LOCAL MEMORY

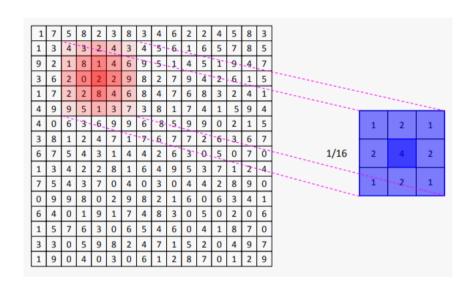
LEARNING OBJECTIVES

- Learn about tiling using local memory
- Learn about how to synchronize work-groups

COST OF ACCESSING GLOBAL MEMORY

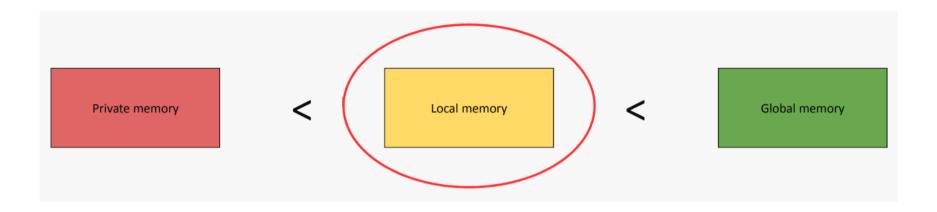
- As we covered earlier global memory is very expensive to access.
- Even with coalesced global memory access if you are accessing the same elements multiple times that can be expensive.
- Instead you want to cache those values in a lower latency memory.

WHY ARE IMAGE CONVOLUTIONS GOOD ON A GPU?



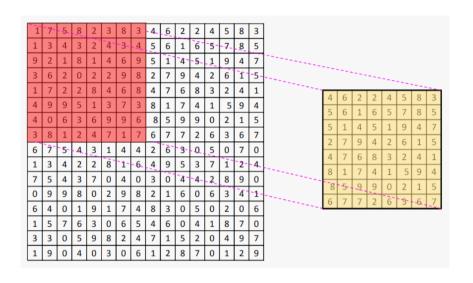
- Looking at the image convolution example.
- For each output pixel we are reading up to NxM pixels from the input image, where N and M are the dimensions of the filter.
- This means each input pixel is being read up to NxM times:
- 3x3 filter: up to 9 ops.
- 5x5 filter: up to 25 ops.
- 7x7 filter: up to 49 ops.
- If each of these operations is a separate load from global memory this becomes very expensive.

USING LOCAL MEMORY



- The solution is local memory.
- Local memory is generally on-chip and doesn't have a cache as it's managed manually so is much lower latency.
- Local memory is a smaller dedicated region of memory per work-group.
- Local memory can be used to cache, allowing us to read from global memory just once and then read from local memory instead, often referred to as a scratchpad.

TILING



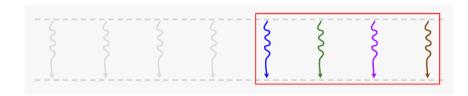
- The iteration space of the kernel function is mapped across multiple work-groups.
- Each work-group has it's own region of local memory.
- You want to split the input image data into tiles, one for each workgroup.

LOCAL ACCESSORS

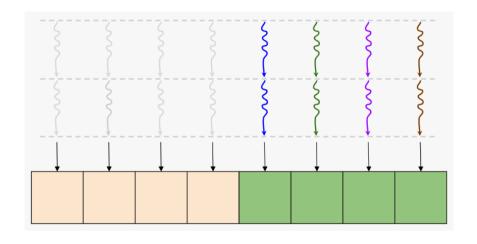
```
auto scratchpad = sycl::accessor<int, 1, sycl::access::target::local>(sycl::range{workGroupSize}, cgh);
```

- Local memory is allocated via an accessor with the access::target::local access target.
- Unlike regular accessors they are not created with a buffer, they allocate memory per work-group for the duration of the kernel function.
- The range provided is the number of elements of the specified type to allocate per work-group.

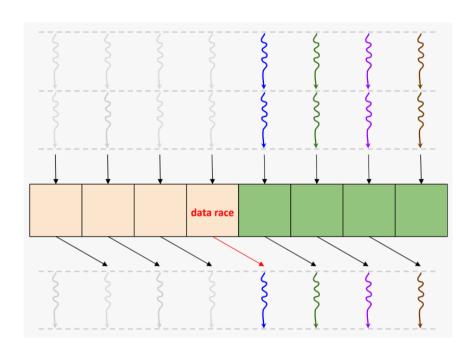
- Local memory can be used to share partial results between work-items.
- When doing so it's important to synchronize between writes and read to memory to ensure all work-items have reached the same point in the program.



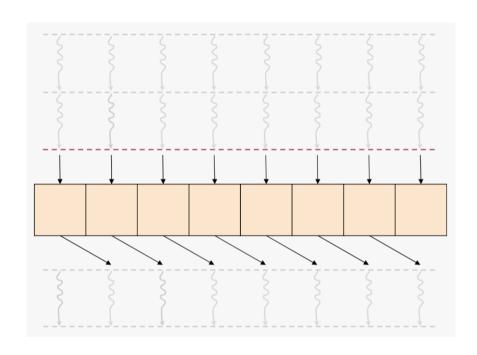
• Remember that work-items are not all guaranteed to execute concurrently.



