Secure Federated Feature Selection



BACKGROUND: we propose the first secure federated $\chi 2$ -test protocol FED- $\chi 2$. To minimize both the privacy leakage and the communication cost, we recast χ2 -test to the second moment estimation problem and thus can take advantage of stable projection to encode the local information in a short vector. As such encodings can be aggregated with only summation, secure aggregation can be naturally applied to hide the individual updates. We formally prove the security guarantee of FED-χ2 that the joint distribution is hidden in a subspace with exponential possible distributions. Our evaluation results show that FED-χ2 achieves negligible accuracy drops with small client-side computation overhead. In several real-world case studies, the performance of FED- $\chi 2$ is comparable to the centralized $\chi 2$ -test.

BUILDING BLOCK: STABLE PROECTION

```
Function Encode (P, u_i):

return P \times u_i

Function Geometric Mean Estimator (e):

\ell: Encoding size.

\hat{d}_{(2),gm} \leftarrow \frac{\prod_{k=1}^{\ell} |\mathbf{e}_k|^{2/\ell}}{(\frac{2}{\pi}\Gamma(\frac{2}{\ell})\Gamma(1-\frac{1}{\ell})\sin(\frac{\pi}{\ell}))^{\ell}}

return \hat{d}_{(2),gm}

Function Decode (e):

return Geometric Mean Estimator (e)
```

METHODOLOGY

STEP 1: Recast $\chi 2$ –test as a second frequency moments estimation problem.

STEP 2: Encode the second frequency moments information using stable projection.

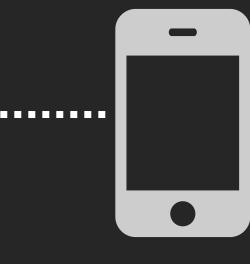
STEP 3: Use secure aggregation to sum the clients' updates.

we propose the first secure federated $\chi 2$ -test protocol FED- $\chi 2$.

Algorithm 2: FED- χ^2 : secure federated χ^2 -test. SECUREAGG is a remote procedure that receives inputs from the clients and returns the summation to the server. INITSECUREAGG is the corresponding setup protocol deciding the communication graph and other hyper-parameters.

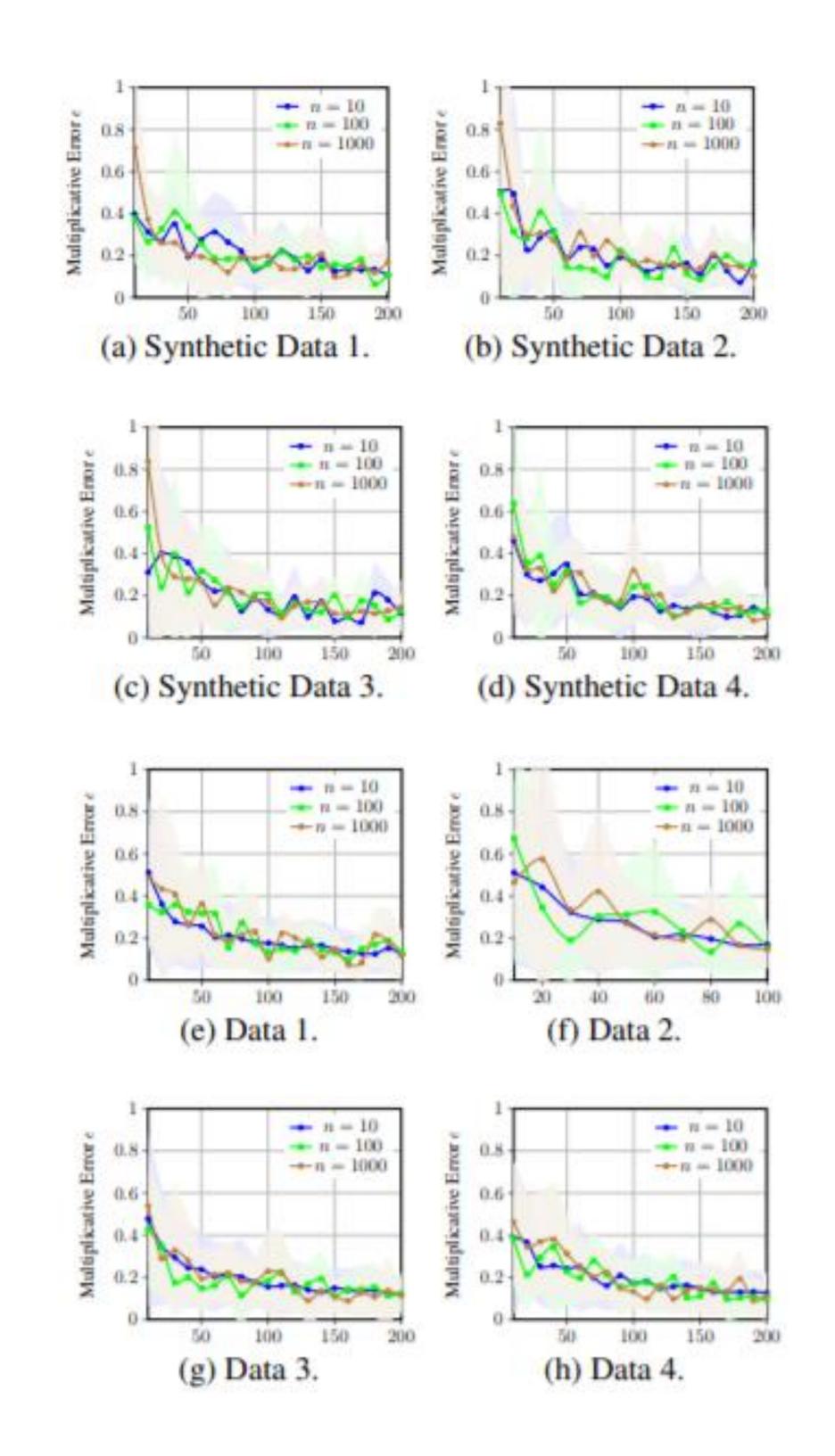
```
Round 1: Reveal the marginal distribution
                                                                                                  // n is the client number.
         INITSECUREAGG(n)
        for x \in [m_x] do v_x = SecureAgg(\{v_x^{(i)}\}_{i \in [n]})
        for y \in [m_y] do v_y = SecureAgg(\{v_y^{(i)}\}_{i \in [n]})
         Server
              Calculate v = \sum_{x} v_x and broadcast v, \{v_x\} and \{v_y\} to all the clients.
   Round 2: Approximate the statistics
        Server
              Sample the projection matrix P from Q_{2,0,1}^{\ell \times m}
              Broadcast the projection matrices to the clients
         Client c_i, i \in [n]
              Calculate \bar{v}_{xy} = \frac{v_x v_y}{v_y}
             Prepare \mathbf{u}_i s.t. \mathbf{u}_i[\mathbb{I}(x,y)] = \frac{v_{xy}^{(i)} - \overline{v_{xy}}/n}{\sqrt{\overline{v_{xy}}}}
13
              Calculate \mathbf{e}_i = \text{Encode}(P, u_i)
14
          i = \mathtt{SecureAgg}(\{e_i\}_{i \in [n]})
        Server
              \hat{s}_{v^2} = \text{Decode}(e)
```





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Experiment



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