

BLG 335E

Homework 2 Report

150130032 – Baran Kaya

a) Worst case:

Cost of Partition: $\Theta(n)$

Recurrence for Quicksort: $T(n) = T(n-1) + T(0) + \Theta(n) = T(n-1) + \Theta(n)$

Total : $T(n) = \Theta(n) T(n-1)$

$$T(n) = \Theta(n) + \Theta(n-1) + \Theta(n-2) + \dots + \Theta(1)$$

$$T(n) = \Theta(n^2) \rightarrow \text{Upperbound} = O(n^2)$$

Best/Average case:

Recurrence for Quicksort: $T(n) \leq 2T(n/2) + \Theta(n)$

Solving Recurrence: $T(n) = O(n \log n) < \text{Upperbound}$

$$T(n) \leq 2T(n/2) + \Theta(n)$$

Master Theorem case 2:

$$a = 2, b = 2, f(n) = n$$

$$n n^{\log_b(a)} = n^{\log_2 2} = n$$

$$f(n) = n \rightarrow f(n) = n^{\log_b(a)}$$

$$T(n) = \Theta(n^{\log_b(a)} \cdot \log^{k+1}(n))$$

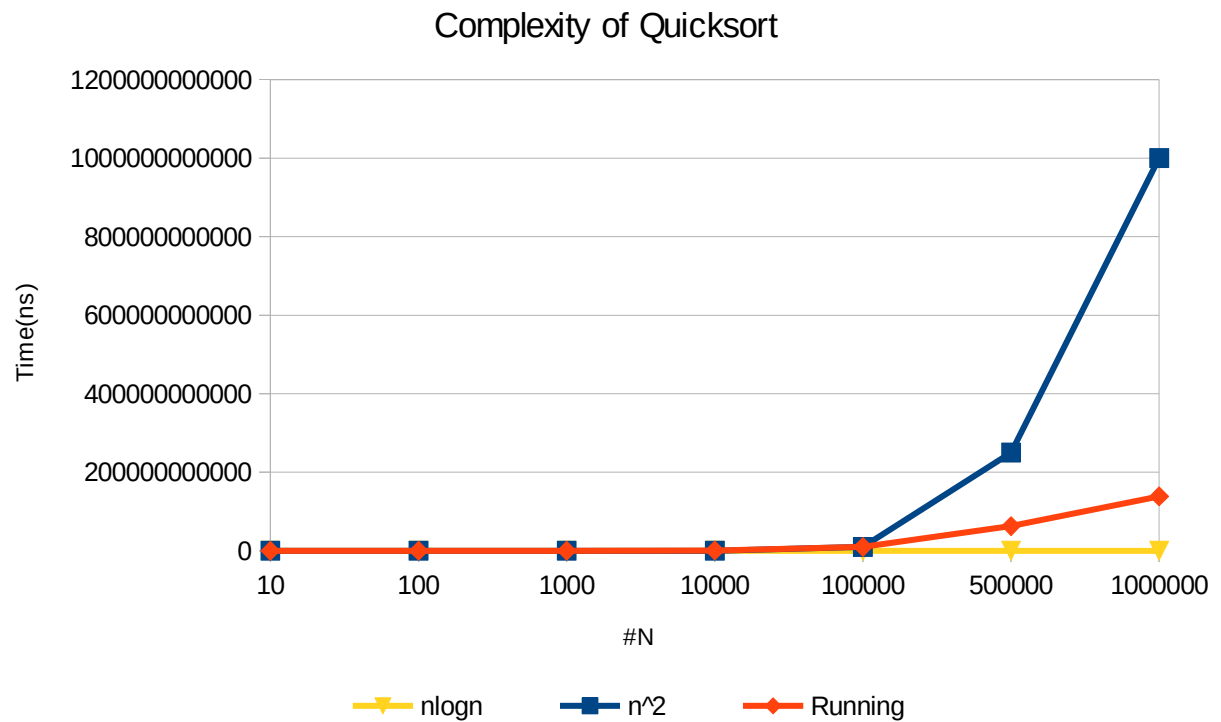
$$T(n) = \Theta(n^{\log_2 2} \cdot \log^{k+1}(n))$$

$$T(n) = \Theta(n \log n)$$

$$T(n) = O(n \log n) \rightarrow \text{Upperbound}$$

b)

