





# <u>Desenvolvimento de</u> <u>Aplicações Distribuídas</u>

Publish-Subscribe Systems





# <u>Agenda</u>

- Introduction
- Elements in Pub-Sub
- Subcription Models
- Architecture Model
  - transport, infrastructure, subscription matching, routing
- Distributed Event Routing
  - flooding, selective (filtering/rendez-vous), gossip,...
- Example Systems
  - Scribe, Siena, ...





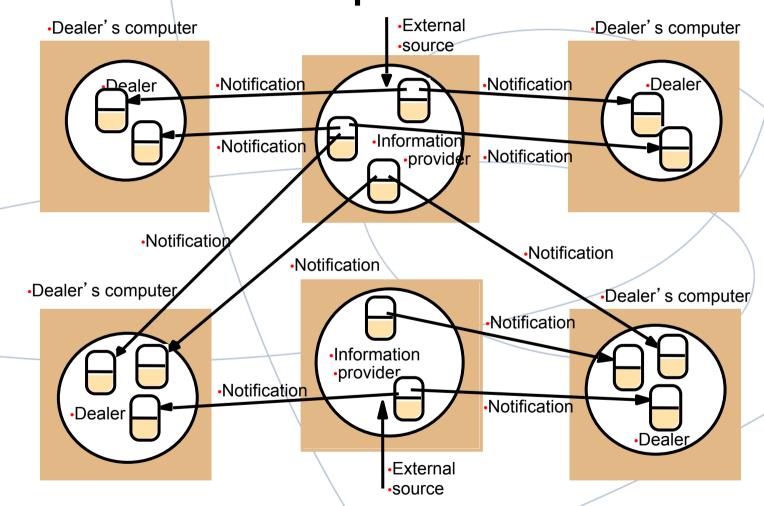
#### Publish-Subscribe Paradigm

- intuitive, application-driven
  - \* stock tickers, news feed, air traffic control, location events, application and Internet monitoring.
- full decoupling
  - \* senders (publishers) and receivers (subscribers)
  - \* no direct knowledge, possibly anonymous
  - \* no synchronous interaction
- indirect communication
  - via events, no explicit addressing, routing according to subscriptions
- scalability
  - large-scale networks, large numbers of publishers, subscribers, subscriptions and events





#### Publish-subscribe Example





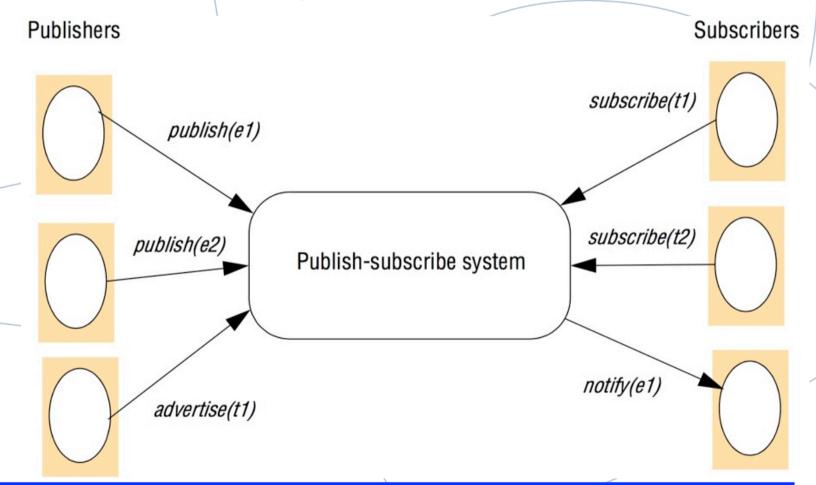


- Elements of a Pub-Sub System
  - clients: publishers, subscribers
  - events: set of (named) attribute-value pairs
  - **subscription** (f): filter, set of constraints
  - basic operations on events
    - publish(e),
    - subscribe(f), unsubscribe(f) (also σ for subscription)
    - advertise(f), unadvertise(f)
  - (matching): test event against subscription filter
  - notification: event delivery
    - notify(e)





