

## UNIT -I

### INTRODUCTION TO DBMS:

#### What is data?

- Data is nothing but facts and statistics stored or free flowing over a network, generally it's raw and unprocessed.
- Data becomes information when it is processed, turning it into something meaningful.
- What is database: The database is a collection of inter-related data which is used to retrieve, insert and delete the data efficiently.
- It is also used to organize the data in the form of a table, schema, views, and reports, etc.
- Using the database, you can easily retrieve, insert, and delete the information.
- For example: The college Database organizes the data about the admin, staff, students and faculty etc.

#### What is dbms?

DBMS	File System
DBMS is a collection of data. In DBMS, the user is not required to write the procedures.	File system is a collection of data. In this system, the user has to write the procedures for managing the database.
Searching data is easy in Dbms	Searching is difficult in File System
Dbms is structured data	Files are unstructured data
No data redundancy in Dbms	Data redundancy is there in file system
Memory utilisation well in dbms	Memory utilisation poor in file system
No data inconsistency in dbms	Inconsistency in file system

DBMS gives an abstract view of data that hides the details.	File system provides the detail of the data representation and storage of data.
DBMS provides a crash recovery mechanism, i.e., DBMS protects the user from the system failure.	File system doesn't have a crash mechanism, i.e., if the system crashes while entering some data, then the content of the file will be lost.
DBMS provides a good protection mechanism.	It is very difficult to protect a file under the file system.
DBMS contains a wide variety of sophisticated techniques to store and retrieve the data.	File system can't efficiently store and retrieve the data.
DBMS takes care of Concurrent access of data using some form of locking.	In the File system, concurrent access has many problems like redirecting the file while deleting some information or updating some information.

- A DBMS is software that allows creation, definition and manipulation of database, allowing users to store, process and analyse data easily.
- DBMS provides us with an interface or a tool, to perform various operations like creating database, storing data in it, updating data, creating tables in the database and a lot more.
- DBMS also provides protection and security to the databases.
- It also maintains data consistency in case of multiple users.

Here are some examples of popular DBMS used these days:

- MySQL
- Oracle
- SQL Server
- IBM DB2

#### **DATABASE APPLICATIONS – DBMS:**

- Applications where we use Database Management Systems are:
- Telecom: There is a database to keep track of the information regarding calls made, network usage, customer details etc.

- Industry: Where it is a manufacturing unit, warehouse or distribution centre, each one needs a database to keep the records of ins and outs
- Banking System: For storing customer info, tracking day to day credit and debit transactions, generating bank statements etc.
- Sales: To store customer information, production information and invoice details.
- Airlines: To travel through airlines, we make early reservations; this reservation information along with flight schedule is stored in database.
- Education sector: Database systems are frequently used in schools and colleges to store and retrieve the data regarding student details, staff details, course details, exam details, payroll data, attendance details, fees details etc.

## **PURPOSE OF DATABASE SYSTEMS**

- The main purpose of database systems is to manage the data. Consider a university that keeps the data of students, teachers, courses, books etc. To manage this data we need to store this data somewhere where we can add new data, delete unused data, update outdated data, retrieve data, to perform these operations on data we need a Database management system that allows us to store the data in such a way so that all these operations can be performed on the data efficiently.

## **Characteristics of DBMS**

- Data stored into Tables: Data is never directly stored into the database. Data is stored into tables, created inside the database.
- Reduced Redundancy: In the modern world hard drives are very cheap, but earlier when hard drives were too expensive, unnecessary repetition of data in database was a big problem. But DBMS follows Normalisation which divides the data in such a way that repetition is minimum.
- Data Consistency: On Live data, i.e. data that is being continuously updated and added, maintaining the consistency of data can become a challenge. But DBMS handles it all by itself.
- Support Multiple user and Concurrent Access: DBMS allows multiple users to work on it(update, insert, delete data) at the same time and still manages to maintain the data consistency.

- Query Language: DBMS provides users with a simple Query language, using which data can be easily fetched, inserted, deleted and updated in a database.

