

WEEK 5

# TREES, TRIES, HEAPS & GRAPHS

CSE 492 J

# GRAPHS & TREES

## GRAPHS & TREES

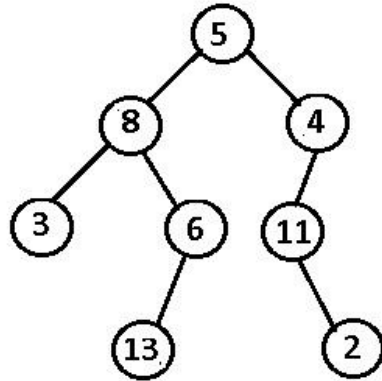
- > Trees:
  - Binary Tree, Binary Search Tree, Balanced Binary Tree(AVL), MST, Heaps, Tries.
- > Graph Algorithms:
  - Breadth-First Search, Depth-First Search, Dijkstra's, Bellman-Ford, Network Flow.



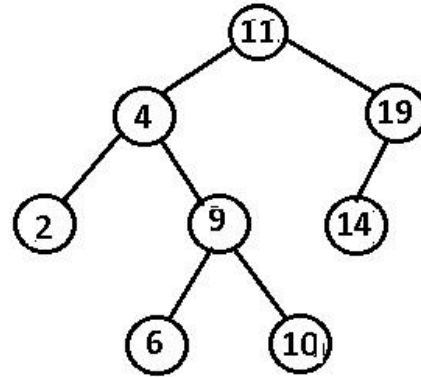
## GRAPHS & TREES

- > Binary Tree, Binary Search Tree, Balanced Binary Tree(AVL)
- > MST, Heaps, Tries
- > Breadth-First Search, Depth-First Search
- > Dijkstra's
- > Bellman-Ford, Network Flow

## BINARY TREE VS BINARY SEARCH TREE



Binary Tree



Binary Search Tree

# AVL Trees

Operation	Case	Runtime
containsKey(key)	best	$\Theta(1)$
	worst	$\Theta(\log n)$
insert(key)	best	$\Theta(\log n)$
	worst	$\Theta(\log n)$
delete(key)	best	$\Theta(\log n)$
	worst	$\Theta(\log n)$

## PROS

All operations on an AVL Tree have a logarithmic worst case

- Because these trees are always balanced!

The act of rebalancing adds no more than a constant factor to insert and delete

➤ Asymptotically, just better than a normal BST!

## CONS

- Relatively difficult to program and debug (so many moving parts during a rotation)
- Additional space for the height field
- Though asymptotically faster, rebalancing *does* take some time
  - Depends how important every little bit of performance is to you

# Review: Dictionaries

## Dictionary ADT

### state

Set of items & keys  
Count of items

### behavior

put(key, item) add item to collection indexed with key  
get(key) return item associated with key  
containsKey(key) return if key already in use  
remove(key) remove item and associated key  
size() return count of items

Why are we so obsessed with Dictionaries?

When dealing with data:

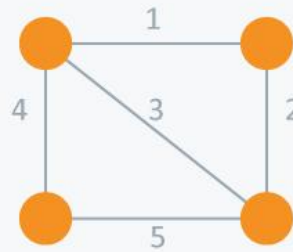
- Adding data to your collection
- Getting data out of your collection
- Rearranging data in your collection

Operation		ArrayList	LinkedList	HashTable	BST	AVLTree
put(key,value)	best	$\Theta(1)$	$\Theta(1)$	$\Theta(1)$	$\Theta(1)$	$\Theta(1)$
	worst	$\Theta(n)$	$\Theta(n)$	$\Theta(n)$	$\Theta(n)$	$\Theta(\log n)$
get(key)	best	$\Theta(1)$	$\Theta(1)$	$\Theta(1)$	$\Theta(1)$	$\Theta(1)$
	worst	$\Theta(n)$	$\Theta(n)$	$\Theta(n)$	$\Theta(n)$	$\Theta(\log n)$
remove(key)	best	$\Theta(1)$	$\Theta(1)$	$\Theta(1)$	$\Theta(1)$	$\Theta(\log n)$
	worst	$\Theta(n)$	$\Theta(n)$	$\Theta(n)$	$\Theta(n)$	$\Theta(\log n)$

