

**CS570**  
**Analysis of Algorithms**  
**Fall 2011**  
**Exam II**

Name: \_\_\_\_\_

Student ID: \_\_\_\_\_

\_\_\_\_\_ 12:30 PM Session      \_\_\_\_\_ 2:00 PM Session

	Maximum	Received
Problem 1	20	
Problem 2	10	
Problem 3	10	
Problem 4	20	
Problem 5	20	
Problem 6	20	
Total	100	

2 hr exam  
Close book and notes

If a description to an algorithm is required please limit your description to within 150 words, anything beyond 150 words will not be considered.

1) 20 pts

Mark the following statements as **TRUE** or **FALSE**. No need to provide any justification.

[ **TRUE**/FALSE ]

Maximum value of an s-t flow could be less than the capacity of a given s-t cut in a flow network.

[ **TRUE**/FALSE ]

In a flow network, if we increase the capacity of an edge that happens to be on a minimum cut, this will increase the max flow in the network .

[ **TRUE**/FALSE ]

The best worst-case time complexity to solve the max flow problem is  $O(Cm)$  where C is the total capacity of the edges leaving the source and m is the number of edges in the network.

[ **TRUE**/FALSE ]

Bellman-Ford algorithm cannot solve the shortest path problem in graphs with negative cost edges in polynomial time.

[ **TRUE**/FALSE ]

If a dynamic programming solution is set up correctly, i.e. the recurrence equation is correct and the sub-problems are always smaller than the original problem, then the resulting algorithm will always find the optimal solution in polynomial time.

[ **TRUE**/FALSE ]

If a problem can be solved by dynamic programming, then it can always be solved by exhaustive search.

[ **TRUE**/FALSE ]

Fractional knapsack problem can be solved in polynomial time.

[ **TRUE**/FALSE ]

0/1 knapsack problem can be solved in polynomial time.

[ **TRUE**/FALSE ]

A flow network with unique edge capacities has a unique min cut.

[ **TRUE**/FALSE ]

The number of iterations it takes Bellman-Ford to converge can vary depending on the order of nodes updated within each iteration.

2) 10 pts

You are given a network with  $n$  vertices,  $m$  edges, a source  $s$  and a sink  $t$ . Suppose your friend presents to you an  $s$ - $t$  flow for the network by assigning for every edge  $e$ , a flow  $f(e)$ . Describe an  $O(n + m)$  algorithm to test if your friend's flow assignment is a maximum  $s$ - $t$  flow.

Run BFS on the residual graph. If we cannot find a path from  $s$  to  $t$ , then the flow is maximum.

Discussion and previous exams

3) 10 pts

You have successfully computed a maximum  $s$ - $t$  flow  $f$  for a network  $G = (V; E)$  with integer edge capacities. Your boss now gives you another network  $G'$  that is identical to  $G$  except that the capacity of exactly one edge is increased by one. You are also explicitly given the edge whose capacity was changed. Describe how you can compute a maximum flow for  $G'$  in  $O(|V| + |E|)$  time, where  $|V|$  is the number of vertices and  $|E|$  is the number of edges.

4) 20 pts

Given a weighted directed graph, describe an algorithm that decides if there exists a negative weight cycle in the graph. The running time should be polynomial in the number of vertices.

**Bellman – Ford Algorithm**

5) 20 pts

We are given a string  $S = s_1s_2s_3\dots s_n$ , and we want to delete some characters such that what remains is a palindrome. Give an efficient dynamic programming algorithm to determine the length of the longest palindrome, and analyze the running time.

A palindrome is a string that reads the same backward or forward such as ACCBCCA.

Solution:

Let  $X[0..n-1]$  be the input sequence of length  $n$  and  $L(0, n-1)$  be the length of the longest palindromic subsequence of  $X[0..n-1]$ .

If last and first characters of  $X$  are same, then  $L(0, n-1) = L(1, n-2) + 2$ .

Else  $L(0, n-1) = \text{MAX}(L(1, n-1), L(0, n-2))$ .

6) 20 pts

Determine if there is a feasible circulation in the following network. Demand values are in circles. Edge capacities are marked on each edge. Show all your steps.