### **Advantages of DBMS**

- Controls database redundancy: It can control data redundancy because it stores all the data in one single database file and that recorded data is placed in the database.
- Data sharing: In DBMS, the authorized users of an organization can share the data among multiple users.
- Easily Maintenance: It can be easily maintainable due to the centralized nature of the database system.
- Reduce time: It reduces development time and maintenance need.
- Backup: It provides backup and recovery subsystems which create automatic backup of data from hardware and software failures and restores the data if required.
- multiple user interface: It provides different types of user interfaces like graphical user interfaces, application program interfaces

### **Disadvantages of DBMS**

- Cost of Hardware and Software: It requires a high speed of data processor and large memory size to run DBMS software.
- Size: It occupies a large space of disks and large memory to run them efficiently.
- Complexity: Database system creates additional complexity and requirements.
- Higher impact of failure: Failure is highly impacted the database because in most of the organization, all the data stored in a single database and if the database is damaged due to electric failure or database corruption then the data may be lost forever.

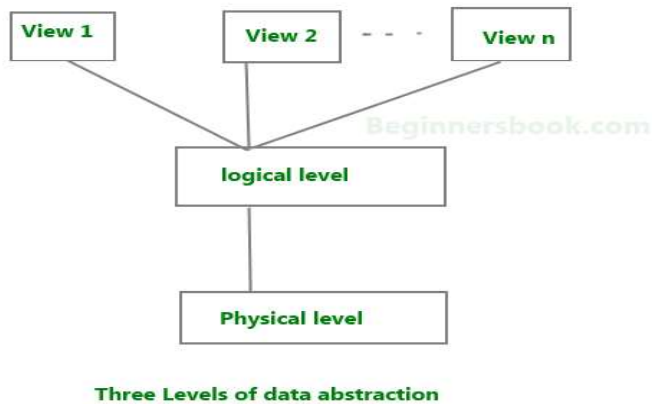
### **View of Data in DBMS**

- Abstraction is one of the main features of database systems.
- Hiding irrelevant details from user and providing abstract view of data to users, helps in easy and efficient user-database interaction.
- the three level of DBMS architecture, The top level of that architecture is “view level”. The view level provides the “view of data” to the users and hides the irrelevant

details such as data relationship, database schema, constraints, security etc from the user.

### Data Abstraction in DBMS

Database systems are made-up of complex data structures. To ease the user interaction with database, the developers hide internal irrelevant details from users. This process of hiding irrelevant details from user is called data abstraction.



We have three levels of abstraction:

**Physical level:** This is the lowest level of data abstraction. It describes how data is actually stored in database. You can get the complex data structure details at this level.

**Logical level:** This is the middle level of 3-level data abstraction architecture. It describes what data is stored in database.

**View level:** Highest level of data abstraction. This level describes the user interaction with database system.

### Instance and schema in DBMS

- Definition of schema: Design of a database is called the schema. Schema is of three types: Physical schema, logical schema and view schema.

- The design of a database at physical level is called physical schema, how the data stored in blocks of storage is described at this level.
- Design of database at logical level is called logical schema, programmers and database administrators work at this level, at this level data can be described as certain types of data records gets stored in data structures, however the internal details such as implementation of data structure is hidden at this level (available at physical level).
- Design of database at view level is called view schema. This generally describes end user interaction with database systems.

#### **Definition of instance:**

The data stored in database at a particular moment of time is called instance of database. Database schema defines the variable declarations in tables that belong to a particular database; the value of these variables at a moment of time is called the instance of that database.

#### **DBMS ARCHITECTURE:**

- Database management systems architecture will help us understand the components of database system and the relation among them.
- The architecture of DBMS depends on the computer system on which it runs.
- the basic client/server architecture is used to deal with a large number of PCs, web servers, database servers and other components that are connected with networks.
- The client/server architecture consists of many PCs and a workstation which are connected via the network.
- DBMS architecture depends upon how users are connected to the database to get their request done.