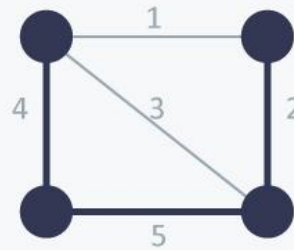


# MST(MINIMUM SPANNING TREE)

## Kruskal's vs Prim's

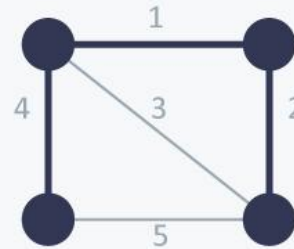


Undirected  
Graph



Spanning  
Tree

Cost =  $11(=4+5+2)$

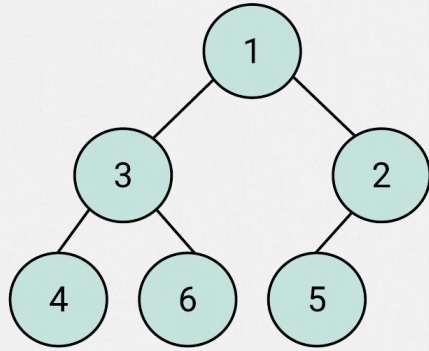


Minimum Spanning  
Tree

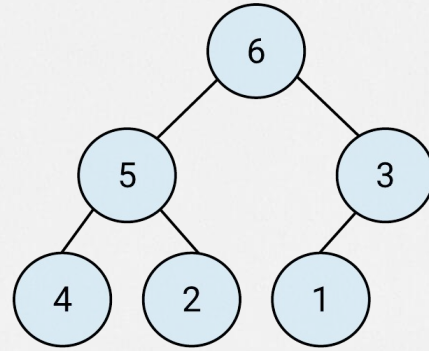
Cost =  $7(=4+1+2)$



# HEAPS

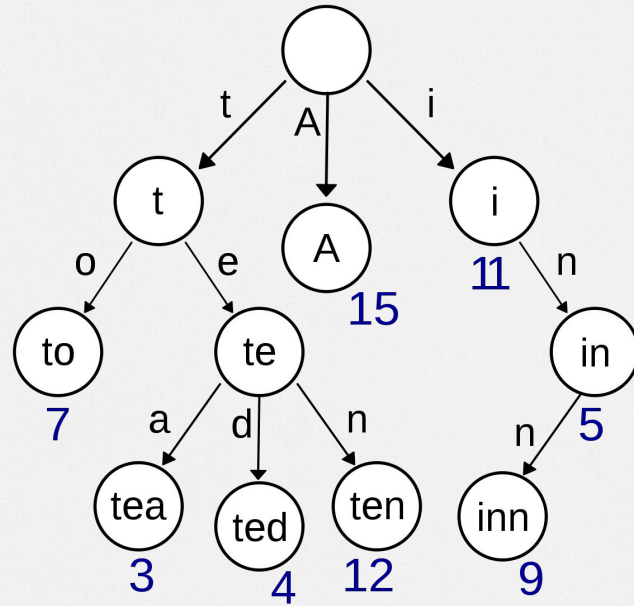


Min heap



Max Heap

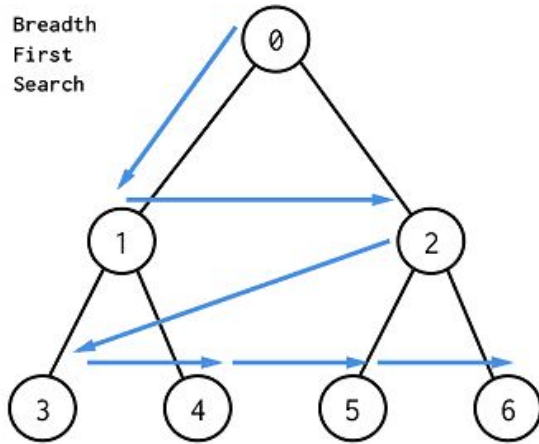
# TRIE



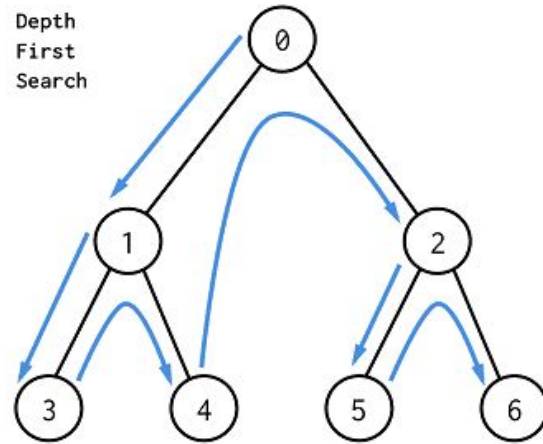


## BFS vs DFS

Breadth  
First  
Search

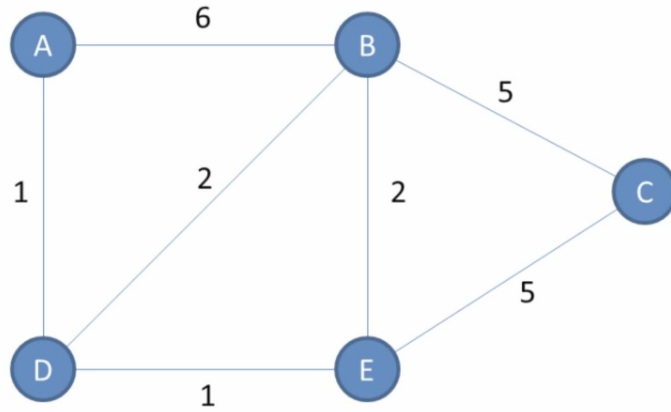


Depth  
First  
Search





# DIJKSTRA'S



Vertex	Shortest distance from A	Previous vertex
A	0	
B	6	A
C	$\infty$	
D	1	A
E	$\infty$	

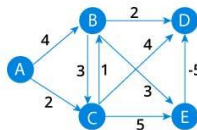
Visited = [A]

Unvisited = [B, C, D, E]

# BELMAN-FORD

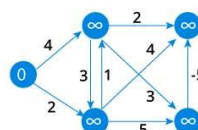
1

Start with a weighted graph



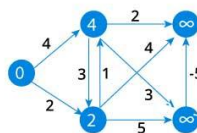
2

Choose a starting vertex and assign infinity path values to all other vertices



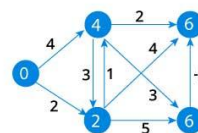
3

Visit each edge and relax the path distances if they are inaccurate



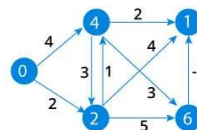
4