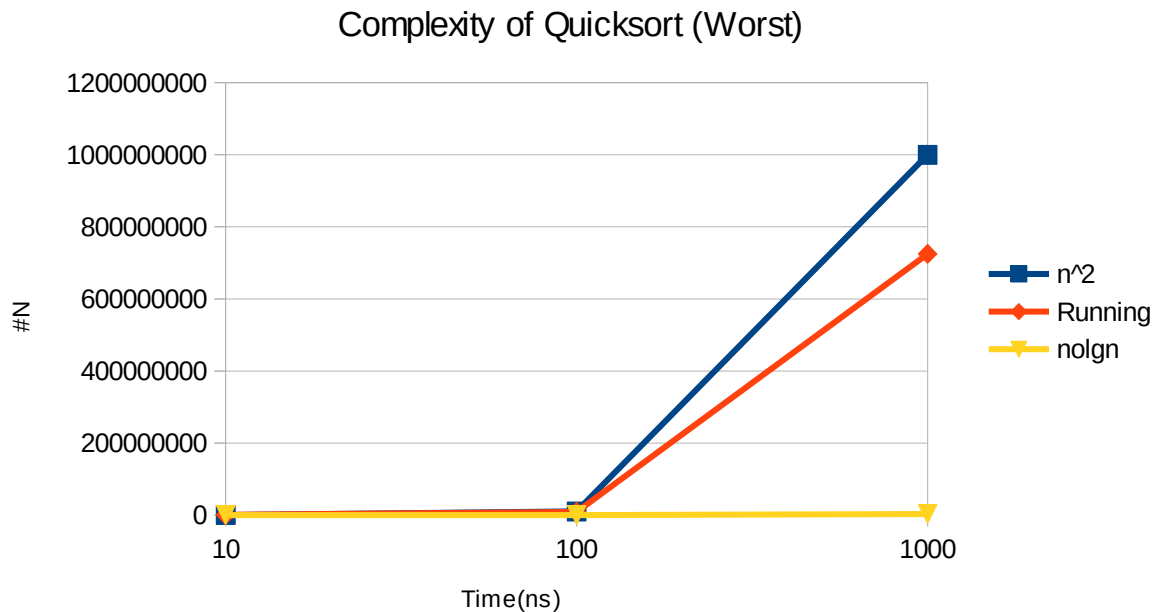
ns	10	100	1000	10000	100000	500000	1000000
1	156419	2214900	44304567	713745435	9998538864	62264656931	138872147521
2	152313	2218595	44264333	710636775	9847431373	62510915618	142297856617
3	156007	2207100	42574111	705008999	9830114510	62473075110	138143475837
4	153545	2261292	45814969	709628882	9762085677	62316047535	137827167605
5	153545	2240354	43598837	701725446	9879535310	62507054425	138534734310
6	169145	2333959	41681583	717297896	9865521706	65048414961	137646938389
7	149850	2228448	43102075	711981315	9831696346	62675197144	137891664504
8	153544	2211616	42704256	707775263	9853215977	62735865284	137756525641
9	150670	2331905	42366785	729653301	9894270245	63438940198	138326167513
10	159292	2240354	42391419	712414442	9881496903	62672317980	138569564689
Average	155433 ns	2248852.3 ns	43280293.5 ns	711986775.4 ns	9864390691.1 ns	62864248518.6 ns	138586624263 ns
Ave. Sec.	0.000155433 s	0.0022488523 s	0.0432802935 s	0.7119867754 s	9.8643906911 s	62.8642485186 s	138.586624263 s



c) If pivot is the biggest or the smallest value in the array than this is the worst case for quick sort. For that purpose i choose the pivot as the one of the (population=1) residence so that it can be smaller than nearly the whole array.

ns	10	100	1000	10000	100000	500000	1000000
1	244686	8723709	686306951				
2	243455	8300025	711402855				
3	243454	8570574	730010067				
4	242223	8272928	722209267				
5	238117	9630606	749444530				
6	238528	9675356	719475026				
7	254949	9235660	706984961				
8	244275	9419175	759552192				
9	305447	9419175	735287235				
10	250433	8214630	727295121				
Average	250556.7 ns	8946183.8 ns	724796820.5 ns				
Average Sec.	0.0002505567 s	0.0089461838 s	0.07247968205 s				

After N=1000 i tried N=10000 but the program failed and i couldn't figure out why it failed. So i decided to make this graph with N=10, N=100 and N=1000



d) Quicksort is not stable also my quicksort in code is not stable neither. Because quicksort changes the elements position that its value equals to pivot (if($A[i] \leq \text{pivot}$)).

ex.

```
if (residenceVec[j].get_population() < pivot.get_population() || (residenceVec[j].get_population() ==
pivot.get_population() && residenceVec[j].get_geo_id() <= pivot.get_geo_id()))
```

-In this code if these two residences populations are the same it will change their position.

-If their populations and geo_ids are the same then it will swap them too.

```
1 - 5,85,0,male,64120,86000000US64120,
2 - 5,55,60,male,52127,86000000US52127 ,
3 - 5,34,40,female,52127,86000000US52127 ,
```

-When comparing 1 and 2, first it compares their population and if they are equal it compares their geo_id. After comparing it will swap 1 and 2's position.

-When comparing 2 and 3, first it compares their population and they are equal so it compares their geo_id. Geo_id's are equal too so it will swap 2 and 3's position.

-Therefore quicksort is not a stable algorithm because it swaps equal elements.

Compiling on SSH: `g++ -std=c++11 ZipcodeQ.cpp -o z -->` It need C++11 for 'stoi' function.

150130032

Baran KAYA