Side-channel attacks on Ascon's S-box

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Introduction

Side-Channel Attacks (SCA): observation of computation time, power consumption, electromagnetic radiation, ... to discover a secret

Goal: Study the leaks from the winner for lightweight cryptography Ascon to theorize a SCA attack

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 - Encryption and decryption phases
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What is Ascon-AEAD?

Authenticated Encryption with Associated Data (AEAD): encrypt, check authentication of content and associated data

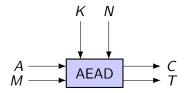


Figure: AEAD algorithm from [1]

Ascon's State

byte 0	byte 1	byte 2	byte 3	byte 4	byte 5	byte 6	byte 7	
IV								S_0
first half of K, K ₀								S_1
Second half of K, K_1								S_2
first half of N, N ₀								S_3
Second half of N, N_1								S_4

Encryption and decryption phases

4 phases: Initialization, Associated data process, plaintext/ciphertext process, Finalization

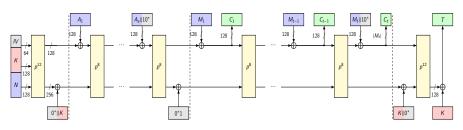


Figure: Ascon-AEAD mode, from [1]

Ascon's permutation

$$p = p_{L} \circ \underbrace{p_{S}} \circ p_{C}$$

$$N \xrightarrow{\downarrow} 1 \xrightarrow{\downarrow} 0 \xrightarrow{\downarrow} 0 ut_{0}$$

$$K_{1} \xrightarrow{\downarrow} 1 \xrightarrow{\downarrow} 0 \xrightarrow{\downarrow} 0 ut_{1}$$

$$N_{1} \xrightarrow{\downarrow} 1 \xrightarrow{\downarrow} 0 \xrightarrow{\downarrow} 0 ut_{4}$$

Figure: Circuit to compute the S-box, from 1 , bijection $[1;31] \rightarrow [1;31]$

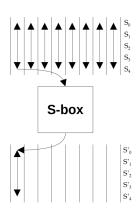


Figure: S-box computation for the first byte of each word

https://extgit.isec.tugraz.at/meichlseder/tikz

Table linking the output of the S-box and the key

(n_0^j, n_1^j, IV^j)	S_4^j
(0,0,0)	k_0^j
(0,0,1)	0
(0,1,0)	1
(0,1,1)	$1 \oplus k_0^j$
(1,0,0)	$1 \oplus k_0^j$
(1,0,1)	1
(1,1,0)	0
(1, 1, 1)	k_0^j

Figure: Link between k_0^j and S_4^j depending on IV and N, from [3]

ChipWhisperer-Lite

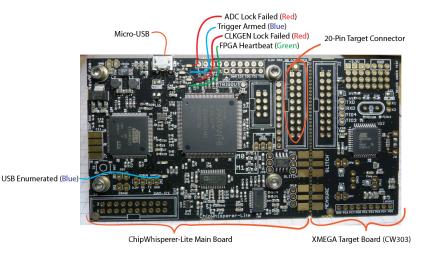


Figure: ChipWhisperer Lite board, from [4]

Steps for a CPA attack

- Compute the algorithm multiple times to gain traces
- Find model for the consumption
- Deduce the hypothesis that correlates best

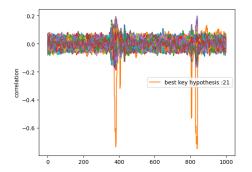


Figure: CPA on another encryption standard



Analyses done

- Finding the best model
 - Vertical vs horizontal
 - HW vs value
- Attack: finding the vertical output and deduct the key

Results vertical vs horizontal and HW vs value

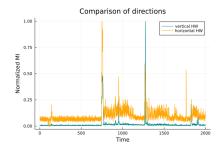


Figure: Mutual information for the horizontal and the vertical value

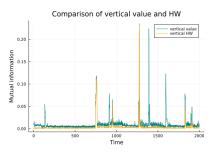


Figure: Mutual information between power consumption and HW or value

Results attack

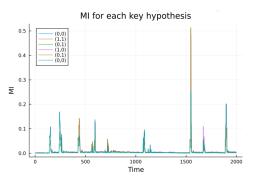


Figure: Mutual information between the Hamming weight of the outputs and the power consumption, for each of the possible outputs for the first nonce

Conclusion

- Good leaks compared to random values
- Though apparent weaknesses, unsuccessful attempts
- Not enough randomness with false key hypotheses
- Leads to follow: belief propagation



L. B. H., T. G., and A. D., "Théorie de la cryptographie."



S. T. M., M. K., C. D., K. J., and K. J., "Ascon-based lightweight cryptography standards for constrained devices: Authenticated encryption, hash, and extendable output functions," *NIST Special Publication 800, NIST SP 800-232 ipd*, 2024. [Online]. Available: https://nvlpubs.nist.gov/nistpubs/SpecialPublications/NIST.SP. 800-232.ipd.pdf