#### **TYPES OF DBMS ARCHITECTURE**

There are three types of DBMS architecture:

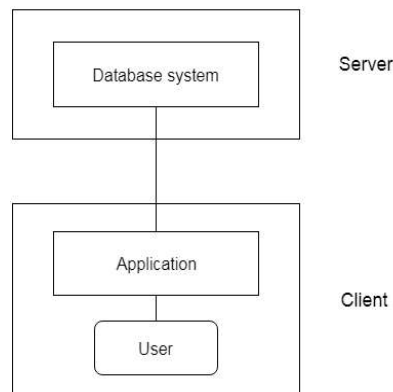
1. Single tier architecture
2. Two tier architecture
3. Three tier architecture

### 1-Tier Architecture

- In this type of architecture, the database is readily available on the client machine, any request made by client doesn't require a network connection to perform the action on the database.
- Any changes done here will directly be done on the database itself. It doesn't provide a handy tool for end users.
- The 1-Tier architecture is used for development of the local application, where programmers can directly communicate with the database for the quick response.

### 2. Two tier architecture

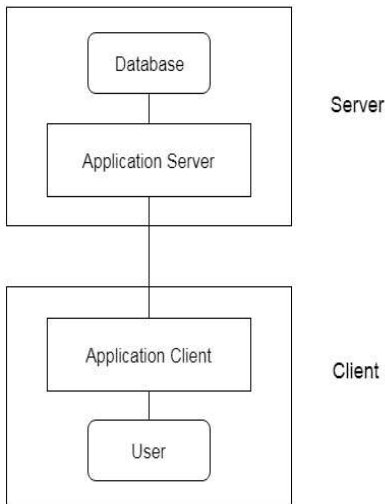
- In two-tier architecture, the Database system is present at the server machine and the DBMS application is present at the client machine, these two machines are connected with each other through a reliable network.
- Whenever client machine makes a request to access the database present at server using a query language like sql, the server perform the request on the database and returns the result back to the client.
- The application connection interface such as JDBC, ODBC are used for the interaction between server and client.



### 3-Tier Architecture

- In three-tier architecture, another layer is present between the client machine and server machine.
- In this architecture, the client application doesn't communicate directly with the database systems present at the server machine, rather the client application

communicates with server application and the server application internally communicates with the database system present at the server.



### **DATA MODELS:**

- Data Model is the modeling of the data description, data semantics, and consistency constraints of the data.
- It provides the conceptual tools for describing the design of a database at each level of data abstraction.
- Therefore, there are following four data models used for understanding the structure of the database:

### **Four Types of DBMS systems are:**

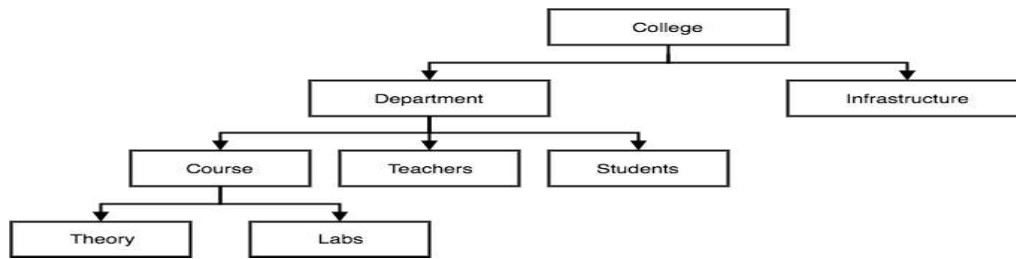
- Hierarchical database
- Network database
- Relational database
- ER model database

### **Hierarchical DBMS**

In a Hierarchical database, model data is organized in a tree-like structure. Data is Stored Hierarchically (top down or bottom up) format. Data is represented using a parent-child

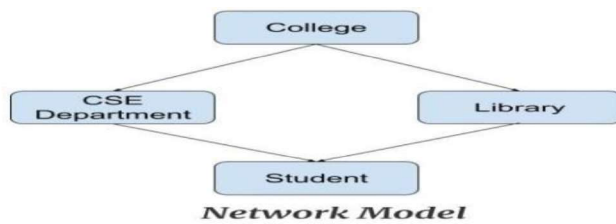


relationship. In Hierarchical DBMS parent may have many children, but children have only one parent.



