# Advanced Algorithm's Project Twitter Followers Project

## **Team members:**

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**Code Explanation and Time Complexity** 

#### **Input section: -**

```
ios_base::sync_with_stdio(0), cin.tie(0), cout.tie(0);
ifstream infile("twitter.gsy");
if (infile.fail())
    cerr << "Error Opening file" << endl;
   exit(1);
string c;
while (!infile.eof())
   string a, b;
   infile >> c;
   a = "", b = "";
   int idx = 0;
    for (int i = 0; i < (int) c.size(); i++)
        if (c[i] == ',')
            idx = i + 1;
           break;
        a.push back(c[i]);
    for (int i = idx; i < c.size(); i++)
        b.push_back(c[i]);
    int x = str_int(a);
   int y = str_int(b);
    adj_1[x].insert(y);
    adj 2[y].insert(x);
infile.close();
```

- Assume that the number of nodes that occurs in the input file = n.
- Assume that the number of edges = m.
- Assume that the length of string s = |s|.

Time complexity: m \* 2 |s| \* (log(n) + log(n)) = O(m \* |s| \* log(n)).

### **Build-tsk-1 function: -**

```
void build()
{
    for (auto x : adj_2)
    {
        int ft = x.first, sc = x.second.size();
        influencers.push_back({sc, ft});
    }
    sort(influencers.begin(), influencers.end());
    reverse(influencers.begin(), influencers.end());
}
```

- In the worst-case scenario adj\_2 will contain all the nodes.
- Time complexity of sort (built-in function) in C++ is n log(n), where n is the size of data, I want to sort it.
- Time complexity of reverse (built-in function) in C++ is n.

#### Time complexity is:

```
O(n + n \log(n) + n) = O(2n + n \log(n)) = O(n \log(n)).
```

## **Task-1 function: -**

```
void tsk_l()
{
   int number;
   cout << "Enter the rank of influencer you want : " << endl;
   cin >> number;
   number--;
   cout << "The influencer ID = " << influencers[number].second;
   cout << " and the number of followers of the influencer = " << influencers[number].first << endl;
}</pre>
```

Time complexity: O (1).

## Task-2 function: -



- In the worst-case scenario adj\_1 and adj\_2 will contain all the nodes.
- we use lower bound built-in function in  $C++(\log(n))$  to check that ID I will increment its frequency is a valid one or not.
- Dealing with map (data structures) is cost log(n).

## Time complexity is: O ( $n^2 \log(n)$ ).

## Task-3 function: -

- We will need to use two pointers (one forward to begin and the other forward to the end) to go into adj\_1 of each ID and it will cost time of log(n) for initialization of two pointers and 2n to move into the adjacency of each ID.

Time complexity is: O (n).