S. M., "Side channel analysis against aead." [Online]. Available: https://theses.hal.science/tel-04816066v1



Chipwhisperer documentation. [Online]. Available: https://chipwhisperer.readthedocs.io/en/latest/getting-started.html

Permutation (1), p_C

Constant for the round i: $const_{16-nb_{rounds}+i}$

i	const;	i	$const_i$
0	0x000000000000003c	8	0x00000000000000b4
1	0x0000000000000002d	9	0x000000000000000a5
2	0x000000000000001e	10	0x00000000000000096
3	0x000000000000000000000000000000000000	11	0x0000000000000087
4	0x000000000000000000000000000000000000	12	0x0000000000000078
5	0x000000000000000001	13	0x00000000000000069
6	0x00000000000000d2	14	0x000000000000005a
_7	0x0000000000000001	15	0x000000000000004b

Figure: Constant-addition layer, constants

Permutation (2), p_C

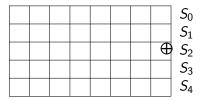


Figure: Constant-addition layer, each box representing a byte of one of the 64-bit words

Permutation (3), p_S

X	0	1	2	3	4	5	6	7
S - box(x)	4	b	1f	14	1a	15	9	2
X	8	9	a	b	С	d	е	f
S - box(x)	1b	5	8	12	1d	3	6	1c
X	10	11	12	13	14	15	16	17
S - box(x)	1e	13	7	е	0	d	11	18
X	18	19	1a	1b	1c	1d	1e	1f
S - box(x)	10	С	1	19	16	a	f	17

Figure: Lookup table for the 5-bit S-box

Permutation (4), p_S

```
state [0] ^= state [4];
             state [4] ^= state [3];
             state [2] ^= state [1];
4
             uint64 t t0 = "state[0];
5
             uint64 t t1 = "state[1];
6
             uint64 t t2 = "state[2];
             uint64 t t3 = "state[3];
8
             uint64 t t4 =  state [4];
             t0 &= state [1];
             t1 &= state [2];
             t2 &= state [3];
             t3 &= state [4];
             t4 &= state [0];
14
             state[0] ^= t1
             ; state[1] ^= t2;
15
             state[2] ^= t3;
17
             state[3] ^= t4;
18
             state [4] ^= t0;
             state [1] ^= state [0];
             state [0] ^= state [4];
             state[3] ^= state[2];
             state[2] = state[2]:
```

Figure: Equations to compute the S-box

Permutation (5), p_L

Diffusion: $S_i \leftarrow \Sigma_i(S_i)$:

$$\Sigma_{0}(S_{0}) = S_{0} \oplus (S_{0} >>> 19) \oplus (S_{0} >>> 28)$$

$$\Sigma_{1}(S_{1}) = S_{1} \oplus (S_{1} >>> 61) \oplus (S_{1} >>> 39)$$

$$\Sigma_{2}(S_{2}) = S_{2} \oplus (S_{2} >>> 1) \oplus (S_{2} >>> 6)$$

$$\Sigma_{3}(S_{3}) = S_{3} \oplus (S_{3} >>> 10) \oplus (S_{3} >>> 17)$$

$$\Sigma_{4}(S_{4}) = S_{4} \oplus (S_{4} >>> 7) \oplus (S_{4} >>> 41)$$

Finding this table (1)

$$S_{4}^{j} = n_{o}^{j} \oplus n_{1}^{j} \oplus k_{0}^{j} \times \left(1 \oplus IV^{j} \oplus n_{1}^{j}\right)$$

$$S_{4}^{j} = \begin{cases} k_{0}^{j} \times \left(1 \oplus IV^{j}\right) & if \ (n_{0}^{j}, n_{1}^{j}) = (0, 0) \\ k_{0}^{j} \times IV^{j} & if \ (n_{0}^{j}, n_{1}^{j}) = (1, 1) \\ 1 \oplus k_{0}^{j} \times IV^{j} & if \ (n_{0}^{j}, n_{1}^{j}) = (0, 1) \\ 1 \oplus k_{0}^{j} \times (1 \oplus IV^{j}) & if \ (n_{0}^{j}, n_{1}^{j}) = (1, 0) \end{cases}$$

Finding this table (2)

Then if $IV^j = 0$:

$$S_4^j = \begin{cases} k_0^j & if \ (n_0^j, n_1^j) = (0, 0) \\ 0 & if \ (n_0^j, n_1^j) = (1, 1) \\ 1 & if \ (n_0^j, n_1^j) = (0, 1) \\ 1 \oplus k_0^j & if \ (n_0^j, n_1^j) = (1, 0) \end{cases}$$

Finding this table (3)

Otherwise, if $IV^{j} = 1$:

$$S_4^j = \left\{ \begin{array}{ll} 0 & if \ (n_0^j, n_1^j) = (0,0) \\ k_0^j & if \ (n_0^j, n_1^j) = (1,1) \\ 1 \oplus k_0^j & if \ (n_0^j, n_1^j) = (0,1) \\ 1 & if \ (n_0^j, n_1^j) = (1,0) \end{array} \right.$$