#### Elements of a Pub-Sub System







- Channel-based
- Topic-based
- Content-based
- Type-based
- Context/Location-aware
- Concept-based
- ...and complex subscriptions on several events/conditions





#### Channel-based

- primitive scheme
- subscription to given communication channel
  - \* e.g., CORBA event service
- all events sent (published) to channel are received (notified)
- advantages: easy implementation
- drawbacks: inflexible, no filtering, scalability issues





#### Topic-based

- events group around topics
- topic as a logical channel
- coarse-grained
  - \* subscription implies all events with specific topic
- e.g., Scribe, Bayeux, iBus...
- advantages: leverage multicast, diffusion trees
- drawbacks: limited expressiveness of the subscription language
- solution: hierarchical topic space
  - \* topic subscription includes all of its sub-topics





#### Content-based

- conditions (predicates) over event content
- query filter
  - \* conjunction, disjunction, equality, comparison operators
  - regular expressions and more complex matching
- fine-grained
  - correspondence between publisher and subscribers event by event basis
- e.g., Gryphon, SIENA, JEDI,...
- advantages: higher expressive power
- drawbacks: higher resource consumption
  - \* calculating set of interested subscribers





#### Type-based

- typed events, type hierarchy/graph
  - events belong to a specific type
  - encapsulated attributes and methods
- middle-ground approach
  - coarse-grained as in topic-based
    - a type subscription is a channel
  - fine-grained as in content-based
    - constraints over attributes or methods

#### advantages:

- \* type-safety at the pub-sub system not the application,
- \* robustness





#### Context/Location-Awareness

- suited to mobile environments
- subscriptions based on/filtered by
  - user locations
  - vicinity of service
  - \* usually, close-by event publisher

#### drawbacks:

requires monitoring on user mobility



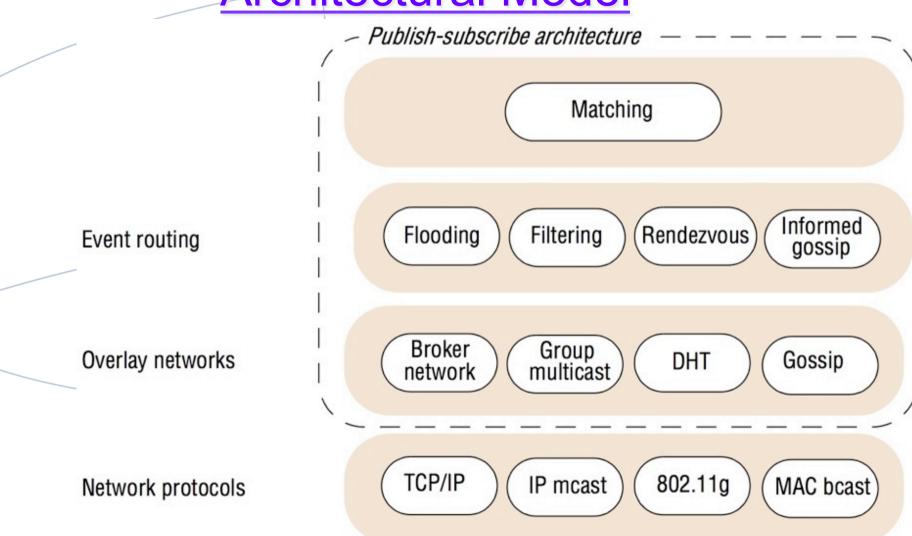


#### Concept-based

- release from syntatical restrictions
  - \* number, name, type of attributes
- release from semantical assumptions
  - (implied) meaning of each attribute
- event schema using ontologies
  - Metadata, mapping functions, XML
- advantages:
  - interoperability, extensibility, implementation-independence
- drawbacks:
  - \* heavier XML-based subscription processing











