Huffman Coding

DCTI – IT Companies Seminary

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Short Introduction

About Huffman Coding:

- Entropy encoding algorithm
- Used for lossless data compression
- Is using a variable-length code for encoding a source symbol (such as a character in a file)
- Is using a specific method for choosing the representation for each symbol, resulting in a prefix code ("prefix-free codes")

Given the input text: To be or not to be. This is not a question.

Step 1: The algorithm starts by reading the text and counting how many times each characters is found. As a result we get a Character – Frequency table:

11 11	9
"."	2
" T "	2
" a "	1
" b "	2
" e "	3
" h "	1
" i "	3

" n "	3
" o "	6
" q "	1
" r "	1
" s "	3
" t "	4
" u "	1

Step 2 : - Creating Leaf Nodes

- Creating Internal Nodes



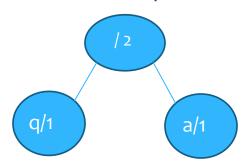
The Tree

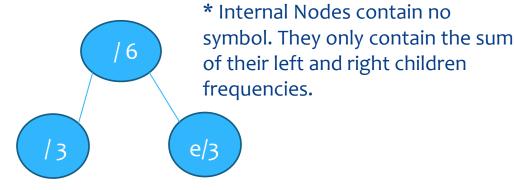
Leaf-Node example: -



* Leaf-Nodes contain Symbols and their frequency.

Internal-Node examples:





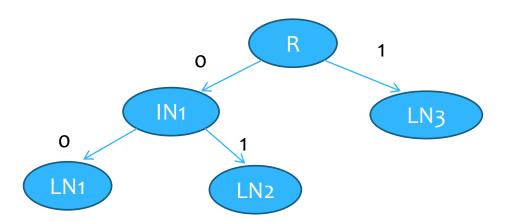
Every LeafNode is inserted into a priority queue ordered by their frequency.



Internal Nodes are created by getting Nodes from the priority queue. We get it's left and it's right child and we form it. After that we insert it in the queue and re order the queue.

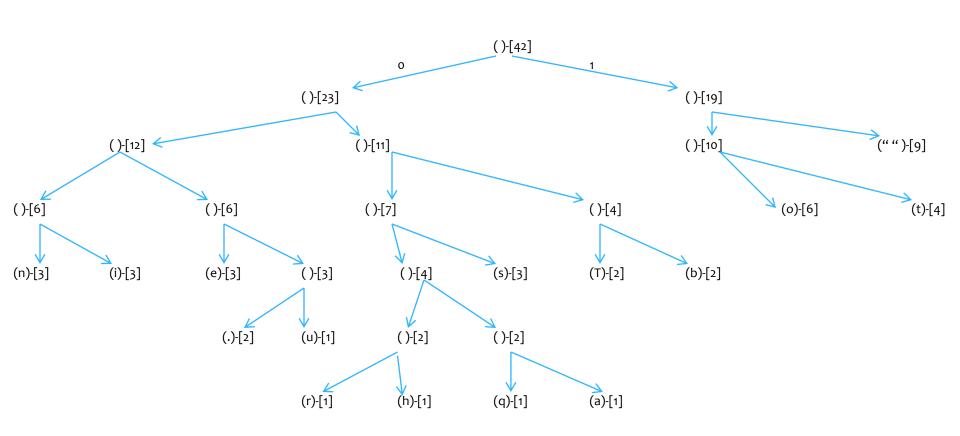


Step 3: After the tree is created it's time to fill it's branches in order to encode the symbols. Every left branch is coded with 'o' while every right branch is coded with '1'.



As a result for the Symbol in LeafNode 1 we have the following code: '00', for LN2: '01' and etc.

The Huffman Tree



As a result we get the **Variable-Length Prefix-Free Codes** for every symbol found in the tree. The code is computed while walking the tree until a symbol is reached.

Prefix-Free Code - the bit string representing some particular symbol is never a prefix of the bit string representing any other symbol.

11 11	11
"."	00110
"T"	0110
" a "	010011
" b "	0111
" e "	0010
" h "	010001
" i "	0001

" n "	0000
" o "	100
" q "	010010
"r"	010000
" s "	0101
" t "	101
" u "	00111

Step 4: Encoding

We replace every symbol in the text with it's Huffman Code.

Input:

To be or not to be. This is not a question.

Output:

T o ''b e ''o

Step 5: Decoding

We walk down the tree while reading the Encoded Sequence. Reading 'o' means going to the left while reading '1' until we find a symbol. The sequence read is replaced by that symbol. This repeats until we finish reading the sequence.

Result: Compression Ratio

Input	42 Symbols	336 Bits	
Output	151 Symbols	151 Bits	2.23

Other Results:

Input	592 Symbols	4736 bits	1.74
Output	2717 Symbols	2717 bits	1.74
Input	1159 Symbols	9272 bits	
Output	5627 Symbols	5627 bits	1.65
Input	2002 Symbols	16736 bits	
Input	2092 Symbols	10/30 DILS	4.54
Output	11095 Symbs	11095 bits	1.51
Input	4095 Symbs	32760 bits	
Output	21707 bits	21707 bits	1.51

Resources

Links:

en.wikipedia.org stackoverflow.com www.youtube.com http://ocw.mit.edu/courses/electrical-engineering-and-computer-science/ http://mathworld.wolfram.com/

Books:

- **1.** Burdescu, Dan Dumitru, Mihaescu, Marian Cristian, Algorithms and Data Structures, Academica Griefswald, 2012.
- **2.**Kernighan, W. Brian, Ritchie, Dennis M., The C Programming Language 2nd edition.
- **3.**Drozdek, Adam, Data Structures Algorithms in C++ Second edition.
- 4. Cormen, Thomas H., Leiserson Charles E., Rivest, Ronald
- R,Stein,Clifford,Introduction to Algorithms 3rd edition

Thank you for your time!

The algorithm can be found at: https://github.com/mihaescu/ADS/tree/master/Data-Compression