# **GROUP 5**

# **DBMS For Art Gallery Management**

### **ABSTRACT**

The Art Gallery Management System aims to efficiently manage and organize the database of an art gallery, providing accurate tracking and storage of art and paintings details. The system enhances safety and efficiency in managing exhibitions, gallery operations, and art stocks. Here's an abstract:

This project focuses on the design and implementation of an Art Gallery Management System, utilizing a MYSQL database platform with PHP and WAMP Server support. The system is developed to ensure effective management of art-related activities within the gallery. The primary functionalities include:

- 1. **Exhibition Management:** Tracks details of ongoing and upcoming exhibitions, including exhibit themes, dates, and participating artists.
- 2. **Gallery Management:** Manages gallery operations such as inventory management, artwork categorization, and gallery space utilization.
- 3. **Art Stocks:** Maintains a comprehensive database of art pieces and paintings, including details such as artist information, artwork descriptions, and pricing.

The application interface is developed using HTML5 and CSS3, with PHP facilitating dynamic content generation and interaction with the MYSQL database. Through this Art Gallery Management System, gallery administrators can streamline administrative tasks, enhance visitor experience, and effectively showcase the gallery's collection to art enthusiasts and patrons.

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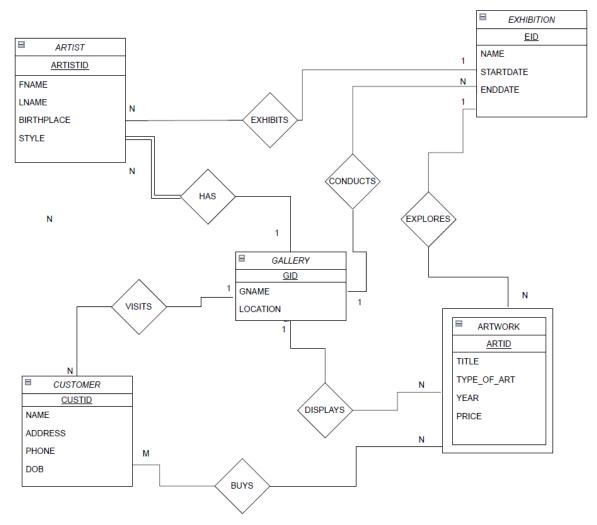
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## **UNVERSAL TABLE**

Attribute Name	Data Type	Description	
GID	INT	Gallery ID (Primary Key)	
GNAME	VARCHAR(100)	Gallery Name	
LOCATION	VARCHAR(200)	Gallery Location	
EID	INT	Exhibition ID (Primary Key)	
STARTDATE	DATE	Exhibition Start Date	
ENDDATE	DATE	Exhibition End Date	
	ARTISTID	INT	
Artist ID (Primary Key)			
FNAME	VARCHAR(50)	Artist First Name	
LNAME	VARCHAR(50)	Artist Last Name	
BIRTHPLACE	VARCHAR(100)	Artist Birthplace	
STYLE	VARCHAR(100)	Artist Style	
ARTID	INT	Artwork ID (Primary Key)	
TITLE	VARCHAR(200)	Artwork Title	
TYPE_OF_ART	VARCHAR(100)	Type of Art	
YEAR	INT	Year of Creation	
PRICE	DECIMAL(10, 2)	Artwork Price	
CUSTID	INT	Customer ID (Primary Key)	
STREET	VARCHAR(200)	Customer Street Address	
CITY_ID	INT	City ID (Foreign Key)	
DOB	DATE	Customer Date of Birth	
PHONE	VARCHAR(15)	Customer Phone Number	
CITY	VARCHAR(100)	City Name	

STATE	VARCHAR(100)	State Name
ZIPCODE	VARCHAR(20)	Zip Code

### **ER-DIAGRAM**



# MAPPING OF ER DIAGRAM TO RELATIONS

# **STEP 1: Mapping of Regular Entities**

For each regular entity type E in the ER schema, create relation R that includes all simple attributes of E.

