### Augmenting Dual Decomposition for MAP Inference

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- Still parallelizable, much faster to reach consensus
- Handles global structural constraints efficiently
- Experiments: Ising models and natural language parsing

#### Outline

- Problem Formulation
- 2 Dual Decomposition
- 3 Augmented Lagrangian Method
- 4 Experiments
- Conclusions

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Let  $\mathbf{X} \triangleq (X_1, \dots, X_n) \in \mathcal{X}$  be a vector of *discrete* random variables

$$P_{\theta,\phi}(\mathbf{x}) \propto \exp\left(\sum_{i=1}^n \frac{\theta_i(x_i)}{\theta_i(x_i)} + \sum_{a \in \mathcal{A}} \phi_a(\mathbf{x}_a)\right)$$

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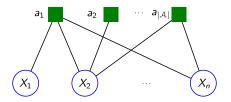
$$P_{\boldsymbol{\theta}, \boldsymbol{\phi}}(\mathbf{x}) \propto \exp\left(\sum_{i=1}^{n} \frac{\theta_{i}(x_{i})}{\theta_{i}(x_{i})} + \sum_{a \in \mathcal{A}} \frac{\phi_{a}(\mathbf{x}_{a})}{\theta_{a}(\mathbf{x}_{a})}\right)$$

 $\mathcal{A}$  is a set of factors,  $\theta_i$  and  $\phi_a$  are unary and higher-order log-potentials

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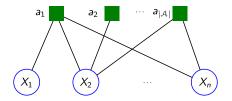
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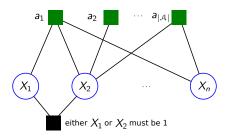
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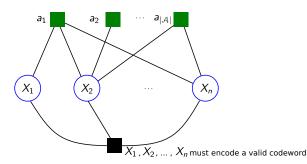
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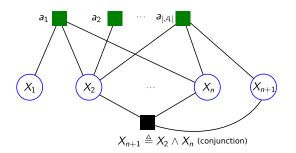
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Goal: compute  $\hat{\mathbf{x}}$  which maximizes  $P_{\theta,\phi}(\mathbf{x})$ 

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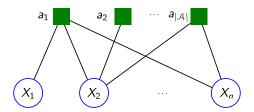
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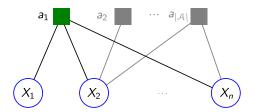
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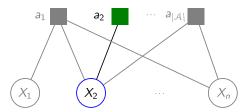
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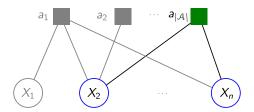
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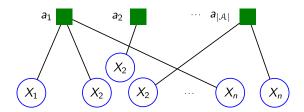
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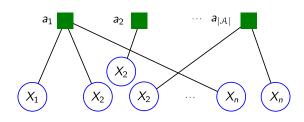








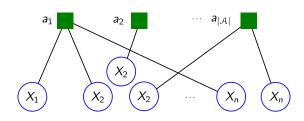




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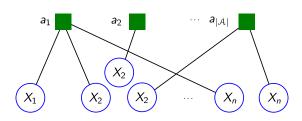
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8 / 22

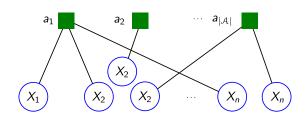


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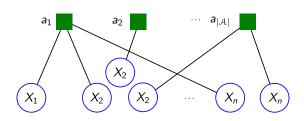
$$\begin{aligned} \max_{\boldsymbol{\mu}, \boldsymbol{\nu}} & & \sum_{i=1}^{n} \boldsymbol{\theta}_{i}^{\top} \boldsymbol{\nu}_{i}^{a} + \sum_{a \in \mathcal{A}} \boldsymbol{\phi}_{a}^{\top} \boldsymbol{\nu}_{a} = \sum_{a \in \mathcal{A}} \left( \sum_{i \in N(a)} d_{i}^{-1} \boldsymbol{\theta}_{i}^{\top} \boldsymbol{\nu}_{i}^{a} + \boldsymbol{\phi}_{a}^{\top} \boldsymbol{\nu}_{a} \right) \\ \text{s.t.} & & & & & & & & & \\ (\boldsymbol{\nu}_{N(a)}^{a}, \boldsymbol{\nu}_{a}) \in \mathcal{M}(\boldsymbol{\mathfrak{I}}_{a}), & \forall a \in \mathcal{A} \\ & & & & & & & & & \\ \boldsymbol{\nu}_{i}^{a} = \boldsymbol{\mu}_{i}, & \forall a \in \mathcal{A}, i \in N(a) \end{aligned}$$



