# OR 3: Chapter 14 - Stochastic games

### Recap

In the previous chapter:

- We considered games of incomplete information;
- Discussed some basic utility theory;
- Considered the principal agent game.

In this chapter we will take a look at a more general type of random game.

## Stochastic games

Definition			
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A stochastic game is defined by:

- X a set of states with a stage game defined for each state;
- A set of strategies  $S_i(x)$  for each player for each state  $x \in X$ ;
- A set of rewards dependant on the state and the actions of the other players:  $u_i(x, s_1, s_2)$ ;
- A set of probabilities of transitioning to a future state:  $\pi(x'|x, s_1, s_2)$ ;
- Each stage game is played at a set of discrete times t.

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We will make some simplifying assumptions in this course:

- 1. The length of the game is not known (infinite horizon);
- 2. The rewards and transition probabilities are not dependent;
- 3. We will only consider strategies called Markov strategies.

#### Definition

A strategy is call a **Markov strategy** if the behaviour dictated is not time dependent.

#### Example

Consider the following game with  $X = \{x, y\}$ :

- $S_1(x) = \{a, b\}$  and  $S_2(x) = \{c, d\}$ ;
- $S_1(y) = \{m\} \text{ and } S_2(x) = \{n\};$

We have the stage game corresponding to state x:

$$\begin{pmatrix} (8,4) & (5,3) \\ (1,5) & (2,6) \end{pmatrix}$$

The stage game corresponding to state y:

The transition probabilities corresponding to state x:

$$\begin{pmatrix} (.8,2) & (1,0) \\ (1,0) & (1,0) \end{pmatrix}$$

The transition probabilities corresponding to state y:

Here is a concise way of representing all this:

We see that the Nash equilibrium for stage game corresponding to x is (a, c) however as soon as the players play that strategy profile they will go to state y which is an absorbing state at which players gain no further utility.

To calculate utilities for players in infinite horizon stochastic games we use a discount rate. Thus without loss of generality if the game is in state x and we

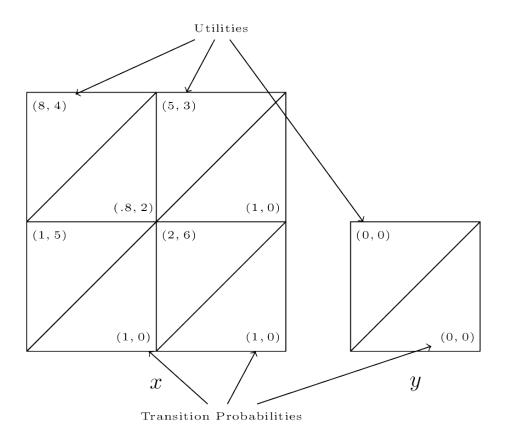


Figure 1:

assume that both players are playing  $\sigma_i^*$  then player 1 would be attempting to maximise future payoffs:

$$U_1(r,s) = \left(u_1(x,r,s) + \delta \sum_{x' \in X} \pi(x'|x,r,s) U_1^*(x')\right)$$

where  $U_1^*$  denotes the expected utility to player 1 when both players are playing the Nash strategy profile.

Thus a Nash equilibrium satisfies:

$$U_1^*(x) = \max_{r \in S_1(x)} (u_i(x, r, s^*) + \delta \sum_{x' \in X} \pi(x'|x, r, s^*) U_1^*(x')$$

$$U_2^*(x) = \max_{s \in S_2(x)} (u_i(x, r^*, s) + \delta \sum_{x' \in X} \pi(x'|x, r^*, s) U_1^*(x')$$

Solving these equations is not straightforward. We will take a look at one approach by solving the example we have above.

### Finding equilibria in stochastic games

Let us find a Nash equilibrium for the game considered above with  $\delta = 2/3$ .

State x gives no value to either player so we only need to consider state x. Let the future gains to player 1 in state x be u, and the future gains to player 2 in state x be v. Thus the players are facing the following game:

$$\begin{pmatrix}
(8 + \frac{1}{3}v, 4 + \frac{1}{3}u) & (5 + \frac{2}{3}v, 3 + \frac{2}{3}u) \\
(1 + \frac{2}{3}v, 5 + \frac{2}{3}u) & (2 + \frac{2}{3}v, 6 + \frac{2}{3}u)
\end{pmatrix}$$

We consider each strategy pair and state the condition for Nash equilibrium:

- 1. (a, c):  $v \le 21$  and  $u \le 3$ .
- 2. (a, d): u > 3.
- 3. (b, c):  $v \ge 21$  and  $u \le 3$ .
- 4. (b,d):  $5 \ge 2$ .

Now consider the implications of each of those profiles being an equilibrium:

1.  $8 + v/3 = v \Rightarrow v = 12$  and  $4 + u/3 = u \Rightarrow u = 6$  which contradicts the corresponding inequality.

- 2.  $3 + 2u/3 = u \Rightarrow u = 9$ .
- 3.  $1+2v/3=v\Rightarrow v=3$  and  $5+2u/3=u\Rightarrow u=15$  which contradicts the corresponding inequality.
- 4. The inequality cannot hold.

Thus the unique Markov strategy Nash equilibria is (a, d) which is not the stage Nash equilibria!