

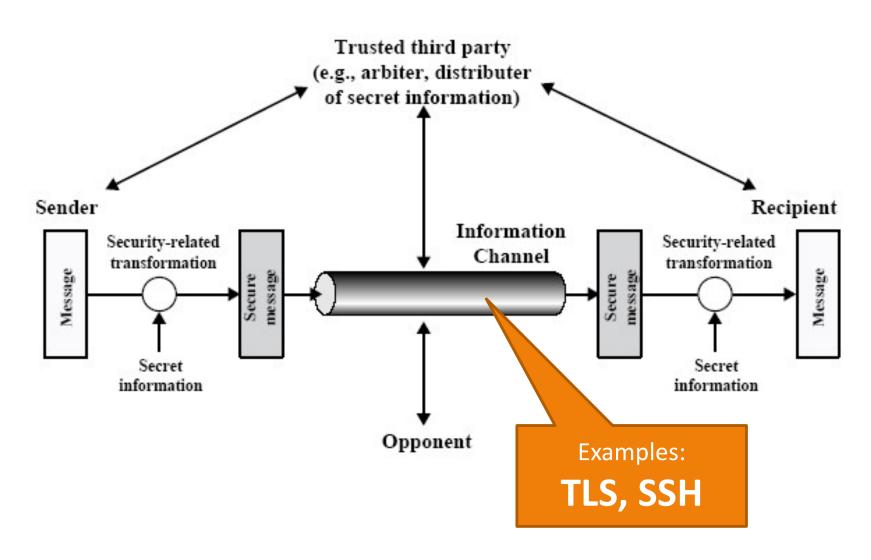
TLS and SSH overview

Segurança Informática em Redes e Sistemas 2022/23

Miguel Pardal

Ack: Carlos Ribeiro, André Zúquete, Miguel P. Correia, Ricardo Chaves

Secure communication channel



Roadmap

- TLS Transport Layer Security
- SSH Secure Shell
- Key management

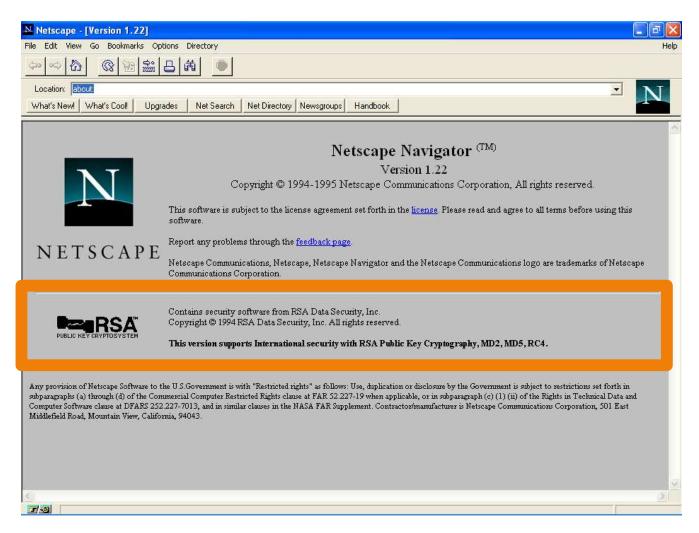
Roadmap

- TLS Transport Layer Security
 - HTTPS
- SSH Secure Shell
- Key management

HTTPS

- HTTPS = HTTP over SSL (now TLS)
 - In fact, SSL was originally designed to be used with HTTP
 - Netscape Navigator (1994)
 - Forefather of Mozilla Firefox
 - Version history
 - SSL 1.0 1994
 - SSL 2.0 1995
 - SSL 3.0 1996
 - TLS 1.0 1999
 - TLS 1.1 2006
 - TLS 1.2 2008
 - TLS 1.3 2018

Netscape Navigator



 A few months later, in July 1995, a computer programmer ordered the first book ever sold by an online bookstore named after a river in South America...

