

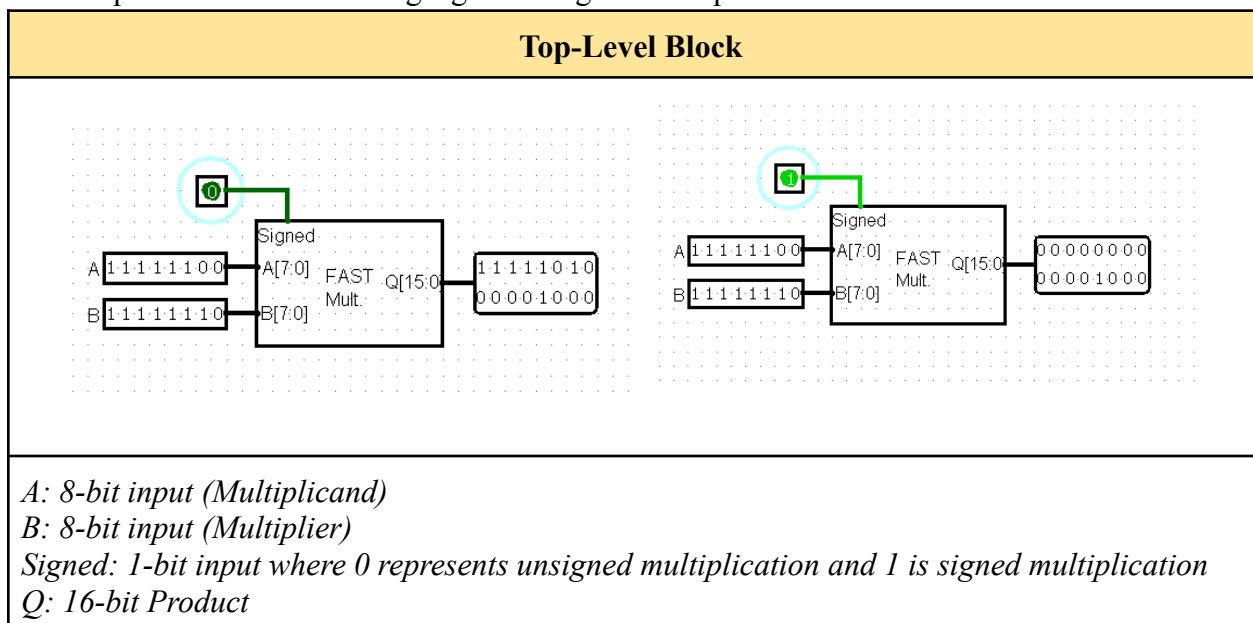
## Fast Multiplier

### Introduction

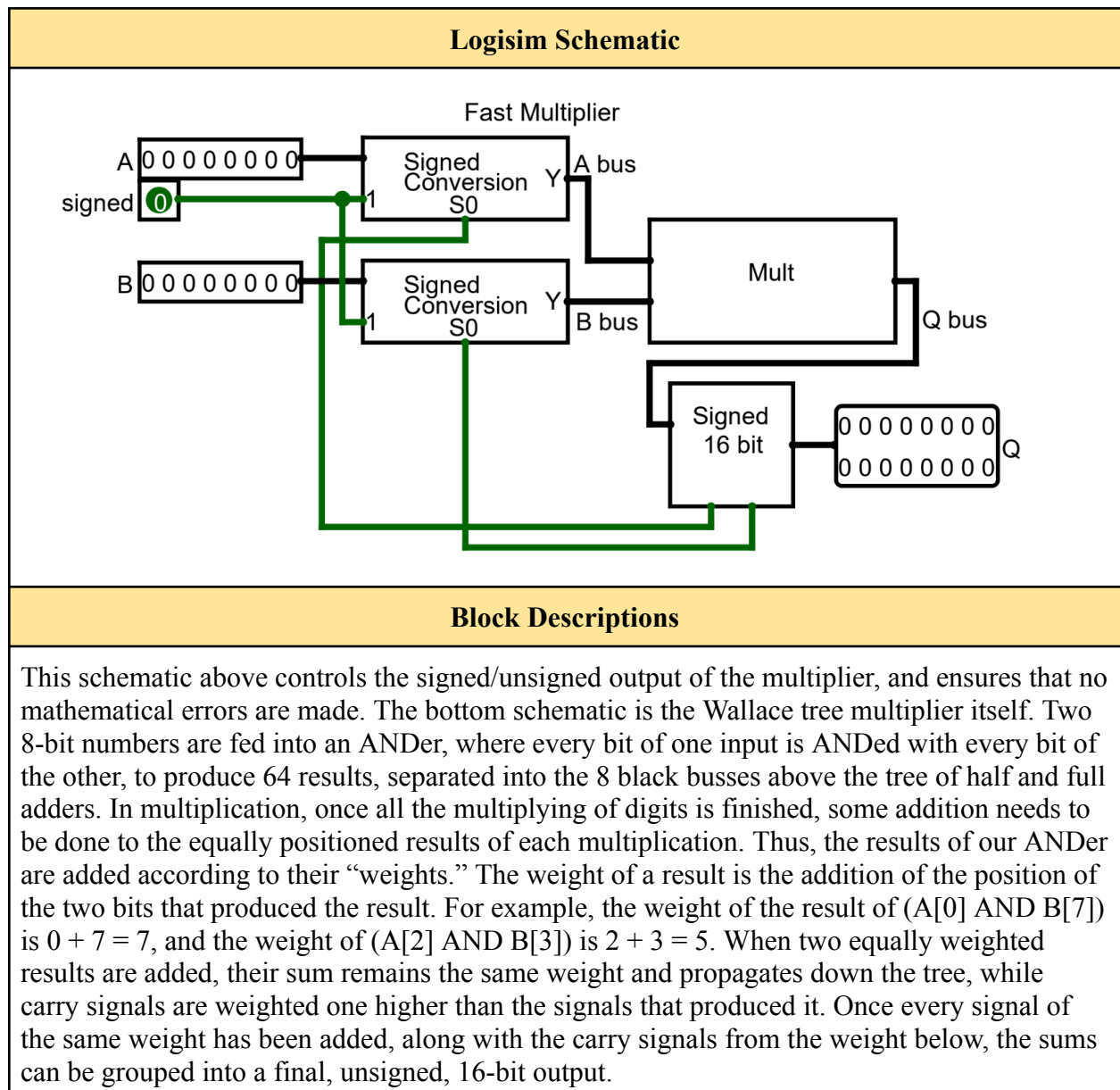
Computer chips were made to perform mathematical calculations for the ease of humans, and they were built to do it fast. So it is an area of intrigue for Computer Engineers to consider how best to do an operation like multiplication. Because multiplication is repeated addition, it can be considerably slower than any of the other operations a CPU may perform, up to 32 times for a standard MIPS processor if one addition is performed every clock cycle. A much faster method of multiplication is the Wallace tree, or a flow-through multiplier, which completes the whole process in one clock cycle, “flowing through” the whole system like an unmitigated waterfall.

### Specification

Our goal was to implement a Wallace tree multiplier in Magic, the VLSI software. For simplicity, we wanted our multiplier to multiply two eight-bit numbers together, to have an output of a maximum of 16 bits. We also wanted to control whether the multiplier understood the sign of an input, in two’s complement form, or if it treated the two inputs as unsigned integers. Because a Wallace tree performs solely unsigned multiplication, we need to manipulate the data before it enters the tree, and then manipulate the data that comes out of it depending on the state of the inputs and if we are doing signed/unsigned multiplication.