### Network Model

The network database model allows each child to have multiple parents. It helps you to address the need to model more complex relationships like as the orders/parts many-to-many relationship. In this model, entities are organized in a graph which can be accessed through several paths.



### Relational model

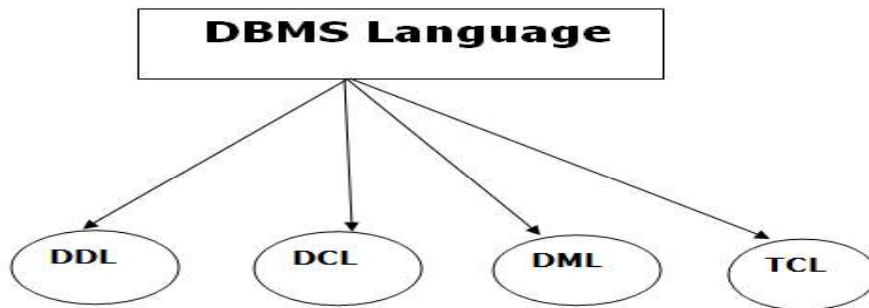
Relational DBMS is the most widely used DBMS model because it is one of the easiest. This model is based on normalizing data in the rows and columns of the tables. Relational model stored in fixed structures and manipulated using SQL.

### Entity-Relationship Model

Entity-Relationship (ER) Model is based on the notion of real-world entities and relationships among them. While formulating real-world scenario into the database model, the ER Model creates entity set, relationship set, general attributes and constraints.

## DBMS languages

Database languages are used to read, update and store data in a database. There are several such languages that can be used for this purpose; one of them is SQL (Structured Query Language).



- DDL – Data Definition Language:  
(CREATE,DROP,ALTER,TRUNCATE,COMMENT,RENAME)
- DML – Data Manipulation Language: (INSERT, UPDATE,DELETE)
- DCL – Data Control Language: (GRANT,REVOKE)
- TCL-Transaction Control Language: (COMMIT,ROLLBACK)

**1. DDL(Data Definition Language)** : DDL or Data Definition Language actually consists of the SQL commands that can be used to define the database schema. It simply deals with descriptions of the database schema and is used to create and modify the structure of database objects in the database.

**CREATE** – it is used to create the database or its objects (like table, index, function, views, store procedure and triggers).

There are two CREATE statements available in SQL:

- CREATE DATABASE
- CREATE TABLE

### **CREATE DATABASE**

A Database is defined as a structured set of data. So, in SQL the very first step to store the data in a well structured manner is to create a database. The CREATE DATABASE statement is used to create a new database in SQL.

Syntax:

**CREATE DATABASE database\_name;**

Example:

```
SQL> CREATE DATABASE Employee;
```

In order to get the list of all the databases, you can use SHOW DATABASES statement.

Example –

```
SQL> SHOW DATABASES;
```

### **CREATE TABLE:**

The CREATE TABLE statement is used to create a table in SQL. We know that a table comprises of rows and columns. So while creating tables we have to provide all the information to SQL about the names of the columns, type of data to be stored in columns, size of the data etc. Let us now dive into details on how to use CREATE TABLE statement to create tables in SQL.

Syntax:

```
CREATE TABLE table_name  
(  
column1 data_type(size),  
column2 data_type(size),  
column3 data_type(size),  
....  
);
```

### **Example Query:**

This query will create a table named Students with three columns, ROLL\_NO, NAME and SUBJECT.

```
CREATE TABLE Students  
(  
ROLL_NO int(3),  
NAME varchar(20),  
SUBJECT varchar(20),  
);
```

### **DROP:**

DROP is used to delete a whole database or just a table. The DROP statement destroys the objects like an existing database, table, index, or view.

A DROP statement in SQL removes a component from a relational database management system (RDBMS).

Syntax:

DROP object object\_name;

Examples:

DROP TABLE table\_name;

table\_name: Name of the table to be deleted.

DROP DATABASE database\_name;

database\_name: Name of the database to be deleted.

