

Street for pedestrians

Sina Hazratpour*
sinahazratpour@gmail.com

November 29, 2017

Abstract

The singular aim of these notes is to give a less condensed and more clear explanation of Ross Street's significant paper (Street 1974). We highly encourage reader to also refer to (Kock 1995) and (Lack 2000).

1 Pseudo algebras for strict 2-monads

DEFINITION 1.1. Let \mathcal{K} be a 2-category and $(T: \mathcal{K} \rightarrow \mathcal{K}, i: 1 \Rightarrow T, m: T^2 \Rightarrow T)$ a strict 2-monad on \mathcal{K} . A *pseudo-algebra* of T consists of

- i. a 0-cell A in \mathcal{K}
- ii. a 1-cell $a: TA \rightarrow A$
- iii. invertible 2-cells $\zeta: 1_A \Rightarrow a \circ i_A$ and $\theta: a \circ Ta \Rightarrow a \circ m_A$

$$\begin{array}{ccc} A & & T^2 A \xrightarrow{Ta} TA \\ i_A \downarrow & \searrow 1 & m_A \downarrow \quad \theta \Downarrow \quad a \downarrow \\ TA & \xrightarrow{a} A & TA \xrightarrow{a} A \end{array} \quad (1)$$

subject to the following coherence axioms:

$$(\theta \cdot m_{TA}) \circ (\theta \cdot T^2 a) = (\theta \cdot Tm_A) \circ (a \cdot T\theta)$$

*School of Computer Science, University of Birmingham, Birmingham, UK.

expressed by equality of pasting diagrams:

$$\begin{array}{c}
 \begin{array}{ccccc}
 T^3 A & \xrightarrow{T^2 a} & T^2 A & \xrightarrow{T a} & T A \\
 m_{TA} \downarrow & & \downarrow m_A & & \downarrow a \\
 T^2 A & \xrightarrow{T a} & T A & \xleftarrow{\theta} & T A \\
 m_A \downarrow & \theta \Downarrow & \downarrow a & & \downarrow a \\
 T A & \xrightarrow{a} & A & \xleftarrow{a} & A
 \end{array} \\
 \end{array} = \begin{array}{c}
 \begin{array}{ccccc}
 T^3 A & \xrightarrow{T^2 a} & T^2 A & \xrightarrow{T a} & T A \\
 m_{TA} \downarrow & T m_A \downarrow & T \theta \Downarrow & \downarrow T a & \downarrow a \\
 T^2 A & \xrightarrow{T a} & T A & \xleftarrow{\theta} & T A \\
 m_A \downarrow & m_A \downarrow & \theta \Downarrow & \downarrow a & \downarrow a \\
 T A & \xrightarrow{a} & A & \xleftarrow{a} & A
 \end{array}
 \end{array} \quad (2)$$

and

$$(\theta \cdot T i_A) \circ (a \cdot T \zeta) = id_a = (\theta \cdot i_{TA}) \circ (\zeta \cdot a)$$

expressed by equality of pasting diagrams:

$$\begin{array}{c}
 \begin{array}{ccccc}
 T A & \xrightarrow{1_{TA}} & T A & & \\
 T i_A \downarrow & T \zeta \Downarrow & \downarrow 1_{TA} & & \\
 T^2 A & \xrightarrow{T a} & T A & & \\
 m_A \downarrow & \theta \Downarrow & \downarrow a & & \\
 T A & \xrightarrow{a} & A & &
 \end{array} \\
 \end{array} = \begin{array}{c}
 \begin{array}{ccc}
 T A & \xrightarrow{a} & A \\
 1_{TA} \downarrow & & \downarrow 1_A \\
 T A & \xrightarrow{a} & A
 \end{array}
 \end{array} = \begin{array}{c}
 \begin{array}{ccccc}
 T A & \xrightarrow{a} & A & & \\
 i_{TA} \downarrow & & \downarrow i_A & & \\
 T^2 A & \xrightarrow{T a} & T A & \xleftarrow{\zeta} & T A \\
 m_A \downarrow & \theta \Downarrow & \downarrow a & & \downarrow a \\
 T A & \xrightarrow{a} & A & &
 \end{array}
 \end{array} \quad (3)$$

