\mathbf{x})_j$
	J (24, y, 11)	$\sum_{j=2}^{nun}  min( y_j - LE(\mathbf{x})j , c) \text{ for } y_j \leq LE(\mathbf{x})j$ 0 otherwise
		((v u )2   (v u )2
	D 1 .	$(x_{L_x} - y_{L_y})^2 + (x_{L_x-1} - y_{L_y-1})^2$
	Boundaries	$(x_1 - y_1)^2 + min \left\{ (x_{L_x} - x_{L_{x-1}})^2 + \lambda \right\}$
TWED	TIP	$(y_{L_y} - y_{L_y-1})^2 + \lambda$
	UE LE	$UE(\mathbf{x})_{j} = max(x_{max(1,j-w)} : x_{min(L_{\mathbf{x}},j+w)}); UE(\mathbf{y})_{i} = max(y_{max(1,i-w)} : y_{min(L_{\mathbf{y}},i+w)})$
	LE	$LE(\mathbf{x})_{j} = min(x_{max(1,j-w)} : x_{min(L_{\mathbf{x}},j+w)}); LE(\mathbf{y})_{i} = min(y_{max(1,i-w)} : y_{min(L_{\mathbf{y}},i+w)})$
		$min((x_i - UE(\mathbf{y})_i)^2 + (x_{j-1} - UE(\mathbf{y})_{i-1})^2)$
		for $x_i \ge UE(\mathbf{y})_i, x_{i-1} \ge UE(\mathbf{y})_i$
	$\delta(\mathbf{y}, \mathbf{x}, w)$	$\sum_{i=2}^{L_X-1} \left\{ \min((x_i - LE(\mathbf{y})_i)^2 + (x_{i-1} - LE(\mathbf{y})_{i-1})^2) \right.$
		for $x_i \le LE(\mathbf{y})_i, x_{i-1} \le LE(\mathbf{y})_i$
		0 otherwise
		$min((y_j - UE(\mathbf{x})_j)^2 + (y_{j-1} - UE(\mathbf{x})_{j-1})^2)$
		for $y_j \ge UE(\mathbf{x})_j, y_{j-1} \ge UE(\mathbf{x})_j$
	$\delta(\mathbf{x}, \mathbf{y}, w)$	$\sum_{j=2}^{L_Y-1} \left\{ \min((y_j - LE(\mathbf{x})_j)^2 + (y_{j-1} - LE(\mathbf{x})_{j-1})^2) \right.$
		for $y_j \le LE(\mathbf{x})_j, y_{j-1} \le LE(\mathbf{x})_j$
		0 otherwise
		$M_{LCSS}(x_1, y_1, \epsilon) + M_{LCSS}(x_{L_x}, y_{L_y}, \epsilon)$ where
	Boundaries	$M_{LCSS}(A, B, \epsilon) = \begin{cases} 1 \text{ if }  A - B  \le \epsilon \\ 0 \text{ otherwise} \end{cases}$
		$M_{LCSS}(A, B, \epsilon) = \begin{cases} 0 \text{ otherwise} \end{cases}$
LCSS	UE	$\mathrm{UE}(\mathbf{x})_j = \max(x_{\max(1,j-w)}: x_{\min(L_\mathbf{x},j+w)}); UE(\mathbf{y})_i = \max(y_{\max(1,i-w)}: y_{\min(L_\mathbf{y},i+w)}); UE(\mathbf{y})_i = \max(y_{\min(1,i-w)}: y_{\min(1,i-w)}: y_{\min(1,i-w)}); UE(\mathbf{y})_i = \max(y_{\min(1,i-w)}: y_{\min(1,i-w)}: y_{\min(1,i$
	LE	$LE(\mathbf{x})_j = min(x_{max(1,j-\mathbf{w})} : x_{min(L_{\mathbf{x}},j+\mathbf{w})}); LE(\mathbf{y})_i = min(y_{max(1,i-\mathbf{w})} : y_{min(L_{\mathbf{y}},i+\mathbf{w})})$
		$M_{LCSS}(x_i, UE(\mathbf{y})_i) \text{ for } x_i \ge UE(\mathbf{y})_i$
	$\delta(\mathbf{y}, \mathbf{x}, w)$	$\sum_{i=2}^{L_{x}-1} \left\{ M_{LCSS}(x_{i}, LE(\mathbf{y})_{i}) \text{ for } x_{i} \leq LE(\mathbf{y})_{i} \right.$
		0 otherwise
		$M_{LCSS}(y_j, UE(\mathbf{x})_j) \text{ for } y_j \ge UE(\mathbf{x})_j$
	$\delta(\mathbf{x}, \mathbf{y}, w)$	$\sum_{j=2}^{Ly-1} \begin{cases} M_{LCSS}(y_j, UE(\mathbf{x})_j) \text{ for } y_j \ge UE(\mathbf{x})_j \\ M_{LCSS}(y_j, UE(\mathbf{x})_j) \text{ for } y_j \le UE(\mathbf{x})_j \end{cases}$
		(0 otherwise
	GL	B Variants for Elastic Measures without Existing Lower Bounds