### **TRUNCATE**

It is used to remove all records from a table, including all spaces

The TRUNCATE TABLE mytable statement is logically (though not physically) equivalent to the DELETE FROM mytable statement (without a WHERE clause).

Syntax:

TRUNCATE TABLE table\_name;

### **DROP vs TRUNCATE**

- Truncate is normally ultra-fast and its ideal for deleting data from a temporary table.
- Truncate preserves the structure of the table for future use, unlike drop table where the table is deleted with its full structure.
- Table or Database deletion using DROP statement cannot be rolled back, so it must be used wisely.

To delete the whole database

DROP DATABASE student\_data;

After running the above query whole database will be deleted.

To truncate Student\_details table from student\_data database.

TRUNCATE TABLE Student\_details;

### **ALTER (ADD, DROP, MODIFY)**

ALTER TABLE is used to add, delete/drop or modify columns in the existing table. It is also used to add and drop various constraints on the existing table.

ALTER TABLE – ADD:

ADD is used to add columns into the existing table. Sometimes we may require to add additional information, in that case we do not require to create the whole database again, ADD comes to our rescue.

Syntax:

```
ALTER TABLE table_name
    ADD (Columnname_1 datatype,
        Columnname_2 datatype,
        ...
        Columnname_n datatype);
```

### **ALTER TABLE – DROP**

DROP COLUMN is used to drop column in a table. Deleting the unwanted columns from the table.

Syntax:

```
ALTER TABLE table_name
DROP COLUMN column_name;
```

### **ALTER TABLE-MODIFY**

It is used to modify the existing columns in a table. Multiple columns can also be modified at once.

```
ALTER TABLE table_name
MODIFY column_name column_type;
```

QUERY:

**To ADD 2 columns AGE and COURSE to table Student.**

```
ALTER TABLE Student ADD (AGE number(3),COURSE varchar(40));
```

**MODIFY column COURSE in table Student**

```
ALTER TABLE Student MODIFY COURSE varchar(20);
```

### **Comments**

As is any programming languages comments matter a lot in SQL also. In this set we will learn about writing comments in any SQL snippet.

Comments can be written in the following three formats:

- Single line comments.
- Multi line comments
- In line comments

**DML(Data Manipulation Language)** : The SQL commands that deals with the manipulation of data present in the database belong to DML or Data Manipulation Language and this includes most of the SQL statements.

### **SELECT Statement**

select statement is used to fetch data from relational database. A relational database is organized collection of data. As we know that data is stored inside tables in a database. SQL select statement or SQL select query is used to fetch data from one or more than one tables.

#### **SELECT Syntax**

##### **One column:**

Here column\_name is the name of the column for which we need to fetch data and table\_name is the name of the table in the database.

SELECT column\_name FROM table\_name;

##### **More than one columns:**

SELECT column\_name\_1, column\_name\_2, ... FROM table\_name;

##### **To fetch the entire table or all the fields in the table:**

SELECT \* FROM table\_name;

Example:

SELECT EMP\_NAME FROM EMPLOYEES;

To fetch the entire EMPLOYEES table:

SELECT \* FROM EMPLOYEES;

Query to fetch the fields ROLL\_NO, NAME, AGE from the table Student:

SELECT ROLL\_NO, NAME, AGE FROM Student;

### **INSERT INTO Statement**

The INSERT INTO statement of SQL is used to insert a new row in a table. There are two ways of using INSERT INTO statement for inserting rows:

Only values: First method is to specify only the value of data to be inserted without the column names.

INSERT INTO table\_name VALUES (value1, value2, value3,...);

Column names and values both: In the second method we will specify both the columns which we want to fill and their corresponding values as shown below:

```
INSERT INTO table_name (column1, column2, column3,...) VALUES ( value1, value2, value3,...);
```

**Example:**

Method 1 (Inserting only values) :

```
INSERT INTO Student VALUES ('5','HARSH','WEST BENGAL','XXXXXXXXXX','19');
```

Method 2 (Inserting values in only specified columns):

```
INSERT INTO Student (ROLL_NO, NAME, Age) VALUES ('5','PRATIK','19');
```

**UPDATE Statement**

The UPDATE statement in SQL is used to update the data of an existing table in database.