#### Publish-Subscribe Functional Layers

- system
  - \* set of distributed nodes coordinated to dispatch events
  - \* to all (preferrably only) interested subscribers
- overlay infrastruture
  - \* node organization, (de-)centralization aspects, dedicated servers,
  - groups in GCS can represent channels or topics
- event routing
  - \* dispatching of events, exploit infrastructure, enhance with routing information to achieve scalbility
- matching
  - testing events against subscriptions





#### Overlay Infrastructure: Broker Overlay

- set of independent servers (brokers)
- logically connected, not permanent links
- each broker knows (some of) the others, (mostly) inherently static
  - topology changes rare
  - addition of new broker, failure handling
- \* clients access system via any broker

#### Hierarchical:

- brokers as trees, subscribers at leaf nodes, publishers at rootnodes (or vice-versa)
- Flat: broker connected to any other brokers, less load on highernodes
- ♣ E.g., TIB/RV, Gryphon, SIENA, JEDI, REDS





# Overlay Infrastructure:Peer-to-Peer Structured Overlay (DHT)

- ★ structured graph over virtual key space (key→node)
- \* efficient discovery of data, granted correspondence
  - between any address (key) and an active node

#### advantages:

- \* allows efficient unicast and multicast
- better handling of dynamic aspects (joins, faults)
- E.g., usage of P2P structured overlays
  - Pastry, Chord, Tapestry, CAN,
- ♣ E.g., systems
  - Scribe, Bayeux (topic-based),
  - \* Hermes, Rebecca, Meghdoot (content-based)





- Overlay Infrastructure: Peer-to-Peer UnStructured Overlay (Gossip)
  - \* absence of structure or direct mapping of keys
  - use flooding, gossiping, random walks for event difusion and information (subscriptions) retrieving
  - probabilistic in nature
  - fewer examples (mainly gossip-based)





#### Matching

- checking events against subscriptions
- performed on massive data sizes
- routing algorithms need matching phase
  - \* distributed over nodes/brokers
  - \* avoid propagating all events or all subscriptions





#### Event Routing

- flooding algorithms: event, subscription
- selective algorithms: rendez-vous, filter-based
  - reduce message and/or memory overhead
- event gossiping: basic, informed gossip
- issues:
  - message overhead, memory overhead, subscription language limitations, subscription changes, joins, leaves, failures (churn)

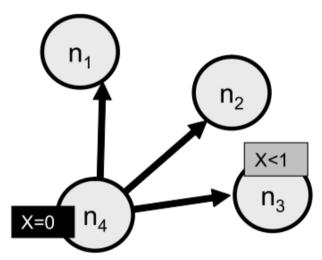
		Routing	Filtering	Nodes storing Subs	Nodes handling Events
Flooding	Event flooding	Det.	Subscribers	None	All
	Subs Flooding	Det.	Publishers	All	None
Selective	Filter-Based	Det.	Intermediaries	Subset	Subset
	Rendezvous-based	Det.	Intermediaries	Subset	Subset
Gossiping	Basic gossiping	Prob.	Subscribers	None	All 6.2
	Informed gossiping	Prob.	Intermediaries	Subset	Subset





#### Event Flooding

- each <u>event</u> fully propagated
  - from publisher to all nodes in the system
- subcriptions stored locally at subscribers
- broadcast may be used
- no message scalability
- minimal memory overhead
- no language limitations



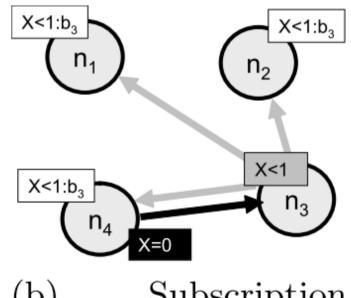
(a) Event Flooding





#### Subscription Flooding

- each <u>subscription</u> fully propagated
  - sent to all (publisher) nodes in the system
- each node has complete knowledge of entire system
- single-hop, optimal filtering
- message overhead
  - subscription flood
- memory overhead
  - all subscriptions
- high coupling regarding
  - overlay membership