DEFINITION 1.2. Suppose $(a, \zeta_A, \theta_A) : TA \rightarrow A$ and $(b, \zeta_B, \theta_B) : TB \rightarrow B$ are pseudo-algebras of a 2-monad T . A **lax morphism** from a to b consists of a 1-cell $f : A \rightarrow B$ and a 2-cell \check{f}

$$\begin{array}{ccc}
 T A & \xrightarrow{T f} & T B \\
 a \downarrow & \check{f} \Downarrow & \downarrow b \\
 A & \xrightarrow{f} & B
 \end{array}$$

in such a way that

- $f \cdot \zeta_A = (\check{f} \cdot i_A) \circ (\zeta_B \cdot f)$ and
- $(f \cdot \theta_A) \circ (\check{f} \cdot T a) \circ (b \cdot T \check{f}) = (\check{f} \cdot m_A) \circ (\theta_B \cdot T^2 f)$

DEFINITION 1.3. A 2-monad $T : \mathcal{K} \rightarrow \mathcal{K}$ is said to be **lax idempotent** if given any two (pseudo) T -algebras $a : TA \rightarrow A$, $b : TB \rightarrow B$ and a 1-cell $f : A \rightarrow B$, there exists a unique 2-cell $\check{f} : b \circ T f \Rightarrow f \circ a$ rendering (f, \check{f}) a lax morphism of pseudo T -algebras.

$$\begin{array}{ccc}
 T A & \xrightarrow{T f} & T B \\
 a \downarrow & \check{f} \Downarrow & \downarrow b \\
 A & \xrightarrow{f} & B
 \end{array}$$

REMARK 1.4. Dually, reverse the direction of \check{f} in definition 1.3, then we get the notion of *co-lax idempotent* monad.

2 KZ-monads

DEFINITION 2.1. A 2-monad $T : \mathcal{K} \rightarrow \mathcal{K}$ is said to be **KZ-monad**¹ if $m \dashv i \cdot T$ in the 2-category $[\mathcal{K}, \mathcal{K}]$ with identity counit.

REMARK 2.2. Dual to the definition above, we define a monad T to be a **co-KZ-monad** by requiring $i \cdot T \dashv m$ with identity unit.

Suppose T is a co-KZ-monad and $i \cdot T \dashv m$. In particular unit of this adjunction is identity since $m \circ (i \cdot T) = 1$. Moreover, the identity 2-cell

$$\begin{array}{ccc} T & \xrightarrow{1} & T \\ 1 \uparrow & id \Downarrow & \uparrow m \\ T & \xrightarrow{T.i} & T^2 \end{array}$$

has a mate

$$\begin{array}{ccc} T & \xrightarrow{1} & T \\ 1 \downarrow & \tau \Downarrow & \downarrow i.T \\ T & \xrightarrow{T.i} & T^2 \end{array} \quad (4)$$

with property that $m \cdot \tau = id_{1_T}$. Suppose $a : TA \rightarrow A$ is a pseudo algebra. We now would like to calculate the composite 2-cell

$$\begin{array}{ccccc} TA & \xrightarrow{i_{TA}} & T^2 A & \xrightarrow{a \circ Ta} & TA \\ \tau_A \Downarrow & & \Downarrow \theta & & \\ TA & \xrightarrow{Ti_A} & T^2 A & \xrightarrow{a \circ m_A} & TA \end{array}$$

In the diagram below, since $m_A \circ \tau = id$, the left column of 2-cells collapses to identity, and therefore we have

$$\begin{array}{ccccc} TA & \xrightarrow{1} & TA & \xrightarrow{a} & A \\ 1 \downarrow & \tau \Downarrow & \downarrow i_{TA} & \downarrow i_A & \\ TA & \xrightarrow{Ti_A} & T^2 A & \xrightarrow{Ta} & TA \\ 1 \downarrow & & \downarrow m_A & \theta \Downarrow & \downarrow a \\ TA & \xrightarrow{1} & TA & \xrightarrow{a} & A \end{array} \begin{array}{c} \curvearrowright 1_A \\ \zeta \end{array} = \begin{array}{ccc} TA & \xrightarrow{a} & A \\ 1_{TA} \downarrow & & \downarrow 1_A \\ TA & \xrightarrow{a} & A \end{array}$$

¹KZ: short for ‘Kock-Zöberlein’