Complementary graph (1)

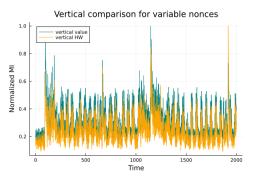


Figure: Mutual information between power consumption and Hamming weight of the concatenation of the first bit of each of the word of S and its value like 9 but for random nonces

Complementary graph (2)

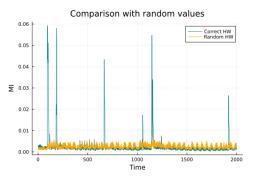


Figure: Mutual information between power consumption and vertical HW or random possible HW

```
1 #include "ascon.h"
2 #include "permutation.h"
3 #include <stdint.h>
4 #include < stdlib.h>
5 #include <stdbool.h>
6
7 //-
8 // ASCON implementation following the NIST Special Publication 800, NIST SP
       800-232 ipd
11 //unsure of the generalization of this IV, as t=128 is not defined,
12 //only given for this specific algorithm
14
                                                     NB RNDS A)<<16) + VERSION;
16 //main functions
17
 void ASCONEncrypt (uint8 t* C, uint8 t* T, uint8 t* P, int len p,
18
                  uint8 t* A, int len a, uint8 t* n, uint8 t* key){
      State t* state = \overline{i} nitialization (P, len p, A, len a, \overline{n}, key, true);
19
      associated data process(state);
23
      plaintext process(state, len p%RATIO);
24
      finalization (state, key, T);
      slice(state->blocktext out, C, len p);
29
      free state(state);
30
      return:
31
32
33
^{34} void ASCONDecrypt(uint8 tst C, uint8 tst P, int len p, uint8 tst A, int len a,
```

Side-channel attacks on Ascon's S-box

10 / 10

```
uint8 t* n, uint8 t* key, uint8 t* T, bool* fail){
      State t* state = initialization(C, len p, A, len a, n, key, false);
      associated data process(state);
      ciphertext process(state, len p%RATIO);
      uint8 t* Tprim = malloc(KEY LENGTH/BYTE LENGTH*sizeof(uint8 t));
      finalization (state, key, Tprim);
      if (same tag(T, Tprim)){
          *fail = false:
          slice(state->blocktext out, P, len p);
      else{
          *fail = true;
      free (Tprim);
      free state(state);
      return:
  //initialzation functions
  uint64 t get IV(void){
      return IV:
64 uint8 t** combine(uint8 t* text, int len, bool padding){
      //new len: has to be incremented either for the added padding
      //or because the value has been floored
      //if padding=false and len%RATIO=0, adding an empty block
      int new len = len/RATIO + 1;
      uint8 t** new text = malloc(new len*sizeof(uint8 t*));
                              Side-channel attacks on Ascon's S-box
```

