Discussion 11

1. In the *Min-Cost Fast Path* problem, we are given a directed graph *G=(V,E)* along with positive integer times t_e and positive costs c_e on each edge. The goal is to determine if there is a path *P* from *s* to *t* such that the total time on the path is at most *T* and the total cost is at most *C* (both *T* and *C* are parameters to the problem). Prove that this problem is **NP**-complete.

Solution:

1- Prove that Min-Cost Fast Path is in NP
  Certificate: an s-t path with total cost ≤C and total time ≤T
  Certifier: Can easily check in polynomial time that
  a- Set of edges given are in fact a path from s-t
  b- Total time is ≤T and total cost is ≤C
  a and b can be easily done in polynomial time. → Min-Cost Fast Path ϵ NP

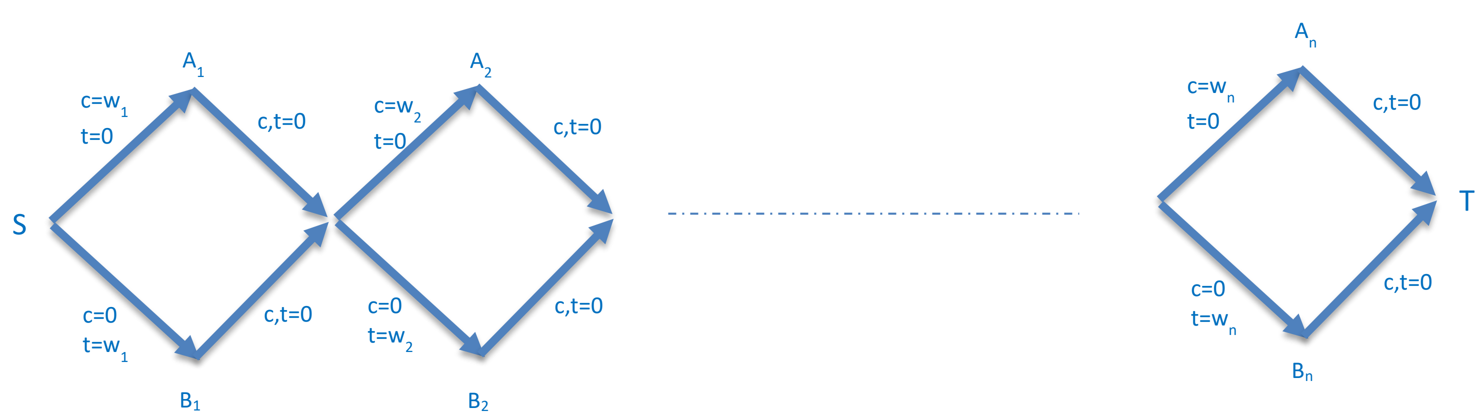
2- Choose Subset Sum for our reduction

3- Will show that Subset Sum $\leq_p$ Min-Cost Fast Path

  Background: Decision version of the Subset Sum problem asks whether given n items where item i has weight w_i, if there is a subset of them whose total weight is less than W and greater than M.

  Plan: We build a graph G such that it has an s-t path with total cost ≤W and total time $\leq \Sigma w_i$ – M iff there is a subset of items whose total weight is between M and W.

  We use gadgets to represent each item. Each gadget will offer two paths through the gadget. If the s-t path in G goes through the A node of the gadget we can interpret that as item being selected as part of the set. If the s-t path goes through the B node of the gadget we interpret that as the item not being selected as part of the set. We string up the gadgets and set time te and costs ce to edges as follows:



Proof:

A- If we are given a set of items with total weight between M and W, we can find a path from S to T with total cost of at most W and total time of at most Σw_i – M by choosing the path through each gadget based on whether the item is part of the set (Go through the A node) or not (go through the B node). If we go through the A node for object i, the object contributes w_i to the cost of the path and if we go through the B node, the object contributes w_i to the total time for the path. So the total cost for the path will be the total weight of the objects selected which we know if ≤ W and the total time of the path is total weight of the objects that are not selected which we know is Σw_i – M (because we know the total weight of the objects selected is ≥ M)

B- If we are given a path from S to T which has a total cost of at most W and a total time of at most Σw_i – M, we can find a set of objects with total weight between M and W. The S-T path can easily select the objects that belong to the set. If the path goes through the A node for an object, we place that object in the set, otherwise not. Since the total cost of the path is at most W then the total weight of the objects selected will be at most W and since the total time for the path is at most Σw_i – M, the total weight of the objects not selected will be Σw_i – M, which means that the total weight of objects selected will be at least M.