	Boundaries	$0 + \min \begin{cases} 1 \\ M_{EDR}(x_{L_x}, y_{L_y}, \epsilon) \end{cases}, \text{ where } M_{EDR}(A, B, \epsilon) = \begin{cases} 0 \text{ if }  A - B  \le \epsilon \\ 1 \text{ otherwise} \end{cases}$
		$M_{EDR}(x_{L_x}, y_{L_y}, \epsilon)$ , where $M_{EDR}(x_L, b, \epsilon)$ 1 otherwise
EDR	UE	$UE(\mathbf{x})j = max(x_{max}(1,j-w) : x_{min}(L_{\mathbf{x}},j+w)); UE(\mathbf{y})i = max(y_{max}(1,i-w) : y_{min}(L_{\mathbf{y}},i+w))$
	LE	$\mathrm{LE}(\mathbf{x})_j = \min(x_{max(1,j-w)} : x_{\min(L_{\mathbf{x}},j+w)}); LE(\mathbf{y})_i = \min(y_{max(1,i-w)} : y_{\min(L_{\mathbf{y}},i+w)})$
		$M_{EDR}(y_j, UE(\mathbf{x})_j, \epsilon)$ for $y_j \ge UE(\mathbf{x})_j$
	$\delta(\mathbf{x}, \mathbf{y}, w)$	$\sum_{j=2}^{L_Y-1} \left\{ M_{EDR}(y_j, LE(\mathbf{x})_j, \epsilon) \text{ for } y_j \leq LE(\mathbf{x})_j \right.$
		0 otherwise
		$M_{EDR}(x_j, UE(\mathbf{y})_i, \epsilon)$ for $x_j = i \ge UE(\mathbf{y})_i$
	$\delta(\mathbf{y}, \mathbf{x}, w)$	$\sum_{i=2}^{L_X-1} \left\{ M_{EDR}(x_i, LE(\mathbf{y})_i, \epsilon) \text{ for } x_i \le LE(\mathbf{y})_i \right.$
		0 otherwise
		$r \text{ if }  A - B  < \varepsilon$
	Boundaries	$M_{SWALE}(x_{L_x}, y_{L_y}, \epsilon)$ , where $M_{SWALE}(A, B, \epsilon) = \begin{cases} r &  A - B  \le \epsilon \\ p & \text{otherwise} \end{cases}$
SWALE	UE	$UE(\mathbf{x})_j = max(x_{max(1,j-\mathbf{w})} : x_{min(L_{\mathbf{x}},j+\mathbf{w})}); UE(\mathbf{y})_i = max(y_{max(1,i-\mathbf{w})} : y_{min(L_{\mathbf{y}},i+\mathbf{w})})$
	LE	$LF(\mathbf{x}) := min(\mathbf{x}_{max}(1,j-\mathbf{w}) \cdot \mathbf{x}_{min}(L_{\mathbf{x}},j+\mathbf{w})) \cdot UF(\mathbf{y}) := min(\mathbf{x}_{max}(1,i-\mathbf{w}) \cdot ymin(L_{\mathbf{y}},i+\mathbf{w})) \cdot UF(\mathbf{y}) := min(\mathbf{x}_{max}(1,i-\mathbf{w}) \cdot ymin(L_{\mathbf{y}},i$
		$LE(\mathbf{x})_j = min(x_{max(1,j-w)} : x_{min(L_{\mathbf{x}},j+w)}); LE(\mathbf{y})_i = min(y_{max(1,i-w)} : y_{min(L_{\mathbf{y}},i+w)})$ $M_{CM-M-P}(y, LE(\mathbf{x}); \varepsilon) \text{ for } y \ge LE(\mathbf{x}).$
	8(*******	$\sum_{L_Y = 1}^{L_Y = 1} \begin{cases} M_{SWALE}(y_j, UE(\mathbf{x})_j, \epsilon) \text{ for } y_j \ge UE(\mathbf{x})_j \\ M_{SWALE}(y_j, UE(\mathbf{x})_j, \epsilon) \text{ for } y_j \le UE(\mathbf{x})_j \end{cases}$
	$\delta(\mathbf{y}, \mathbf{x}, w)$	$\sum_{j=2}^{L_Y-1} M_{SWALE}(y_j, LE(\mathbf{x})_j, \epsilon) \text{ for } y_j \le LE(\mathbf{x})_j$
		(0 otherwise
		$M_{SWALE}(x_j, UE(\mathbf{y})_i, \epsilon) \text{ for } x_j = i \ge UE(\mathbf{y})_i$
		$1 \rightarrow 2 \rightarrow 44 - \cdots - (v_r + V(v_r), c) \text{ for } v_r < 1 E(v_r)$
	$\delta(\mathbf{x}, \mathbf{y}, w)$	$\sum_{i=2}^{L_{\mathbf{x}}-1} M_{SWALE}(x_i, LE(\mathbf{y})_i, \epsilon) \text{ for } x_i \leq LE(\mathbf{y})_i$ 0 otherwise