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Problem would be separable, if not for the coupling constraints



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- Dualize them out by adding Lagrange multipliers  $\lambda_i^a$

#### Dual formulation:

#### Slave subproblem at factor a

$$\min_{\boldsymbol{\lambda}} \quad L(\boldsymbol{\lambda}) \triangleq \sum_{\boldsymbol{a} \in \mathcal{A}} \underbrace{\max_{(\boldsymbol{\nu}_{N(\boldsymbol{a})}^{\boldsymbol{a}}, \boldsymbol{\nu}_{\boldsymbol{a}}) \in \mathcal{M}(\boldsymbol{\beta}_{\boldsymbol{a}})} \sum_{i \in N(\boldsymbol{a})} \left( d_{i}^{-1} \boldsymbol{\theta}_{i} + \boldsymbol{\lambda}_{i}^{\boldsymbol{a}} \right)^{\top} \boldsymbol{\nu}_{i}^{\boldsymbol{a}} + \boldsymbol{\phi}_{\boldsymbol{a}}^{\top} \boldsymbol{\nu}_{\boldsymbol{a}}}$$
s.t. 
$$\boldsymbol{\lambda} \in \boldsymbol{\Lambda} \triangleq \left\{ \boldsymbol{\lambda} \mid \sum_{\boldsymbol{a} \in N(\boldsymbol{i})} \boldsymbol{\lambda}_{i}^{\boldsymbol{a}} = 0, \ \forall \boldsymbol{i} \right\}$$

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Slave subproblems: one per each factor  $a \in A$ —parallelizable!

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Master problem: minimize  $L(\lambda)$  via the projected subgradient algorithm

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Subgradient given by the slaves:

$$abla_{\lambda_i^a} L(\lambda) = \hat{
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 (solution of a *local* MAP subproblem at  $\mathfrak{G}_a$ )

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• Projection onto  $\Lambda$ : a simple *centering* operation



9 / 22

### DD-Subgradient Algorithm

```
input: \mathcal{G}, \theta, \phi, number of iterations \mathcal{T}, sequence (\eta_t)_{t=1}^I
Initialize \lambda = 0
for t = 1 to T do
   for each factor a \in A do
       Set unary potentials \omega_i^a = d_i^{-1}\theta_i + \lambda_i^a, for i \in N(a)
       Compute (\hat{\nu}_{N(a)}^a, \hat{\nu}_a) = MAP(\omega_{N(a)}^a, \phi_a)
   end for
   Compute average \bar{\nu}_i = d_i^{-1} \sum_{a:i \in N(a)} \hat{\nu}_i^a
   Update \lambda_i^a \leftarrow \lambda_i^a - \eta_t (\hat{\nu}_i^a - \bar{\nu}_i)
end for
output: \lambda
```

- Converges for a suitable stepsize sequence  $(\eta_t)_{t\in\mathcal{T}}$
- Slow when the number of slaves is large

## Outline

- Problem Formulation
- 2 Dual Decomposition
- 3 Augmented Lagrangian Method
- 4 Experiments
- Conclusions

# $A_{\eta}(\mu, \nu, \lambda) \triangleq \sum_{a \in \mathcal{A}} \left( \sum_{i \in \mathcal{N}(a)} \left( d_i^{-1} \theta_i + \lambda_i^a \right)^\top \nu_i^a + \phi_a^\top \nu_a \right) - \sum_{a \in \mathcal{A}} \sum_{i \in \mathcal{N}(a)} \lambda_i^{a\top} \mu_i - \frac{\eta}{2} \sum_{a \in \mathcal{A}} \sum_{i \in \mathcal{N}(a)} \| \nu_i^a - \mu_i \|^2 \right)$ Residual Term

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Problem: the quadratic term breaks the separability!

