



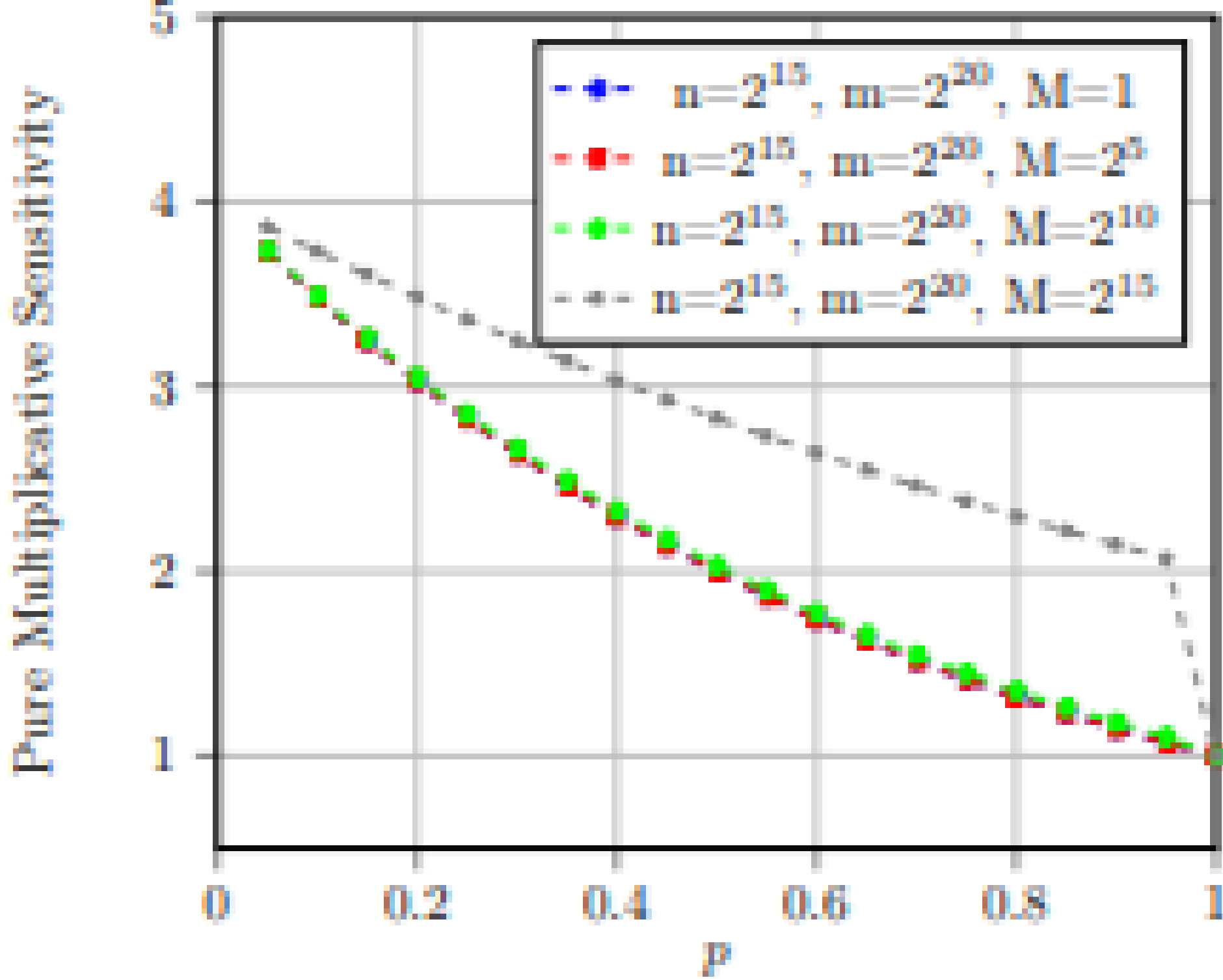
BACKGROUND: In the present work, we make the first customized effort towards DP estimation of fractional frequency moments, i.e. $p \in (0, 1]$ with low space complexity. We show that a well-known streaming algorithm, namely Fp sketch [21], preserves differential privacy as is. With its small space complexity, Fp sketch elegantly solves the trilemma between efficiency, accuracy, and privacy.

Fp SKETCH REVIEW

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Algorithm 1:  $\mathbb{F}_p$  sketch.  