# **GLB Variants of Other Elastic Measures**

Having introduced GLB for DTW, we now focus on six alternative elastic measures. Table 5 summarizes all LBs produced by GLB.

4.2.1 Edit Distance with Real Penalty (ERP). Similar to GLB\_DTW, GLB\_ERP uses  $(x_i - y_j)^2$  as  $initial\_distance(\mathbf{x}, \mathbf{y})$ . On the other hand, as shown in Table 4, while DTW has the same distance functions for diagonal and horizontal/vertical movements, ERP has  $(x_i - g)^2$  or  $(y_j - g)^2$  for horizontal/vertical movements, but  $(x_i - y_j)^2$  for diagonal movements; having different distance functions results in two differences of GLB\_ERP from GLB\_DTW: (1) GLB\_ERP takes the minimum of the two possible distance functions in the construction of  $\delta(\mathbf{y}, \mathbf{x}, w)$  and  $\delta(\mathbf{x}, \mathbf{y}, w)$ , as shown in Table 5 (2) GLB\_ERP takes the minimum of possible distance functions when computing the ending boundary distance.

4.2.2 Longest Common Subsequence (LCSS). As a threshold-based elastic measure, LCSS employs a matching function  $M_{LCSS}$  (See Table 2) to determine if two elements are close enough to be "matched". GLB\_LCSS uses the same matching function  $M_{LCSS}$  to compute boundary distances by checking if the two initial elements and two ending elements from two time series "match" with each other (see Table 5). In constructing  $\delta(\mathbf{y}, \mathbf{x}, w)$  and  $\delta(\mathbf{x}, \mathbf{y}, w)$  (see Table 5), the same matching function  $M_{LCSS}$  is used to compute distances arising from the distance between one time series and envelopes of the other time series. A unique feature of LCSS is that LCSS first calculates similarity and computes distance based on similarity; accordingly, GLB\_LCSS also calculates similarity ( $LCSS\_Sim$ ) before converting similarity to distance (see Table 5).

4.2.3 Move-Split-Merge (MSM). As shown in Table 3, the alignment distance between each pair of elements  $x_i$  and  $y_i$  using MSM depends on both distance between  $x_i$  and  $y_i$  and  $x_{i-1}$  and  $y_{i-1}$ , GLB\_MSM defines a new distance function C to capture the distance for ending boundary alignment (see Table 5). Since the distance associated with each movement to the next cell in the distance matrix can be either  $|x_i - y_i|$  or c, GLB\_MSM takes the minimum of the two to ensure the distance is less than the actual distance.

4.2.4 Time Warp Edit Distance (TWED). Since the distance function for diagonal movements includes pairwise comparison between  $x_i$  and  $y_j$  as well as  $x_{i-1}$  and  $y_{j-1}$  in TWED, GLB\_TWED only captures the distance associated with cases when two consecutive elements both fall outside of envelopes. Since there are three possible distances associated with ending boundary alignment, GLB\_TWED takes the minimum among the three alternatives.

4.2.5 EDR. Similar to LCSS, EDR is a threshold-based elastic measure that employs a match function  $M_{EDR}$  (see Table 5) to decide if two elements mismatch and add to distance between time series. As shown in Table 5, GLB\_EDR adopts the same  $M_{EDR}$  in defining the distance when computing  $\delta(\mathbf{y}, \mathbf{x}, w)$  and  $\delta(\mathbf{x}, \mathbf{y}, w)$ .