We need to do this V times because in the worst case, a vertex's path length might need to be readjusted V times



5

Notice how the vertex at the top right corner had its path length adjusted



6

After all the vertices have their path lengths, we check if a negative cycle is present.

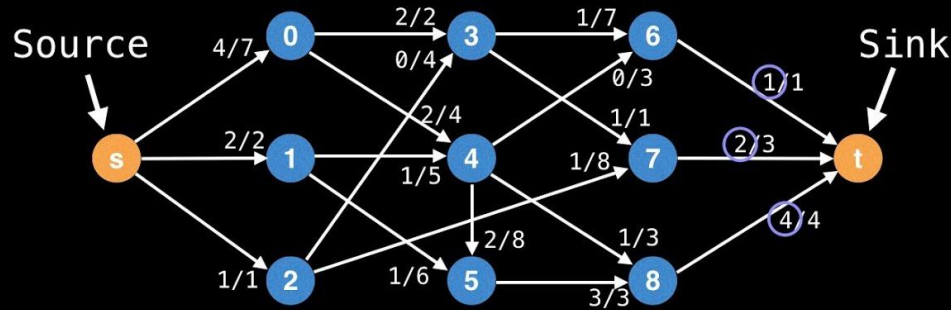
A	B	C	D	E
0	∞	∞	∞	∞
0	4	2	∞	∞
0	3	2	6	6
0	3	2	1	6
0	3	2	1	6



# NETWORK FLOW

## Max flow

Q: With an infinite input source, how much “flow” can we push through the network given that each edge has a certain capacity?







LET'S PRACTICE  
WITH CEO-IT!

## QUESTION

Given an integer array `nums` and an integer `k`, return the `k`th largest element in the array.

What would we use to solve this problem?

## CLARIFY

1. Does  $k$  refer to the  $k$ th largest distinct element?
  - a. No, it is the  $k$ th largest element \*after\* sorting the input array.
2. Can we use additional data structures?
  - a. Yes, assume we want the fastest runtime.
3. Can there be duplicate/non-positive integers in the input array `nums`?
  - a. Yes.
4. Can the value of  $k$  be greater than the length of the input array `nums`?
  - a. No. The length of `nums` is greater than  $k$ , which is greater than 1.



