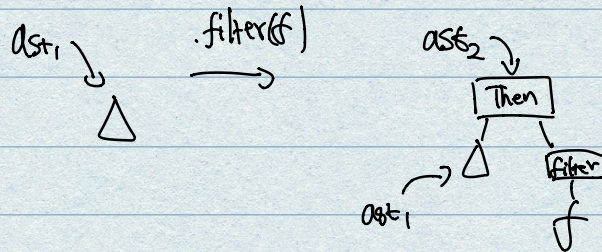


Lecture #5.

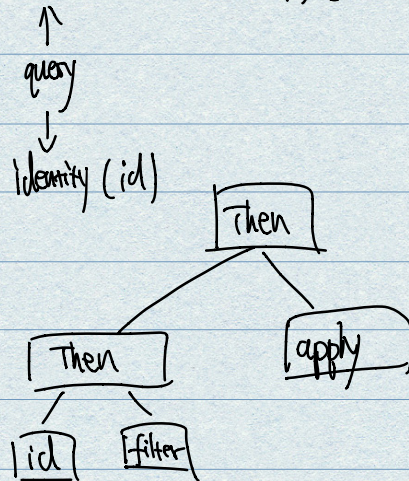
Implementation of fluent AST construction.

$Q.\text{apply}(\text{Math.sqrt}) \dots \text{filter}(f)$
└──────────┘
ast₁



Adjust for call chaining

$Q.\text{filter}(f).\text{apply}(g)$



Q has a type of `ASTNode` (subtype `IdNode`)

Optimization.

inefficiency of ' $q.\text{filter}(f).\text{count}()$ '

- ①. stores the data
- ②. two iterations.

ThenNode(ThenNode(x, FilterNode(f)), count)
→ ThenNode(x, countIf(f))
" tree rewrite rule "

Optimization traversal order.

- Bottom up / Top down ?
- post-order optimization

{ 1 → left child
2 → right child
3 → this node

(Lecture #5)

Global Optimization,

- control flow.
- multiple execution paths
- Control Flow Graph.

⇒ • Global Flow Analysis

- Constant propagation
- Liveness analysis (eliminate redundant stmts)

Local Optimization.

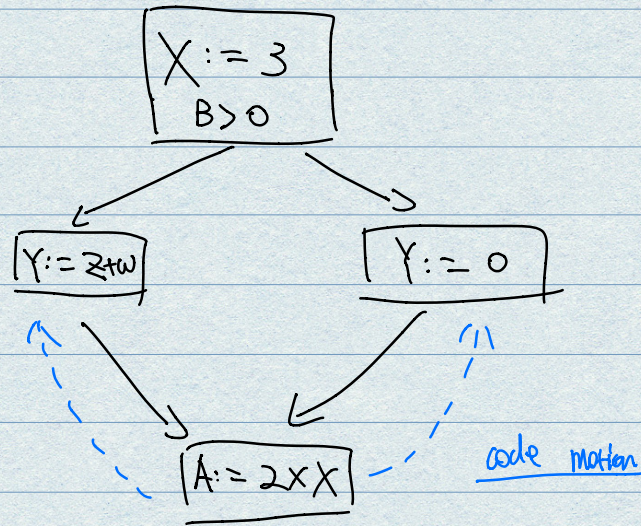
- Basic block: code sequence with no jumps
- Basic block optimization:
 - Constant propagation.
 - Dead code elimination.

- etc

$X := 3$

$Y := Z \times W$

$Q := X + Y$



To Replace a use of x by a constant k .

→ every path to the use of x

the last assignment to x is $x := k$.

Global Analysis

• The optimization depends on the property Prop at a particular point in program execution,

⌋ Prop is definitely true

⌋ don't know whether Prop is true.