4.2.6 SWALE. Similar to LCSS and EDR, SWALE is a threshold-based elastic measure that employs a match function  $M_{EDR}$ , with the difference being SWALE parameterizes the penalty for a mismatch and the reward for a match with p and r respectively, instead of using 1 and 0 for match and mismatch. GLB\_SWALE uses the same matching function to compute the initial and ending boundary distances, as well as in  $\delta(\mathbf{x}, \mathbf{y}, w)$  and  $\delta(\mathbf{y}, \mathbf{x}, w)$  to compute the distance between one time series and envelopes of the other.

# 5 EXPERIMENT SETTINGS

In this section, we report our experimental settings. We aim to provide the first comprehensive study of LBs for a diverse set of elastic measures. We aim to show in each case how GLB improves the current state of the art. Therefore, first, we evaluate the performance of GLB against state-of-the-art LBs for DTW, which is the most widely used elastic measure. Then, we compare GLB with state-of-the-art LBs of alternative elastic measures. Finally, for elastic measures without known LBs, we present the performance of GLB, which could serve as the baseline for further research.

**Datasets:** Our experiments are performed on 128 datasets from the UCR archive [26]. The UCR archive includes datasets from various domains and is the largest public collection of labeled time series datasets. Datasets are normalized and split into training and test sets. For datasets with varying lengths and missing values, we resort to pre-processed versions in [70], which used standardized resampling and interpolation methods to fix these issues.

**Platforms:** We ran our experiments on a server with the following configuration: Dual Intel(R) Xeon(R) Silver 4116 (12-core with 2-way SMT), 2.10 GHz, 196GB RAM. The server ran Ubuntu 18.04.3 LTS (64-bit) with Python 3.7.5, Numba 0.53.1, and GCC 8.4.0 compiler. **Implementations:** We implemented all methods in Python for consistency. We employ Numba [53], which translates Python functions to optimized machine code at runtime to accelerate the computation of lower bounds and elastic measures. For reproducibility purposes, we make our source code available.<sup>2</sup>

Choice of Parameters: We choose parameters that demonstrated optimal performance in empirical evaluation [78] or recommended by authors who proposed these elastic measures (Table 6). For the choice of warping window, w, searching for the optimal window size would require a laborious process [90]. For consistency with previous evaluation efforts (see references in Table 1), we use a window size of 5% for all lower bounds and elastic measures. This window size has also achieved state-of-the-art 1-NN classification accuracy [74, 78] for unsupervised settings.

**Baselines:** We compare GLB variants against state-of-the-art LBs for elastic time series distances. For DTW, we compare GLB\_DTW against the following strong LBs:

- LB\_Yi: a simple LB that compares every element of the data with the min and max values of the query time series
- LB\_Keogh: the LB that introduced query envelopes
- LB\_New: a tight LB with boundary dependence
- LB\_Improved: a LB that builds on LB\_Keogh and exploits data time series characteristics to increase its tightness

We focus on key existing LBs for DTW (as identified in Section 2, Table 4) and omit comparisons with variants that combine techniques or introduce adaptations to trade off pruning power for efficiency. These combinations or adaptations are often orthogonal contributions that can be incorporated into GLB as well with similar effects. Instead, we include comparisons with state-of-the-art LBs of 4 additional elastic measures, namely, ERP, MSM, TWED, and LCSS, overall 11 LBs spanning 5 elastic measures (see Table 1 for references and Tables 2, 3, and 4 for their exact mathematical formulations). More specifically, we omitted LB\_Enhanced [89] because it requires a time-consuming parameter tuning process and, therefore, it is not directly comparable. Similarly, we omitted LB\_Petitjean and LB\_Webb, which are variants of LB\_Improved.

 $<sup>^2</sup> www.timeseries.org/glb \\$ 

Elastic Measures	Parameters
Dynamic Time Warping (DTW)	-
Edit Distance on Real Sequences (EDR)	$\epsilon = 0.1$
Longest Common Subsequence (LCSS)	$\epsilon = 0.2$
Sequence Weighted Alignment (SWALE)	$\epsilon = 0.2, p = 5, r = 1$
Edit Distance with Real Penalty (ERP)	gap = 0
Move-Split-Merge (MSM)	c = 0.5
Time Warp Edit Distance (TWE)	$\lambda = 1, \nu = 0.0001$

Tighter lower bounds, like LB\_Improved and LB\_New, typically involve higher costs, resulting in poor speedup relative to their impressive pruning power. Therefore, we evaluate LBs in isolation but also, as noted in [54], in a cascade, where the expensive LBs are only computed for the cases that cannot be pruned by cheaper LBs. To take advantage of such property, we evaluate the performance of LB\_Improved and LB\_New running together with LB\_Keogh.