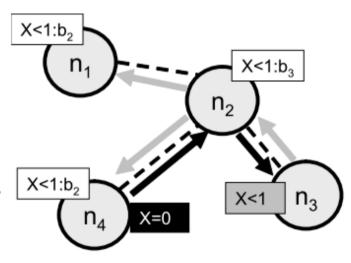
(b) Subscription Flooding 6.22





#### Filtering-based

- events forwarded only along
  - paths leading to interested subscribers
- eliminate as soon as possible
  - events without subscribers
  - stop forwarding
- diffusion paths
  - routing information at nodes
  - neighbour brokers and
    - reachable subscriptions
- (c) Filtering-based
- reverse paths followed by events published







#### Filtering-based

- reduced mesage overhead
- optimizations: subscription containment
  - brokers modify routing
    - only when new subscriptions change the...
    - ...overall condition to be tested
- natural architecture: brokers, tree or acyclic graph
  - more neighbors,
    - smaller network diameter
    - memory overhead increases
- no subscription language limitations





#### Filtering-based

- data structures:
  - neighbor list, routing table, subscription list
- routing table maps neighbors with subscription sets

```
matchlist \leftarrow match(e,subscriptions);

send notify(e) to <math>matchlist;

fwdlist \leftarrow match(e,routing);

send publish(e) to <math>fwdlist - x;

upon receive subscribe(subscription s) from node x

if x is client then

add \ s to subscriptions;

else add \ (x,s) to routing

send \ s to neighbors - x;
```

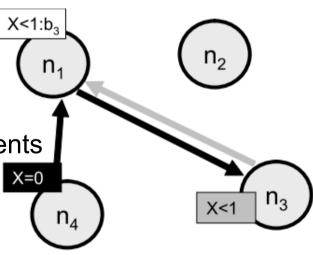
**upon receive** publish(event e) from node x





#### Rendez-vous based

- mapping functions: SN, EN
- $\blacksquare$  SN( $\sigma$ ) → set of rendez-voud nodes for subscription  $\sigma$ 
  - store σ and forward matching events matching to all subscribers
- EN(e) → set of rendez-voud nodes for event
  - match event against subscriptions
- two phases:
  - publisher sends events to EN(e)
  - EN(e) match e against subscriptions
  - forward to subscribers
- $\blacksquare$  EN(e) and SN( $\sigma$ ) must have common elements
  - mapping intersection rule
  - not trivial with content-based
  - clustering of subscription space



(d) Rendezvous- $_{6.26}$ 

based





#### Rendez-vous based

- not suited to broker overlay
  - dynamism not efficiently handled
  - new node joins in, node leaving, crashing: similar problems
    - implies whole repartitioning among nodes
    - (many) subscriptions moved among nodes
- unstructured P2P networks also not easily suited
- highly-efficient on P2P structured overlays
  - mapping functions, fixed-size address space used as target
  - no need to know globally individual node in charge of each key
- limitations: subscription language
  - mostly topic-based,
  - content-based: attribute mapping, numerical-only, limited
- memory overhead, balanced subscriptions over keys





#### Rendez-vous based

- routing over neightbour lists
- directed towards rendez-vous nodes

```
upon receive publish(event e) from node x at node i

rvlist \leftarrow EN(e);

if i \in rvlist then

matchlist \leftarrow match(e,subscriptions);

send notify(e) to matchlist;

else

send(e) to rvlist;

upon receive subscribe(subscription s) from node x at node i

rvlist \leftarrow SN(s);

if i \in rvlist then

add s to subscriptions;

else

send(s) to rvlist;
```





#### Gossip-based

- random choice of nodes
- driven by local information
  - acquired during execution
  - fine-tuning routing tables
- aggregation of subscriptions
- organized into groups (super-peers)
  - delegate actions on other group leaders
  - \* outside vicinity
- fully distributed, resilient to churn
- fully probabilistic





#### Guaranteed Delivery

- non-zero delay between
  - \* subscription changes
  - event routing data capturing changes
- problems
  - \* event loss
  - unnecessary event forwarding