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10/10

```
70
       for (int i=0; i < new len; i++){
            new text[i] = calloc(RATIO, size of (uint8 t));
73
74
       for (int i=0; i<len; i++)\{// little endian for each 64 bit block
75
            new text[i/RATIO][8*((i%RATIO)/8) + 7 +(-i)%8] = text[i];
76
       //adding a 1 at the end of the block if a padding is needed
       if (padding)
79
            new text [new len -1][8*((len%RATIO)/8) + 7 +(-len)%8] += 0 \times 01;
80
       return new text;
81 }
82
83
  State t* state initialization (uint8 t* textin, int len p, uint8 t* A,
                                   int len a, uint8 t* n, uint8 t* key, bool plain){
86
       State t* state = (State t*) malloc(sizeof(State t));
88
       state -> words [0] = get IV();
89
90
       //|K| + |n| + 64 = 320, so K U n has 32 bytes,
91
       //adds K then n to the four last words of state
       state \rightarrow words[1] = 0:
93
       state \rightarrow words[2] = 0;
94
       state \rightarrow words[3] = 0;
       state \rightarrow words[4] = 0:
96
       for (int i = 0: i < 32: i++){
97
            uint64 t curr byte = (i<KEY LENGTH/BYTE LENGTH)?
                                        kev[i] : n[i—KEY LENGTH/BYTE LENGTH]:
            state \rightarrow words [1 + i/8] += curr byte <math><<(((i*BYTE LENGTH)\%64);
99
00
       }
       state -> Ai = combine (A, len a, true);
03
       state \rightarrow s = len a/RATIO + 1;
04
```

```
05
       state -> blocktext in = combine(textin, len p, plain);
       state \rightarrow t = len p/RATIO + 1;
08
       state \rightarrow blocktext out = malloc(state \rightarrow t*sizeof(uint8 t*));
09
       for (int i=0; i < state -> t; i++){
10
            state -> blocktext out[i] = calloc(RATIO, size of (uint8 t));
11
12
13
       return state:
14 }
15
16 State t* initialization(uint8 t* textin, int len p, uint8 t* A, int len a,
                               uint8 t* n, uint8 t* key, bool plain){
17
18
       State t* state = state initialization(textin, len p, A, len a, n, key, plain);
       permutation (state, NB RNDS A);
       for(int i=0; i<KEY LENGTH/BYTE LENGTH; i++){</pre>
            //finds the first modified word to add the key at the end of the state
24
            int ind = 5 - \text{KEY LENGTH}/64 - (\text{KEY LENGTH}\%64 \stackrel{!}{=} 0) + i/8:
            state \rightarrow words [ind] \hat{} = (uint64 t) key [i] < ((i*BYTE LENGTH)%64);
26
       return state:
29
30
  //processing associated data
31
  void associated data process(State t* state){
       for (int i=0; i < state -> s; i++){
            for (int j=0; j < RATIO; j++){
                 int shift = (RATE-((j+1)*BYTE LENGTH))\%64;
34
35
                 state \rightarrow words[i/8] ^= (uint64 t) state \rightarrow Ai[i][i] < shift;
36
37
            permutation(state, NB RNDS B):
38
39
       //update for the least significant byte in little endian
                                  Side-channel attacks on Ascon's S-box
                                                                                               10/10
```

```
state -> words [4] \hat{} = (uint64 t) 1<<63;
       return:
42 }
44
  //processing plaintext/ciphertext
45
  void plaintext process(State t* state, int 1){
       for (int i=0; i < state -> t-1; i++){
           for (int i=0; i < RATIO; i++){
                int shift = (RATE-((i+1)*BYTE LENGTH))%64:
                state \rightarrow words [j/8] \hat{t} state \rightarrow blocktext in [i][j] << shift;
                //too much bits for a uint8 \overline{t}, only takes the smallest byte
                state -> blocktext out[i][j] = state -> words[j/8] >> shift;
           }
           permutation (state, NB RNDS B);
       //C I is going to have its RATIO-I least significant bytes uninitialized.
       //but it is unimportant because C will be computed only with its beginning
       for (int i=0; i<l+1; i++){
           int i = state \rightarrow t - 1:
           //not every byte will be explored so the order matters
           int ind = 8*((i\%RATIO)/8) + 7 + (-i)\%8;
           int shift = (RATE-((ind+1)*BYTE LENGTH))%64;
           state -> words [ind /8] ^= (uint64 t) state -> blocktext in [i] [ind] << shift;
           state -> blocktext out [i] [ind] = state -> words [(ind -1)/8] >> shift;
       return:
68
  void ciphertext process(State t* state, int 1){
       for (int i=0; i < state -> t-1; i++){
           for (int j=0; j < RATIO; j++){
                uint8 t wbyte = state->words[j/8]>>((RATE-((j+1)*BYTE LENGTH))%64);
                state -> blocktext out[i][j] = wbyte^state -> blocktext in[i][j];
                char* word = (char*) &(state -> words[i/8]);
                word[(RATIO-1-j)\%8] = state \rightarrow blocktext in[i][j];
                                 Side-channel attacks on Ascon's S-box
                                                                                         10/10
```

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74

```
permutation (state, NB RNDS B);
       }
79
       for (int i=0; i<1; i++){
           int i = state \rightarrow t - 1:
81
           //not every byte will be explored so the order matters
82
           int ind = 8*((j\%RATIO)/8) + 7 + (-j)\%8;
           uint8 t wbyte = state->words[ind/8]>>((RATE-((ind+1)*BYTE LENGTH))%64);
84
           state -> blocktext out[i][ind] = wbyte^state -> blocktext in[i][ind];
           char* word = (char*) &(state->words[j/8]);
85
86
           word[(RATIO-1-ind)%8] = state->blocktext in[i][ind];
87
       }
88
       state \rightarrow words [1/8] \hat{} = (uint64 t) 1<<((BYTE LENGTH*1)%64);
89
       return:
90 }
91
92
  //finalization functions
93
  void slice (uint8 t** blocktextout, uint8 t* textout, int len){
94
       for (int i=0; i<len; i++){
95
           textout[i] = (blocktextout[i/RATIO][8*((i%RATIO)/8) + 7 + (-i)%8]);
96
97
       return:
98 }
99
00
  void finalization (State t* state, uint8 t* key, uint8 t* T){
       for(int i=0; i<KEY LENGTH/BYTE LENGTH; i++){</pre>
           //adding the key after the rate
           int ind = (RATIO-1)/8 + 1 + i/8;
04
           state \rightarrow words[ind] ^= (uint64 t) key[i] < ((i*BYTE LENGTH)%64);
05
       }
06
07
       permutation (state, NB RNDS A);
08
09
       for(int i=0; i<TAG LENGTH/BYTE LENGTH; i++){</pre>
                                                                                          10/10
```

```
10
           //xors the last TAG LENGTH bits of the key and the state
           int ind = 5 - TAG LENGTH/64 - (TAG LENGTH%64 != 0) + i/8;
           uint8 t key byte = key[i + (KEY LENGTH-TAG LENGTH)/BYTE LENGTH];
13
           T[i] = state->words[ind]>>((i*BYTE LENGTH)\%64) ^ key byte;
14
15
       return:
16 }
17
18
  bool same tag(uint8 t* T.uint8 t* Tprim){
19
      bool res = true:
       for(int i=0; i<TAG LENGTH/BYTE LENGTH; i++){</pre>
           //to have the same computation time no matter the value of T and T'
           res = res \&\& (T[i] == Tprim[i]);
       return res;
27
  void free state(State t* state){
       for (int i=0; i < state -> s; i++)
           free(state->Ai[i]);
30
       free(state->Ai);
31
       for (int i=0; i < state ->t; i++){
           free(state->blocktext in[i]);
33
           free(state -> blocktext out[i]);
34
       free(state->blocktext in):
       free(state -> blocktext out);
       free(state);
       return:
39 }
```

