

Priority Queues

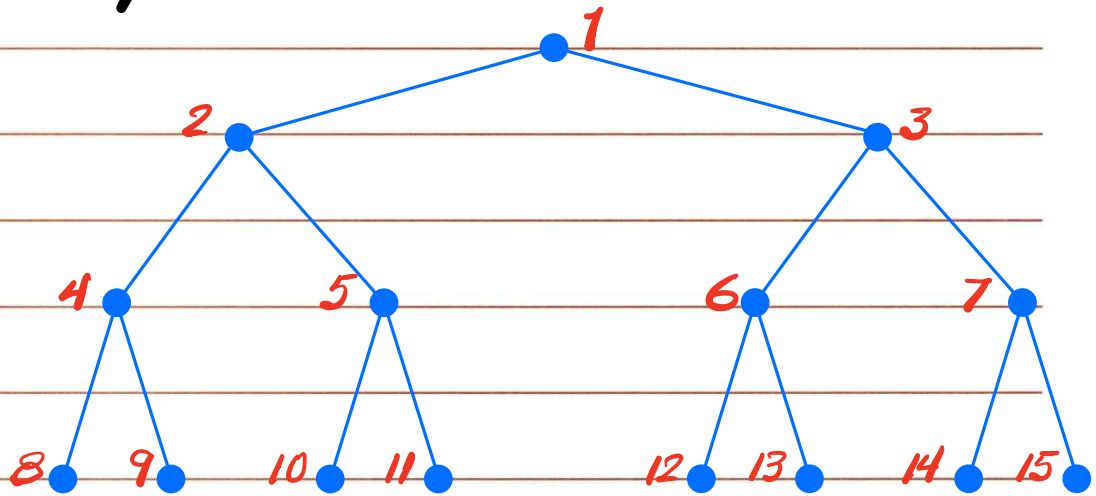
A priority queue has to perform these two operations fast!

1. Insert an element into the set

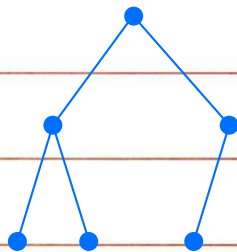
2. Find the smallest element in the set

	<u>Insert</u>	<u>Find Min</u>
Array Implementation	$O(1)$	$O(n)$
Sorted array "	$O(n)$	$O(1)$
Linked List "	$O(1)$	$O(n)$
Sorted Linked List "	$O(n)$	$O(1)$

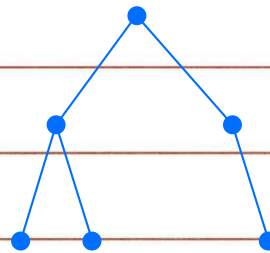
Def. A binary tree of depth k which has exactly $2^k - 1$ nodes is called a full binary tree.



Def. A binary tree with n nodes and of depth k is complete iff its nodes correspond to the nodes which are numbered 1 to n in the full binary tree of depth k .



Complete



Not Complete

Traversing a complete binary tree stored as an array.

Parent(i) is at $\lfloor i/2 \rfloor$ if $i \neq 1$
if $i=1$, i is the root

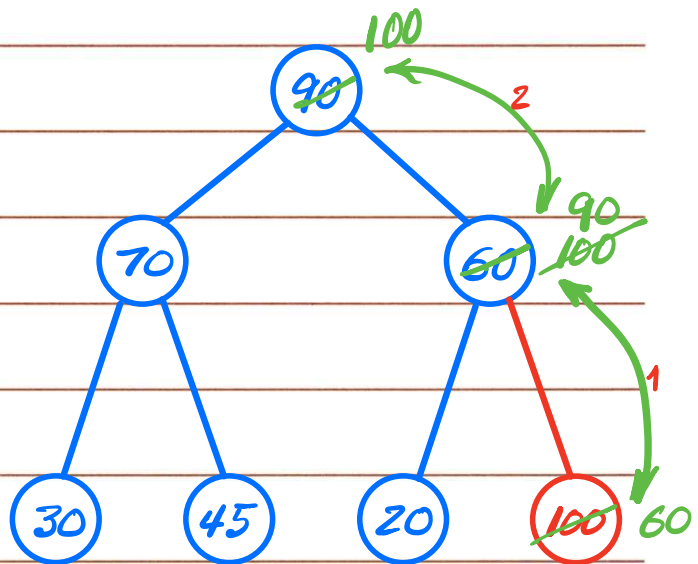
Lchild(i) is at $2i$ if $2i \leq n$
otherwise, it has no left child

