

Let $f : \mathbb{Z}_{10} \rightarrow \mathbb{Z}_{10}$ be a random function. What is the probability that $f(0) = 0$ and $f(1) = 1$? Express your answer as a reduced fraction without any spaces (eg, 1/3 and not 12/36), or as 0 or 1 if appropriate.

$$\begin{aligned} p(A \text{ and } B) &= p(A) * p(B) \\ &= \frac{1}{10} * \frac{1}{10} \\ &= \frac{1}{100} \end{aligned}$$

Let $f : \mathbb{Z}_{10} \rightarrow \mathbb{Z}_{10}$ be a random permutation. What is the probability that $f(0) = 0$ and $f(1) = 1$? Express your answer as a reduced fraction without any spaces (eg, 1/3 and not 12/36), or as 0 or 1 if appropriate.

$$\begin{aligned} p(A \text{ and } B) &= p(A) * p(B|A) \\ &= \frac{1}{10} * \frac{1}{9} \\ &= \frac{1}{90} \end{aligned}$$

You are given a black box $f : \mathbb{Z}_{10} \rightarrow \mathbb{Z}_{10}$ that contains either a random permutation or a random function. Your distinguisher is allowed to invoke f twice. What is the best advantage you can achieve? Express your answer as a reduced fraction without any spaces (eg, 1/3 and not 12/36), or as 0 or 1 if appropriate.

$$\begin{aligned} \text{Advantage} &= \Pr[\text{right}] - \Pr[\text{wrong}] \\ &= \Pr[\text{output func} \mid f \text{ is func}] \\ &\quad - \Pr[\text{output func} \mid f \text{ is perm}] \\ &= \frac{1}{10} - 0 \\ &= \frac{1}{10} \end{aligned}$$

You are given a black box $f()$ that contains either a fair coin or a pair of six-sided dice. If $f()$ is a pair of dice, then each invocation of $f()$ rolls the dice, sums the die faces, and reports 0 if the sum is even and 1 if the sum is odd. If $f()$ is a coin, then each invocation of $f()$ flips the coin and reports 0 if it's heads and 1 if it's tails. What is the advantage of the following distinguisher?

```
if f() == 0
    output "dice"
else
    output "coin"
```

The intuition behind this distinguisher is that there are 6 possible even dice outcomes and only five odd ones. Enter your answer as a reduced fraction without any spaces (eg, $1/3$ and not $12/36$), or as 0 or 1 if appropriate. Note that the probability that a pair of dice sum to 2, 3, 4, 5, 6, 7, 8, 9, 10, 11, 12 is $1/36, 2/36, 3/36, 4/36, 5/36, 6/36, 5/36, 4/36, 3/36, 2/36, 1/36$.

$$\text{Advantage} = \Pr[\text{output dice} \mid f \text{ is dice}] - \Pr[\text{output dice} \mid f \text{ is coin}]$$

pair dice

Even sum : 2 , 4 , 6 , 8 , 10 , 12

$\frac{1}{36}, \frac{3}{36}, \frac{5}{36}, \frac{5}{36}, \frac{3}{36}, \frac{1}{36}$

$$\left(\frac{1}{36} + \frac{3}{36} + \frac{5}{36} + \frac{5}{36} + \frac{3}{36} + \frac{1}{36} \right) - \frac{1}{2}$$

$$\text{Advantage} = 0$$

5

2/2 points

Let's say the following code is executed on a little-endian computer.

```
uint32_t *p = malloc(8);
p[0] = 0x12345678;
p[1] = 0x23456789;
```

What are the 8 bytes in memory that begin at the address that's in p ? Express as 8 two-digit hexadecimal values with a single space between each (eg, ab cd ef 01 02 03 04 50).

1) Starting at $p[0]$, read each byte in Little Endian

78 56 34 12

2) Now from $p[1]$, read each byte in little Endian

89 67 45 23

3) Write all bytes starting from $p[0]$ results

78 56 34 12 89 67 45 23