#### solutions

- resend lost events
- \* keep separate streams for failed subscribers
- \* store events with persistence







System (and further reading)	Subscription model	Distribution model	Event routing
CORBA Event Service (Chapter 8)	Channel-based	Centralized	-
TIB Rendezvouz [Oki et al. 1993]	Topic-based	Distributed	Ffiltering
Scribe [Castro et al. 2002b]	Topic-based	Peer-to-peer (DHT)	Rendezvous
TERA [Baldoni et al. 2007]	Topic-based	Peer-to-peer	Informed gossip
Siena [Carzaniga et al. 2001]	Content-based	Distributed	Filtering
Gryphon [www.research.ibm.com]	Content-based	Distributed	Filtering
Hermes [Pietzuch and Bacon 2002]	Topic- and content-based	Distributed	Rendezvous and filtering
MEDYM [Cao and Singh 2005]	Content-based	Distributed	Flooding
Meghdoot [Gupta et al. 2004]	Content-based	Peer-to-peer	Rendezvous
Structure-less CBR [Baldoni et al. 2005]	Content-based	Peer-to-peer	Informed gossip





#### Scribe

- topic-based routing over structured overlay (Pastry)
  - \* each broker has unique ID with efficient routing to it
- event routing: rendez-vous over structured overlay
  - each topic assigned unique ID (hashing)
  - broker of topic is the node with ID closest to topic's
  - \* multicast tree for each topic (broker as root)
  - \* EN(e) = h(e), SN(s) = h(S), i.e., EN=SN for a given topic
- subscriptions: routed to broker
  - updating tree structure to include new subscriber
  - \* higher nodes in tree seldom changed
- publishing event routed to broker directly
  - \* broker starts diffusion of notification through the topic tree
  - until all leaf nodes (subscribers) are reached

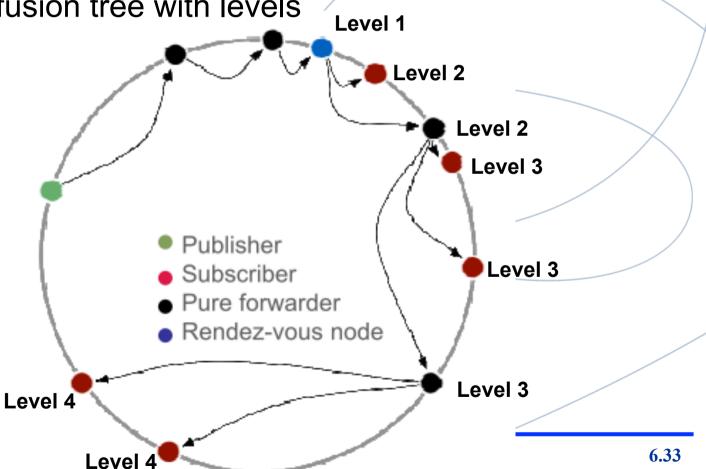


#### inesc id lisboa

# Pub-Sub Systems

#### Scribe

- topic rendez-vous node
- event diffusion tree with levels





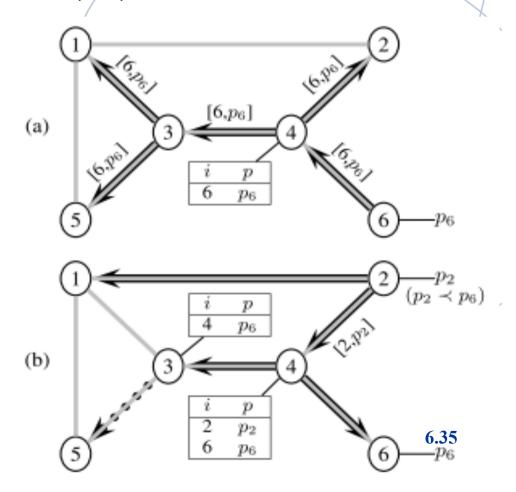


- uses broker network for routing over TCP or UDP layer
- filtering-based content-based routing (influential work)
  - for subcription propagation and event delivery
- optimizations:
  - \* subscription covering/containment
    - s1 contains s2 iff all events matching s2 also match s1
    - subscription updates only forwarded
      - if broker has not already propagated a containing subscription
  - \* subscription refreshing protocol triggered by brokers
  - \* advertisements (publisher advertisements)
    - publishers declare set of events they are going to produce
    - considered in routing paths, futher reduce set of involved brokers