Listing: Implementation for ascon.c

```
2 #ifndef ASCON H
3 #define ASCON H
5 #define KEY LENGTH 128
6 #define TAG LENGTH 128
7 #define NONCE LENGTH (256 - KEY LENGTH)
8 #define RATE 128
9 #define BYTE LENGTH 8
10 #define NB RNDS A 12
11 #define NB RNDS B 8
12 #define VERSION 1
13 #define RATIO (RATE/BYTE LENGTH)
15 #include <stdint.h>
16 #include <stdbool.h>
17
18 // defining a type State \, t in order to implement the state with the data
19 typedef struct State s {
20
      uint64 t words [5]:
      uint8 t** Ai: //asociated data in blocks of size RATE
      int s; //number of blocks in Ai
23
      uint8 t** blocktext in; //blocks of the plaintext when encryption,
24
                               // when decryption blocks of ciphertext
25
      int t; //number of blocks in the given text, whether plain or ciphered
      uint8 t** blocktext out;// output blocks (resp. ciphered or plain)
27 }
    State t:
30 //main functions
31 void ASCONEncrypt(uint8 t* C, uint8 t* T, uint8 t* P, int len p,
                   uint8 t* A, int len a, uint8 t* n, uint8 t* key);
33 //takes a pointer C pointing towards allocated memory of size len p.
34 //and a pointer T for the tag for a space of TAG LENGTH
```

```
36 void ASCONDecrypt(uint8 t* C, uint8 t* P, int len p, uint8 t* A, int len a,
                   uint8 t* n, uint8 t* key, uint8 t* T, bool* fail);
\frac{1}{2} //same thing but for \overline{P}
39 //modifies fail thanks to the pointer if authentication failed
40
41
42 //initialization functions
43 uint64 t get IV(void);//getter for the IV
44
45 State t* state initialization(uint8 t* textin, <mark>int</mark> len p, uint8 t* A,
46
                               int len a, uint8 t* n, uint8 t* key, bool plain);
47 //creates the state with the five words
48 //and the blocks for the plaintext, ciphertext and associated data
49
50 uint8 t** combine(uint8 t* text, int len, bool padding);
51 //transforms list of uint8 t in list of lists of RATE uint8 t.
52 //gives the responsability to free the list .adds a padding if padding=true
54 State t* initialization(uint8 t* textin, int len p, uint8 t* A, int len a,
55
                           uint8 t* n. uint8 t* kev. bool plain):
56 //creates the state and does the initialization phase, if plain=true encryption
57 //therefore there will be a padding, otherwise not
//gives responsability to free the state
60 //processing associated data
61 void associated data process(State t* state):
64 //processing plaintext/ciphertext
65 void plaintext process(State t* state, int 1);
66 //I, the length in bytes of the last block in the original text
67 void ciphertext process(State t* state, int 1);
70 //finalization functions
```

```
void slice(uint8_t** blocktextout, uint8_t* textout, int len);
//flattens the list of list blocktextout in a list textout
// //flattens text of length len

void finalization(State_t* state, uint8_t* key, uint8_t* T);

bool same_tag(uint8_t* T, uint8_t* Tprim); //checks T=T'

void free_state(State_t* state);

#endif
```

