1 Homework sheet 1 - Normal form games

1. Represent the following game in normal form:

Alice, Bob and Celine are childhood friends that would like to communicate online. Alive likes facebook, Bob likes twitter and Celine like G+.

Clearly state the players, strategy sets and interpretations of the utilities.

Solution

The set of players is $\{A, B, C\}$. The strategy space is $S_A = S_B = S_C = \{\text{fb}, t, G\}$.

The payoff functions are given below:

$$u_i(s_1, s_2, s_3) = |\{r \in s \setminus s_i \mid r = s_i\}|$$

(ie the utility of player i is equal to the number of players that have picked the same strategy as player i.)

Here is the normal form representation:

• $s_A = \text{fb}$:

$$\begin{pmatrix} (2,2,2) & (1,1,0) & (1,1,0) \\ (1,0,1) & (0,1,1) & (0,0,0) \\ (1,0,1) & (0,0,0) & (0,1,1) \end{pmatrix}$$

• $s_A = t$:

$$\begin{pmatrix} (0,1,1) & (1,0,1) & (0,0,0) \\ (1,1,0) & (2,2,2) & (1,1,0) \\ (0,0,0) & (1,0,1) & (0,1,1) \end{pmatrix}$$

• $s_A = G$:

$$\begin{pmatrix} (0,1,1) & (0,0,0) & (1,0,1) \\ (0,0,0) & (0,1,1) & (1,0,1) \\ (1,1,0) & (1,1,0) & (2,2,2) \end{pmatrix}$$

2. Represent the following game in normal form:

Assume two neighbouring countries have at their disposal very destructive armies. If both countries attack each other the countries' civilian population will suffer 10 thousand casualties. If one country attacks whilst the other remains peaceful, the peaceful country will lose 15 thousand casualties but would also retaliate causing the offensive country 13 thousand casualties. If both countries remain peaceful then there are no casualties.

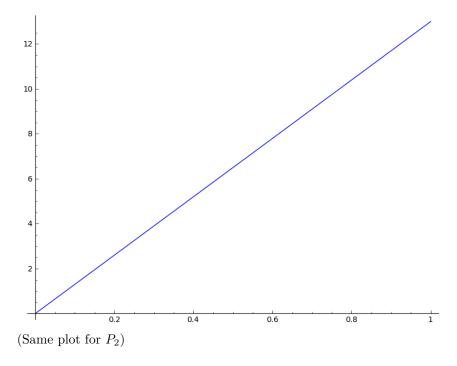
- Clearly state the players and strategy sets.
- Plot the utilities to both countries assuming that they play a mixed strategy while the other country remains peaceful.
- Plot the utilities to both countries assuming that they play a mixed strategy while the other country attacks.

Solution

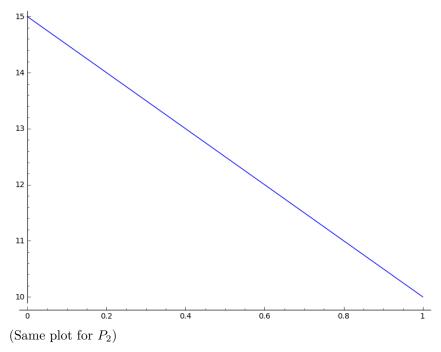
Players: $\{P_1, P_2\}$. The strategy space is $\{A, P\}$. We take the utility to be the number of casualties suffered (players aim to minimize their utility):

$$\begin{pmatrix} (10,10) & (13,15) \\ (15,13) & (0,0) \end{pmatrix}$$

$$u_1((x, 1-x), (0,1)) = 13x$$



 $u_1((x, 1-x), (1,0)) = 10x + 15(1-x) = 15 - 5x$



3. Dominance

Attempt to predict rational behaviour using iterated elimination of dominated strategies for the following:

 $\begin{pmatrix} (2,1) & (1,1) \\ (1,1) & (1,3) \end{pmatrix}$

Solution

We see that r_1 weakly dominates r_2 so we have:

Thus a predicted strategy profile is (r_1, s_1) .