## EXAMPLE

**Middle Case:**

`nums = [3, 2, 1, 5, 6, 4]; k = 2`

`Output = 5`

**Edge Case 1:**

`nums = [-45]; k = 1`

`Output = -45`

## EXAMPLE

Edge Case 2:

`nums = [3, 2, 3, 1, 2, 4, 5, 5, 6]; k = 4`

Output = 4

Notice the duplicates!

(This is why clarifying is important, interviewers may or may not care about duplicates...Ask, do not read minds although that would be cool)

## APPROACH

Brute Force:

1. Use an  $O(n * \log n)$  sorting algorithm
2. Sort the entire array
3. Traverse through the array and return the  $k$ th element

There is a faster way using a data structure we discussed today...



## OPTIMIZE

- > Optimized: Use a Heap  $O(n)$  for adding elements of input array to Heap +  $O(\log n)$  for finding Kth largest element in Heap =  $O(n)$
- > Optimized no additional data structure: use an in-place sorting algorithm with  $O(n * \log n)$  runtime
- > Note: When you hear "find the largest/smallest.." in an interview, think about using a heap if applicable.

## IMPLEMENT

- > There are no sorting algorithms that run in  $O(n)$  runtime. So we create a MinHeap using a **Priority Queue** to have  $O(n)$  runtime.
- > Iterate through `nums` and add its elements into our heap. For every iteration, make sure our heap size is not greater than `k`, otherwise remove root of MinHeap.
- > Our result will be the final root of our MinHeap after iterating through all of `nums`.



## IMPLEMENT - JAVA CODE

```
public int findKthLargest(int[] nums, int k) {  
    PriorityQueue<Integer> minHeap = new PriorityQueue<Integer>();  
    for (int i: nums) {  
        minHeap.add(i);  
        if (minHeap.size() > k) {  
            minHeap.remove();  
        }  
    }  
    return minHeap.remove();  
}
```

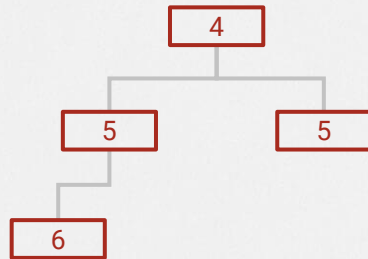


## TEST

Test with Edge Case with Duplicates:

`nums = [3, 2, 3, 1, 2, 4, 5, 5, 6]; k = 4`

Resulting MinHeap after for loop ends:



Remove root element to  
get answer

Output  
= 4

## EXTRA PRACTICE PROBLEMS FOR PATTERN-BUILDING!

### Binary Trees (Easy):

- > <https://leetcode.com/problems/validate-binary-search-tree/>
- > <https://leetcode.com/problems/path-sum/>
- > <https://leetcode.com/problems/maximum-difference-between-node-and-ancestor/>

### Binary Trees (Medium):

- > <https://leetcode.com/problems/path-sum-ii/>
- > <https://leetcode.com/problems/binary-search-tree-iterator/>



## EXTRA PRACTICE PROBLEMS FOR PATTERN-BUILDING!

### Graphs (Easy):

- > <https://leetcode.com/problems/number-of-islands/>
- > <https://leetcode.com/problems/rotting-oranges/>
- > <https://leetcode.com/problems/accounts-merge/>
- > <https://leetcode.com/problems/evaluate-division/>
- > <https://leetcode.com/problems/word-ladder/>
- > <https://leetcode.com/problems/all-paths-from-source-to-target/>



# EXTRA PRACTICE PROBLEMS FOR PATTERN-BUILDING!

## Graphs (Medium):

- > <https://leetcode.com/problems/reconstruct-itinerary/>
- > <https://leetcode.com/problems/binary-tree-right-side-view/>
- > <https://leetcode.com/problems/pacific-atlantic-water-flow/>
- > <https://leetcode.com/problems/clone-graph/>
- > <https://leetcode.com/problems/course-schedule/>
- > <https://leetcode.com/problems/number-of-provinces/>
- > <https://leetcode.com/problems/symmetric-tree/>

## EXTRA PRACTICE PROBLEMS FOR PATTERN-BUILDING!

### Heaps (Easy):

- > <https://leetcode.com/problems/kth-largest-element-in-a-stream>
- > <https://leetcode.com/problems/last-stone-weight/>
- > <https://leetcode.com/problems/maximum-product-of-two-elements-in-an-array>

### Tries (Medium):

- > <https://leetcode.com/problems/search-suggestions-system/>
- > <https://leetcode.com/problems/map-sum-pairs/>