On the other hand, we can compose row-wise instead, and we get

Thus, in the end, we have

LEMMA 2.3. Let T be a KZ-monad, and A an object of \mathcal{K} . Then any pseudo T -algebra is left adjoint to i_A . Conversely, if i_A has a left adjoint with invertible counit then this left adjoint is a pseudo T -algebra.

$$\begin{array}{ccc}
& \overset{i_A}{\curvearrowright} & \\
TA & \overset{\perp}{\rightleftarrows} & A \\
& \underset{a}{\curvearrowright} &
\end{array}$$
$$\begin{array}{ccccc}
& & A & & \\
& \xrightarrow{a} & & \xrightarrow{i_A} & \\
TA & \xrightarrow{i_{TA}} & T^2A & \xrightarrow{Ta} & TA \\
& \tau_A \Downarrow & & & \\
& \xrightarrow{Ti_A} & T\zeta^{-1}\Downarrow & & \\
& & 1 & &
\end{array}
\quad (6)$$

²The dual of this situation, i.e. unit in the case of KZ-monad, is calculated in page 112 of (Street 1974).

4 Fibrations as pseudo-algebras of a co-KZ-monad

Let \mathcal{K} be a representable 2-category. Define \mathcal{K}/B to be the strict slice 2-category over B , meaning the morphism triangles commute up to equality. (Street 1974) constructs KZ-monads $L, R: \mathcal{K}/B \rightrightarrows \mathcal{K}/B$. The idea is, for a morphism $p: E \rightarrow B$, an algebra $R(p) \rightarrow p$ (resp. $L(p) \rightarrow p$) if it exist, corresponds to the fibration structure on p (resp. opfibration structure)⁴ and thus, we will mainly concern ourselves with 2-monad R . However, when necessary, we will comment on the dual results for the case of opfibrations. We now define 2-monad R : It takes an object (E, p) to $(B/p, R(p))$ where

$$\begin{array}{ccc} B/p & \xrightarrow{\hat{d}_1} & E \\ R(p) \downarrow & \phi_p \uparrow & \downarrow p \\ B & \xrightarrow{1} & B \end{array} \quad (8)$$

is a comma square.

REMARK 4.1. 2-cell ϕ_p can be constructed as follows:

$$\begin{array}{ccc} B/p & \xrightarrow{\hat{d}_1} & E \\ R(p) \downarrow & \phi_p \uparrow & \downarrow p \\ B & \xrightarrow{1} & B \end{array} = \begin{array}{ccc} B/p & \xrightarrow{\hat{d}_1} & E \\ \hat{p} \downarrow \lrcorner & & \downarrow p \\ B & \xrightarrow{d_1} & B \\ d_0 \downarrow & \phi \uparrow & \downarrow 1 \\ B & \xrightarrow{1} & B \end{array}$$

The action of R on morphisms is given as follows:

If $f: (E', p') \rightarrow (E, p)$ is a 1-cell in \mathcal{K}/B , then define $R(f)$ to be the unique 1-cell with $\hat{d}_1 \circ R(f) = f \circ \hat{d}'_1$ and $\hat{p} \circ R(f) = p'$.

$$\begin{array}{ccc} B/p' & \xrightarrow{\hat{d}'_1} & E' \\ R(f) \downarrow \lrcorner & & \downarrow f \\ B/p & \xrightarrow{\hat{d}_1} & E \\ \hat{p} \downarrow \lrcorner & & \downarrow p \\ B & \xrightarrow{d_1} & B \end{array}$$

(Curved arrows labeled p' and p connect B/p' to B and E' to B respectively.)

⁴Unlike Street's paper whereby he works with opfibration structures and thus chooses to work with 2-monad L .

Similarly if $\sigma: f \Rightarrow g$ is a 2-cell in \mathcal{K}/B , then we have a unique induced 2-cell $R(\sigma): R(f) \Rightarrow R(g)$ with $\hat{d}_1 \cdot R(\sigma) = \sigma \cdot \hat{d}_1'$ and $\hat{p} \cdot R(\sigma) = id_{\hat{p}}$.

PROPOSITION 4.2. 2-functor $R: \mathcal{K}/B \rightarrow \mathcal{K}/B$ is a 2-monad.