Input : Data stream:  $\mathcal{S} = \{(k_1, v_1), \dots, (k_n, v_n)\}$   
1 Construct:  
2   Initialize  $\mathbf{a} = \{0\}^r$ ,  $\mathbf{P} \sim \mathcal{D}_{p,1}^{r \times m}$ ;  
3 Update:  
4   for  $i \in [n]$  do Let  $e_{k_i}$  be the one-hot encoder of  $k_i$ ,  $\mathbf{a} = \mathbf{a} + \mathbf{P} \times v_i \mathbf{e}_{k_i}$ ;  
5 Query:  
6   return scale_estimator( $\mathbf{a}$ );
```

METHODOLOGY & RESULTS

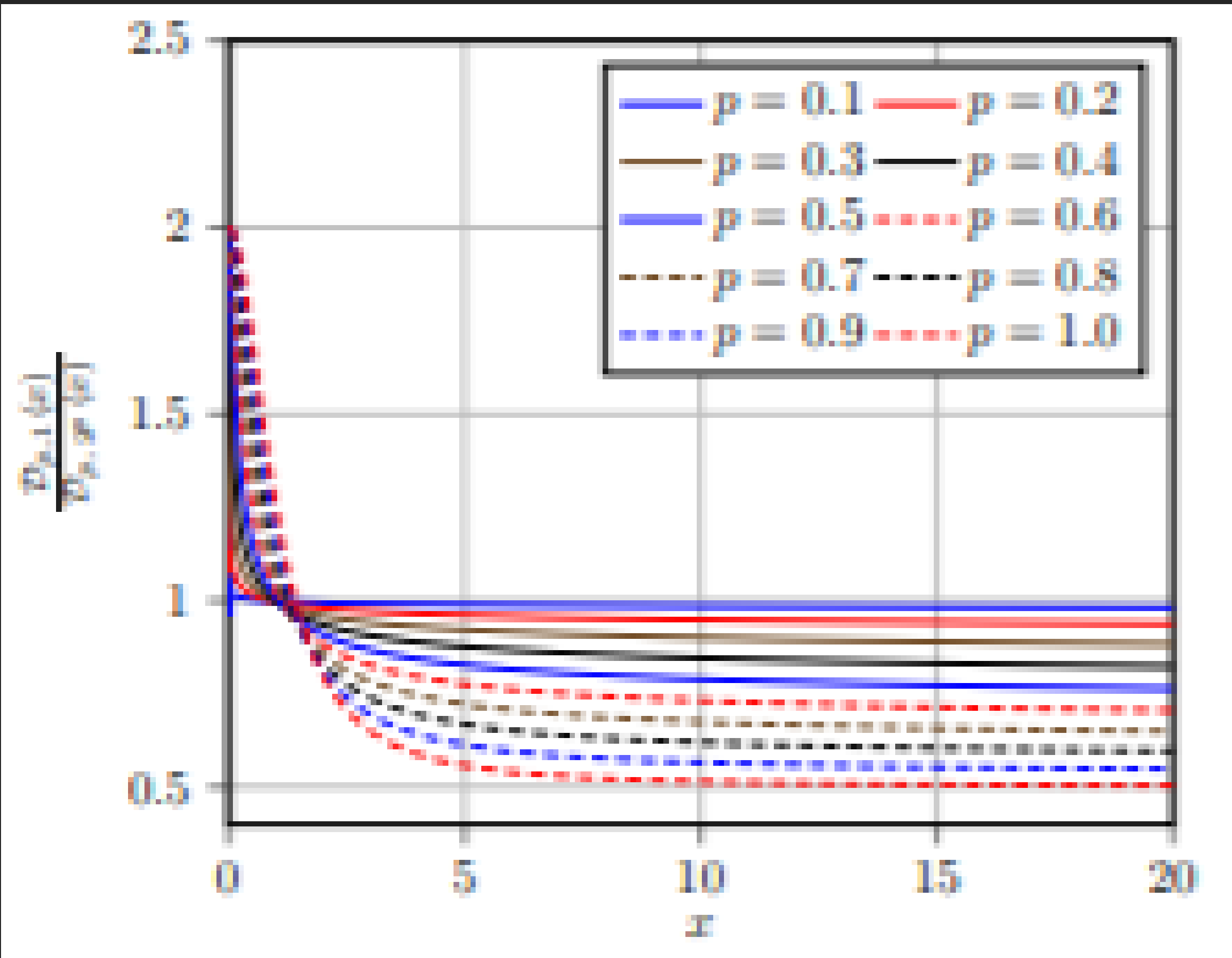
STEP 1: Bound the difference between scales of neighboring datasets



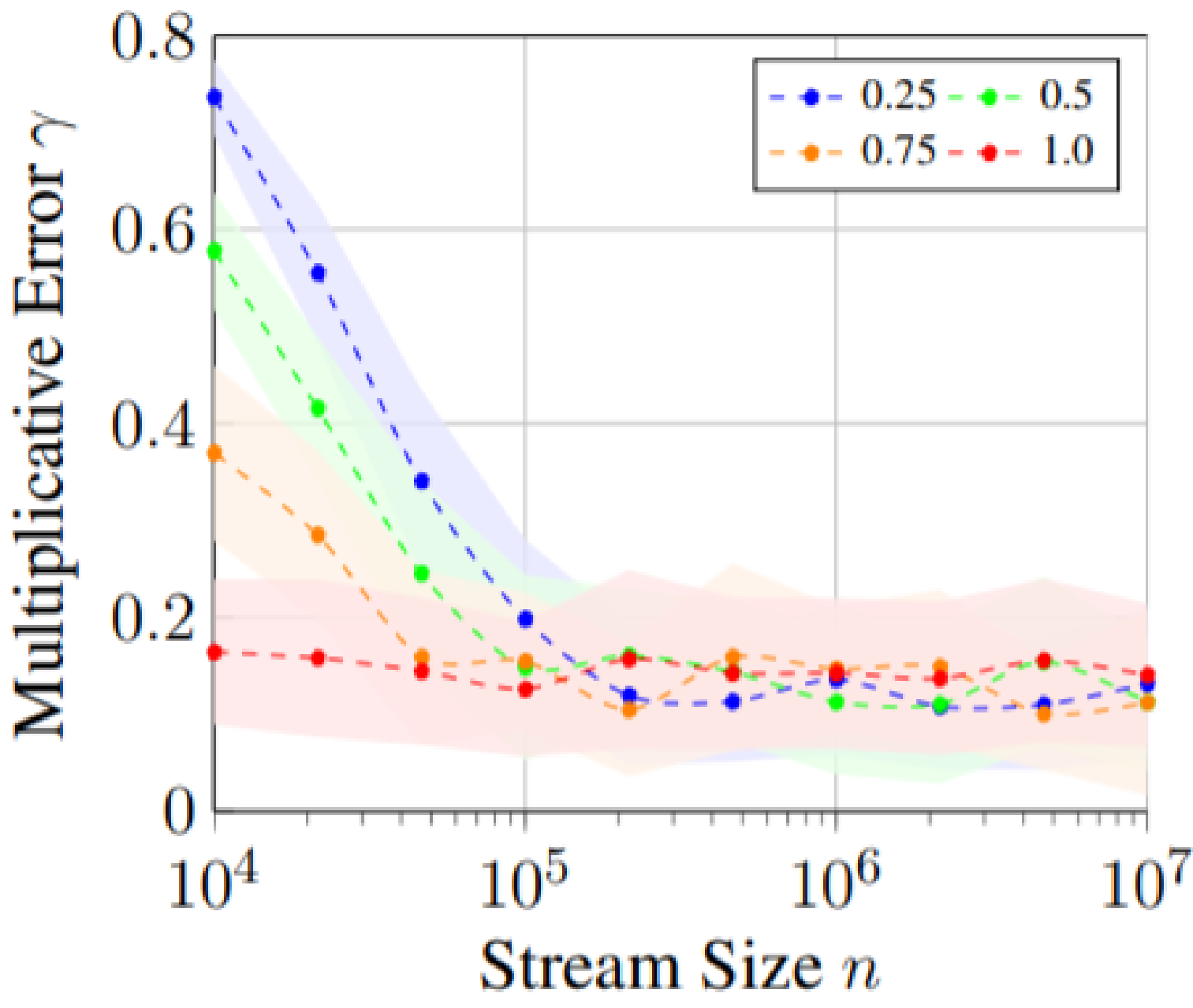
STEP 2: Bound the privacy loss

Theorem 3 (ϵ -DP for Algorithm 1). Let ρ_p represent the multiplicative sensitivity of the p -th frequency moments. When $r = 1$ and $p \in (0, 1]$, moments. When $r = 1$ and $p \in (0, 1]$, \mathbb{F}_p sketch is $\frac{1}{p} \ln \rho_p$ -differentially private.