#### **GALLERY**

GID	GNAME	LOCATION

#### **EXHIBITION**

<b>EID</b> STARTDATE		ENDDATE

## ARTIST

ARTISTID	FNAME	LNAME	BIRTHPLACE	STYLE

### **CUSTOMER**

CUSTID	ARTID	FNAME1	LNAME1	ADDRESS	PHONE	DOB

# **STEP 2: Mapping of Weak Entity Types**

#### **ARTWORK**

ARTID	ARTISTID	TITLE	TYPE_OF_ART	YEAR	PRICE
	FK	I			

# **STEP 3: Mapping of 1:1 Relationship**

Identify the relation S that represents the participating entity type at the 1-side of the relationship type.

Include as foreign key in S the primary key of the relations T that represents the other entity type participating in R.

For each binary 1:1 relationship type R in ER schema, identify the relations S and T that correspond to the entity types participating in R if any.

There are **no** 1:1 relationship.

## STEP 4: Mapping of 1:N Relationship

#### **EXHIBITION**

EID	STARTDATE	ENDDATE	GID

FK

#### **ARTIST**

ARTISTID	FNAME	LNAME	BIRTHPLACE	STYLE	EID	GID	CUSTID	1
					1		FK I	īΚ

#### **CUSTOMER**

CUSTID	ARTID	FNAME1	LNAME1	STREET	CITY_ID	DOB	GID
	FK				FK	FK	

#### **ARTWORK**

ARTID	ARTISTID	TITLE	TYPE_OF_ART	YEAR	PRICE	EID	GID
	FK					F	K FK

### CITY\_ID

CITY ID	CITY	STATE	ZIPCODE

## **STEP 5: Mapping of M:N Relationship**

Create a new relation S to represent R.

Include as foreign key attributes in S the primary key of the relations that represents the participating entity types their combination will form the primary key of S.

Also, include any simple attributes of the M:N relationship type as attributes of S.

## **STEP 6: Mapping of Multi-Valued Attributes**

For each multivalued attributes A, create a new relation R. This relation R will include an attribute corresponding to A, plus the primary key attribute K-as a foreign key in R-

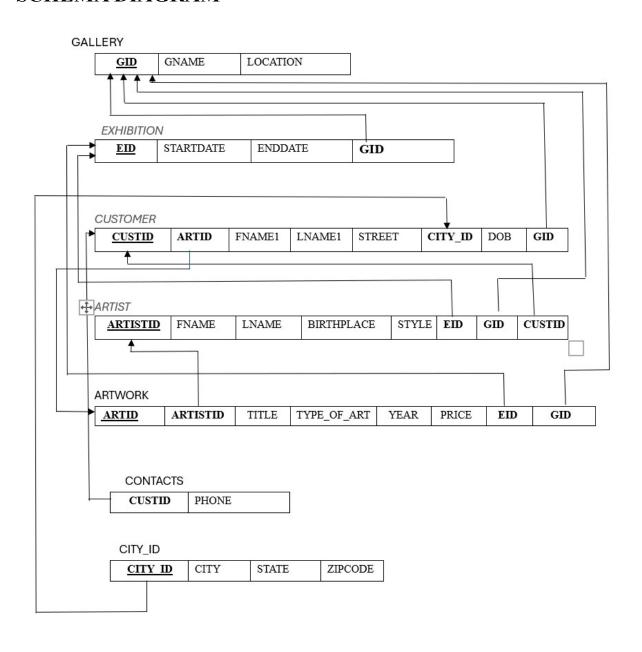
of the relation that represents the entity type of relationship type that has A as an attribute.

The Primary Key of R is the combination of A and K. If the multivalued attribute is composite, we include its simple components.

#### **CONTACTS**

CUSTID	PHONE

### **SCHEMA DIAGRAM**



## NORMALIZE THE RELATIONS

Database normalization, or simply normalization, is the process of organizing the columns(attributes) and tables(relations) of a relational database to reduce data redundancy and improve data integrity. Normalization involves arranging attributes in relations based on dependencies between attributes.

