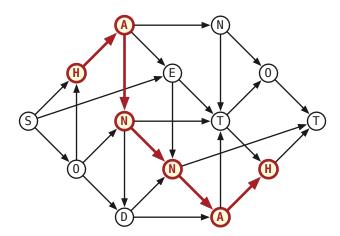
1. Suppose we are given a directed acyclic graph *G* with labeled vertices. Every path in *G* has a label, which is a string obtained by concatenating the labels of its vertices in order. Recall that a *palindrome* is a string that is equal to its reversal.

Describe and analyze an algorithm to find the length of the longest palindrome that is the label of a path in *G*. For example, given the graph below, your algorithm should return the integer 6, which is the length of the palindrome **HANNAH**.



2. Let *G* be a connected directed graph that contains both directions of every edge; that is, if $u \rightarrow v$ is an edge in *G*, its reversal $v \rightarrow u$ is also an edge in *G*. Consider the following non-standard traversal algorithm.

SPAGHETTITRAVERSAL(G):
for all vertices v in Gunmark vfor all edges $u \rightarrow v$ in Gcolor $u \rightarrow v$ white $s \leftarrow$ any vertex in GSPAGHETTI(s)

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\frac{\text{SPAGHETTI}(v)}{\text{mark } v} : \qquad & \langle \langle \text{"visit } v \text{"} \rangle \rangle
if there is a white arc v \rightarrow w
if w is unmarked
\text{color } w \rightarrow v \text{ green}
\text{color } v \rightarrow w \text{ red} \qquad & \langle \langle \text{"traverse } v \rightarrow w \text{"} \rangle \rangle
\text{SPAGHETTI}(w)
else if there is a green arc v \rightarrow w
\text{color } v \rightarrow w \text{ red} \qquad & \langle \langle \text{"traverse } v \rightarrow w \text{"} \rangle \rangle
\text{SPAGHETTI}(w)
\langle \langle \text{else every arc } v \rightarrow w \text{ is red, so halt} \rangle \rangle
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We informally say that this algorithm "visits" vertex ν every time it marks ν , and it "traverses" edge $\nu \rightarrow w$ when it colors that edge red. Unlike our standard graph-traversal algorithms, Spaghetti may (in fact, *will*) mark/visit each vertex more than once.

The following series of exercises leads to a proof that Spaghetti traverses each directed edge of *G* exactly once. Most of the solutions are very short.

- (a) Prove that no directed edge in *G* is traversed more than once.
- (b) When the algorithm visits a vertex v for the kth time, exactly how many edges into v are red, and exactly how many edges out of v are red? [Hint: Consider the starting vertex s separately from the other vertices.]

- (c) Prove each vertex ν is visited at most $\deg(\nu)$ times, except the starting vertex s, which is visited at most $\deg(s)+1$ times. This claim immediately implies that SpaghettiTraversal(G) terminates.
- (d) Prove that when SpaghettiTraversal(G) ends, the last visited vertex is the starting vertex s.
- (e) For every vertex v that SpaghettiTraversal(G) visits, prove that all edges incident to v (either in or out) are red when SpaghettiTraversal(G) halts. [Hint: Consider the vertices in the order that they are marked for the first time, starting with s, and prove the claim by induction.]
- (f) Prove that SPAGHETTITRAVERSAL(G) visits every vertex of G.
- (g) Finally, prove that SpaghettiTraversal(G) traverses every edge of G exactly once.