2. We saw in lecture that finding a Hamiltonian Cycle in a graph is **NP**-complete. Show that finding a Hamiltonian Path -- a path that visits each vertex exactly once, and isn't required to return to its starting point -- is also **NP**-complete.

Solution:

4- Prove that Hamiltonian Path is in NP
   Certificate: an ordering of all nodes in G that forms a Hamiltonian Path
   Certifier: Can easily check in polynomial time that
   a- There is an edge between each pair of adjacent vertices in the given order
   b- All nodes in G are visited by the path
   a and b can be easily done in polynomial time. → Hamiltonian Path ϵ NP

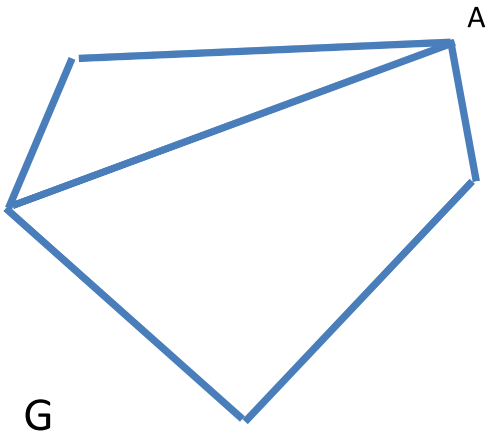
5- Choose Hamiltonian Cycle for our reduction

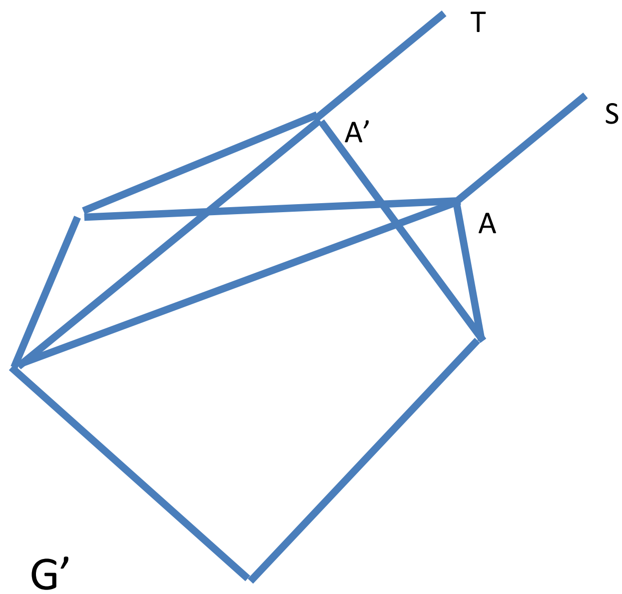
6- Will show that Hamiltonian Cycle $\leq_p$ Hamiltonian Path

Plan: Given graph G—an instance of the Hamiltonian Cycle problem, we will construct G' such that G' has a Hamiltonian Path iff G has a Hamiltonian Cycle.

Construction of G'. We will split one of the nodes in G, say node A. Nodes A and A' will have the same connections as the original node A in G. We will then add two node nodes S and T and connect one with A and the other with A'.

Now G' has a Hamiltonian Path from S to T iff there is a Hamiltonian Cycle in G.





Proof:

A- If we are given a Hamiltonian Path in G’, since S and T have a degree of 1, and can only be the beginning or the end of the path, the path must go from S to T. Ignoring the two new edges SA and TA’, this path will give us a Hamiltonian Cycle in G since A and A’ are the same node in G, i.e. the path will start and end at the same node (A).

B- If we are given a Hamiltonian Cycle in G, we can create a Hamiltonian Path in G’ by splitting the Cycle at node A and creating a path from A’ to A. We can then form a Hamiltonian Path in G’ by starting at S going to A, following the Hamiltonian Cycle to A’ and end the Path at T.