#### 1. First Normal Form

As per First normal form, no two rows of data must contain repeating group of information. Each set of columns must have a unique value, such that multiple columns cannot be used to fetch the same row. Each table should be organized into rows, and each row should have a primary key that will distinguishes it as unique.

#### **Example:**

#### **GALLERY**

GID	GNAME	LOCATION		

All the tables in the database are normalized to 1NF as all the attributes are atomic.

## 2. Second Normal Form (2NF)

A table is in 2NF if it is in 1NF and if all non-key attributes are fully functionally dependent on all of the key.

#### **Example:**

### **CUSTOMER**

CUSTID	ARTID	FNAME1	LNAME1	STREET	CITY_ID	DOB	GID
FD1							

#### FD1

CUSTID	FNAME1	LNAME1	DOB

### 3. Third Normal Form(3NF):

A table is in 3NF if it is in 2NF and if it has no transitive dependency. X->Y, Y->Z, X>Z

According to CODD's definition a relation schema R is in 3NF. It satisfies 2NF and no non-prime attribute of R is transitively dependent on the primary key. All tables of database satisfies upto 3NF.

### **CREATION OF TABLES**

### **Table 1: GALLERY**

```
CREATE TABLE GALLERY (
GID INT PRIMARY KEY,
GNAME VARCHAR(100) NOT NULL,
LOCATION VARCHAR(200) NOT NULL
);
```

### **Table 2: EXHIBITION**

```
CREATE TABLE EXHIBITION (
EID INT PRIMARY KEY,
STARTDATE DATE NOT NULL,
ENDDATE DATE NOT NULL,
GID INT,
FOREIGN KEY (GID) REFERENCES GALLERY(GID)
);
```

### **Table 3: ARTIST**

```
CREATE TABLE ARTIST (
ARTISTID INT PRIMARY KEY,
FNAME VARCHAR(50) NOT NULL,
LNAME VARCHAR(50) NOT NULL,
BIRTHPLACE VARCHAR(100),
STYLE VARCHAR(100)
);
```

### **Table 4: ARTWORK**

```
CREATE TABLE ARTWORK (
ARTID INT PRIMARY KEY,
ARTISTID INT,
TITLE VARCHAR(200) NOT NULL,
TYPE_OF_ART VARCHAR(100),
YEAR INT,
PRICE DECIMAL(10, 2),
EID INT,
GID INT,
FOREIGN KEY (ARTISTID) REFERENCES ARTIST(ARTISTID),
FOREIGN KEY (EID) REFERENCES EXHIBITION(EID),
FOREIGN KEY (GID) REFERENCES GALLERY(GID)
);
```

### **Table 5: CUSTOMER**

```
CREATE TABLE CUSTOMER (
CUSTID INT PRIMARY KEY,
FNAME VARCHAR(50) NOT NULL,
LNAME VARCHAR(50) NOT NULL,
STREET VARCHAR(200),
CITY_ID INT,
DOB DATE,
GID INT,
FOREIGN KEY (GID) REFERENCES GALLERY(GID)
);
```

### **Table 6: CITY**

```
CREATE TABLE CITY (
CITY_ID INT PRIMARY KEY,
CITY VARCHAR(100) NOT NULL,
STATE VARCHAR(100) NOT NULL,
ZIPCODE VARCHAR(20) NOT NULL
);
```

### **Table 7: CONTACTS**

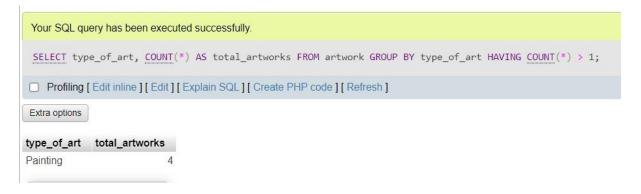
```
CREATE TABLE CONTACTS (
CUSTID INT,
PHONE VARCHAR(15),
PRIMARY KEY (CUSTID, PHONE),
FOREIGN KEY (CUSTID) REFERENCES CUSTOMER(CUSTID)
);
```

