# DIVIDE AND CONQUER

#### **DIVIDE AND CONQUER**

#### Concept

a recursive method of breaking down a problem into unique sub problems to be solved and combined into a solution for the original larger problem

at first this definition may sound like Dynamic Programming where a large problem is solved by resolving smaller sub problems

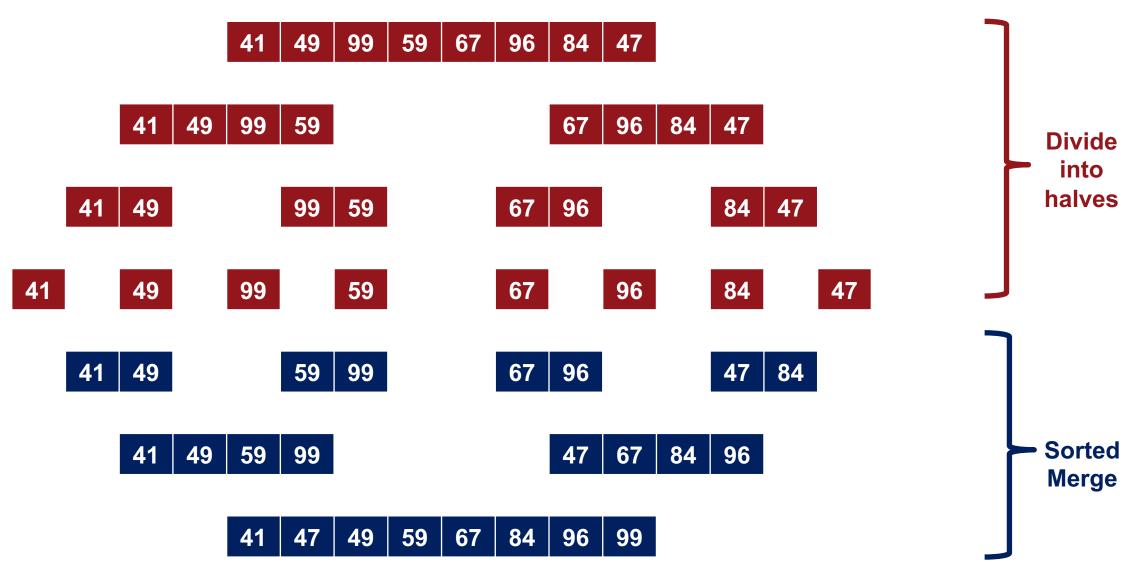
Dynamic Programming is specifically used to optimize overlapping sub problems by replacing repetitive sub problem calculations with memory storage and lookup

Divide and Conquer divides a large problem into unique sub problems which are not overlapping

#### Merge Sort

```
Concept
            recursively divide array into halved subsets which are merged into sorted order
            O(n log n) time complexity
Example
            void sort(int *a, int start, int end) {
                                                        // recursively halve array into subsets
                                                        // recurse until out of values
                if( start < end) {</pre>
                    int middle = (start + end) / 2;
                                                        // find middle
                                                 // sort lower half
                    sort(a, start, middle);
                    sort(a, middle+1, end);
                                                        // upper half
                    merge(a, start, middle, end);
                                                        // merge
            void merge(int *a, int start, int middle, int end) {
                    // store left and right halves of the array a into new left/right arrays
                    // sort the smaller values of left/right arrays into the array a
                    // sort any remaining values from left/right arrays into array a
```

# Merge Sort



#### **Quick Sort**

Concept recursively select a pivot value from the array, then partition the array such that: all values less than the pivot value have a lower index then the pivot value all values greater than the pivot value have a higher index then the pivot value  $O(n^2)$  time complexity but can be optimized to compete with  $O(n \log n)$  algorithms

```
// recursively partition the array
            void sort(int *a, int low, int high) {
Example
                if( low < high) {</pre>
                                                         // if values remain in the array
                    int pivot = partition(a, low, high); // find new pivot value
                    sort(a, low, pivot-1);
                                                         // sort partition before pivot
                    sort(a, pivot+1, high);
                                                             sort partition after pivot
            int partition(int *a, int low, int high) {
                // set pivot to last value
                // place all values relative to the pivot (less than or greater than)
                // swap pivot into its relative position
```

# **Quick Sort**

	22	62	48	91	12	15
PIVOT	15					
	22	62	48	91	12	15
	12	62	48	91	22	15
	12	15	48	91	22	62
PIVOT	62					
	48	91	22	62		
	48	22	91	62		
	48	22	62	91		
PIVOT	22					
	48	22				
	22	48				
	4	32	43	56	74	75

### Merge Sort vs. Quick Sort

Merge Sort

requires more space than quick sort because it creates new arrays

optimal for large data structures stored on disk due to less read operations

**Quick Sort** 

requires less space then merge sort because it sorts in place

optimal for smaller data structures stored in memory

randomized pivot value improves average time complexity making it potentially faster then merge sort

considered to be the fastest sorting method for most applications