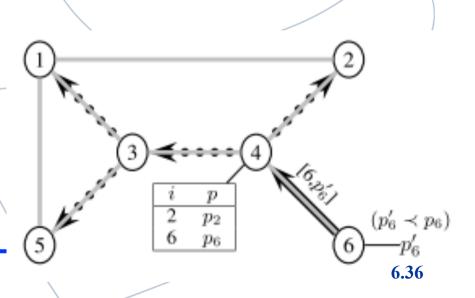
- subscription propagation
  - \* Receiver Advertisement Protocol (RA)
- subscription containment
  - (subscription covering)
  - p6 covers p2
  - \* no need to propagate \$2
    - over to b3 and beyond







- subscription/adress inflation
- when updated subscriptions by one client
  - are covered by its early subscriptions
  - no forwarding takes place
- in the long run
  - increasing number of events
  - wastefully forwarded
  - via broker routes directing to
  - no-longer subscribing nodes







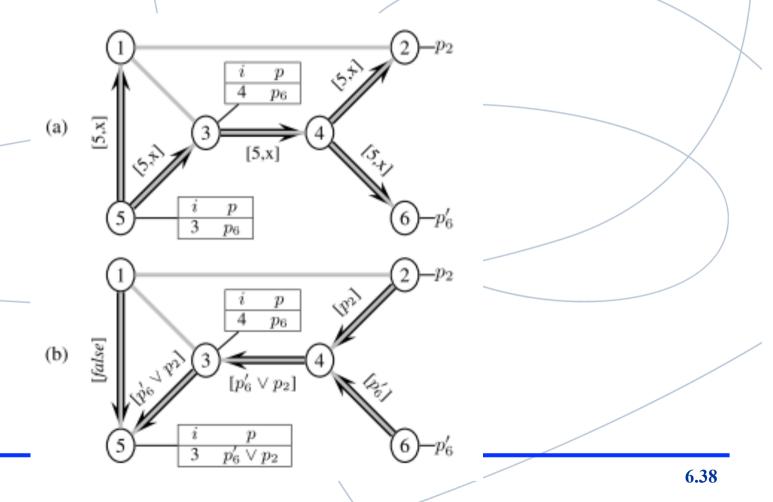
- Sender Request / Update Replies (SR/UR) protocol
  - periodical subscription update
  - triggered by broker nodes to pull subscription information
  - tackles address inflation
- receiving brokers:
  - leaf-broker:
    - replies immediately with subscriptions
  - non-leaf:
    - assembles UR
    - combining its subscriptions with
    - those received from downstream
    - waits until all replies or timeout and replies upstream





#### SIENA

Sender Request / Update Replies (SR/UR) protocol







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- Sender Request / Update Replies (SR/UR) protocol
- Opportunistic Update Reply Processing
  - allows caching of UR replies to previous SR
  - brokers may reuse URs to SRs sent by other brokers iff
    - set of downstream nodes from caching broker...
    - equal to downstream nodes for initial requester...
    - at the same link
  - e.g., b4 caches UR from b6
  - reuses as reply to SR from b2
    - nodes from b5 to b6 via b4
    - same nodes as
    - from b2 via b4 (i.e., b6)

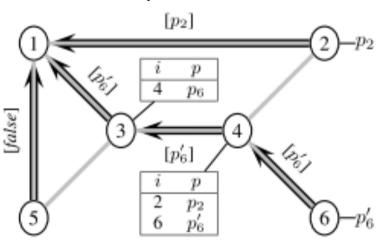


Fig. 9. Opportunistic UR Processing