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## Tutorial - 5

Ques 1

BFS

- Stands for Breadth first search

- BFS uses queue to find the shortest path

- BFS is better when target is closer to source

- If BFS consider all neighbours so it is not suitable for decision tree used in puzzle games

- BFS is slower than DFS

DFS

- Stands for Depth first Search

- DFS uses stack to find shortest path

- DFS is better when target is far from source

- DFS is more suitable for decision tree. If with one decision we need to traverse further to argument the decision. If we search the conclusion

- DFS is faster than BFS

## Application of DFS

- Using DFS we can find path between two vertices.
- We can perform topological sorting which is used to scheduling jobs.
- We can use DFS to detect cycles.
- Using DFS, we can find strongly connected components of a graph.

## Application BFS:

- BFS may also used to detect cycles.
- Finding shortest path and minimal Spanning tree in unweighted graphs.
- In networking finding a route for packet transmission.
- Finding a route through GPS navigation system.

Ques 2

Sol: Breadth First Search (BFS) uses Queue data structure. In BFS you mark any node in the graph as source node and start traversing. BFS visited an adjacent node, marks it as done and insert it into Queue.

DFS uses stack data structure because DFS traverse a graph in a depth first motion and uses stack to remember upto get the next vertex to start a search, when a dead end occurs in any iteration.

Ques 3

Sol: Sparse Graph:

A graph in which the number of edges is much less than the possible number of edges.

Dense Graph:

A dense graph is a graph in which the number of edges is close to the maximal no. of edges of edges.

→ If the graph is sparse, we should store it as list of edges.

Alternatively if a graph is dense, we should store it as a adjacency matrix.

#### Ques 4

Soln → DFS can be used to detect cycle in a Graph. DFS for a Connected Graph produces a tree. Three back edge are present in the graph. A back edge is an edge that is from a node to itself or one of its ancestor in the tree produced by DFS.

BFS can also be used to detect cycles. Just perform BFS while keeping a list of previous nodes to at each node visited or else construct a tree from the starting node. If I visit a node that is already marked by BFS, I found a cycle.

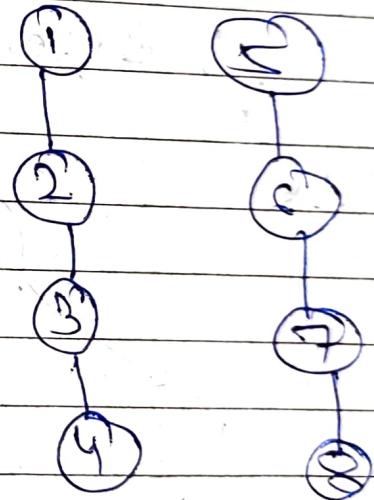
Ques

## Soln: Disjoint Set Data Structure:

- It allows to find out whether the two elements are in the same set or not efficiently.
- A disjoint set can be defined as the subsets when there is no common element between the two sets.

E.g.  $S_1 = \{1, 2, 3, 4\}$

$S_2 = \{5, 6, 7, 8\}$



## Operations performed:

(i) find:

```
g int find (int v)
```

```
if (v == parent[v])
```

```
    return v;
```

```
    return parent[v] = find (parent[v]);
```

```
}
```

Union :

void union (int a, int b)

{

a = find (a)

b = find (b)

if (a != b)

{

if (size[a] &lt; size[b])