Listing: Implementation for ascon.h

```
2 #include "permutation.h"
4
  // permutation implementation for Ascon, following the NIST Special Publication
       800. NIST SP 800-232 ipd
7
8 void permutation (State t* state, int nb rounds, bool attack) { // modifies the
       current state by applying nb rounds times the permutation
      for (int r=0; r< nb rounds; r++){
           uint8 t c = round const(nb rounds, r);
          perm const(state, c);
          perm sub(state);
14
15
           if (attack && r==0){
16
               trigger high();
17
               perm lin(state);
```

```
trigger low();
            else
                perm lin(state);
       return:
24 }
26 uint8 t round const(int nb rounds, int round){
       // finds the right index, the table of consts defining for 12 rounds c 0 = 0
        ×f0
28
        int i = 12 - nb rounds + round;
30
       //relationship between index and value defined for i in (-4.11)
        return 0 \times f0 - i * 0 \times 10 + (0 \times 10 + i) \% 0 \times 10;
32 }
33
34 void perm const(State t* state, uint8 t c){
35
       state -> words [2] ^= (uint64 t)c;
36
        return:
37
39 void perm sub(State t* state){
40
       //applies the bitsliced implementation from figure 4
41
       state -> words [0] ^= state -> words [4]; state -> words [4] ^= state -> words [3];
        state -> words [2] ^= state -> words [1]:
42
       uint64 t t0 = state\rightarrowwords[0]; uint64 t t1 = state\rightarrowwords[1]; uint64 t t2 =
          state -> words [2];
43
       uint64 t t3 = state->words[3]; uint64 t t4 = state->words[4];
       t0 = \overline{t0}; t1 = \overline{t1}; t2 = \overline{t2}; t3 = \overline{t3}; t4 = \overline{t4};
45
       t0 &= state->words[1]; t1 &= state->words[2]; t2 &= state->words[3]; t3
        &= state -> words [4]; t4 &= state -> words [0];
46
       state \rightarrow words[0] ^= t1 ; state \rightarrow words[1] ^= t2 ; state \rightarrow words[2] ^= t3 ;
        state \rightarrow words[3] \stackrel{}{}= t4 ; state \rightarrow words[4] \stackrel{}{}= t0;
       state -> words [1] ^= state -> words [0]; state -> words [0] ^= state -> words [4];
    Alexane Boldo (OCIF, IRISA)
                                    Side-channel attacks on Ascon's S-box
```

Main code

state -> words [3] ^= state -> words [2]; state -> words [2] = state -> words [2];

```
return:
50 }
  uint64 t circular shift(uint64 t word, int nb){
53
       uint64 t shifted = word>>nb:
54
       uint64 t extracted = word << (RATE-nb):
55
       return shifted+extracted;
56 }
58
  void perm lin(State t* state){
59
       //calculates the sigma functions
       state -> words [0] = state -> words [0] ^ circular shift (state -> words [0], 19) ^
        circular shift(state->words[0],28);
61
       state -> words [1] = state -> words [1] ^ circular shift (state -> words [1],61) ^
        circular shift (state -> words [1], 39);
       state -> words [2] = state -> words [2] ^ circular shift (state -> words [2], 1) ^
        circular shift(state->words[2],6);
63
       state -> words [3] = state -> words [3] ^ circular shift (state -> words [3], 10) ^
        circular shift (state -> words [3], 17);
64
       state -> words [4] = state -> words [4] ^ circular shift (state -> words [4],7) ^
        circular shift(state -> words[4],41);
65
       return:
66 }
                         Listing: Implementation for permutation.c
```

```
3 #ifndef PERMUTATION H
4 #define PERMUTATION H
6 //functions for the permutation
```

1 #include "ascon.h"

48

Listing: Implementation for permutation.h

```
1 #include "ascon.h"
2 #include "hal/hal.h"
3 #include "hal/simpleserial.h"
4 #include <stdint.h>
5 #include <stdlib.h>
7 uint8 t key [KEY LENGTH/BYTE LENGTH];
8 uint8 t A[MAX DATA SIZE];
9 uint8 t P[MAX DATA SIZE];
10 uint8 t C[2*MAX DATA SIZE]:
11 uint8 t * T = &(C[MAX DATA SIZE]);
12 uint8 t n[NONCE LENGTH];
13
14 uint8 t set key(uint8 t*k, uint8 t len)
15
16
      if (len != KEY LENGTH/BYTE LENGTH) {
17
           led error (\overline{1});
18
           return 0x01;
19
```

```
for(int i = 0; i < KEY LENGTH/BYTE LENGTH; i++) {</pre>
           kev[i] = k[i]:
24
    return 0x00:
25
  uint8 t set assodata(uint8 t* ad, uint8 t len){
28
       if (len != MAX DATA SIZE) {
           led error(1):
           return 0x01:
31
33
    for (int i = 0: i < MAX DATA SIZE: i++) {
34
           A[i] = ad[i]:
35
36
    return 0x00:
37 }
38
39
  uint8 t set nonce(uint8 t* nonce, uint8 t len){
40
       if (len != NONCE LENGTH/BYTE LENGTH){
41
           led error(1);
42
           return 0x01;
43
44
45
       for(int i=0; i < NONCE LENGTH/BYTE LENGTH; i++){</pre>
46
           n[i] = nonce[i];
48
       return 0x00;
49 }
51
  uint8 t set pt(uint8 t* pt, uint8 t len)
53
       if (len != MAX DATA SIZE) {
54
           led error(1);
           return 0x01;
```

10/10

```
for (int i = 0; i < MAX DATA SIZE; i++) {
           P[i] = pt[i];
61
       /*
63
       Encrypt here using ciphertext, plaintext, tag, associated data and key
       variables.
64
      */
65
    ASCONEncrypt (C, T, P, len, A, len, n, key);
66
    simpleserial put('r', 16, C);//sends ciphertext
68
    return 0x00;
69 }
70
71 int
      main (void)
72 {
       platform init();
74
       init uart();
       trigger setup();
76
       led ok(1);
78
       led error(0);
79
       simpleserial init():
81
       simpleserial addcmd('k', 16, set key);
       simpleserial addcmd('a', 16, set assodata);
       simpleserial addcmd('n', 16, set nonce);
84
       simpleserial addcmd('p', 16, set pt);
85
86
87
       while(1) {
           simpleserial get();// get next command and react to it
                                                                      May-July 2025
```

```
90 }
```