Metrics: We compare the effectiveness lower bounds on (1) pruning power and (2) speedup. Pruning power is the percentage of the true distance computation avoided in the 1-NN search due to adopting a lower bound. Speedup is calculated by dividing the runtime of traditional 1-NN search by the runtime of 1-NN search using a lower bound. We compute all reusable components of lower bounds as a pre-processing step to avoid repetitive calculation and include pre-processing time in our runtime measurement.

## **6 EXPERIMENTAL RESULTS**

In this section, we report the results of our experiments. First, we evaluate GLB\_DTW against existing lower bounds for DTW, which is the most widely used and studied elastic measure (Section 6.1). Then, we compare GLB variants for elastic measures other than DTW (including LCSS, ERP, MSM, and TWED) with their existing lower bounds (Section 6.2). Additionally, we showcase the performance of novel lower bounds we propose for EDR and SWALE (Section 6.3). Next, we present an exploratory breakdown analysis of GLB\_DTW to understand the contributions of its three components to its performance (Section 6.4). Finally, we demonstrate trade-offs for GLB by varying relevant parameters (Sections 6.5 and 6.6).

# 6.1 Evaluation of GLB DTW

To understand if GLB\_DTW (Section 3) is an effective LB for DTW, we evaluate it against existing DTW LBs using their pruning power and speedup in 1-NN classification task across 128 datasets. Table 7 reports the performance of GLB\_DTW against DTW LBs.

Comparison against LB\_Kim, LB\_Yi, LB\_Keogh: As shown in Table 7, GLB\_DTW's pruning power is higher than LB\_Keogh by 11.7%, LB\_Kim by 58.5%, and LB\_Yi by 22.2% on average. Across 128 datasets, the difference in pruning power between GLB\_DTW and LB\_Keogh goes up to 62.3% and for LB\_Kim up to 99.0%. The advantage in pruning power is consistent as out of a total of 128 datasets, GLB\_DTW obtains higher pruning power than LB\_Kim, LB\_Keogh in 127 datasets (Figure 6, Part (a)) and higher than LB\_Yi in 110 datasets. GLB's high pruning power shows the effectiveness of query/data dependencies and boundary distances in establishing a tighter lower bound. In terms of speedup, GLB\_DTW outperforms LB\_Kim in 127 datasets, LB\_Keogh in 115 datasets (Figure 6 Part (b)),

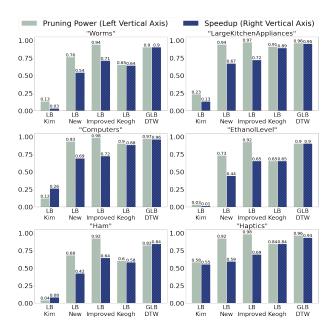


Figure 5: Average pruning power and speedup of DTW LBs.

and LB\_Yi in 109 datasets. As summarized in Table 6, GLB\_DTW's is faster than LB\_Keogh by 1.6x, faster than LB\_Kim by 8x on average, and faster than LB\_Yi by 4.14x. The advantage of GLB\_DTW over LB\_Keogh goes up to 3.4x, for LB\_Kim up to 36.6x, and for LB\_Yi up to 17.44x. GLB\_DTW's superior speedup results from its higher tightness over LB\_Kim, LB\_Keogh, and LB\_Yi.

Comparison against LB New and LB Improved: We compare the pruning power and speedup of GLB DTW against LB New and LB\_Improved both when they are used in isolation and running in cascade (with LB\_Keogh). Firstly, we compare the pruning power and speedup of GLB\_DTW, LB\_New, and LB\_Improved when used in isolation. Table 7 shows the GLB\_DTW's pruning power is slightly better than LB New by 0.6%, and worse than LB Improved by 2.8% on average. Interestingly, despite the similarity in pruning power among three lower bounds, LB\_Improved and LB\_New achieved significantly less speedup than GLB\_DTW. Across 128 datasets, GLB\_DTW outperforms LB\_New in 127 datasets and LB\_Improved in all 128 datasets in terms of speedup. It is worth noting that the fact that GLB\_DTW has less pruning power than LB\_Improved is a result of purposeful design. As we mentioned in Section 3, GLB can integrate LB Improved into GLB DTW by replacing current envelopes with the tighter LB\_Improved and thus achieve higher pruning power. Nonetheless, the inability to precompute and reuse LB\_Improved results in a drastic increase in runtime. In rare situations where even a slight increase in pruning power would compensate for large overheads, GLB\_DTW can integrate LB\_Improved and meet the needs.

