

# CS 1550

Week 5 – Synchronization with xv6

Teaching Assistant

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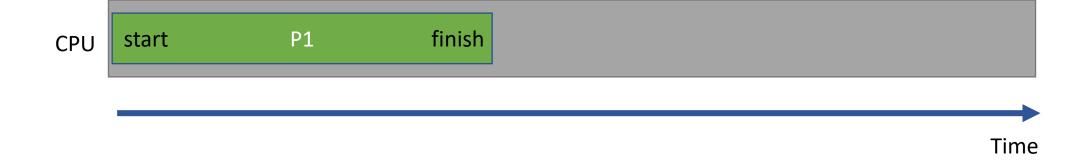
(Slides credited to Henrique Potter)

#### Keep in mind the different qemu

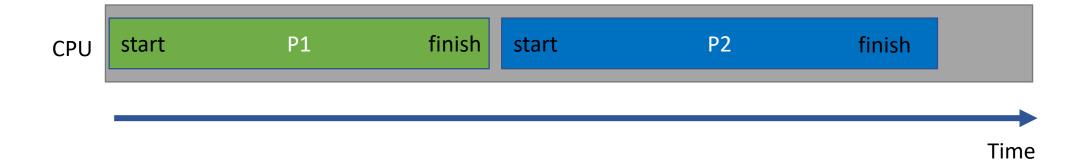
- qemu with xv6 (Labs) Refer to Lab 1 if needed!
- qemu-x86 i386 (Project 1 and 2)







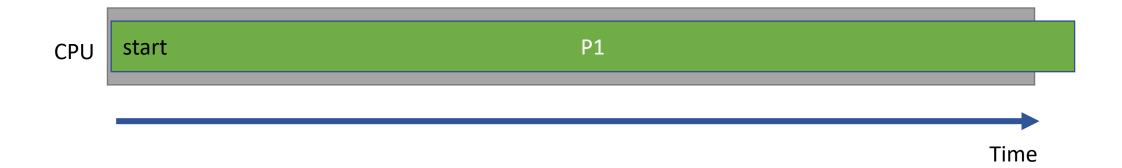
• OS chooses another processes to execute once the first finishes



• OS chooses another processes to execute once the first finishes

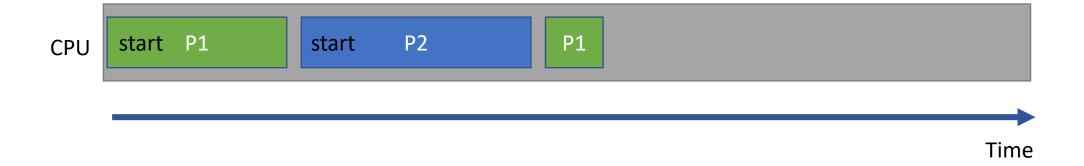


What if P1 is a big process?