3. Some **NP**-complete problems are polynomial-time solvable on special types of graphs, such as bipartite graphs. Others are still **NP**-complete.
Show that the problem of finding a Hamiltonian Cycle in a bipartite graph is still **NP**-complete.

Solution:

7- Prove that Hamiltonian Cycle in a bipartite graph is in NP
   Certificate: an ordering of all nodes in G that forms a Hamiltonian Cycle
   Certifier: Can easily check in polynomial time that
   a- There is an edge between each pair of adjacent vertices in the given order
   b- All nodes in G are visited by the path
   c- There is an edge between the last node in the order and the first node
   a, b and c can be easily done in polynomial time. → Hamiltonian Cycle in a bipartite graph ϵ NP

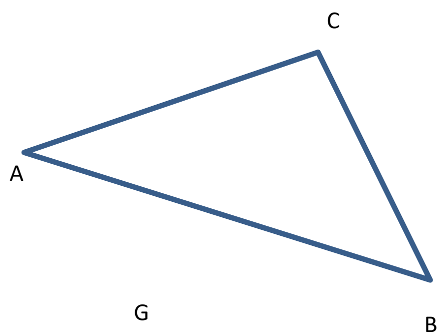
8- Choose Hamiltonian Cycle for our reduction

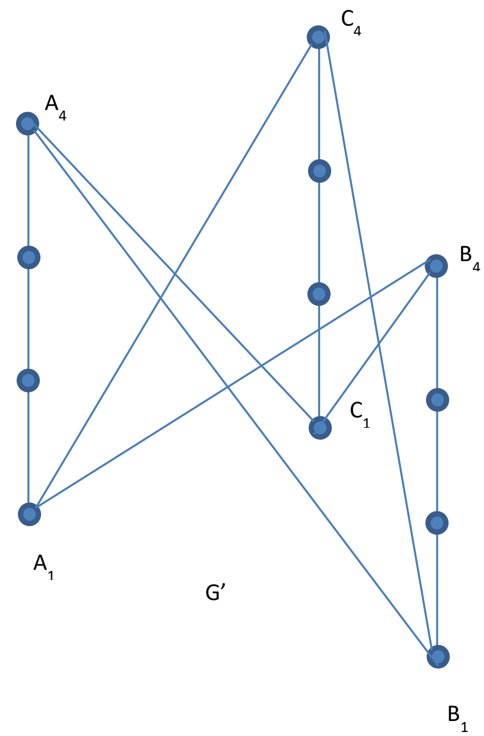
9- Will show that Hamiltonian Cycle $\leq_p$ Hamiltonian Cycle in a bipartite graph

Plan: Given graph G—an instance of the Hamiltonian Cycle problem, we will construct G' such that G' is bipartite and has a Hamiltonian Cycle iff G has a Hamiltonian Cycle.

Construction of G': For each node A in G we will use a gadget with four nodes as shown below. If A and B are connected in G, we connect nodes A_1 and B_4 and nodes B_1 and A_4 in G'.

Gadget in G' representing the node A in G





G' is bipartite since we place all nodes V_1 and V_3 into the set X and nodes V_2 and V_4 into the set Y, all edges in G' go between sets X and Y. And G' has a Hamiltonian Cycle iff G has a Hamiltonian Cycle.

Proof:

A- If we are given a Hamiltonian Cycle in G' it must be of the form $V_1 V_2 V_3 V_4 U_1 U_2 U_3 U_4 \ldots V_1$ since there is no other way for the Hamiltonian Cycle to go through the nodes of each gadget. We can then use the same sequence of nodes V, U, …., V to form a Hamiltonian Cycle in G, since if there is a connection between V_4 and U_1 in G', there must be an edge between V and U in G.

B- If we are given a Hamiltonian Cycle in G that goes through nodes V, U, …., V we can form a Hamiltonian Cycle in G' by going through the gadgets corresponding to nodes V, U, …., V i.e. nodes $V_1 V_2 V_3 V_4 U_1 U_2 U_3 U_4 \ldots V_1$ since if there is a connection between nodes V and U, there must be a connection between nodes V_4 and U_1 in G'.