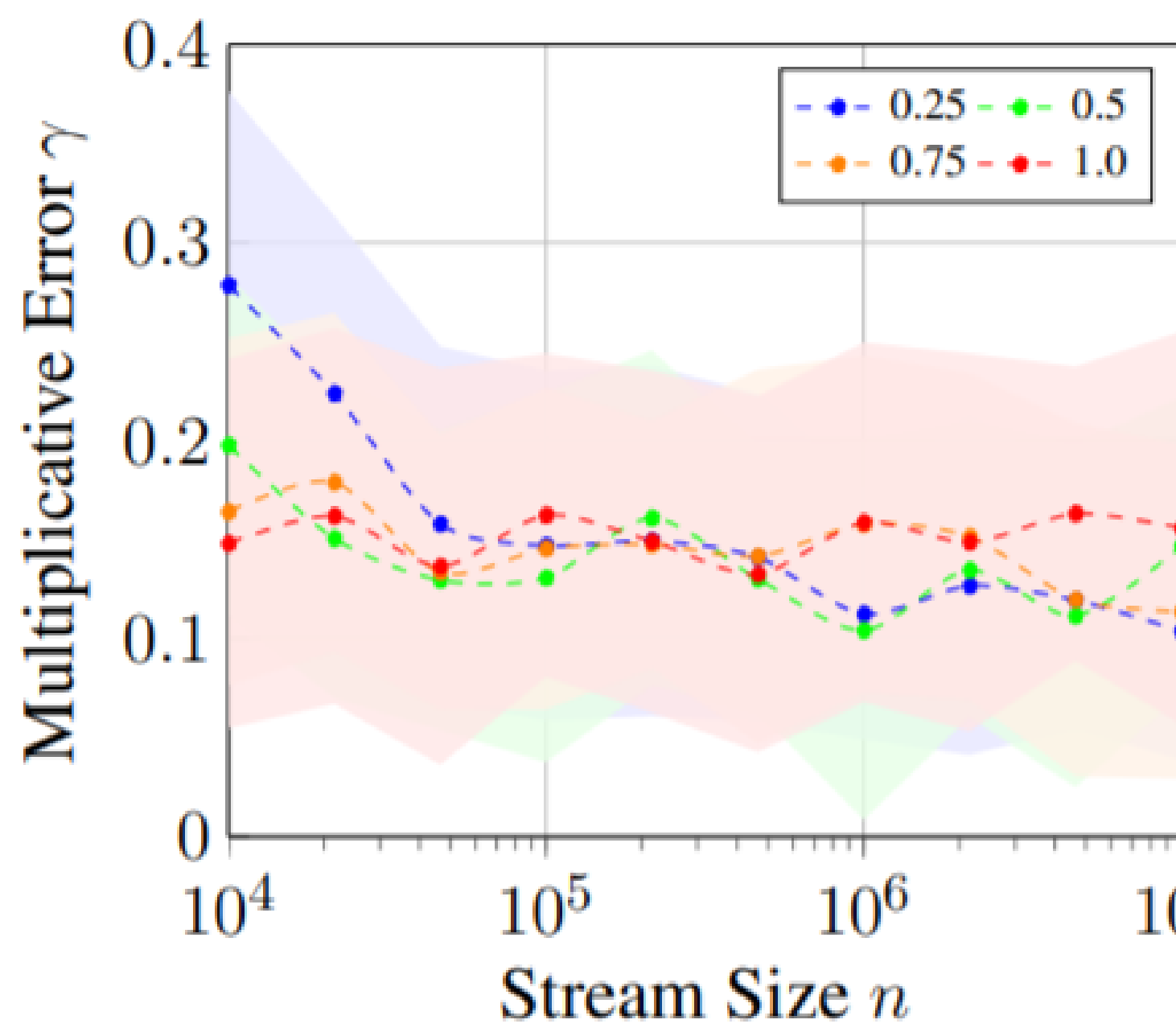
We prove that **Fp sketch**, a streaming algorithm for **frequency moments estimation**, is **differentially private** when $p \in (0, 1]$.