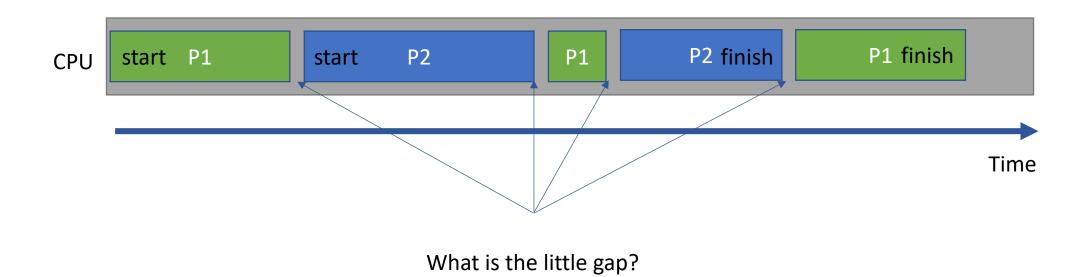


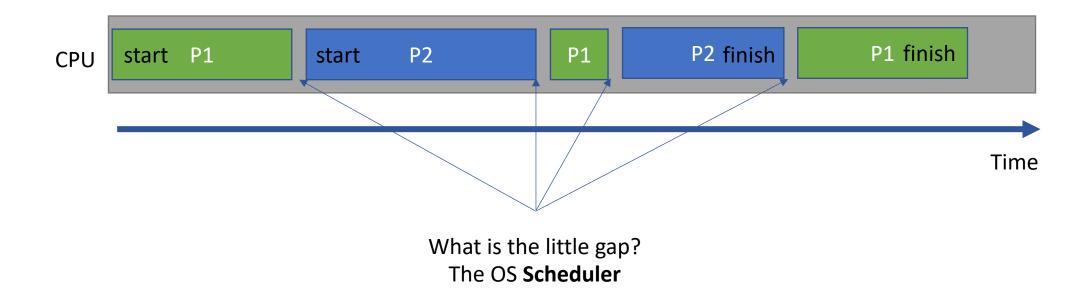








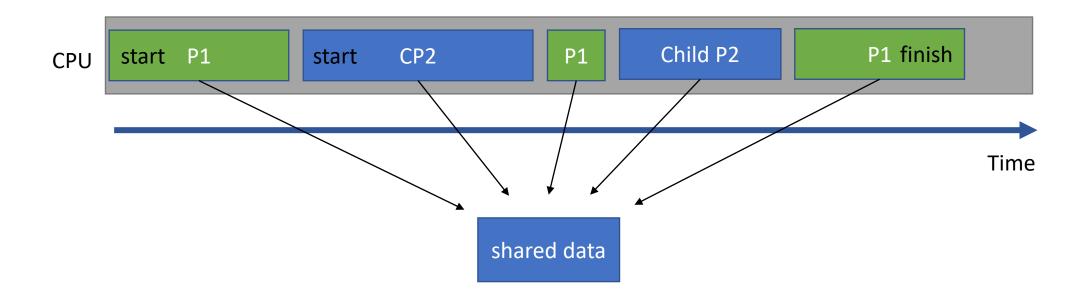




What happens in Parent-Child Process scenario?



- What happens in Parent-Child Process scenario?
- How to keep integrity/correctness on race conditions?



```
struct list {
  int data;
  struct list *next;
};
```

```
struct list {
   int data;
   struct list *next;
};

struct list *list = 0;
```

```
struct list {
  int data;
  struct list *next;
};
struct list *list = 0;
void
insert(int data) {
      struct list *1;
      1 = malloc(sizeof *1);
      1->data = data;
      1->next = list;
      list = 1;
```

```
struct list {
  int data;
  struct list *next;
};
                                       CPU
                                              P1
struct list *list = 0;
void
insert(int data) {
      struct list *1;
                                       P1 stops here the
      1 = malloc(sizeof *1);
                                       OS switches to P2
      1->data = data;
      l->next = list;
      list = 1;
```

```
struct list {
         int data;
         struct list *next;
       };
                                                                  CP2
                                                 CPU
                                                         P1
       struct list *list = 0;
       void
       insert(int data) {
              struct list *1;
                                                   P2 gets the same
              1 = malloc(sizeof *1);
                                                   reference to the
P1 stopped
              1->data = data;
                                                   same block of
              l->next = list;
                                                   data of list and
              list = 1;
                                                   overwrites it
```

```
struct list {
        int data;
        struct list *next;
      };
      struct list *list = 0;
      void
      insert(int data) {
            struct list *1;
             1 = malloc(sizeof *1);
             l->data = data;
CP2 stopped
             l->next = list;
             list = 1;
```

CPU P1 CP2 P1

When P1 comes back it will have written the wrong data

- Sharing CPU among processes
- Ensuring data integrity/correctness
- Ensure that a critical section of your code is only executed by one process

```
struct list *list = 0;
struct lock listlock;
void
insert(int data)
      struct list *1;
      acquire(&listlock);
      1 = malloc(sizeof *1);
      1->data = data;
      1->next = list;
      list = 1;
      release(&listlock);
```

CPU P1 CP2 P1

```
struct list *list = 0;
struct lock listlock;
void
insert(int data)
      struct list *1;
      acquire(&listlock);
      1 = malloc(sizeof *1);
      l->data = data;
      l->next = list;
      list = 1;
      release(&listlock);
```

