

DB实验1

本次实验在java + jdbc拼sql完成实验中的导入部分，在索引部分使用命令行排除jdbc的干扰和可能的优化[项目地址](#)，截图为在mac OSX 16G环境下的运行时间

步骤二

jdbc直接从txt导入到数据库表，对应语句

```
private static String LOAD_ORDERS_DATA =
    "load data local infile \"/Users/py/Downloads/索引作业数据/data1.txt\"\\n" +
    "into table orders(id, name, age, sex, amount);";

private static String LOAD_PRODUCTS_DATA =
    "load data local infile \"/Users/py/Downloads/索引作业数据/data2.txt\"\\n" +
    "into table products(id, pid, nums);";

private void loadOrdersData() {
    try {
        Statement statement = connection.createStatement();
        long startTime=System.currentTimeMillis();
        statement.execute(LOAD_ORDERS_DATA);
        long endTime=System.currentTimeMillis();
        System.out.println("Load orders data costs: " + (endTime - startTime)/
1000 + " s.");
    }catch (Exception ex) {
        ex.printStackTrace();
    }
}

private void loadProductsData() {
    try {
        Statement statement = connection.createStatement();
        long startTime=System.currentTimeMillis();
        statement.execute(LOAD_PRODUCTS_DATA);
        long endTime=System.currentTimeMillis();
        System.out.println("Load products data costs: " + (endTime - startTime)/10
00 + " s.");
    }catch (Exception ex) {
        ex.printStackTrace();
    }
}
```

```
/Library/Java/JavaVirtualMachines/jdk1.8.0_121.jdk/Contents/Home/bin/java ...  
Succeed to connect...  
Load orders data costs: 29 s.  
Load products data costs: 0 s.  
  
Process finished with exit code 0
```

步骤三

问题1:在 **orders** 表中找出购买人年龄小于20岁的order列表。

未加入索引前sql语句:

```
select * from orders where age < 20
```

```
571196 rows in set (1.96 sec)
```

加入索引:

```
create index age_idx on orders(age)
```

```
mysql> create index age_idx on orders(age)  
->  
->  
->  
-> ;
```

```
Query OK, 0 rows affected (9.86 sec)  
Records: 0 Duplicates: 0 Warnings: 0
```

再

次执行查询:

```
select * from orders where age < 20
```

```

+-----+-----+-----+-----+-----+
571196 rows in set (1.89 sec)

mysql>

```

发现在age列建立B+树索引并没有对查询效率有明显帮助

```
explain select * from orders where age < 20;
```

```

mysql> explain select * from orders where age < 20;
+----+-----+-----+-----+-----+-----+-----+-----+
| id | select_type | table | partitions | type | possible_keys | key | key_len |
| ref | rows      | filtered | Extra      |      |               |     |         |
+----+-----+-----+-----+-----+-----+-----+-----+
| 1  | SIMPLE     | orders | NULL       | ALL  | age_idx       | NULL | NULL    |
| NULL | 4982784    | 33.33  | Using where |      |               |     |         |
+----+-----+-----+-----+-----+-----+-----+-----+
1 row in set, 2 warnings (0.00 sec)

```

发现这依旧是一个全表查询，type为ALL

究其原因在查询过程中集合index value的大小十分重要，总结来说是当列越长，越少的值能被放入B+树的结点，因此树的深度会变大，当树的深度越大时，需要更多磁盘访问，磁盘访问越多效率越低

这是援引ovasitariq的一则回答翻译理解过来

在这个例子中，age列集合显然不会超过200与整个order表无法相比，此时磁盘开销巨大，需要更多的索引，而这些具体的开销高昂以至于MySQL直接全表查询

问题2:在 orders 表中找出所有姓王的人的order列表。

```
select * from orders where name like '王%';
```

```

+-----+-----+-----+-----+-----+
11160 rows in set (1.62 sec)

```

继续建B+索引在name上

```
create index name_idx on orders(name);
```

```
mysql> create index name_index on orders(name);  
Query OK, 0 rows affected (10.83 sec)  
Records: 0 Duplicates: 0 Warnings: 0
```

继续查询

```
1 4999928 | 王忆 | 42 | 女 | 72 |  
+-----+-----+-----+-----+-----+  
11160 rows in set (0.11 sec)
```

效率大大增加

可以理解，因为名字集合本身很大，统计一下

```
select count(distinct(name)) from orders;
```

```
mysql> select count(distinct(name)) from orders  
-> ;  
+-----+  
| count(distinct(name)) |  
+-----+  
| 3477290 |  
+-----+  
1 row in set (1.84 sec)
```

总: 3477290确实很巨大，考量表的其他列，名字的区别性确实让本身建立索引查询占优

问题3:统计 orders 表中所有男性的人的数量。

```
select count(*) from orders where sex = '男';
```

```
+-----+
| 2499997 |
+-----+
1 row in set (1.17 sec)
```

心里有数然而继续实践一下

```
create index sex_idx on orders(sex);
```

```
mysql> create index sex_idx on orders(sex);
Query OK, 0 rows affected (6.01 sec)
Records: 0 Duplicates: 0 Warnings: 0
```

```
select count(*) from orders where sex = '男';
```

```
mysql> select count(*) from orders where sex = '男'
+-----+
| count(*) |
+-----+
| 2499997 |
+-----+
1 row in set (0.44 sec)
```