EXPERIMENT



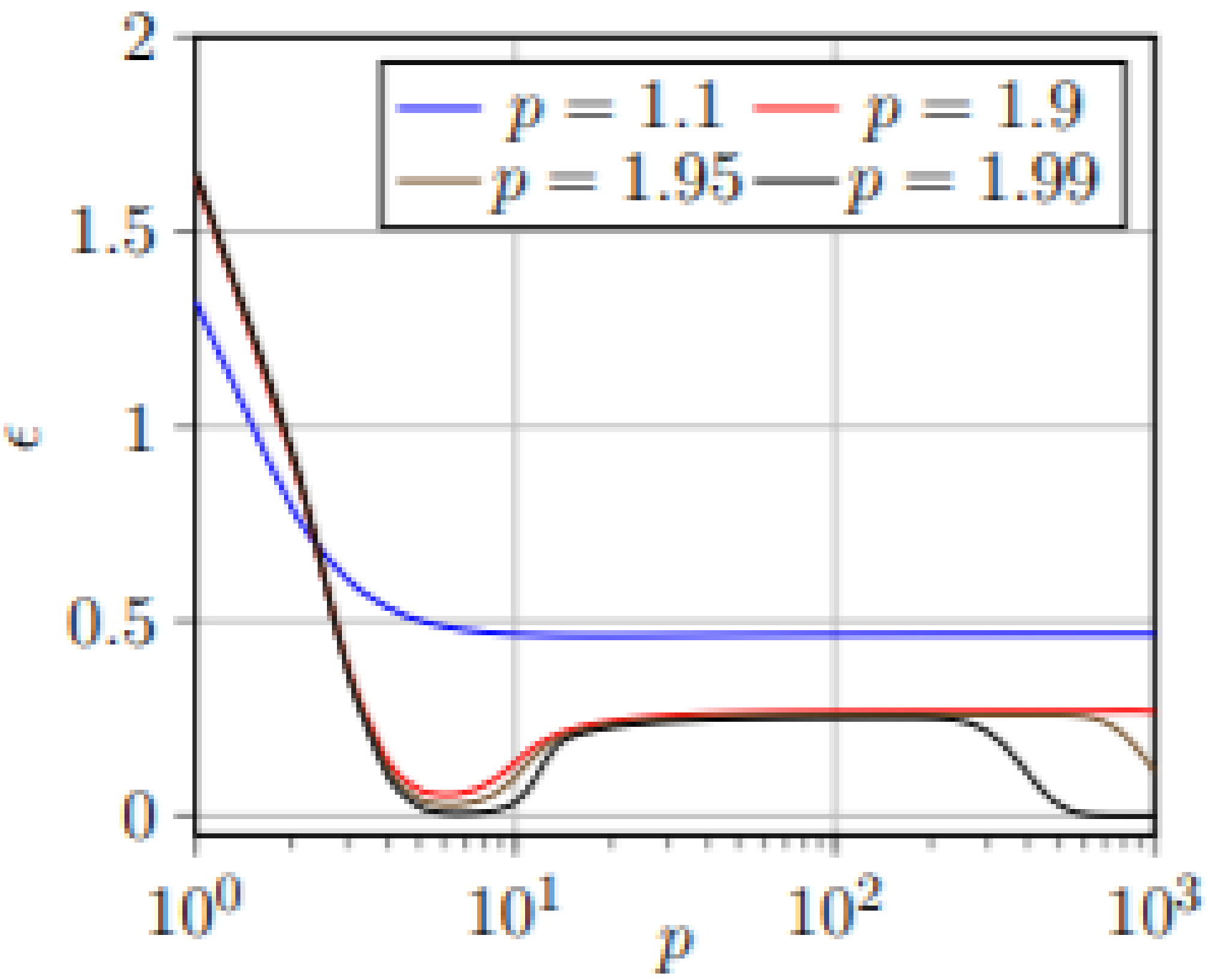
(a) Uniformly Distributed Stream.



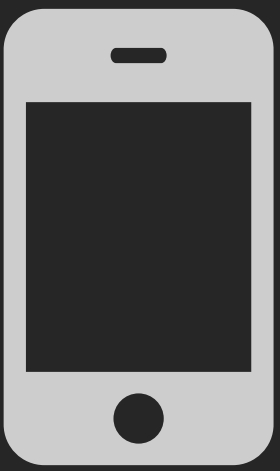
(b) Binomially Distributed Stream.

OPEN PROBLEMS

The proof does not easily extend to $p \in (1, 2)$.



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