Now we focus on the comparison between GLB\_DTW running in a cascade of two envelopes against the cascade of LB\_Keogh + LB\_New and the cascade of LB\_Keogh + LB\_Improved. This setting is important for cases where the precomputation of the GLB envelopes is not possible or parameters may need to change dynamically and, therefore, caching of envelopes is not useful. In

Table 7: Summary of pruning power and speedup for exist-
ing DTW lower bounds and GLB_DTW.

Evaluation of DTW Lower Bounds							
Metrics	Pr	Pruning Power			Speedup		
LBs	Average	Max	Std	Average	Max	Std	
LB_Kim	23.41%	87.44%	0.1166%	2.456x	7.591x	1.0752x	
LB_Yi	59.72%	83.28%	21.28%	1.321x	4.331x	0.5325x	
LB_Keogh	70.19%	98.49%	27.37%	6.314x	25.071x	6.0849x	
LB_New	81.27%	99.42%	22.12%	1.110x	2.696x	0.5178x	
LB_Improved	84.71%	99.69%	20.02%	1.492x	3.077x	0.6814x	
GLB_DTW	81.93%	99.69%	22.44%	10.176x	41.725x	10.3284x	
	LBs	when ru	nning in	a cascade			
LB_New (Cascade)	80.70%	99.29%	21.92%	4.779x	19.838x	4.8138x	
LB_Improved (Cascade)	84.0%	99.61%	20.10%	7.393x	32.192x	8.1032x	
GLB_DTW (Cascade)	80.35%	99.62%	23.08%	10.381x	43.491x	10.1959x	

terms of speedup, GLB\_DTW enjoys a significant advantage by outperforming LB\_Keogh + LB\_New in 123 datasets and LB\_Keogh + LB\_Improved in 124 datasets. On average, GLB\_DTW is 1.4x faster than LB\_Keogh + LB\_Improved and 2.1x faster than LB\_Keogh + LB\_New. The evaluation of GLB\_DTW, LB\_New, and LB\_Improved highlights the fact that high pruning power doesn't necessarily translate into high speedup. As "HandOutlines", "AllGestureWiimoteX" and "ScreenType" datasets in Figure 5 demonstrate, although LB\_Keogh + LB\_Improved has higher pruning power than GLB\_DTW, their speedups are significantly less GLB\_DTW.

# 6.2 Evaluation of GLB on Elastic Measures other than DTW

Having demonstrated the advantage of GLB\_DTW over existing DTW LBs, we now turn our attention to the comparison between GLB and state-of-the-art LBs of other common elastic measures, including LCSS, ERP, MSM, and TWED, in terms of pruning power and speedup. Next, we report the performance of GLB against three ERP LBs and the LBs of LCSS, MSM, and TWED, respectively.

Comparison against LB\_Kim-ERP, LB\_ERP, and LB\_Keogh-ERP: Out of a total of 128 datasets, GLB\_ERP obtains higher pruning power than LB\_Kim-ERP, LB\_ERP, and LB\_Keogh-ERP in 127 datasets, as shown in Figure 6 Part (a). As shown in Table 8, GLB\_ERP's pruning power is higher than LB\_Keogh-ERP by 66.9%, LB\_Kim-ERP by 53.7%, and LB\_Keogh-ERP by 11.7% on average. In terms of speedup, GLB\_DTW outperforms LB\_Kim-ERP in 124 datasets, LB\_ERP in 126 datasets, and LB\_Keogh-ERP in 118 datasets, as shown in Figure 6 Part (b). As summarized in Table 8, GLB\_DTW's is faster by LB\_Keogh-ERP by 1.3x, faster than LB\_Kim-ERP by 4.3x on average, and faster than LB\_ERP by 4.8x. GLB\_ERP outperforms all ERP lower bounds in speedup.