### **Additional Considerations**

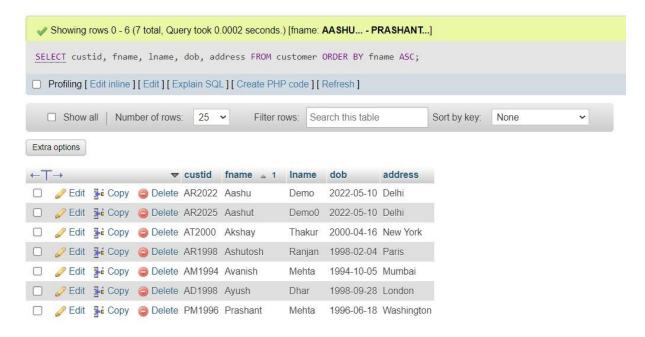
- All primary key constraints ensure that each record is uniquely identifiable.
- Foreign key constraints establish the relationships between tables, ensuring referential integrity.

# **Execution of Queries**

# 1) Aggregate functions, Group by...having



## 2) Order by



## 3) Join, Outer Join



# 4) Query having Boolean operators



## 5) Query having arithmetic operators



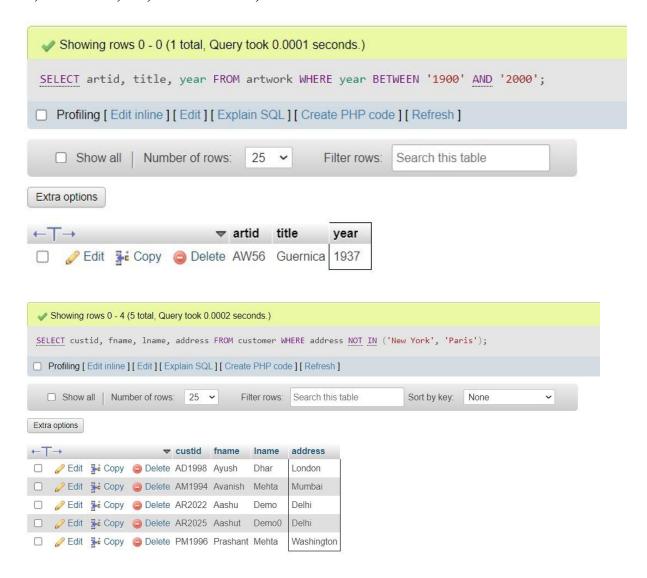
## 6) A search query using string operators