TLS goals

- Secure communication channels over TCP/IP
 - Current version: TLS 1.3 august 2018
 - Standard based on the deprecated SSL (Secure Sockets Layer)
 - Sometimes called SSL for that reason
 - Manages secure sessions over TCP/IP per application
 - Initially designed for the HTTP protocol (HTTPS)
 - Currently used by other protocols, e.g., SMTP, IMAP, POP3
- Security mechanisms
 - Authentication of the communicating parties
 - Confidentiality and integrity of the communication
 - Key distribution

TLS utilization

- TLS is just a protocol; not a standard API
- Common APIs:
 - Reference API for SSL: SSLref (Netscape)
 - Public implementations: OpenSSL, GnuTLS, SSLeay,
 Mbed TLS, Java Secure Socket Extension (JSSE)
- Remote interfaces:
 - Conventional: protocol/port
 - e.g., TCP/443 for HTTPS
 - STARTTLS: allows upgrading text connection to TLS
 - Defined for SMTP, IMAP, POP, SMTP, FTP, IRC,...

TLS operation management

- Client-Server model as in TCP
- The applications (e.g., web, mail) define the strategy
- Authentication
 - If it is needed and how it is performed
- Cryptographic algorithms
 - The client presents the cipher suites it supports
 - The server selects one
- Session key management
 - Lifetime of the master secret = lifetime of the TLS session
 - · Used whenever the client and the server decide to communicate
 - Maximum of 24 hours recommended
 - Lifetime of the session keys: at most lifetime of the TCP connection

TLS Cipher Suites

Cipher suite

- Public-key algorithm
- Symmetric encryption algo.
- MAC algorithm
- TLS supports several

Negotiation:

- Client offers choice
- Server picks 1
 - e.g., the most secure

Common algorithms

- Public-key encryption
 - RSA, ECDSA
- Symmetric ciphers
 - AES: block
 - ChaCha20: stream
- Hash functions
 - SHA-2
 - SHA-3

TLS protocols

Handshake Protocol

- Exchange of identity and supported cipher suites
- Authentication of the communicating parties
 - Client challenges the server, that proves its identity with certificate(s) [Optional, but mandatory if the next exists]
 - Server challenges the client, that proves its identity with certificates [Optional]
- Key distribution

Record Protocol

- Creation and verification of secure messages
 - Compression, cipher, integrity control



TLS handshake (1)

Purpose

- 1. Server authentication
- 2. Negotiation: agree on crypto algorithms
- 3. Establish keys
- 4. Client authentication (optional)

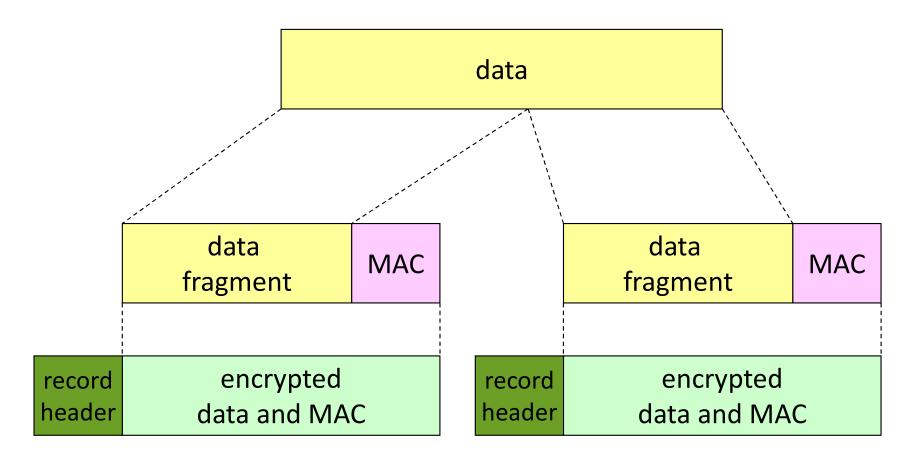
TLS handshake (2)

- Client sends list of algorithms it supports, along with client nonce
- 2. Server chooses algorithms from list; sends back: choice + certificate + server nonce
- 3. Client verifies certificate, extracts server's public key, generates pre_master_secret, encrypts with server's public key, sends to server
- 4. Client and server independently compute encryption and MAC keys from pre_master_secret and nonces
- 5. Client sends a MAC of all the handshake messages
- 6. Server sends a MAC of all the handshake messages

TLS handshake (3)

- Last two steps protect handshake from tampering:
 - Client typically offers range of algorithms, some strong, some weak
 - Man-in-the-middle could delete best algorithms from list
 - Last two message MACs allows detecting such an attack
- Why not simply sign or add MACs to the handshake messages?
 - Not possible: it is the handshake that sets up the keys!