We can update single columns as well as multiple columns using UPDATE statement as per our requirement.

Basic Syntax:

```
UPDATE TableName
```

```
SET column_name1 = value, column_name2 = value....
```

```
WHERE condition;
```

EX1:

```
SQL> UPDATE EMPLOYEES
```

```
SET EMP_SALARY = 10000
```

```
WHERE EMP_AGE > 25;
```

EX2;

```
SQL> UPDATE EMPLOYEES
```

```
SET EMP_SALARY = 120000
```

```
WHERE EMP_NAME = 'Apoorv';
```

**DELETE Statement**

The DELETE Statement in SQL is used to delete existing records from a table. We can delete a single record or multiple records depending on the condition we specify in the WHERE clause.

Basic Syntax:

```
DELETE FROM table_name WHERE some_condition;
```

Deleting single record: Delete the rows where NAME = 'Ram'. This will delete only the first row.

```
DELETE FROM Student WHERE NAME = 'Ram';
```

Deleting multiple records: Delete the rows from the table Student where Age is 20. This will delete 2 rows(third row and fifth row).

DELETE FROM Student WHERE Age = 20;

Delete all of the records: There are two queries to do this as shown below,

query1: "DELETE FROM Student";

query2: "DELETE \* FROM Student";

**DCL(Data Control Language)** : DCL includes commands such as GRANT and REVOKE which mainly deals with the rights, permissions and other controls of the database system.

Examples of DCL commands:

**GRANT-gives user's access privileges to database.**

**REVOKE-withdraw user's access privileges given by using the GRANT command.**

**TCL(transaction Control Language)** : TCL commands deals with the transaction within the database.

Examples of TCL commands:

**COMMIT- commits a Transaction.**

**ROLLBACK- rollbacks a transaction in case of any error occurs.**

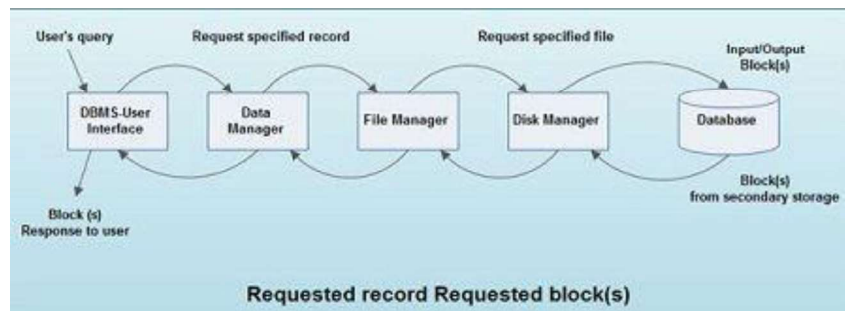
**SAVEPOINT-sets a savepoint within a transaction.**

**SET TRANSACTION-specify characteristics for the transaction.**

### **What is the Procedure for Database Access?**

- Any access to the stored data is done by the data manager. A user's request for data is-received by the data manager, which determines the physical record required. The decision as to which physical record is needed may require some preliminary consultation of the database and/or the data dictionary prior to the access of the actual data itself.
- The data manager sends the request for a specific physical record to the file manager. The file manager decides which physical block of secondary storage devices contains the required record and sends the request for the appropriate block to the disk manager. A block is a unit of physical input/output operations between primary and secondary storage. The disk manager retrieves the block and sends it to the file manager, which sends the required record to the data manager.





## DATA BASE USERS AND ADMINISTRATORS:

Database users are the persons who interact with the database and take the benefits of database.

They are differentiated into different types based on the way they expect to interact with the system.

**Naive users:** They are the unsophisticated users who interact with the system by using permanent applications that already exist. Example: Online Library Management System, ATMs (Automated Teller Machine), etc.

**Application programmers:** They are the computer professionals who interact with system through DML. They write application programs.

**Sophisticated users:** They interact with the system by writing SQL queries directly through the query processor without writing application programs.