Listing: Implementation for main.c

```
1 import chipwhisperer as cw
2 import matplotlib.pyplot as plt
3 import numpy as np
4 from tgdm import tgdm
5
  import time
6
  def reset target(scope):
8
       scope.io.pdic = 'low'
       time. sleep (0.05)
       scope.io.pdic = 'high'
       time.sleep (0.05)
12
13
  def generate nonce()
14
       nonce = np.random.randint(0,high=1<<64-1,size=2)
15
       nonce[1] = 0
       for i in range (64):
17
            nonce [1] += ((IV[i//8] >> (i\%8))\%2) << i
19
       nonce = bytearray(nonce)
       return nonce
  IV = [0 \times 00, 0 \times 00, 0 \times 10, 0 \times 00, 0 \times 80, 0 \times 8c, 0 \times 00, 0 \times 01]
  scope = cw.scope()
  scope.default setup()
  scope.adc.samples = 2000
```

30 target = cw.target(scope)

37

46

54

60

61 62

63

```
32 \text{ ktp} = \text{cw.ktp.Basic()}
33 key, pt = ktp.new pair()
34 \text{ tag} = \text{ktp.next text()}
35 ad = ktp.next text()
36 nonce = generate nonce()
38 target.simpleserial write('k',key)
39 target.simpleserial wait ack()
40 target.simpleserial write('a',ad)
41 target.simpleserial wait ack()
42 target.simpleserial write('n', nonce)
43 target.simpleserial wait ack()
44 target.simpleserial write('t',tag)
45 target.simpleserial wait ack()
47 nonces = []
48 traces = []
49 pts = []
50 \text{ cts} = []
51 \text{ tags} = []
53 for i in tqdm(range(10000)):
       ct = ktp.next text()
       tag = ktp.next text()
       ad = ktp.next text()
       if True :
           nonce = generate nonce()
       target.simpleserial write('a',ad)
       target.simpleserial wait ack()
       target.simpleserial write('t',tag)
   Alexane Boldo (OCIF, IRISA)
```

Main code

target.simpleserial wait ack() target.simpleserial write('n', nonce) target.simpleserial wait ack()

65

69

70

71

73

74

75

77

80

84

87

```
trace = cw.capture trace(scope, target, pt, key)
      traces.append(trace.wave)
      cts.append(trace.textin)
      pts.append(trace.textout)
      tags.append(tag)
      nonces.append(nonce)
76 nptraces = np.asarray(traces)
  np.save("Mesures hf/traces.npy", nptraces)
79 nppts = np.asarray(pts)
 np.save("Mesures hf/pts.npy", nppts)
82 npcts = np.asarray(cts)
83 np.save("Mesures hf/cts.npy", npcts)
npnonces = np.asarray([nonces])
86 np.save("Mesures hf/nonces.npy", npnonces)
88 nptags = np.asarray(tags)
89 np.save("Mesures hf/tags.npy", nptags)
91 plt.plot(traces[0])
92 plt.show()
           Listing: Implementation for trace capture for the ChipWhisperer
```

1 """Using the equations from Modou Sarry during the decryption with a fixed nonce to find the key by finding the constant output of the initialization