因为是单纯计数而不需要访问本身，其索引建立的意义得到体现

```
explain select count(*) from orders where sex = '男';
```

```
mysql> explain select count(*) from orders where sex = '男';
+-----+-----+-----+-----+-----+-----+-----+-----+
| id | select_type | table | partitions | type | possible_keys | key | key_
len | ref | rows | filtered | Extra |
+-----+-----+-----+-----+-----+-----+-----+-----+
| 1 | SIMPLE | orders | NULL | ref | sex_idx | sex_idx | 2
| const | 2491392 | 100.00 | Using index |
+-----+-----+-----+-----+-----+-----+-----+-----+
1 row in set, 1 warning (0.00 sec)
```

此时是索引查询

问题4:在 **orders** 表中计算女性，姓张，年龄大于50，且消费小于100的人数。

```
select count(*) from orders where sex = '女' and name like '张%' and age > 50 and a
mount < 100;
```

```
mysql> select count(*) from orders where sex = '女' and name like '张%' and age
> 50 and amount < 100;
+-----+
| count(*) |
+-----+
|      2586 |
+-----+
1 row in set (0.10 sec)
```

沿用上面两题建立对**sex, name, amount, age**的复合索引而在计数不需要担心全部访问的尴尬

```
create index mul_idx on orders(sex, name, amount, age);
```

```
mysql> create index mul_idx on orders(sex, name, amount, age);
Query OK, 0 rows affected (13.42 sec)
Records: 0 Duplicates: 0 Warnings: 0
```

Once again

```
select count(*) from orders where sex = '女' and name like '张%' and age > 50 and a
mount < 100;
```

```
mysql> select count(*) from orders where sex = '女' and name like '张%' and age
> 50 and amount < 100;
+-----+
| count(*) |
+-----+
|      2586 |
+-----+
1 row in set (0.01 sec)
```

Stack Overflow对于复合列索引有一个类比电话本的解释我很喜欢，写一下：

1. 如果你找一个姓潘的人，你可以很容易找到因为电话本按姓排列
2. 如果你找一个名羽的人，那没办法了，因为电话本又链接不到名，你得查找全部😭
3. 如果你找一个姓潘名羽的肥宅，电话本会很有帮助，因为你能找到潘以后按名的顺序找到羽

以上解释了查找准确值，但是如果很不幸你想找一个范围，比如名是羽姓以p开头，这时候你得先找到那些名羽的人，只有先这样你才能group

我记下来的原因是因为这对理解复合列的索引很有帮助，不要以为你做完了，你只是运气好吧sex放第一个，那么如果你运气不好呢，尝试一把

```
create index test_idx0 on orders(name, amount, age, sex);
```

```
mysql> create index test_idx0 on orders(name, amount, age, sex);
Query OK, 0 rows affected (12.85 sec)
Records: 0 Duplicates: 0 Warnings: 0

mysql> █
```

忽略那些令人尴尬的warning，对multi-index都有数，MySQL还是选择了我第一把建的mul_index，这是有道理的

```
-----+-----+-----+-----+-----+-----+-----+-----+
-----+
| 1 | SIMPLE      | orders | NULL      | range | age_idx,name_index,sex_idx,mu
l_idx,test_idx0 | mul_idx | 140      | NULL | 5715 | 11.11 | Using where; Usin
g index |
+---+-----+-----+-----+-----+-----+-----+-----+
-----+
1 row in set, 4 warnings (0.00 sec)

mysql> █
```

此时心中应该规定优先级 exact > range > like

问题5:统计 orders 表中姓名为三个字的人数。

```
select count(*) from orders where name like '___';
```

```
mysql> select count(*) from orders where name like '___';
+-----+
| count(*) |
+-----+
| 2501252 |
+-----+
1 row in set (1.18 sec)
```

```
create index name_idx on orders(name);
```

```
mysql> create index name_idx on orders(name);
Query OK, 0 rows affected, 1 warning (10.04 sec)
Records: 0 Duplicates: 0 Warnings: 1
```

好了，这个索引好像没用

```
mysql> select count(*) from orders where name like '___';
+-----+
| count(*) |
+-----+
| 2501252 |
+-----+
1 row in set (1.16 sec)
```

```
explain select count(*) from orders where name like '___';
```


id	select_type	table	partitions	type	possible_keys	key	key_len	ref	rows	filtered	Extra
1	SIMPLE	orders		index	NULL				4982784	11.11	Using where; Using index

其实还是用了的，然而这个索引之所以没有提升也不能理解，索引本身的数量决定了性能，而对于name列来说这本身就是一个巨大的挑战，MySQL拿出资源来维护索引本身也是高昂的开销

问题6:在 products 表中查找库存大于150的product列表。

```
select * from products where nums>150;
```

```
2534 rows in set (0.00 sec)
```

其实我已经不想建索引了，0s我也看不出来，但没办法，Once again

```
create index nums_idx on products(nums);
```

```
mysql> create index nums_idx on products(nums);
Query OK, 0 rows affected (0.04 sec)
Records: 0 Duplicates: 0 Warnings: 0
```

```
2534 rows in set (0.01 sec)
```

其实你建了也没用

还是那句话，索引本身是有代价的，MySQL会在资源和查询中作出权衡，具体视数据而定，这道题不应该建