The unit of monad $i: id \Rightarrow R$ at (E, p) is given by the unique arrow $i(p): E \rightarrow B/p$ with property that $R(p) \circ i(p) = p$ and $\hat{d}_1 \circ i(p) = 1_E$, and moreover $\phi_p \cdot i(p) = id_p$, all inferred by universal property of comma object B/p .

$$\begin{array}{ccccc}
 & & 1 & & \\
 & \curvearrowright & & \curvearrowleft & \\
 E & \xrightarrow{i(p)} & B/p & \xrightarrow{\hat{d}_1} & E \\
 & \searrow p & \downarrow R(p) & \uparrow \phi_p & \downarrow p \\
 & & B & \xrightarrow{1} & B
 \end{array}$$

It also follows that $\hat{d}_1 \dashv i(p)$ with identity counit. Indeed, $i(p)$ is v in proposition 3.1, when $f = 1$ and $g = p$. From there, we also get the unit β of adjunction with $R(p) \cdot \beta = \phi_p$.

The multiplication $m: R^2 \Rightarrow R$ of monad at (E, p) is given by the unique arrow $m(p): B/R(p) \rightarrow B/p$

$$\begin{array}{ccccc}
 B/R(p) & \xrightarrow{\widehat{d_1}} & B/p & \xrightarrow{\hat{d}_1} & E \\
 \downarrow \hat{p} & \lrcorner & \downarrow \hat{p} & \lrcorner & \downarrow p \\
 B \rightarrow & \xrightarrow{\widehat{d_1}} & B \rightarrow & \xrightarrow{\hat{d}_1} & B \\
 \downarrow d_0 & \lrcorner & \downarrow d_0 & \lrcorner & \downarrow 1 \\
 B \rightarrow & \xrightarrow{\widehat{d_1}} & B & \xrightarrow{1} & B \\
 \downarrow d_0 & \lrcorner & \downarrow 1 & & \\
 B & \xrightarrow{1} & B & &
 \end{array} \tag{9}$$

with property that $R(p) \circ m(p) = R^2(p)$ and $\hat{d}_1 \circ m(p) = \hat{d}_1 \circ \widehat{d_1}$, and moreover $\phi_p \cdot m(p) = (\phi_p \cdot \widehat{d_1}) \circ (\phi \cdot d_0 \rightarrow \hat{p}) = (\phi_p \cdot \widehat{d_1}) \circ \phi_{R(p)}$, all inferred by universal property of comma object B/p .

PROPOSITION 4.3. 2-monad $R: \mathcal{K}/B \rightarrow \mathcal{K}/B$ is a co-KZ-monad.

Proof. We have to show that $i \cdot T \dashv m$. □

Now, we would like to see what an algebra $a: R(p) \rightarrow p$ in \mathcal{K}/B looks like. The fact that

a is a morphism in \mathcal{K}/B provides us with a morphism a which makes the diagram

$$\begin{array}{ccc} B/p & \xrightarrow{a} & E \\ & \searrow R(p) \quad \swarrow p & \\ & B & \end{array} \quad (10)$$

commute. Moreover, by remark 2.4 R being a co-KZ-monad generates an adjunction $i(p) \dashv a$ whose unit is the invertible 2-cell $\zeta(p): 1 \Rightarrow a \circ i(p)$

$$\begin{array}{ccccc} & & 1 & & \\ & & \Downarrow \zeta(p) & & \\ E & \xrightarrow{i(p)} & B/p & \xrightarrow{a} & E \\ & \searrow p \quad \swarrow R(p) & & \swarrow p & \\ & & \mathcal{A} & & \end{array} \quad (11)$$

such that $p \cdot \zeta(p) = id_p$.

In the example below we investigate how the construction above look like when we choose 2-category of (locally small) categories as our working 2-category.

EXAMPLE 4.4. Let's take $\mathcal{K} = \mathbf{Cat}$ to be the strict 2-category of categories, functors, and natural transformations. First and foremost, for a functor $p: E \rightarrow B$, the comma category B/p is given as a category whose objects are pairs $\langle f: a \rightarrow p(e); e \rangle$ where f is morphism in B :⁵

$$\begin{array}{ccc} & e & \\ & \downarrow p & \\ b_0 & \xrightarrow{f} & b_1 \end{array}$$