TLS record protocol



Record header: content type; TLS version; length

MAC: is function of the sequence number and the MAC key M_x

Fragment: each data fragment has up to 2¹⁴ bytes (~16 Kbytes)

TLS record format

1 byte 2 bytes 2 bytes content length **TLS version** type data (encrypted) MAC (encrypted)

TLS performance

- First implementations were slow
- Currently, when properly configured, can be fast
- "On our production frontend machines,
 TLS accounts for
 less than 1% of the CPU load,
 less than 10 KB of memory per connection and
 less than 2% of network overhead."
 - Adam Langley, Google "Overclocking SSL"
 - https://istlsfastyet.com/

TLS tool: SSLTest

Home

Projects



You are here: Home > Projects > SSL Server Test > www.tecnico.ulisboa.pt

SSL Report: www.tecnico.ulisboa.pt

Assessed on: Wed, 23 May 2018 16:02:59 UTC | Hide | Clear cache

Scan Another >>

Contact

Qualys.com

	Server	Test time	Grade
	2001:690:2100:1:0:0:c03a:216e Ready	Wed, 23 May 2018 15:57:43 UTC Duration: 154.709 sec	В
2	193.136.128.169 Ready	Wed, 23 May 2018 16:00:18 UTC Duration: 161.248 sec	В

SSL Report v1.31.0

Copyright © 2009-2018 Qualys, Inc. All Rights Reserved.

Terms and Conditions

Qualys is the leading provider of integrated infrastructure security, cloud infrastructure security, endpoint security, devsecops, compliance and web app security solutions.

https://www.ssllabs.com/ssltest/

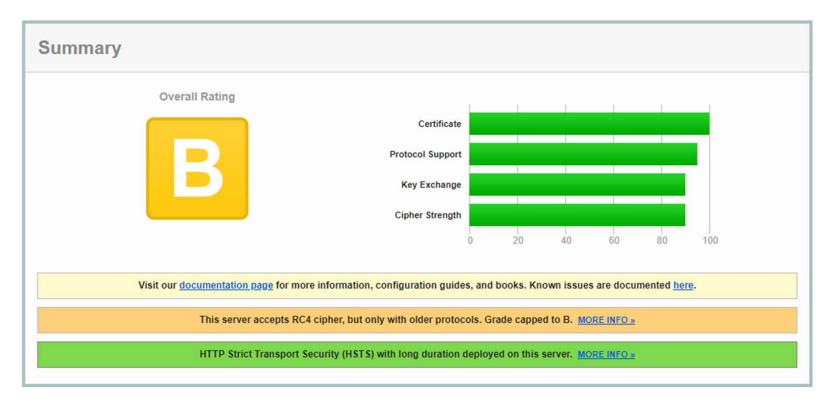
SSLTest output for Técnico



Home Projects Qualys.com Contact

You are here: Home > Projects > SSL Server Test > www.tecnico.ulisboa.pt > 2001:690:2100:1:0:0:c03a:216e

SSL Report: <u>www.tecnico.ulisboa.pt</u> (2001:690:2100:1:0:0:c03a:216e)



https://www.ssllabs.com/ssltest/analyze.html? d=tecnico.ulisboa.pt&s=193.136.128.169

Roadmap

- TLS Transport Layer Security
- SSH Secure Shell
- Key management

SSH

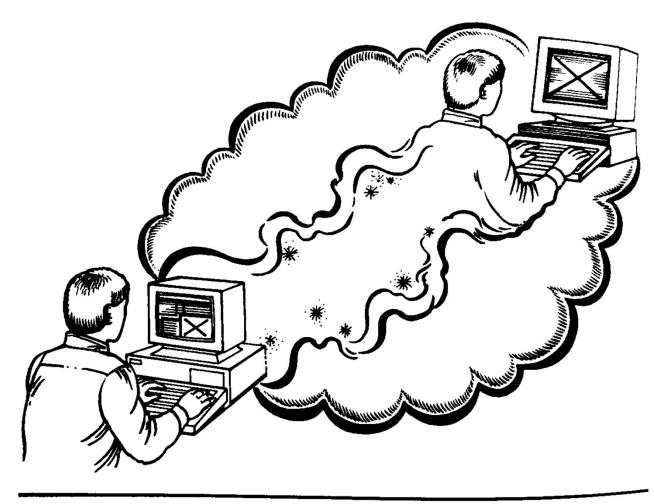


Figure 7.1 Remote login is a lot like astral projection.