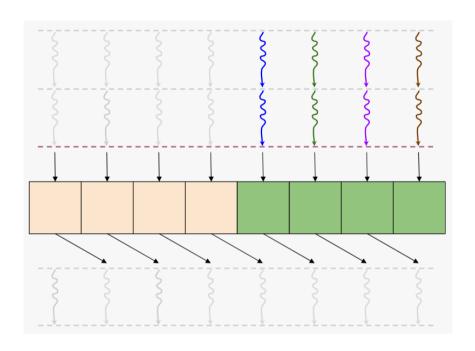
• A work-item can share results with other work-items via local (or global) memory.



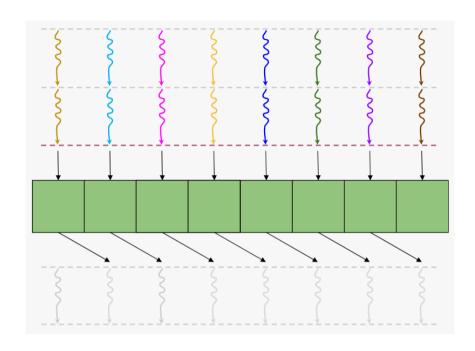
- This means it's possible for a work-item to read a result that hasn't been written to yet.
- This creates a data race.



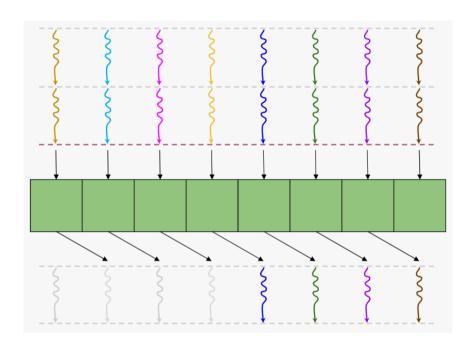
• This problem can be solved with a synchronization primitive called a work-group barrier.



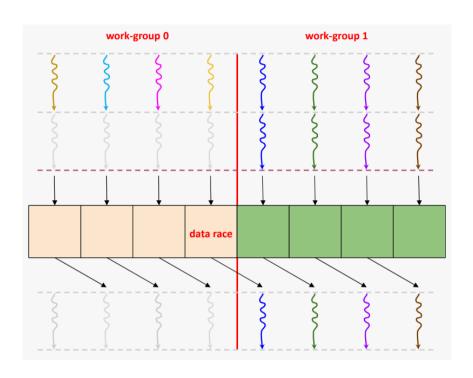
 When a work-group barrier is inserted work-items will wait until all work-items in the work-group have reached that point.



• Only then can any work-items in the work-group continue execution.



• So now you can be sure that all of the results that you want to read have been written to.



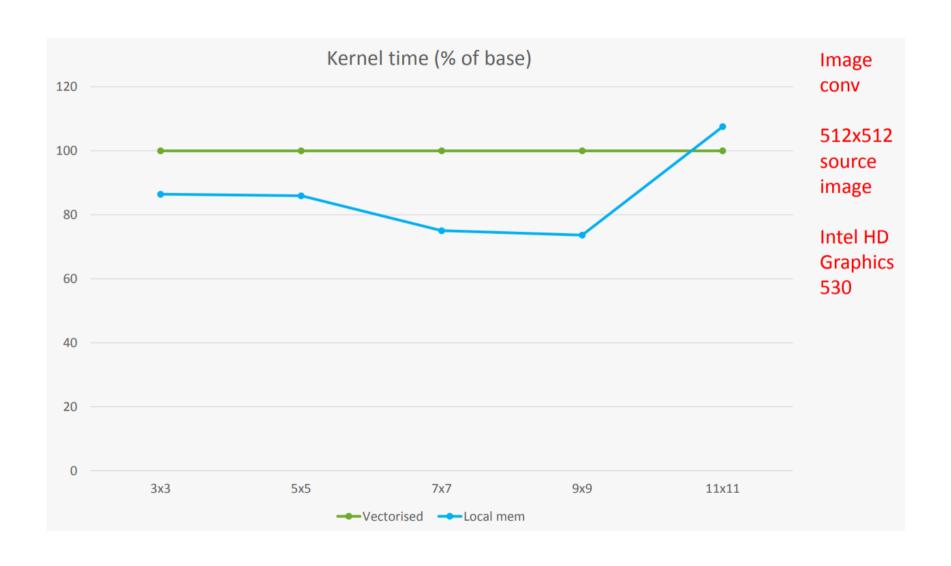
- However note that this does not apply across work-group boundaries.
- So if you write in a work-item of one work-group and then read it in a work-item of another workgroup you again have a data race.
- Furthermore, remember that work-items can only access their own local memory and not that of any other work-groups.

GROUP_BARRIER

```
sycl::group_barrier(item.get_group());
```

- Work-group barriers can be invoked by calling group_barrier and passing a group object.
- You can retrieve a group object representing the current work-group by calling get_group on an nd_item.
- Note this requires the nd_range variant of parallel_for.

LOCAL MEMORY IMAGE CONVOLUTION PERFORMANCE



QUESTIONS

EXERCISE

Code_Exercises/Exercise_18_Local_Memory_Tiling/source

Use local memory to cache a tile of the input image data per work-group.