{

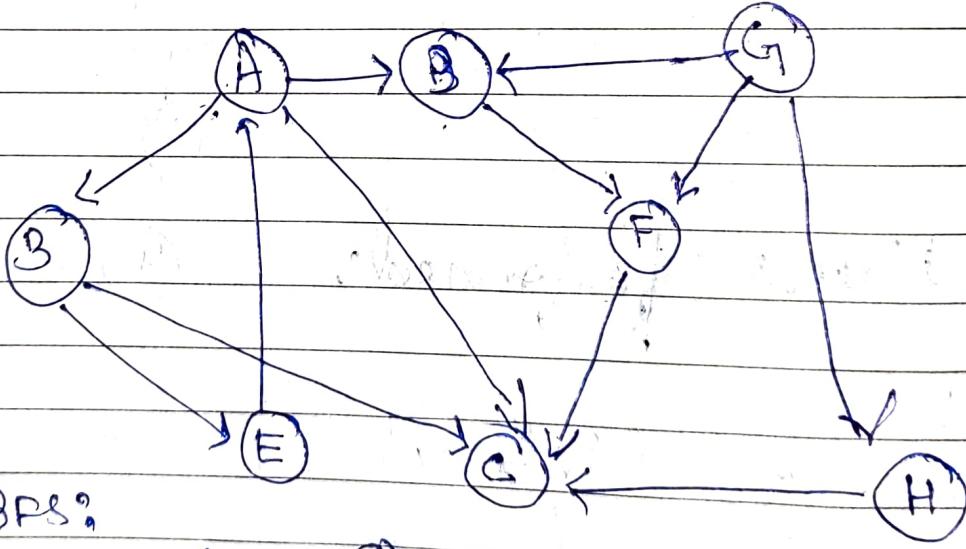
swap (a, b) }

parent [b] = a ;

size [a] += size [b] ;

}

Ques.

Soln

BFS:

Node :	(B)	(E)	(C)	(A)	(G)	(F)
parent :	-	B	B	E	A	D

Path : B → E → A → D → F

DFS:

Node processed  
StackB    B    C    E    A    D    F  
B    CE    EE    AE    DE    FE    E

path: B → C → E → A → D → F

Ques:

$$\text{Soln} \rightarrow V = \{a\} \cup \{b\} \cup \{c\} \cup \{d\} \cup \{e\} \cup \{f\} \cup \{g\} \cup \{h\} \cup \{i\} \cup \{j\}$$

$$E = \{a, b\}, \{a, c\}, \{b, c\}, \{b, d\}, \{c, e\}, \{c, f\}, \{e, g\}, \{h, i\}, \{h, j\}$$

(a,b) {a, b} {c} {d} {e} {f} {g} {h} {i} {j}

(a,c) {a, b, c} {d} {e} {f} {g} {h} {i} {j}

(b,c) {a, b, c} {d} {e} {f} {g} {h} {i} {j}

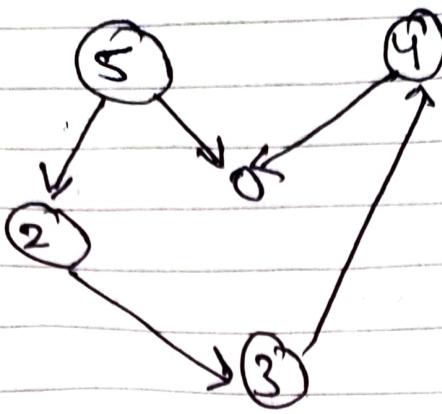
(b,d) {a, b, c, d} {e} {f} {g} {h} {i} {j}

(e,f) {a, b, c, d} {e, f} {g} {h} {i} {j}

(h,j) {a, b, c, d} {e, f, g} {h, i} {j}

No. of Components connected = 3

Ques 8



Adjacency list

Visited

0 →  
1 →  
2 → 3  
3 → 1  
4 → 0, 1  
5 → 2, 0

0	1	2	3	4	5
False	False	False	False	False	False

Stack (empty)

Step 1: Topological sort(0), visited[0] = true

Stack [0]

Step 2: Topological sort(1), visited[1] = true

Stack [0] | 1

Step 3: Topological sort(2), visited[2] = true

Topological sort(3), visited[3] = true.

Stack [0] | 1 | 3 | 2

Step 4:

Stack 

0	1	3	2	4
---	---	---	---	---

Step 5:

Stack 

0	1	3	2	4	5
---	---	---	---	---	---

Step 6: Print all elements of stack from top to bottom.

→ 5, 4, 2, 3, 1, 0

Ques

Algorithms that use priority Queue:

i) Dijkstra's Algorithm -

When a graph is stored in the form of list or matrix, priority queue can be used to extract minimum efficiency when implementing Dijkstra's Algorithm.

ii) Prim's Algorithm:

It is used to implement prim's algorithm to store key of nodes to extract minimum key node at every step.

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iii) Date Compression:

It is used in Huffman's code which is used to compress data.

Ques 10

Sol →

Min Heap

Max Heap

- In min heap the key present at root node must be less than or equal to among the keys present at all its children.
- Uses the ascending priority.
- The minimum key present at the root node.

- In max-heap the key present at root node must be greater or equal to the key present at all its children.
- Uses descending priority.
- The minimum key present at the root node.