SSH goals

- Secure communication application and protocol over TCP
 - Allows secure remote sessions ("telnet) and file transfer ("ftp)
 - Allows tunneling of TCP/IP traffic
- Security mechanisms
 - Confidentiality and integrity of the communication
 - Distribution of keys
 - Communicating parties' authentication
 - The server (or, usually, the server machine)
 - The user

SSH protocols (1/2)

- Transport Layer Protocol
 - Server authentication
 - Signature of the exchanged DH ephemeral values
 - Server public key can be transferred at this time, if not already held by the client
 - → TOFU (Trust on First Use) model
 - Vulnerable to man-in-the-middle
 - No certificates or CA
 - Distribution of keys
 - Diffie-Hellman key exchange
 - The session keys are computed with ephemeral DH values
 - Creation and analysis of secure messages
 - Compression, encryption, integrity control

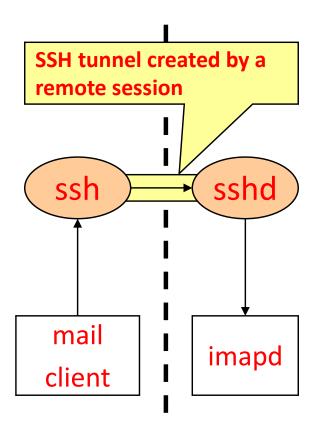
SSH protocols (2/2)

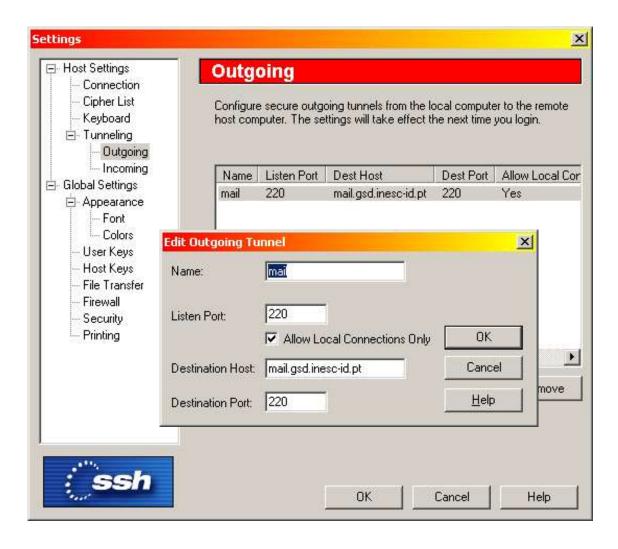
- User authentication towards the remote machine:
 - Password or
 - Signature with private key of client
 - Server knows the public key of client
- Connection Protocol
 - Information/data flow multiplexing over a secure session

Tunneling with SSH

- In the client machine, a mapping is created between a local TCP port and a port in the remote machine
 - e.g., localhost:IMAP → mail.myorg.pt:IMAP
- SSH session / tunnel is created to mail.myorg.pt
- Client application is configured to use the local port
 - When using the port, it will securely interact via SSH with the remote server mail.myorg.pt
 - Client SSH and server SSHd operate as a secure relay mechanism

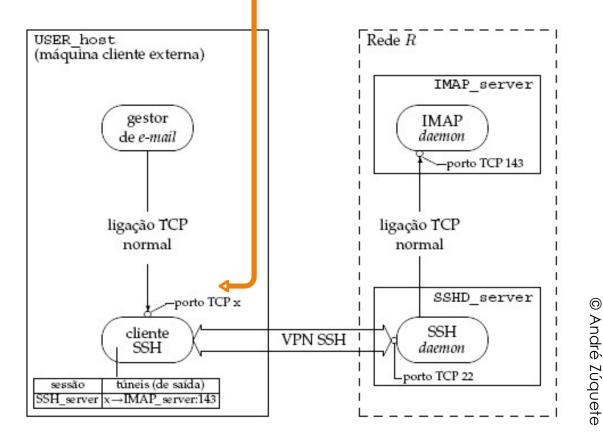
Configuration example of SSH tunnel for IMAP (port 220)