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- $lue{}$  Maximize w.r.t.  $\mu$  (closed form)
- **2** Maximize w.r.t.  $\nu$  (can be carried out in parallel at each factor)
- **3** Adjust the Lagrange multipliers  $\lambda_i^a \leftarrow \lambda_i^a au \eta \left( 
  u_i^a \mu_i \right)$



### DD-ADMM Algorithm

```
input: \mathcal{G}, \theta, \phi, number of iterations T, sequence (\eta_t)_{t=1}^T, parameter \tau Initialize \mu uniformly, \lambda = \mathbf{0} for t=1 to T do

for each factor a \in \mathcal{A} do

Set unary potentials \omega_i^a = d_i^{-1}\theta_i + \lambda_i^a + \eta_t\mu_i, for i \in \mathcal{N}(a)

Update (\boldsymbol{\nu}_{\mathcal{N}(a)}^a, \boldsymbol{\nu}_a) \leftarrow \text{QUAD}_{\eta_t}(\boldsymbol{\omega}_{\mathcal{N}(a)}^a, \boldsymbol{\phi}_a)
end for

Update \mu_i \leftarrow d_i^{-1} \sum_{a:i \in \mathcal{N}(a)} (\boldsymbol{\nu}_i^a - \eta_t^{-1} \lambda_i^a)
Update \lambda_i^a \leftarrow \lambda_i^a - \tau \eta_t (\boldsymbol{\nu}_i^a - \mu_i)
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output: \mu, \nu, \lambda
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### DD-ADMM Algorithm

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• Instead of MAP, we solve a quadratic problem QUAD at each factor:

$$\min_{(\boldsymbol{\nu}_{N(a)}^{a},\boldsymbol{\nu}_{a})\in\mathbb{M}(\boldsymbol{\S}_{a})}\frac{\eta_{t}}{2}\sum_{i\in\boldsymbol{a}}\|\boldsymbol{\nu}_{i}^{a}-\eta_{t}^{-1}\boldsymbol{\omega}_{i}^{a}\|^{2}-\boldsymbol{\phi}_{a}^{\top}\boldsymbol{\nu}_{a},$$

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• Under certain conditions, converges even when QUAD is approximately solved (Eckstein and Bertsekas, 1992).

Binary pairwise factors: closed form solution

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Hard constraint factors (Tarlow et al., 2010; Martins et al., 2010):

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- For many hard factors imposing logical constraints, can be done efficiently with sort operations

Larger slaves (sequences, trees): solved approximately with primal-dual cyclic projection algorithms (work in progress)

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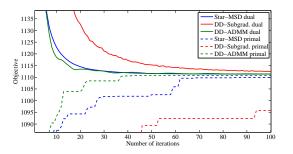
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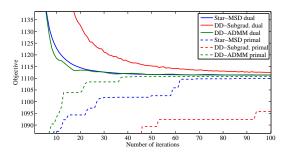
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- DD-subgradient is the slowest (many slaves!)
- DD-ADMM outperforms the others: it approaches a near optimal primal-dual solution in a few tens of iterations

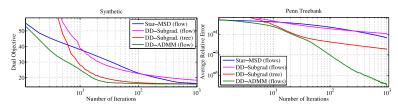
A problem with heavily constrained outputs, to which DD has recently been applied (Rush et al., 2010; Koo et al., 2010)



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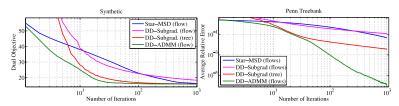
Our model: 2nd-order with  $O(n^3)$  slaves (distinct from Koo et al. 2010)



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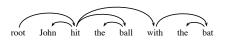


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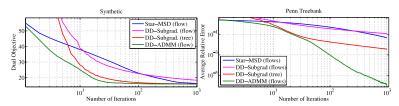


• DD-ADMM wins with both synthetic (10 words) and real sentences

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- DD-ADMM wins with both synthetic (10 words) and real sentences
- Also outperforms DD-Subgradient with a TREE model (fewer slaves)

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#### Conclusions

- DD-ADMM: a new algorithm for LP-MAP inference
- Dual decomposable, hence the slaves can all be solved in parallel
- Allies the simplicity of DD with the effectiveness of AL methods
- Suitable for problems with many slaves, outperforming Komodakis et al. (2007)
- Optimality certificates for LP-MAP (not just MAP)
- A significant amount of computation can be saved by caching and warm-starting the subproblems
- Related work in accelerating DD and in quadratic projections: Jojic et al. (2010); Ravikumar et al. (2010)
- Future work: larger slaves and approximate ADMM steps

## Thank you!

• Questions?

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