

1. Implement the Backtracking algorithm for the n -Queens problem (Algorithm 5.1) on your system, and run it on problem instances in which $n = 4, 8, 10$, and 12 .

The Backtracking Algorithm for the n -Queens Problem

Problem: Position n queens on a chessboard so that no two are in the same row, column, or diagonal.

Inputs: positive integer n .

Outputs: all possible ways n queens can be placed on an $n \times n$ chessboard so that no two queens threaten each other. Each output consists of an array of integers col indexed from 1 to n , where $col[i]$ is the column where the queen in the i th row is placed.

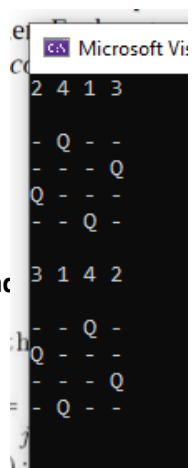
```
void queens (index i)
{
    index j;

    if (promising(i))
        if (i == n)
            cout << col[1] through col [n];
        else
            for (j = 1; j <= n; j++){          // See if queen in
                col[i + 1] = j;                // (i + 1)st row can be
                queens(i + 1);                  // positioned in each of
            }                                   // the n columns.
}

bool promising (index i)
{
    index k;
    bool switch;

    k = 1;
    switch = true;          // Check if any queen threatens
    while (k < i && switch){ // queen in the ith row.
        if (col[i] == col[k] || abs(col[i] - col[k]) == i - k)
            switch = false;
        k++;
    }
    return switch;
}
```

2. Visualize the results in 2D. See Examples for $n = 4$
3. Modify the Backtracking algorithm for then-Queens problem (Algorithm 5.1) so that, instead of generating all possible solutions, it finds only a single solution.



2. Given an $n \times n \times n$ cube containing n^3 cells, we are to place n queens in the cube so that no two queens challenge each other (so that no two queens are in the same row, column, or diagonal). Can then-Queens algorithm (Algorithm 5.1, see AZA lab 5 sheet) be extended to solve this problem? If so, write the algorithm and implement it on your system to solve problem instances in which $n = 4$ and $n = 8$.