Comparing with LB\_Keogh-LCSS, LB\_MSM, and LB\_TWED: Now we focus on LBs of LCSS, MSM, and TWED, which only have one respective LB reported in the literature. For LCSS, Figure 6 part (a) shows that GLB\_LCSS significantly outperforms LB\_Keogh-LCSS in pruning power: out of a total of 128 datasets, GLB\_LCSS achieves higher pruning power than LB\_Keogh-LCSS in 119 datasets. In terms of speedup, GLB\_LCSS outperforms LB\_Keogh-LCSS in 76 datasets, as shown in Figure 6 part (b). For MSM,

Table 8: Summary of average pruning power and speedup for existing LBs and GLB variants. The last column, "win," indicates the number of datasets where the corresponding LB is winning out of a total of 128 datasets in terms of speedup.

Evaluation of Lower Bounds of Other Elastic Measures								
Metrics	Pruning Power Speedup				Win			
LBs	Average	Max	Std	Average	Max	Std	-	
Edit Distance with Real Penalty (ERP)								
LB_ERP	0.61%	36.99%	3.66%	0.967x	1.469x	0.0953x	1	
LB_Kim	13.87%	76.50%	16.57%	1.134x	3.536x	0.3153x	1	
LB_Keogh-ERP	55.84%	96.23%	29.15%	3.222x	10.784x	2.4601x	9	
GLB_ERP	67.55%	97.56%	26.67%	4.780x	19.191x	4.4416x	117	
	Move-Split-Merge (MSM)							
LB_MSM	18.97%	92.92%	23.01%	1.358x	7.495x	0.8324x	15	
GLB_MSM	42.09%	95.90%	30.82%	2.302x	10.217x	1.9056x	113	
	Т	ime Warp	Edit Distar	ice (TWED	)			
LB_TWED	2.42%	56.36%	6.8%	0.988x	2.082x	0.1182x	9	
GLB_TWED	69.75%	99.78%	31.94%	6.517x	26.285x	6.227x	119	
Longest Common Subsequence (LCSS)								
LB_Keogh-LCSS	43.72%	100%	29.48%	2.291x	14.257x	2.0396x	52	
GLB_LCSS	53.58%	100%	31.19%	2.541x	11.775x	2.1548x	76	

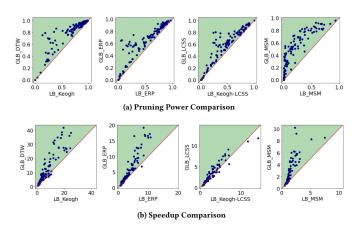


Figure 6: Comparison of GLB variants and state-of-the-art LBs for DTW, ERP, LCSS, and MSM. Part(a) and Part(b) show the pruning power and speedup over 128 datasets, respectively. The blue dots above the diagonal indicate datasets over which GLB outperforms the state of the art.

we illustrate that GLB MSM significantly outperforms LB MSM. GLB MSM achieves higher pruning power than LLB MSM in 120 out of 128 datasets, as shown in Figure 6 Part (a). GLB MSM pruning power is higher than LB\_MSM by 23.13% on average. In terms of speedup, GLB\_MSM outperforms LB\_MSM in 113 datasets, as shown in Figure 6 Part (b), and GLB MSM is faster than LB MSM by 1.57x on average. For **TWED**, we also see that GLB\_TWED significantly outperforms LB TWED. GLB TWED's achieves higher pruning power than LB TWED in 125 out of 128 datasets. GLB TWED pruning power is higher than LB\_TWED by 67.32% on average. In terms of speedup, GLB\_TWED outperforms LB\_TWED in 119 datasets, and GLB\_TWED is faster than LB\_TWED by 6.56x on average. Based on GLB's consistent advantage in pruning power and speedup over existing baselines, we conclude that GLB variants have established new state-of-the-art performance in popular elastic measures, including LCSS, MSM, and TWED.