Importantly, we see however that s_2 weakly dominates s_1 which would give:

$$\begin{pmatrix} (1,1) \\ (1,3) \end{pmatrix}$$

Thus another predicted strategy profile is (r_2, s_2) .

$$\begin{pmatrix} (2,11) & (1,9) & (3,10) & (17,22) \\ (27,0) & (3,1) & (1,1) & (1,0) \\ (4,2) & (6,10) & (7,12) & (18,0) \end{pmatrix}$$

Solution

We see that s_2 is weakly dominated by s_3 so we have:

$$\begin{pmatrix} (2,11) & (3,10) & (17,22) \\ (27,0) & (1,1) & (1,0) \\ (4,2) & (7,12) & (18,0) \end{pmatrix}$$

Now r_3 strictly dominates r_1 so we have:

$$\begin{pmatrix} (27,0) & (1,1) & (1,0) \\ (4,2) & (7,12) & (18,0) \end{pmatrix}$$

Now s_3 strictly dominates s_1 and s_4 so we have:

$$\begin{pmatrix} (1,1) \\ (7,12) \end{pmatrix}$$

Thus the predicted rational behaviour is (r_3, s_3) .

$$\begin{pmatrix} (3,2) & (3,1) & (2,3) \\ (2,2) & (1,3) & (3,2) \end{pmatrix}$$

Solution

There are no dominated strategies in this game.

$$\begin{pmatrix} (3,-3) & (-1,1) \\ (2,1) & (7,-6) \end{pmatrix}$$

Solution

There are no dominated strategies in this game.

Explain when games occur that cannot be handled this way.

4. For all of the games of question 3, identify all best responses and attempt to predict rational behaviour.

Solution

$$\begin{pmatrix} (\underline{2},\underline{1}) & (\underline{1},\underline{1}) \\ (1,1) & (\underline{1},\underline{3}) \end{pmatrix}$$

We have 3 pairs of best responses: $(r_1, s_1), (r_1, s_2), (r_2, s_2)$.

$$\begin{pmatrix} (2,11) & (1,9) & (3,10) & (17,\underline{22}) \\ (\underline{27},0) & (3,\underline{1}) & (1,\underline{1}) & (1,0) \\ (4,2) & (\underline{6},10) & (\underline{7},\underline{12}) & (\underline{18},0) \end{pmatrix}$$

We have a single pair of best responses: (r_3, s_2) .

$$\begin{pmatrix} (\underline{3},2) & (\underline{3},1) & (2,\underline{3}) \\ (2,2) & (1,\underline{3}) & (\underline{3},2) \end{pmatrix}$$

There are no pairs of best responses.

$$\begin{pmatrix} (\underline{3},-3) & (-1,\underline{1}) \\ (2,\underline{1}) & (\underline{7},-6) \end{pmatrix}$$

There are no pairs of best responses.

Explain when games occur that cannot be handled this way.

5. Consider the following game:

$$\begin{pmatrix}
(7,3) & (0,2) \\
(2,1) & (6,1) \\
(4,0) & (4,2)
\end{pmatrix}$$

Compute directly B_1, B_2, UD_1 and UD_2 .

Solution

By definition:

$$UD_1 = \{s \in \{r_1, r_2, r_3\} \mid s \text{ is not strictly dominated}\}$$

Thus $UD_1 = \{r_1, r_2, r_3\}$

Similarly:

$$UD_2 = \{s \in \{s_1, s_2\} \mid s \text{ is not strictly dominated}\}$$

Thus $UD_2 = \{s_1, s_2\}$

By definition

$$B_1 = \{s \in \{r_1, r_2, r_3\} | \exists \sigma \in \Delta S_2 \text{ such that } s \text{ is a best response to } \sigma\}$$

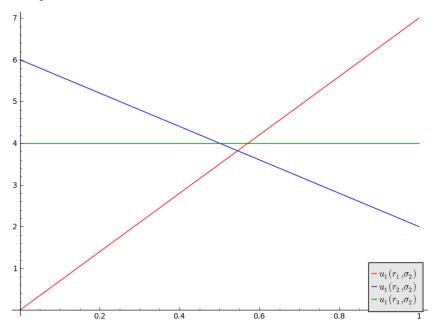
Let us assume that player 2 plays $\sigma_2 = (x, 1 - x)$. This gives:

$$u_1(r_1, \sigma_2) = 7x$$

$$u_1(r_2, \sigma_2) = 2x + 6(1 - x) = 6 - 4x$$

 $u_1(r_3, \sigma_2) = 4x + 4(1 - x) = 4$

If we plot these utilities:



we see that all strategies are a best response to a strategy in ΔS_2 : $B_1 = \{r_1, r_2, r_3\}$.

