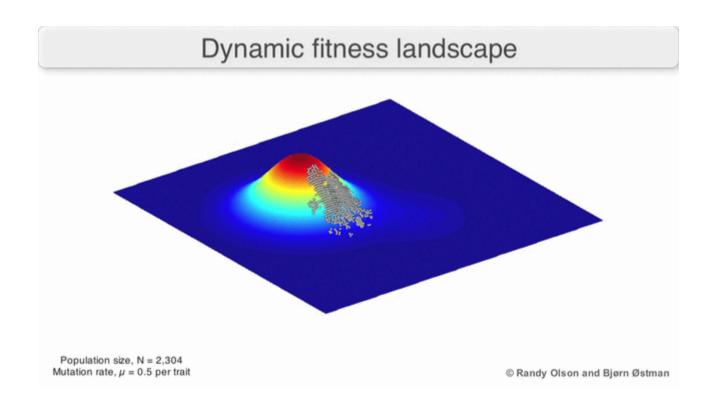
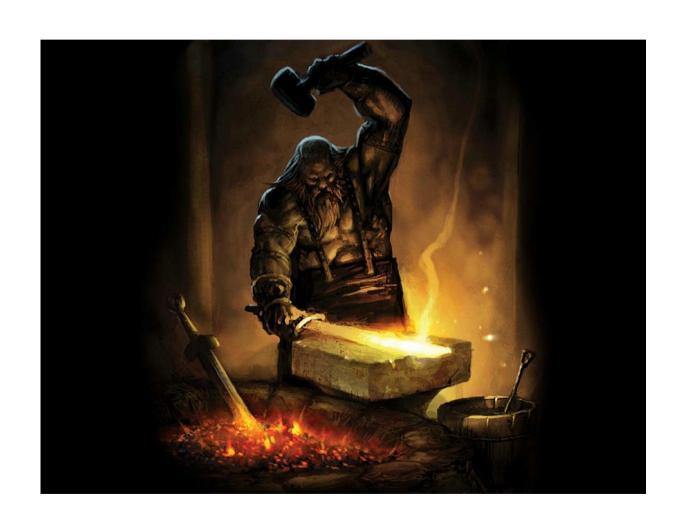
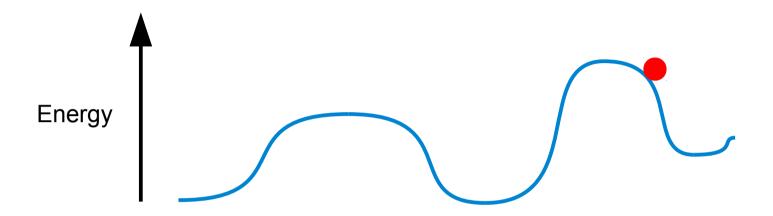
Stochastic Optimization



Simmulated Annealing



Simmulated Annealing

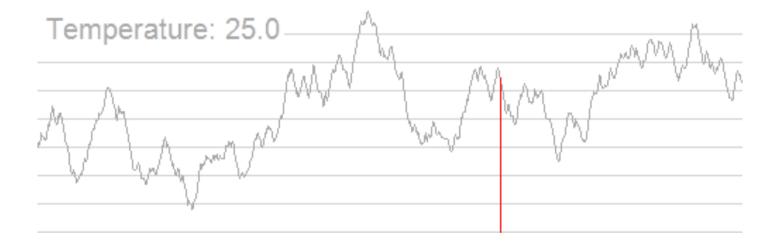


Greedy algorithms can get stuck in local optima, and depend on initial conditions :(

Simmulated Annealing procedure:

- Initial configuration \emph{i} with $E_\emph{i}$
- Choose random (nearby) state \emph{j} with $E_\emph{j}$
- Accept change i o j with probability $P_{i o j} = \min \left[1, \exp \left(\frac{E_j E_i}{T} \right) \right]$
- Gradually decrease T according to a "cooling schedule"

Simmulated Annealing



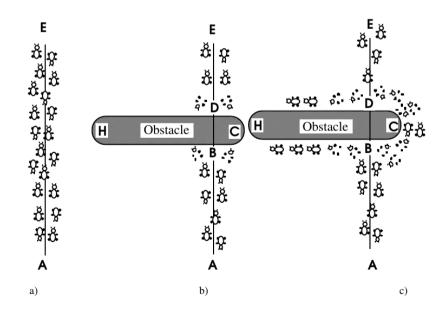


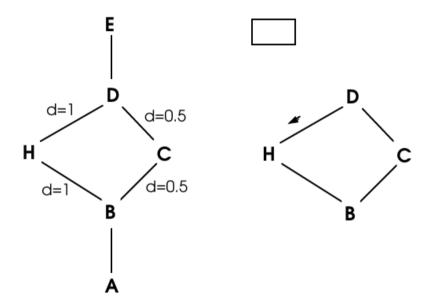
Swarm intelligence

Submitted to IEEE Transactions on Systems, Man, and Cybernetics

The Ant System: Optimization by a colony of cooperating agents

Marco Dorigo, Vittorio Maniezzo, Alberto Colorni





TSP: Which is the shortest path that visits each city exactly once?