Rchild(i) is at $2i+1$ if $2i+1 \leq n$
otherwise, it has no right child

Def. A binary heap is a complete binary tree with the property that the value of the key at each node is at least as large as the key values at its children (Maxheap)

Find-Max

Takes $O(1)$



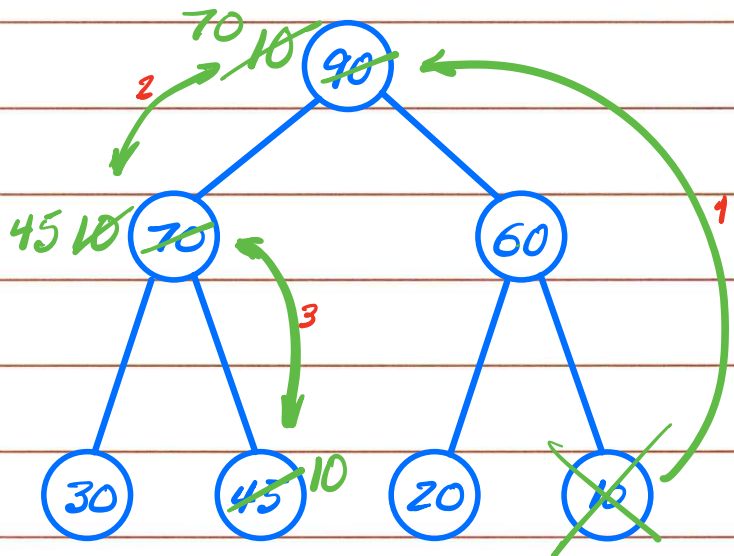
Insert

ex. insert 100

Takes $O(\lg n)$

Extract-Max

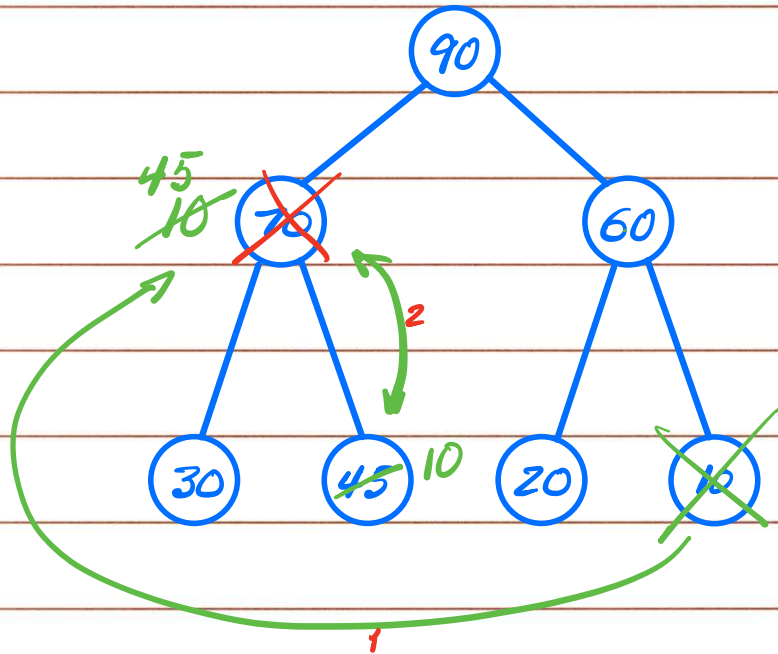
Takes $O(\lg n)$



Delete

ex. Delete 70

Takes $O(\lg n)$

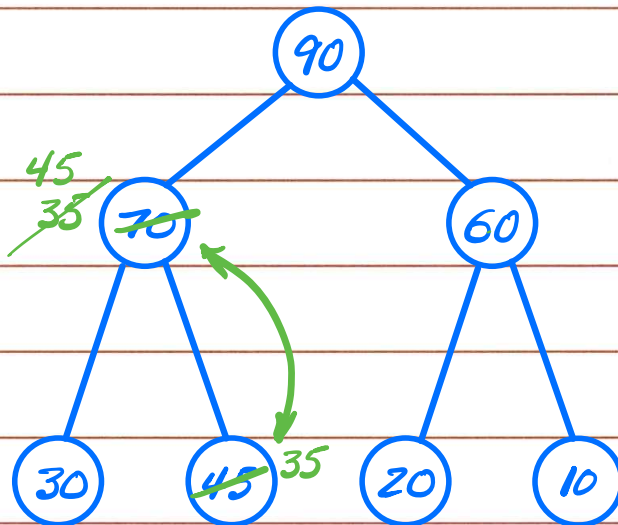


Decrease-key

ex. Decrease-Key

70 to 35