CPU P1 CP2 P1

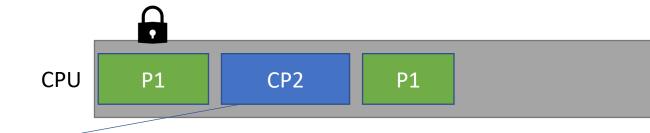
```
struct list *list = 0;
struct lock listlock;
void
insert(int data)
                                       CPU
                                                       CP2
                                              P1
                                                                P1
      struct list *1;
      acquire(&listlock);
      1 = malloc(sizeof *1);
                                              P1 gets locks the lock
      l->data = data;
      l->next = list;
      list = 1;
      release(&listlock);
```

```
struct list *list = 0;
struct lock listlock;
void
insert(int data)
                                       CPU
                                                       CP2
                                              P1
                                                                P1
      struct list *1;
      acquire(&listlock);
      1 = malloc(sizeof *1);
                                              P1 gets locks the lock
      l->data = data;✓
      l->next = list;
      list = 1;
      release(&listlock);
```

release(&listlock);

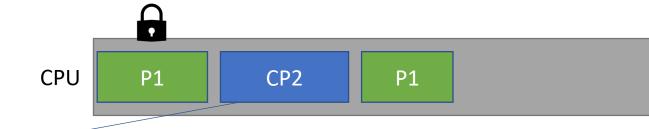
```
struct list *list = 0;
  struct lock listlock;
  void
   insert(int data)
                                                          CP2
                                          CPU
                                                                   P1
         struct list *1; *
         acquire(&listlock);
         1 = malloc(sizeof *1);
                                               When the OS schedule CP2
        1->data = data;
P1 stopped
         1->next = list;
         list = 1;
```

```
struct list *list = 0;
  struct lock listlock;
  void
  insert(int data)
         struct list *1;
         acquire(&listlock);
         l = malloc(sizeof *1);
        l->data = data;
P1 stopped
         1->next = list;
         list = 1;
        release(&listlock);
```



It will try to get the lock but won't.

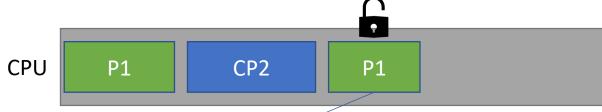
```
struct list *list = 0;
  struct lock listlock;
  void
  insert(int data)
         struct list *1;
         acquire(&listlock);
         1 = malloc(sizeof *1);
P1 stopped
        l->data = data;
         l->next = list;
         list = 1;
         release(&listlock);
```



It will try to get the lock but won't.

It will be constantly try to get it (in a loop). Until the OS switches back to P1

```
struct list *list = 0;
   struct lock listlock;
   void
    insert(int data)
          struct list *1;
          acquire(&listlock);
CP2 stopped
          1 = malloc(sizeof *1);
          l->data = data;
          l->next = list;
          list = 1;
          release(&listlock);
```



P1 release the lock P2 will finally be able to execute, once scheduled

```
struct list *list = 0;
   struct lock listlock;
   void
    insert(int data)
          struct list *1;
CP2 proceeds acquire(&listlock);
          1 = malloc(sizeof *1);
          l->data = data;
          l->next = list;
          list = 1;
          release(&listlock);
```



P1 release the lock P2 will finally be able to execute, once scheduled

#### SpinLock

```
Void
acquire(struct spinlock *lk)
{
    for(;;) {
        if(!lk->locked) {
            lk->locked = 1;
            break;
        }
    }
}
```