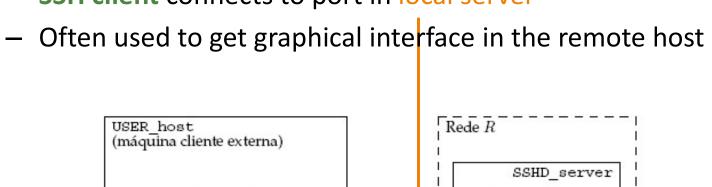
Output tunnel

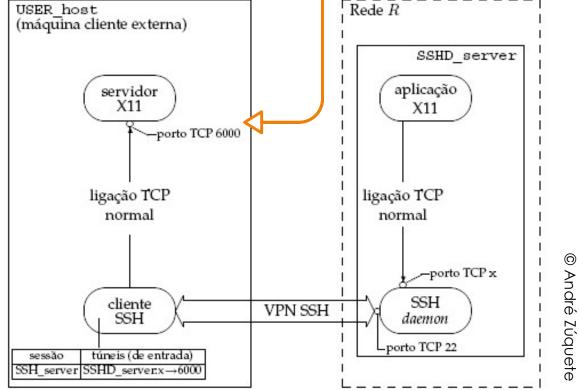
- TCP connections started at the client, so outbound
 - SSH client opens port X and waits for connections from client application
 - Similar situation to remote login



Input tunnel

- TCP connection started at the server, so inbound
 - SSH client connects to port in local server





Roadmap

- TLS Transport Layer Security
- SSH Secure Shell
- Key management

Keys

- TLS and SSH use public-key cryptography
 - (We will study this cryptography later, in more detail)
- Pair of keys
 - One is private, personal, non-transmissible
 - One is public, can/should be widely known
- Allow for
 - Confidentiality with the exchange of secret keys
 - Authentication and Integrity with digital signatures

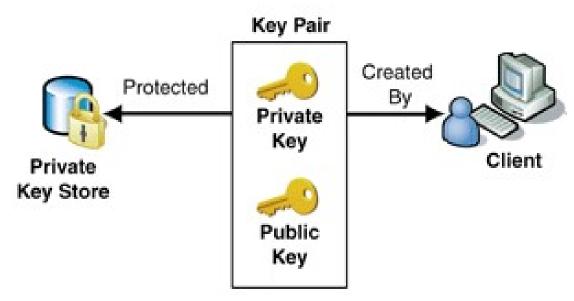
Private key

- The private key represents its owner so:
 - The probability of it being compromised must be minimized
 - Backup copies must be physically secure
- Private key must be protected
 - The access path to the private key must be restricted
 - Password protected, e.g., JKS, PGP
 - Security of applications using the private key must be guaranteed