Morphisms of B/p are of the form

$$\begin{array}{ccccc} & e & & \tilde{h}_1 & \\ & \downarrow p & & \searrow & \\ b_0 & \xrightarrow{f} & b_1 & & e' \\ & \searrow h_0 & \swarrow h_1 & & \downarrow p \\ & c_0 & \xrightarrow{g} & c_1 & \end{array}$$

$R(p)$ as in diagram (8) takes pair $\langle f; e \rangle$ to $b_0 = \text{dom}(f)$, and \hat{d}_1 is simply the second projection; it takes $\langle f; e \rangle$ to e . The unit $i(p): E \rightarrow B/p$ sends an object e of E to the object

$$\begin{array}{ccc} & e & \\ & \downarrow p & \\ p(e) & = & p(e) \end{array}$$

⁵ $e \mapsto b_1$ indicates that $p(e) = b_1$.

We also note that $\widehat{d_1^*}$ (as in diagram 9) is given by the action

$$\begin{array}{ccc} & e & \\ & \downarrow & \\ b_0 & \xrightarrow{f} b_1 & \xrightarrow{g} b_2 \end{array} \mapsto \begin{array}{ccc} & e & \\ & \downarrow & \\ b_1 & \xrightarrow{g} & b_2 \end{array}$$

and multiplication $m(p)$ given by

$$\begin{array}{ccc} & e & \\ & \downarrow & \\ b_0 & \xrightarrow{f} b_1 & \xrightarrow{g} b_2 \end{array} \mapsto \begin{array}{ccc} & e & \\ & \downarrow & \\ b_0 & \xrightarrow{g \circ f} & b_2 \end{array}$$

Now, suppose that $\mathbf{a}: R(p) \rightarrow p$ is a pseudo algebra for 2-monad R . By commutativity of diagram 10 we know that $p(\mathbf{a}\langle f; e \rangle) = \text{dom}(f)$. So we draw

$$\begin{array}{ccc} \mathbf{a}\langle f; e \rangle & & \\ p \downarrow & & \\ b_0 & \xrightarrow{f} & b_1 \end{array}$$

As observed in diagram 11 we get an isomorphism lift of identity in the base:

$$\begin{array}{ccc} e & \xrightarrow{\zeta(p)(e)} & \mathbf{a}\langle 1_{p(e)}; e \rangle \\ p \downarrow & & \downarrow p \\ p(e) & \xlongequal{\quad} & p(e) \end{array}$$

Observe that functors $R(i(p)): B/p \rightarrow B/R(p)$ and $i(R(p)): B/p \rightarrow B/R(p)$ are given as follows:

$$R(i(p)) : \begin{array}{ccc} & e & \\ & \downarrow & \\ b_0 & \xrightarrow{f} & b_1 \end{array} \mapsto \begin{array}{ccc} & e & \\ & \downarrow & \\ b_0 & \xrightarrow{f} b_1 & \xlongequal{\quad} b_1 \end{array}$$

and

$$i(R(p)) : \begin{array}{ccc} & e & \\ & \downarrow & \\ b_0 & \xrightarrow{f} & b_1 \end{array} \mapsto \begin{array}{ccc} & e & \\ & \downarrow & \\ b_0 & \xlongequal{\quad} b_0 & \xrightarrow{f} b_1 \end{array}$$

and there is a natural transformations $\tau: i(R(p)) \Rightarrow R(i(p))$

$$\begin{array}{ccccc}
 & & & e & \\
 & & & \downarrow p & \parallel \\
 b_0 & \xlongequal{\quad} & b_0 & \xrightarrow{f} & b_1 \\
 & \parallel & & \searrow f & \parallel \\
 & & b_0 & \xrightarrow{f} & b_1 \\
 & & & \parallel & \\
 & & & b_1 & \xlongequal{\quad} & b_1
 \end{array}$$

Indeed, τ is the mate 2-cell given in diagram 4. We also keep in mind that $R(\mathfrak{a}) \circ R(i(p))(\langle f; e \rangle) = \langle f; \mathfrak{a}(1_{b_1}; e) \rangle$, and hence $R(\zeta(p))$ is illustrated as in below:

$$\begin{array}{ccc}
& e & \\
& \downarrow p & \searrow \zeta(p)(e) \\
b_0 & \xrightarrow{f} & b_1 \\
\parallel & & \parallel \\
b_0 & \xrightarrow{f} & b_1 \\
& & \downarrow p \\
& & \mathbf{a}\langle 1_{b_1}; e \rangle
\end{array}$$