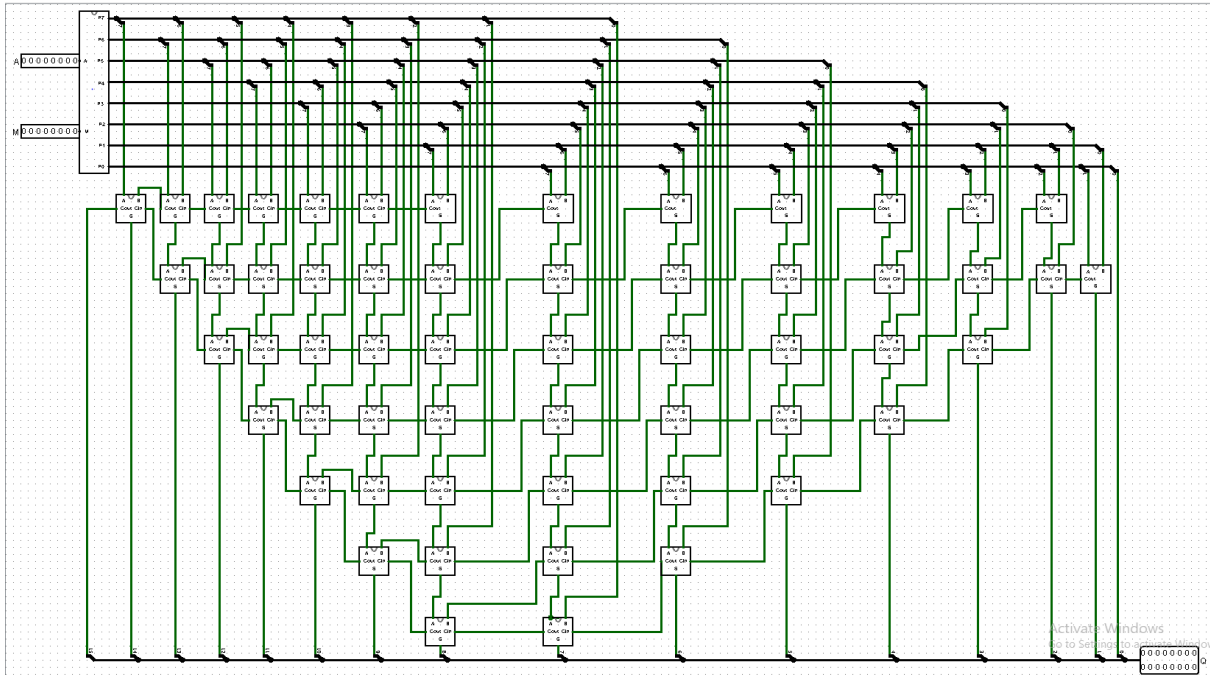


This is the final top-level schematic of our multiplier. The two 8-bit inputs A and B are multiplied together to produce Q, and a “signed” input determines if the block reads those inputs assigned or not. Inside the block is shown below.



The multiplier cannot multiply a “negative” number, so if the leading bit is one, signaling it is negative is two’s complement, and we are doing a signed multiplication, then the input is inverted in two’s complement before entering the multiplier. In a signed multiplication, neither input will be negative entering the tree. The result is then kept the same or negated depending on the state of the two inputs: if one of them is negative, then the result will be re-inverted so that a negative time a positive turns out negative. But if neither or both of them are negative, the result remains positive.