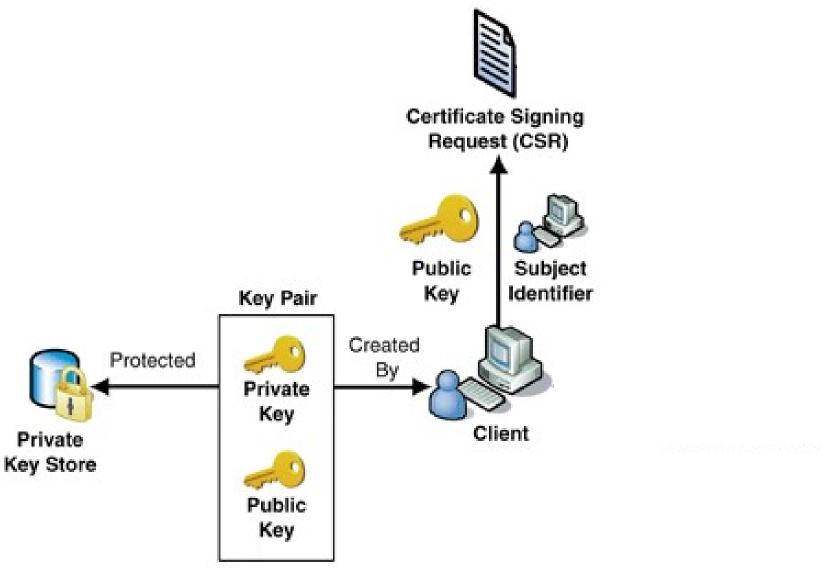
Public key certificate

- Certificates are documents signed by a certification entity
 - Certification Authority (CA), public organization or company
 - Certificates are public documents
 - Certificates have a digital signature to assure authenticity
- Can distribute public key through unsecure channel
 - Receiver can validate the certificate signature using the CA public key
 - If it trusts the CA and the signature is valid,
 then it can trust the public key
- Certificate standard format: X.509 (RFC 3280)

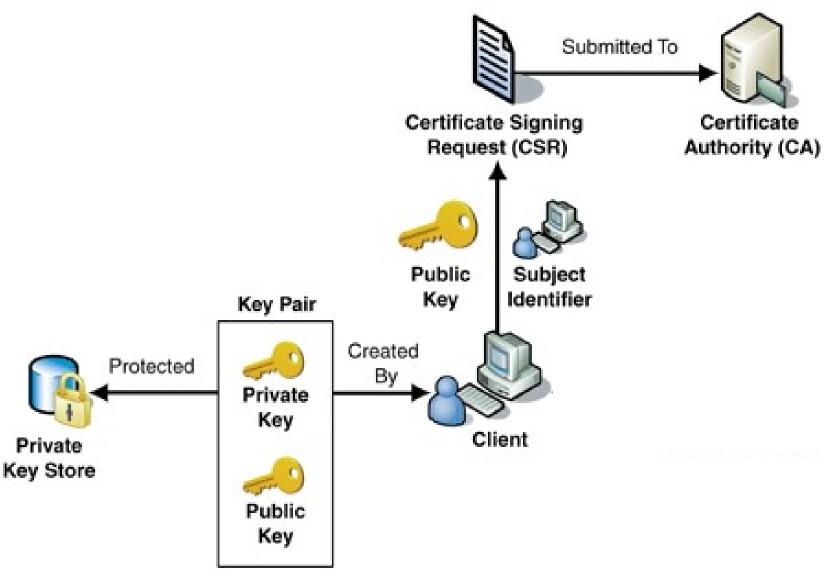
Certificate signing steps 1/4



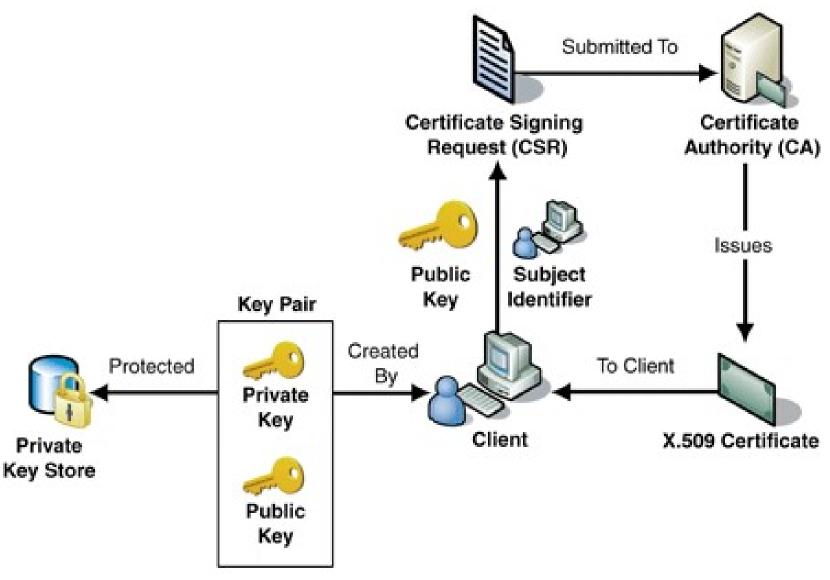
Certificate signing steps 2/4



Certificate signing steps 3/4



Certificate signing steps 4/4



Summary

- TLS Transport Layer Security
- SSH Secure Shell
- Key management