Let $\tau_{ij}(t)$ be the *intensity of trail* on edge (i,j) at time t. Each ant at time t chooses the next town, where it will be at time t+1. Therefore, if we call an *iteration* of the AS algorithm the m moves carried out by the m ants in the interval (t, t+1), then every n iterations of the algorithm (which we call a cycle) each ant has completed a tour. At this point the trail intensity is updated according to the following formula

$$\tau_{ij}(t+n) = \rho \cdot \tau_{ij}(t) + \Delta \tau_{ij} \tag{1}$$

where

 ρ is a coefficient such that $(1 - \rho)$ represents the *evaporation* of trail between time t and t+n,

$$\Delta \tau_{ij} = \sum_{k=1}^{m} \Delta \tau_{ij}^{k} \tag{2}$$

where $\Delta \tau_{ij}^k$ is the quantity per unit of length of trail substance (pheromone in real ants) laid on edge (i,j) by the k-th ant between time t and t+n; it is given by

$$\Delta \tau_{ij}^k = \begin{cases} \frac{Q}{L_k} & \text{if } k \text{ - th ant uses edge } (i,j) \text{ in its tour (between time t and } t+n) \\ \\ 0 & \text{otherwise} \end{cases}$$
 (3)

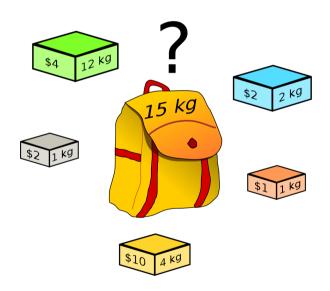
where Q is a constant and Lk is the tour length of the k-th ant.

We call visibility η_{ij} the quantity $1/d_{ij}$. This quantity is not modified during the run of the AS, as opposed to the trail which instead changes according to the previous formula (1).

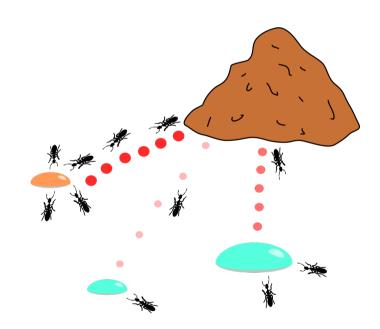
We define the transition probability from town i to town j for the k-th ant as

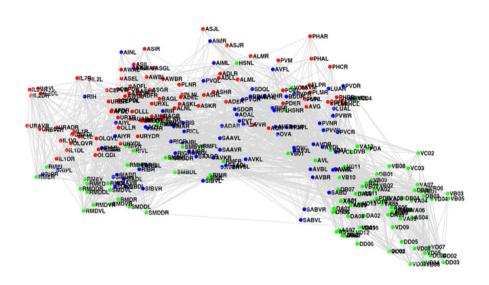
$$p_{ij}^{k}(t) = \begin{cases} \frac{\left[\tau_{ij}(t)\right]^{\alpha} \cdot \left[\eta_{ij}\right]^{\beta}}{\sum_{k \in allowed_{k}} \left[\tau_{ik}(t)\right]^{\alpha} \cdot \left[\eta_{ik}\right]^{\beta}} & \text{if } j \in allowed_{k} \\ 0 & \text{otherwise} \end{cases}$$

(4)



Knapsack problem: Which boxes should be chosen to maximize the amount of money while still keeping the overall weight under or equal to 15 kg?



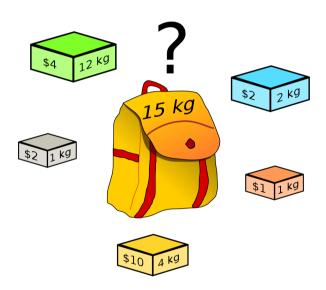


$$Q = \frac{1}{2m} \sum_{vw} \left[A_{vw} - \frac{k_v k_w}{2m} \right] \delta(c_v, c_w)$$

Finding community structure in very large networks

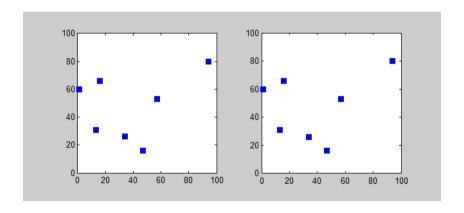
Aaron Clauset, M. E. J. Newman, and Cristopher Moore 1,3

Knapsack problem:



Which boxes should be chosen to maximize the amount of money while still keeping the overall weight under or equal to 15 kg?

Travelling salesman problem:



Which is the shortest path that visits each city exactly once?