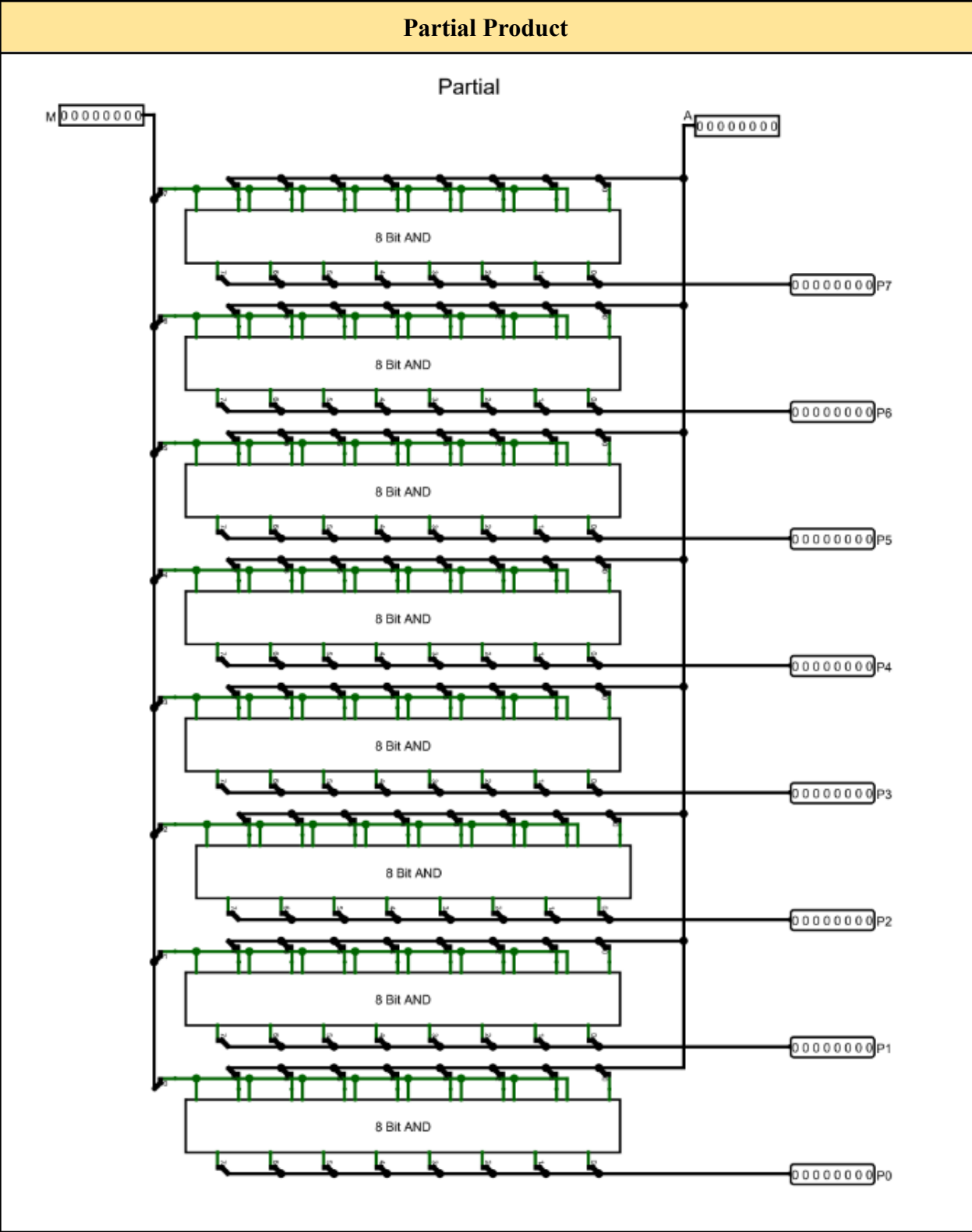
## Multiplier



## Block Descriptions

The schematic below is the circuit that produces the 64 output results of ANDing every bit of the M input with every bit of the A input. So, the bit M[0] is ANDed with A[0], A[1], A[2], ..., and A[7], and the bit M[1] is ANDed with A[0], A[1], A[2], ..., and A[7], etc.

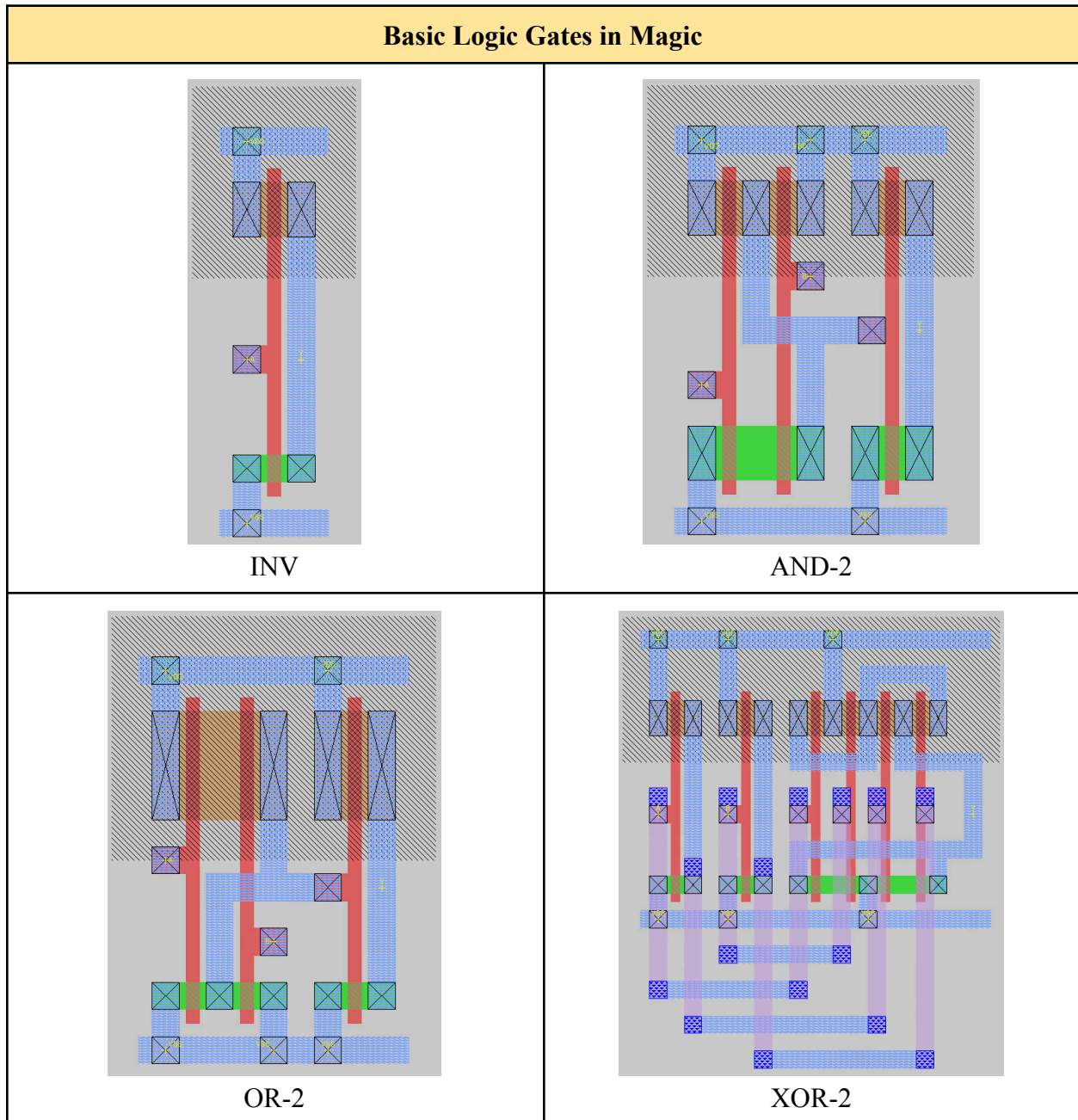
Above is the schematic for the Wallace tree. A Wallace tree takes advantage of the nature of binary multiplication, in that a multiplication between two bits is merely an ANDing of them, and the results of those ANDings need to be added according to their “weight.” The weight of a product is the addition of the position of the bits that produced it. For example, P0\_1 is the product of (A0 AND B1) and its weight is  $(1 + 0 = 1)$ . And P3\_2 is the product of (A3 and B2) and its weight is  $(3 + 2 = 5)$ . Products of the same weight are added to produce a sum bit, which is considered the same weight as its parents, and a carry bit, which is one weight higher than its parents. The reason products of the same weight need to be added is made more obvious in the image below.