In addition, invertible 2-cell $\theta(p) \colon \mathfrak{a} \circ R(\mathfrak{a}) \Rightarrow \mathfrak{a} \circ m(p)$ provides us with an isomorphism $\mathfrak{a}\langle f; \mathfrak{a}\langle g; e \rangle \rangle \rightarrow \mathfrak{a}\langle gf; e \rangle$. Now, we study the coherence equations 1.1 in our case, which state that the following diagrams commute:

$$\begin{array}{ccc}
\alpha\langle f; e \rangle & \xrightarrow{\alpha.R(\zeta(p))} & \alpha\langle f; \alpha(1_{b_1}; e) \rangle \\
\zeta(p) \cdot \alpha\langle f; e \rangle \downarrow & \searrow & \downarrow \theta.R(i(p)) \\
\alpha(1_{b_0}; \alpha\langle f; e \rangle) & \xrightarrow{\theta.i(R(p))} & \alpha\langle f \circ 1_{b_0}; e \rangle
\end{array}
\qquad
\begin{array}{ccc}
\alpha\langle f; \alpha\langle g; \alpha\langle h; e \rangle \rangle \rangle & \longrightarrow & \alpha\langle gf; \alpha\langle h; e \rangle \rangle \\
\downarrow & & \downarrow \\
\alpha\langle f; \alpha\langle gh; e \rangle \rangle & \longrightarrow & \alpha\langle hgf; e \rangle
\end{array}$$

Furthermore, the counit of adjunction $i(p) \dashv \mathfrak{a}$, as computed in diagram 6, gives us the lift $\tilde{f} = \hat{d}_1((Ra \cdot \tau_p) \circ R\zeta(p)^{-1})$ of f :

The diagram illustrates the relationships between various objects and maps. The objects are arranged in a grid-like structure:

- Top row: $a(f; e)$ and $a(1_{b_1}; e)$
- Second row: $a(1_{b_1}; e)$ and e
- Third row: b_0 (double line) and b_0
- Fourth row: b_0 (double line) and b_1
- Fifth row: b_0 (double line) and b_1

The maps are as follows:

- p from $a(f; e)$ to b_0 (double line)
- $\widehat{d_1}(Ra, \tau_p)$ from $a(f; e)$ to $a(1_{b_1}; e)$
- $\widehat{d_1}(R\zeta(p)^{-1})$ from $a(1_{b_1}; e)$ to e
- f from b_0 (double line) to b_1 (double line)
- f from b_0 to b_1
- f from b_0 (double line) to b_1 (double line)
- 1 from b_0 to b_0 (double line)
- p from e to b_1 (double line)

It remains to prove that \tilde{f} as defined is cartesian. One can try to prove this directly. However, we prove this in a more general setting in the next section.

Now, equations (12) and (13), and definition of $R(ai(p))$ altogether prove that $\gamma_1 \lambda_1 = R(ai(p))$. So, we have

$$\gamma_1 \circ \lambda_1 = R(a) \circ R(i(p)) = R(a \circ i(p))$$

and we shall show that counit $\varepsilon: \gamma_1 \circ \lambda_1 \Rightarrow 1$ is given by $R(\zeta(p)^{-1})$ which is invertible.

Now, we will construct the unit and show that counit and unit satisfy triangle equations of adjunction.

To be completed!

$$\begin{array}{ccc}
 E^{\rightarrow} & \xleftarrow{a^{\rightarrow}} & (B/p)^{\rightarrow} \\
 & \searrow \lambda_1 & \nearrow k \\
 & B/p & \\
 p^{\rightarrow} \swarrow & \downarrow \hat{p} & \searrow R(p)^{\rightarrow} \\
 & B^{\rightarrow} &
 \end{array}$$

EXAMPLE 5.2. Let $p: E \rightarrow B$ be a cloven Grothendieck fibration. We will show that p satisfies Chevalley criteria. Let $\mathfrak{s}: B^{\rightarrow} \rightarrow B$ and $\mathfrak{t}: B^{\rightarrow} \rightarrow B$ be the source and target functors, respectively. Note that the objects of $p \downarrow B = \mathfrak{s}^*(E)$ are pairs $\langle e', f: p(e') \rightarrow b \rangle$ where f is morphism in B .