Table 9: Summary of average pruning power and speedup for GLB\_EDR and GLB\_SWALE

Evaluation of Novel EDR and SWALE Lower Bounds							
Metrics	Pr	uning Pow	er	Speedup			
LBs	Average	Max	Std	Average	Max	Std	
Edit Distance for Real Sequences (EDR)							
GLB_EDR	GLB_EDR 52.55% 97.26% 31.99% 4.096x 28.909x 4.6569x					4.6569x	
Sequence Weighted Alignment (SWALE)							
GLB_SWALE	85.07%	99.996%	24.66%	.66% 10.684x 46.764x 9.9162x			

Table 10: Summary of pruning power and speedup for existing DTW lower bounds and GLB\_DTW

Lower Bounds	Pruning Power	Speedup
Query Only	69.08%	6.260x
Query + Boundary	74.40%	6.737x
Query + Data	79.89%	9.437x
Query + Data + Boundary (GLB)	81.93%	10.17x

## 6.3 Case Studies of EDR and SWALE LBs

Having shown the comparison between GLB and existing lower bounds for various elastic measures, we now demonstrate the performance of the novel lower bounds we propose for EDR and SWALE based on GLB. Table 9 shows the pruning power and speedup for EDR and SWALE. We observe that GLB\_EDR and GLB\_SWALE achieve high pruning power and speed up comparable to that of state-of-the-art lower bounds for other elastic measures.

# 6.4 Break Down Analysis

To understand the contribution of various GLB components to its superior performance, we performed a breakdown analysis of GLB\_DTW to find the contributions of four different component combinations: query envelopes only, query envelopes and boundary distances, query envelopes and data envelopes, as well as GLB itself.

As shown in Table 10, adopting data envelopes is the driving force behind the improvement in pruning power and speedup. To further understand the contributions of data envelopes, we have kept a record of the percentage of instances when GLB uses data envelopes and query envelopes. We found that GLB chooses the query envelope 50.63% of the time and the data envelope 49.37% of the time. Such results indicate that using only query envelopes, as in the case of LB\_Keogh, is suboptimal half of the time, and accounting for both data envelopes and query envelopes is essential.

## 6.5 Varying Window and other Parameters

In order to assess the effect of varying parameters and changing window sizes on the performance of elastic measure LBs, we conduct a comprehensive evaluation of their impact on the performance of MSM LBs, which produced exceptional results in recent studies [78]. In the first experiment, we held the parameter c constant at 0.5 while systematically increasing the window size from 5% to 100% of the length of the original time series. As the window size increases, the performance of the LBs deteriorates. The results, depicted in Figure 7(b), indicate a declining speedup of LBs as the window size increases. Despite this trend, it is evident that GLB\_MSM consistently outperforms LB\_MSM across all window sizes.

In the second experiment, we maintained a fixed 5% window size and evaluated ten different parameter values ranging from 0.01 to 500. The results, depicted in Figure 7(a), demonstrate that

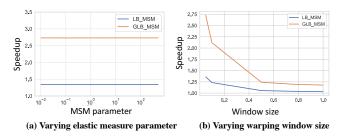


Figure 7: Average speedup of LB\_MSM and GLB\_MSM with varying parameter and warping window size.

variations in this parameter of MSM have limited influence on the performance of MSM LBs with respect to speedup. Furthermore, it is evident that GLB\_MSM consistently surpasses LB\_MSM.

# 6.6 Varying Pre-computed Envelopes

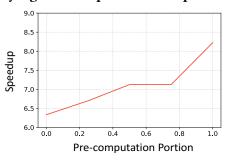


Figure 8: Trade-off of precomputation vs. speedup.

In order to evaluate the trade-off between precomputation of data envelopes and speedup of GLB, we conducted a comprehensive analysis of GLB\_DTW in a cascade with varying degrees of precomputation: 0%, 25%, 50%, 75%, and 100%. As precomputation increases, speedup increases at the cost of higher storage requirements. Our results, as illustrated in Figure 8, confirm this hypothesis, showing a positive correlation between increased precomputation from 0% to 100% and increased speedup.

## 7 CONCLUSION

In this paper, we presented GLB, a generalized framework for deriving lower bounds for elastic measures. Motivated by the disproportionate attention of the research effort to a single elastic measure, we designed an LB to extract cache-friendly summary characteristics, adaptively exploit summaries of both query and target data time series, capture boundary distances in an unsupervised and parameter-free manner, and achieve high LB tightness with low computational cost, resulting in substantial speedup. Based on GLB, we propose new LBs for all elastic measures, including those without existing LBs in the literature. We extensively evaluate GLB against existing LBs, resulting in one of the most comprehensive experimental studies in this area. In particular, we included 11 stateof-the-art lower bounds spanning 5 elastic measures and used 128 datasets in our evaluation. Our findings show that GLB outperforms all existing baseline LBs in terms of speedup. GLB LBs for elastic measures without existing LBs obtain comparable performance to state-of-the-art LBs of other elastic measures. Overall, GLB is a generalizable framework for developing efficient LBs that facilitate accurate and fast time series similarity search.

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