Takes $O(\lg n)$

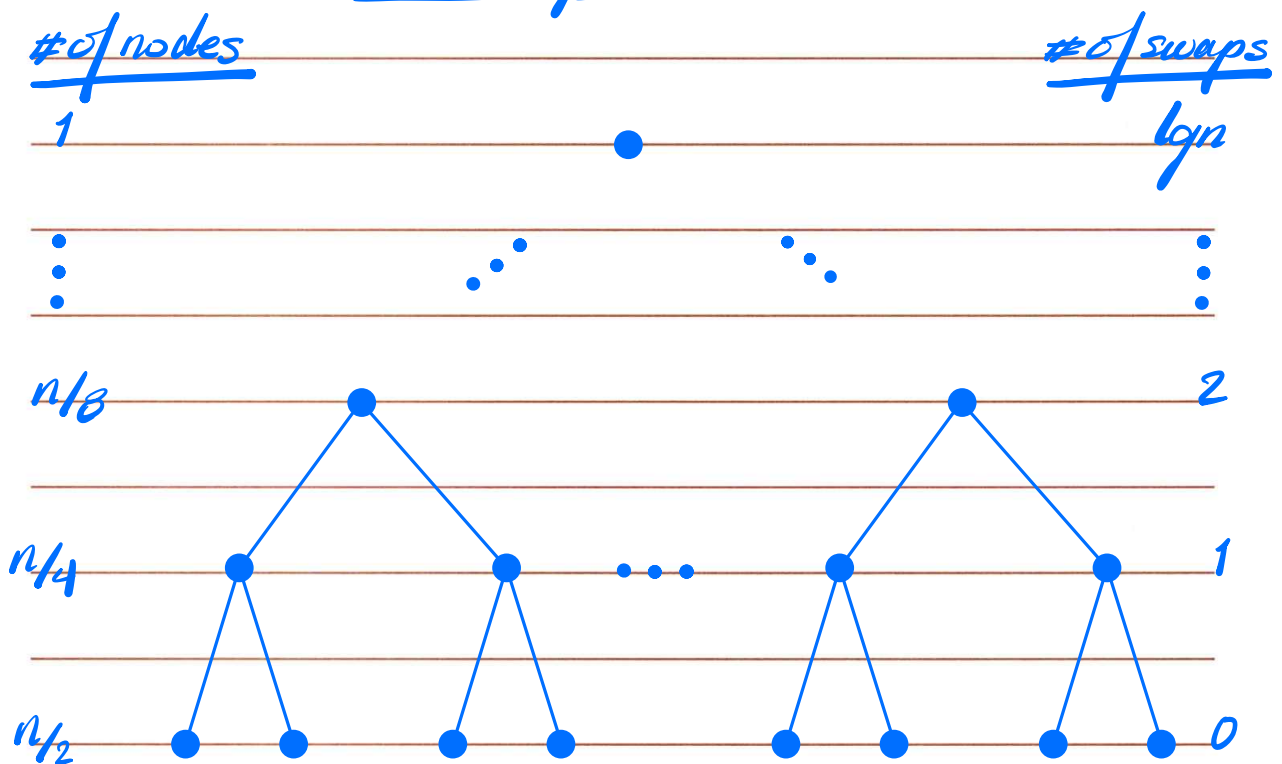


Construction

Can be done in $O(n \lg n)$ time using n insert operations.

Can we do better?

Bottom up construction



$T = \text{Maximum number of swaps needed}$

$$T = n/4 * 1 + n/8 * 2 + n/16 * 3 + \dots$$

$$T/2 = n/8 * 1 + n/16 * 2 + n/32 * 3 + \dots$$

$$T - T/2 = \underbrace{n/4 * 1 + n/8 * 1 + n/16 * 1 + \dots}_{n/2}$$

$$T/2 = n/2$$

$$T = n$$

Bottom up construction takes $O(n)$

Q: What is the best run time to merge two binary heaps of size n ?

A: $O(n)$ using linear time construction