$$\begin{array}{ccc}
 e' & & \\
 \downarrow p & & \\
 a & \xrightarrow{f} & b
 \end{array}$$

($e' \mapsto a$ indicates that $p(e') = a$.) Similarly, the objects of $B \downarrow p = \mathfrak{t}^*(E)$ are pairs $\langle e, f: a \rightarrow p(e) \rangle$ where f is morphism in B .

$$\begin{array}{ccc}
 & e & \\
 & \downarrow p & \\
 a & \xrightarrow{f} & b
 \end{array}$$

Note that the data of a cloven Grothendieck fibration includes structure of a cleavage, that is a choice of cartesian lifts:

$$\rho_{a,b}: \prod_{\text{Hom}(a,b)} \prod_{e \in E_b} \sum_{e' \in E_a} \mathcal{C}art_E(e', e)$$

For all pairs of objects a, b , satisfying:

$$\begin{aligned}
 \mathbf{snd} \rho_{a,c}(g \circ f, e) &\cong \mathbf{snd} \rho_{b,c}(g, e) \circ \mathbf{snd} \rho_{a,b}(f, \mathbf{fst} \rho_{b,c}(g, e)) \\
 \mathbf{snd} \rho_{b,b}(1_{P_e}, e) &\cong 1_e
 \end{aligned} \tag{14}$$

We denote by $\text{Pull}_f(e)$ the domain of cartesian lift i.e. $\mathbf{fst} \rho_{a,b}(f, e)$ and by \tilde{f} the cartesian lift itself i.e. $\mathbf{snd} \rho_{a,b}(f, e)$.

$$\begin{array}{ccc} \text{Pull}_f(e) & \xrightarrow{\tilde{f}} & e \\ p \downarrow & & \downarrow p \\ a & \xrightarrow{f} & b \end{array}$$

The functors $\gamma_0: E^\rightarrow \rightarrow \mathfrak{s}^*(E)$ and $\gamma_1: E^\rightarrow \rightarrow \mathfrak{t}^*(E)$ are defined as follows: for any object $u: e \rightarrow e'$ in E^\rightarrow , we define $\gamma_0(u) = \langle \mathfrak{s}(u), p(u) \rangle$, and $\gamma_1(u) = \langle \mathfrak{t}(u), p(u) \rangle$. Definitions of γ_i ($i=0,1$) on morphisms is rather straightforward: If $u_0: d \rightarrow d'$ and $u_1: e \rightarrow e'$ are in E^\rightarrow and $\langle h, h' \rangle$ is a morphism from u_0 to u_1 in E^\rightarrow , then $\gamma_0(\langle h, h' \rangle) = \langle h, \langle p(h), p(h') \rangle \rangle$. Similarly, $\gamma_1(\langle h, h' \rangle) = \langle h', \langle p(h), p(h') \rangle \rangle$. Moreover, $\lambda_1: \mathfrak{t}^*E \rightarrow E^\rightarrow$ is defined on objects as $\lambda_1 \langle f: a \rightarrow b, e \rangle = \tilde{f}$, and on morphisms by assigning to $\langle u, \langle h, k \rangle \rangle$, morphism $\langle \bar{h}, u \rangle: \tilde{f}_0 \rightarrow \tilde{f}_1$, where \bar{h} is the a unique lift of h which make the upper square commute. \bar{h} is obtained from cartesian property of \tilde{f}_1 . (Note that although \bar{h} is a lift of h it may not be in the cleavage.)

$$\begin{array}{ccc} \text{Pull}_{f_0}(e_0) & \xrightarrow{\tilde{f}_0} & e_0 \\ \tilde{h} \downarrow & & \downarrow u \\ \text{Pull}_{f_1}(e_1) & \xrightarrow{\tilde{f}_1} & e_1 \end{array}$$