2 sbox"""

```
4 using NPZ
5 using Statistics
6 using InformationMeasures
7 using Plots
8
9 plain = npzread("Mesures hf/pts.npy")
10 cipher = npzread("Mesures hf/cts.npv")
11 tags = npzread("Mesures h\bar{f}/tags.npy")
12 traces = npzread("Mesures hf/traces.npy")
13 nonces = npzread("Mesures hf/nonces.npy")[1,:,:]
14
15 function cut(x, len)
16
       [(x << (k-1))>> (len-1) \text{ for } k \text{ in } 1: len]
17
  end
18
19 function assemble(lst,len)
      sum(map(<<,|st|,(|en-1):-1:0))
21 end
23
  function to bits (x, deb, len)
24
       [(UInt8(x[deb + i]) < (k-1)) > 7  for i in 1:len for k in 1:8]
  end
  function to bytes(lst.len)
       [sum(map(<<, |st[j:j+7], 7:-1:0)) for j in 1:8:len]
29
  end
  function sbox 4(k0, k1, n0, n1, v0)
       xor(n0, xor(n1, k0*(xor(1, xor(v0, n1)))))
33 end
34
35
  function sbox 3(k0,k1,n0,n1,v0)
36
       xor(xor(v0,1)*xor(n1, n0), xor(v0, xor(k0,k1)))
```

```
37 end
39
  function sbox 2(k0,k1,n0,n1,v0)
40
       xor(n1*xor(n0,1), xor(1, xor(k0,k1)))
41
  end
42
43
  function sbox 1(k0.k1.n0.n1.v0)
44
       xor(n1, xor(v0, xor(xor(n0,1)*xor(k0,k1), k0*k1)))
45
  end
46
47
  function sbox 0(k0,k1,n0,n1,v0)
48
       xor(n0, xor(v0, xor(k1, k0*(xor(n1, xor(k1, xor(v0,1)))))))
49 end
  function HW(x, nb bits)
       nb = 0
       for i in 0:(nb bits-1)
54
            if (x>>i)\%\overline{2} == 1
                nb += 1
56
           end
       end
58
       nb
59 end
60
61
  function find key (n0, n1, v0, s4)
       if n0 == n1 \&\& n1 == v0
           return s4
       elseif v0 == n1
           return xor(1, s4)
       else
67
            error ("Mauvais nonce")
68
       end
69 end
71 IV =
       cut(0x00001000808c0001,64)
```

```
0 \times 15, 0 \times 88, 09, 0 \times cf, 0 \times 4f, 0 \times 3c, 0, 16)
lst n1 = [to bits(nonces[i,:],8,8)]  for i in 1:size(nonces,1)]
76 function outputs (output, I, jump)
      outputs = zeros(UInt8, size(traces, 1))
       bit output = cut(UInt8(output),8)
       for i in 1: jump
           outputs[i] = output
      end
      m = 0
      for i in 1: size(traces .1)
           if IV[I] == lst n1[i][I]
               m = i
               break
           end
      end
      n0 = |st n0[m][1]
      n1 = |st n1[m][l]
      v0 = IV \overline{1}
      s4 = bit output[8]
      k0 = find key(n0, n1, v0, s4)
      k1 = xor(xor(v0, 1)*xor(n0, n1), xor(v0, xor(k0, bit output[7])))
       for i in (jump+1): jump:(size(traces,1))
           n0 = |st n0[i][1]
           n1 = |st n1[i][l]
           s0 = sbox 0(k0, k1, n0, n1, v0)
           s1 = sbox 1(k0, k1, n0, n1, v0)
           s2 = sbox^{2}(k0, xor(k1, ((l in [57,58,59,60]) ? 1 : 0), n0, n1, v0))
           s3 = sbox^{-}3(k0, k1, n0, n1, v0)
           s4 = sbox^4 (k0, k1, n0, n1, v0)
           output i = assemble([s0, s1, s2, s3, s4], 5)
           for j in 0:(jump-1)
               outputs[i+j] = output i
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                               Side-channel attacks on Ascon's S-box
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                                                                                      10/10
```

72 real key = to bits($[0 \times 2b, 0 \times 7e, 0 \times 15, 0 \times 16, 0 \times 28, 0 \times ae, 0 \times d2, 0 \times a6, 0 \times ab, 0 \times f7,$

74

78

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```
end
       end
       outputs
09 end
10
11
  outs = Array{Dict{Any, Any}}(undef, 64)
12
  for | in 1:64
13
       outs[I] = Dict()
14
       for output in 0:31
           out = outputs(output, 1,1)
           occ = Dict()
17
           for i in 1: size (traces, 1)
               if !(out[i] in keys(occ))
                    occ[out[i]] = 1
               else
                    occ[out[i]] += 1
               end
23
           end
24
           if length (keys(occ)) == 2
25
               outs[]][output] = out
           end
      end
28
  end
31 IM hw = [[[get mutual information(map(x \rightarrow HW(x,5),outs[l][output]),traces[:,t])
        for t in 1: size (traces, 2) for output in keys (outs [1]) for 1 in 1:1]
  IM val = [[[get mutual information(outs[l][output], traces[:,t]) for t in 1:size(
       traces .2) | for output in keys(outs) | for | in 1:64|
33
  plot()
36
  for k in 1:length(keys(outs[1]))
37
       cles = [output for output in keys(outs[1])]
38
       out = cles[k]
```

```
39
      m = 0
       for i in 1: size (traces, 1)
            if IV[1] == lst n1[i][1]
42
                m = i
43
                break
44
           end
45
       end
46
       bit output = cut(UInt8(out),8)
       n0 = |st n0[m][1]
       n1 = |st_n1[m][1]
48
49
       v0 = IV[1]
       s4 = bit output[8]
51
       k0 = find key(n0, n1, v0, s4)
       k1 = xor(xor(v0, 1)*xor(n0, n1), xor(v0, xor(k0, bit output[7])))
       plot!(IM hw[1][k], label = "("*string(k0)*", "*string(k1)*")")
54 end
56
  key guess = Array (UInt8) (undef, 128)
57
  for | in 1:1
58
       possible s = [i for i in keys(outs[1])]
       s max = \overline{0}
       maxi = 0
       for hyp s in 1:length (IM hw[l])
           val max = maximum(IM hw[l][hyp s])
63
           if val max > maxi
64
                s max = possible s[hyp s]
                maxi = val max
66
           end
       end
       println(s max)
69
       m = 0
       for i in 1: size (traces, 1)
            if IV[|] == |st n1[i][|]
               m = i
                break
```

Listing: Analysis in Julia of the traces following the established attack