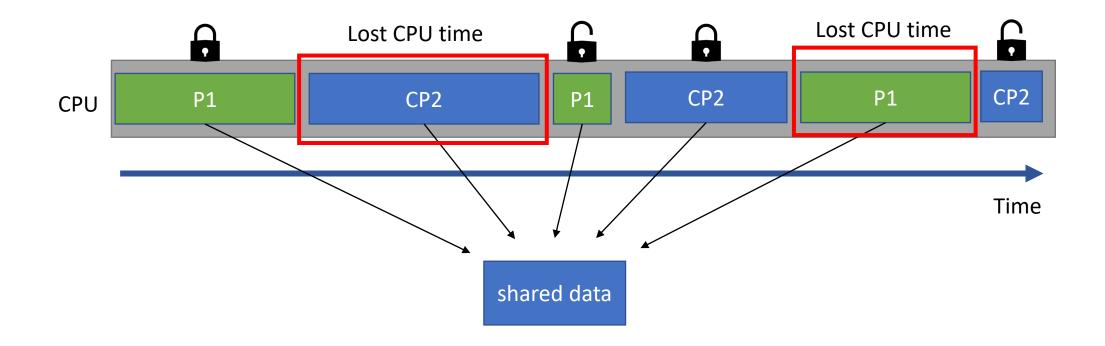
- Keep spinning until find lock is released
- But we can have the same issue as before
- We need to check and lock atomically

- Atomically check and change a register value
  - xchg(&lk->locked, 1)

- Swap a word in memory with the contents of a register
- In acquire function:
  - loop xchg instruction
  - Each round atomically read lock and set the lock to 1

```
void
acquire(struct spinlock *lk)
    pushcli(); // disable interrupts to
avoid deadlock.
// The xchq is atomic.
    while (xchg(&lk->locked, 1) != 0);
// Record info about lock acquisition for
debugging.
    1k->cpu = mycpu();
    getcallerpcs(&lk, lk->pcs);
```

- But the we have another issue
  - Busy waiting



- Sleep Locks
  - For code need to hold a lock for a long time (read/write to disk)
- Avoids the schedule of "spin locked" processes

- Sleep Locks
  - For code need to hold a lock for a long time (read/write to disk)
- Avoids the schedule of "spin locked" processes

```
void
acquiresleep(struct sleeplock *lk)

{
    acquire(&lk->lk);
    while (lk->locked) {
        sleep(lk, &lk->lk);
        lk->locked = 0;
        sleep(lk, &lk->lk);
        lk->pid = 0;
        wakeup(lk);
        release(&lk->lk);
    }
    release(&lk->lk);
}
```

- Sleep Locks
  - For code need to hold a lock for a long time (read/write to disk)
- Avoids the schedule of "spin locked" processes

```
void
acquiresleep(struct sleeplock *lk)
{
    acquire(&lk->lk);
    while (lk->locked) {
        sleep(lk, &lk->lk);
        lk->locked = 0;
        sleep(lk, &lk->lk);
        lk->pid = 0;
        wakeup(lk);
        release(&lk->lk);
    }
    release(&lk->lk);
}
```

 Put one process to sleep waiting for event

Mark current process as sleeping

Call sched() to release the processor

```
void
sleep(void *chan, struct spinlock *lk)
{
   struct proc *p = myproc();
   ...
   p->state = SLEEPING;
   sched();
   ...
}
```

 Put one process to sleep waiting for event

Mark current process as sleeping

Call sched() to release the processor

```
void
sleep(void *chan, struct spinlock *lk)
  struct proc *p = myproc();
  if(p == 0)
    panic("sleep");
  if(lk == 0)
    panic("sleep without lk");
  if(lk != &ptable.lock){
    acquire(&ptable.lock);
    release(lk);
  p->chan = chan;
  p->state = SLEEPING;
  sched();
  p->chan = 0
  if(lk != &ptable.lock){
    release(&ptable.lock);
    acquire(lk);
```

- Wake up process when event happened
- Mark a waiting process as runnable

- Who needs to be a syscall?
  - SpinLocks
  - SleepLocks

#### CS 1550 – Lab exercise 2

#### PROCESS SYNCHRONIZATION IN XV6

- **Due**: Friday, February 22, 2019 @11:59pm
- Part 2 step 5: user.h
  - Add declaration for init\_lock()
    - void init lock(struct spinlock \*);
  - struct condvar;
  - struct spinlock;
- Part 3 step 8: defs.h
  - Add declaration for sleep1()



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