We now show that λ_1 is right adjoint to γ_1 . Notice that the counit of adjunction is identity as it is readily observed that $\gamma_1 \circ \lambda_1 = id_{\mathfrak{t}^*E}$. For obtaining the unit $\eta: Id_{E^\rightarrow} \rightarrow \lambda_1 \circ \gamma_1$, take any object $u: e_0 \rightarrow e_1$ of E^\rightarrow . We have $\lambda_1 \circ \gamma_1(f) = \tilde{p(f)}$, and we define $\eta(f)$ as $\langle \overline{id_{p(e_0)}}, id_{e_1} \rangle: f \rightarrow \lambda_1 \circ \gamma_1(f)$, where $\overline{id_{p(e_0)}}$ is the unique vertical lift of $id_{p(e_0)}$ which makes the following triangle commute:

$$\begin{array}{ccc} e_0 & & \\ & \searrow f & \\ \overline{id_{p(e_0)}} \downarrow & & \text{Pull}_{p(f)}(e_1) \xrightarrow{\tilde{p(f)}} e_1 \end{array}$$

We have to verify that triangle identities of adjunction hold:

$$\begin{array}{ccc} E^\rightarrow & \xlongequal{\quad} & E^\rightarrow \\ \gamma_1 \downarrow & \searrow \eta & \nearrow \lambda_1 \\ \mathfrak{t}^*E & \xlongequal{\quad} & \mathfrak{t}^*E \end{array} = \gamma_1 \left(\begin{array}{c} E^\rightarrow \\ \text{=} \\ \mathfrak{t}^*E \end{array} \right) \gamma_1 \quad \text{and} \quad \begin{array}{ccc} \mathfrak{t}^*E & \xlongequal{\quad} & \mathfrak{t}^*E \\ \lambda_1 \downarrow & \searrow \gamma_1 & \nearrow \eta \\ E^\rightarrow & \xlongequal{\quad} & E^\rightarrow \end{array} = \lambda_1 \left(\begin{array}{c} \mathfrak{t}^*E \\ \text{=} \\ E^\rightarrow \end{array} \right) \lambda_1 \quad (15)$$

Observe that $\gamma_1 = \gamma_1 \circ \lambda_1 \gamma_1$ and $\gamma_1 \cdot \eta = id_{\gamma_1}$, and this proves the first pasting identity. Similarly, $\lambda_1 = \lambda_1 \circ \gamma_1 \circ \lambda_1$ and $\eta \cdot \lambda_1 = id_{\lambda_1}$, and hence we have the second pasting identity.

DEFINITION 5.3. For a category B , define 2-category $\mathbf{Fib}(B)$ of fibrations over B whose 0-cells are Grothendieck fibrations, whose 1-cells are fibred functors over B (i.e. those functors over B which preserve cartesian morphisms), and 2-cells are vertical natural transformations (i.e. transformations over B). Compositions are usual composition of functors and natural transformations.

REMARK 5.4. Example 4.4 can be encapsulated as follows: The forgetful 2-functor $U: \mathbf{Fib}(B) \rightarrow \mathfrak{Cat}/B$ is 2-monadic: the **free fibration** of a functor $p: E \rightarrow B$ is fibration $R(p): B/p \rightarrow B$; cleavage (aka fibration structure) on p is uniquely (in fact unique up to unique isomorphism) determined by a pseudo-algebra structure for 2-monad $R = UF$. Strict algebra structures of R correspond to splitting fibration structures on p .

$$\begin{array}{c}
 \mathbf{Fib}(B) \\
 \begin{array}{c} \nearrow \quad \searrow \\ F \quad \quad U \\ \downarrow \quad \uparrow \\ \mathfrak{Cat}/B \end{array} \\
 \mathfrak{Cat}/B \xleftarrow{\quad} R
 \end{array}$$

We also note that for a category B the domain functor $\text{cod}: B^{\rightarrow} \rightarrow B$ is the free Grothendieck fibration on identity functor $1: B \rightarrow B$; that is $\text{dom} = R(1)$. In more explanatory terms this fact states that

We also note that for a category B with pullbacks the codomain functor $\text{cod}: B^{\rightarrow} \rightarrow B$ is the free Grothendieck fibration *with existential quantifiers* on identity functor $1: B \rightarrow B$;

References

- Kock, A. (1995), ‘Monads for which structures are adjoint to units’, *Journal of Pure and Applied Algebra* **Vol.104**, 41–59.
- Lack, S. (2000), ‘A coherent approach to pseudomonads’, *Advances in Mathematics* **Vol.152** (Issue 2), 179–202.
- Street, R. (1974), ‘Fibrations and yoneda’s lemma in a 2-category’, *Lecture Notes in Math., Springer, Berlin* **Vol.420**, 104–133.