By definition

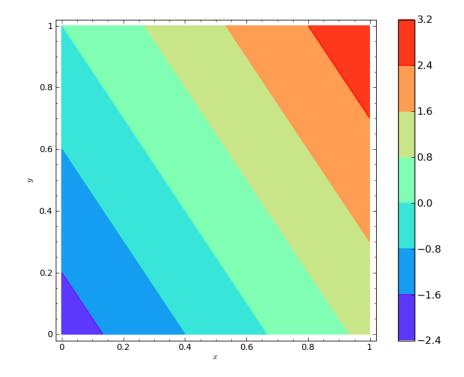
$$B_2 = \{s \in \{s_1, s_2\} | \exists \sigma \in \Delta S_1 \text{ such that } s \text{ is a best response to } \sigma\}$$

Let us assume that player 1 plays $\sigma_1 = (x, y, 1 - x - y)$. This gives:

$$u_2(\sigma_1, s_1) = 3x + y$$
$$u_2(\sigma_1, s_2) = 2x + y + 2 - 2x - 2y = 2 - y$$

We see that $u_2(\sigma_1, s_1) > u_2(\sigma_1, s_2) \Rightarrow 3x + y > 2 - y \Rightarrow x > \frac{2-2y}{3}$. Thus (finding an obvious example) $\sigma_1 = (2/3, 0, 1/3) \in \Delta S_1$ has s_1 as a best response and similarly $\sigma_1 = (1/3, 0, 2/3) \in \Delta S_2$ has s_2 as a best response. This gives $B_2 = \{s_1, s_2\}$.

To illustrate this further here is a contour plot of $u_2(\sigma_1, s_1) - u_2(\sigma_1, s_2)$:



6. In the notes the following theorem is given:

In a 2 player normal form game $B_i = UD_i$ for all $i \in \{1, 2\}$.

Prove the theorem for 2 player games with $|S_1|=|S_2|=2$. I.e. prove the above result in the special case of 2×2 games.

Solution

Let us consider the 2×2 game:

$$\begin{pmatrix} (a_{11}, b_{11}) & (a_{12}, b_{12}) \\ (a_{21}, b_{21}) & (a_{22}, b_{22}) \end{pmatrix}$$

We prove the result for i=1 without loss of generality. Let us consider the two following cases:

$$UD_1 = \{r_1\} \text{ and } UD_1 = \{r_1, r_2\}$$

• If $UD_1 = \{r_1\} \Rightarrow a_{1j} > a_{2j} \text{ for } j \in \{1, 2\}$. This gives:

$$u_1(r_1, \sigma_2) = a_{11}x + a_{12}(1-x) > a_{12}x + a_{22}(1-x) = u_1(r_2, \sigma_2)$$

So r_2 is never a best response giving $B_1 = \{r_1\}$ as required.

• If $UD_1 = \{r_1, r_2\}$ this implies that one of the following must hold:

$$a_{11} > a_{21}$$
 and $a_{12} \le a_{22}$ (1)

$$a_{11} \ge a_{21}$$
 and $a_{12} < a_{22}$ (2)

$$a_{11} < a_{21}$$
 and $a_{12} \ge a_{22}$ (3)

$$a_{11} \le a_{21}$$
 and $a_{12} > a_{22}$ (4)

Consider wlog case (1). As before we have:

$$u_1(r_1, \sigma_2) = a_{11}x + a_{12}(1-x)$$

$$u_1(r_2, \sigma_2) = a_{21}x + a_{22}(1-x)$$

Thus the point at which $u_1(r_1, \sigma_2) - u_1(r_2, \sigma_2) = 0$ is:

$$\frac{a_{22} - a_{12}}{a_{11} - a_{12} + a_{22} - a_{21}}$$

Assuming that equality in (1) does not hold then this point is strictly between 0 and 1 thus a value of 0 < x < 1 can be found for which either r_1 or r_2 is a best response. This gives:

$$B_1 = \{r_1, r_2\}$$

If equality holds in (1) then both strategies are best responses giving the same conclusion.