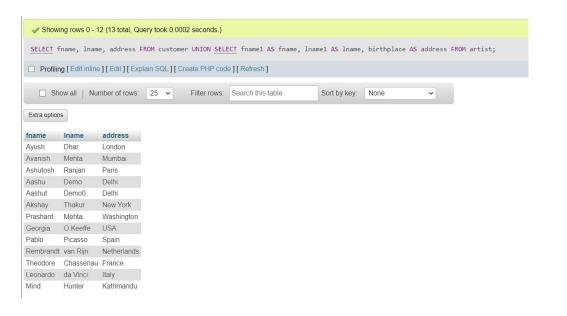
## 7) Usage of to\_char, extract



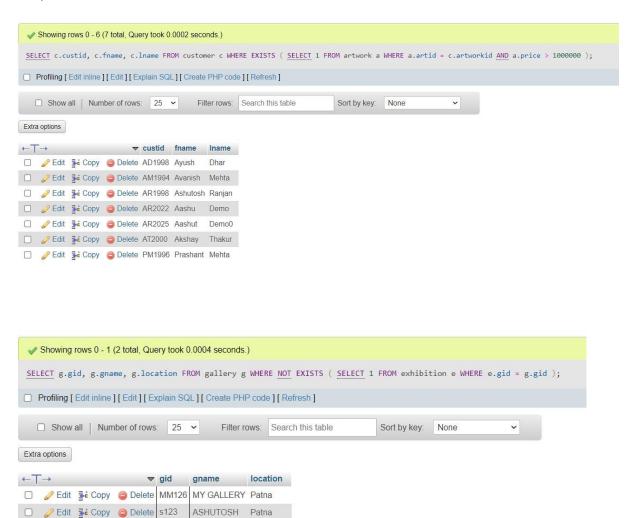
## 8) Between, IN, Not between, Not IN



## 9) Set operations



# 10) Subquery using EXISTS / NOT EXISTS, ANY, ALL

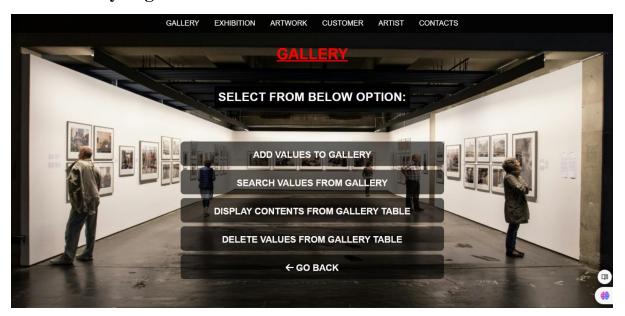


### **User Interface**

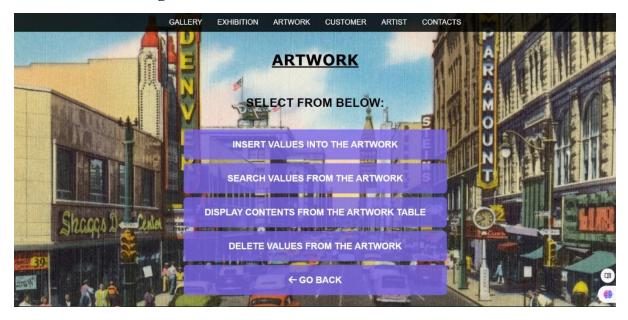
# Landing Page



## • Gallery Page



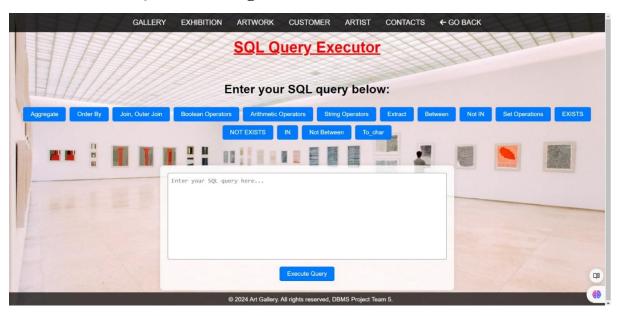
## Artwork Page



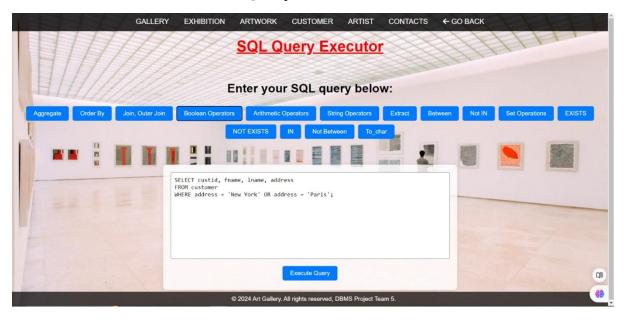
## • Artwork Table Display

					V 100 100		
Artwork ID	Title	Year	Type of Art	Price	E_ID	G_ID	Artist ID
AW12	Mona Lisa	1503	Painting	1000000000.00	G123	NG123	AD11
AW34	Poppies	1873	Painting	15000000.00	H123	MM123	AD22
AW56	Guernica	1937	Painting	25000000.00	I123	TLM123	AD55
AW78	The Night Watch	1642	Painting	9000000.00	J123	BM123	AD88
AW90	Two Sisters	2010	Sculpture	200000.00	K123	JG123	AD00
							8

• Custom SQL Search Page



• Excecution of Boolean Query



custid	fname	lname	address
AR1998	Ashutosh	Ranjan	Paris
AT2000	Akshay	Thakur	New York