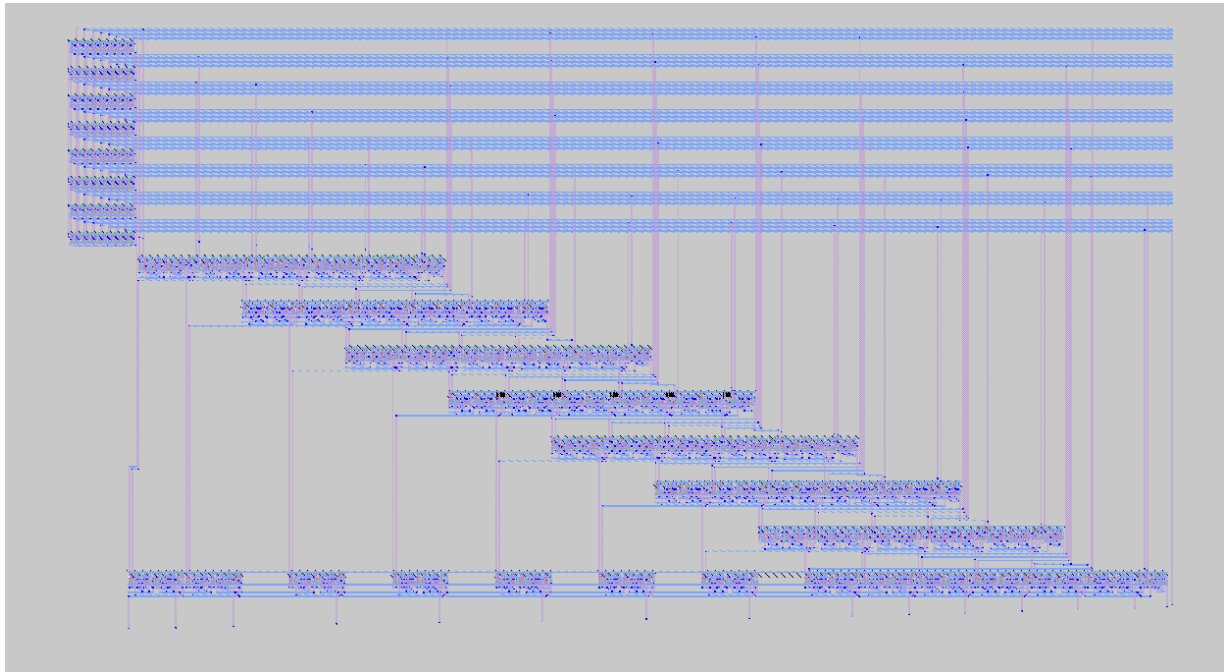
This is the array of AND gates mentioned previously.

### Implementation (magic schematics go under here)

To realize this idea fully, we used Magic, the VLSI simulation software, to lay out the chemical makeup of a microchip that operated this way. We start by making the basic gates that constitute the fundamental building blocks of the massive schematic this will eventually turn into.



## Multiplier



This is what it looks like when the AND signals are pulled out into the ADDers according to their weights. If you look above at Logisim schematics, the similarities in structure are obvious.