**Specialized users:** They are also sophisticated users who write specialized database applications that do not fit into the traditional data processing framework. Example: Expert System, Knowledge Based System, etc.

## Database Administrators

The life cycle of database starts from designing, implementing to administration of it. A database for any kind of requirement needs to be designed perfectly so that it should work without any issues. Once all the design is complete, it needs to be installed. Once this step is complete, users start using the database. The database grows as the data grows in the database. When the database becomes huge, its performance comes down. Also accessing the data from the database becomes challenge. There will be unused memory in database, making the memory inevitably huge. These administration and maintenance of database is taken care

by database Administrator – DBA.

A DBA has many responsibilities. A good performing database is in the hands of DBA.

Database Administrators coordinate all the activities of the database system. They have all the permissions.

### **Tasks of DBA**

- Creating the schema
- Specifying integrity constraints
- Storage structure and access method definition
- Granting permission to other users.
- Monitoring performance
- Routine Maintenance

### **Transaction Management?**

- A Database Transaction is a logical unit of processing in a DBMS which entails one or more database access operation. In a nutshell, database transactions represent real-world events of any enterprise.
- All types of database access operation which are held between the beginning and end transaction statements are considered as a single logical transaction in DBMS. During the transaction the database is inconsistent. Only once the database is committed the state is changed from one consistent state to another.

### **What are ACID Properties?**

ACID Properties are used for maintaining the integrity of database during transaction processing. ACID in DBMS stands for Atomicity, Consistency, Isolation, and Durability.

- **Atomicity:** A transaction is a single unit of operation. You either execute it entirely or do not execute it at all. There cannot be partial execution.
- **Consistency:** Once the transaction is executed, it should move from one consistent state to another.

- **Isolation:** Transaction should be executed in isolation from other transactions (no Locks). During concurrent transaction execution, intermediate transaction results from simultaneously executed transactions should not be made available to each other. (Level 0,1,2,3)
- **Durability:** After successful completion of a transaction, the changes in the database should persist. Even in the case of system failures.

### Storage Manager In DBMS

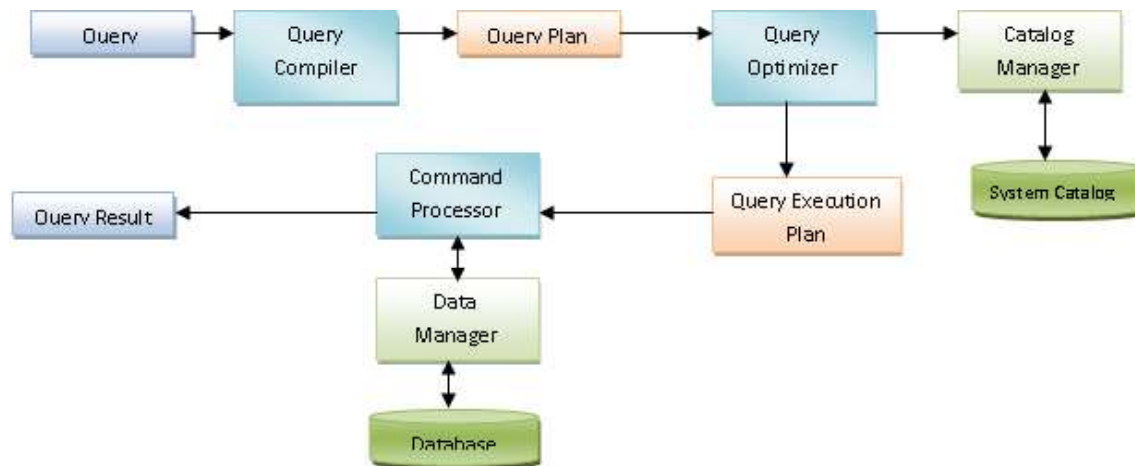
- A storage manager is a program module that provides the interface between the lowlevel data stored in the database and the application programs and queries submitted to the system.
- The storage manager is responsible for the interaction with the file manager.
- The raw data are stored on the disk using the file system, which is usually provided by a conventional operating system.
- The storage manager translates the various DML statements into low-level file-system commands. Thus, the storage manager is responsible for storing, retrieving, and updating data in the database.

The storage manager components include:

1. Authorization and integrity manager, which tests for the satisfaction of integrity constraints and checks the authority of users to access data.
2. Transaction manager, which ensures that the database remains in a consistent (correct) state despite system failures, and that concurrent transaction executions proceed without conflicting.
3. File manager, which manages the allocation of space on disk storage and the data structures used to represent information stored on disk.
4. Buffer manager, which is responsible for fetching data from disk storage into main memory, and deciding what data to cache in main memory. The buffer manager is a critical part of the database system, since it enables the database to handle data sizes that are much larger than the size of main memory. The storage manager implements several data structures as part of the physical system implementation:

## Query Processing in DBMS

A query processor is one of the major components of a relational database or an electronic database in which data is stored in tables of rows and columns. It complements the storage engine, which writes and reads data to and from storage media.



## Parsing and Translation

As query processing includes certain activities for data retrieval. Initially, the given user queries get translated in high-level database languages such as SQL. It gets translated into expressions that can be further used at the physical level of the file system. After this, the actual evaluation of the queries and a variety of query -optimizing transformations and takes place.

## Query Evaluation Plan

- In order to fully evaluate a query, the system needs to construct a query evaluation plan.
- The annotations in the evaluation plan may refer to the algorithms to be used for the particular index or the specific operations.
- Such relational algebra with annotations is referred to as **Evaluation Primitives**. The evaluation primitives carry the instructions needed for the evaluation of the operation.

- Thus, a query evaluation plan defines a sequence of primitive operations used for evaluating a query. The query evaluation plan is also referred to as **the query execution plan**.
- A **query execution engine** is responsible for generating the output of the given query. It takes the query execution plan, executes it, and finally makes the output for the user query.

### **Optimization**

- The cost of the query evaluation can vary for different types of queries. Although the system is responsible for constructing the evaluation plan, the user does need not to write their query efficiently.
- Usually, a database system generates an efficient query evaluation plan, which minimizes its cost. This type of task performed by the database system and is known as Query Optimization.
- For optimizing a query, the query optimizer should have an estimated cost analysis of each operation. It is because the overall operation cost depends on the memory allocations to several operations, execution costs, and so on.

### **What is Relational Model?**

**Relational Model (RM)** represents the database as a collection of relations. A relation is nothing but a table of values. Every row in the table represents a collection of related data values. These rows in the table denote a real-world entity or relationship.

The table name and column names are helpful to interpret the meaning of values in each row. The data are represented as a set of relations. In the relational model, data are stored as tables. However, the physical storage of the data is independent of the way the data are logically organized.

### **Relational Model Concepts**

1. **Attribute:** Each column in a Table. Attributes are the properties which define a relation. e.g., Student\_Rollno, NAME, etc.

2. **Tables** – In the Relational model the, relations are saved in the table format. It is stored along with its entities. A table has two properties rows and columns. Rows represent records and columns represent attributes.
3. **Tuple** – It is nothing but a single row of a table, which contains a single record.
4. **Relation Schema:** A relation schema represents the name of the relation with its attributes.
5. **Degree:** The total number of attributes which in the relation is called the degree of the relation.
6. **Cardinality:** Total number of rows present in the Table.
7. **Column:** The column represents the set of values for a specific attribute.
8. **Relation instance** – Relation instance is a finite set of tuples in the RDBMS system. Relation instances never have duplicate tuples.
9. **Relation key** - Every row has one, two or multiple attributes, which is called relation key.
10. **Attribute domain** – Every attribute has some pre-defined value and scope which is known as attribute domain

### **Keys in DBMS**

**KEYS in DBMS** is an attribute or set of attributes which helps you to identify a row(tuple) in a relation(table). They allow you to find the relation between two tables. Keys help you uniquely identify a row in a table by a combination of one or more columns in that table. Key is also helpful for finding unique record or row from the table. Database key is also helpful for finding unique record or row from the table.

### **Why we need a Key?**

Here are some reasons for using sql key in the DBMS system.

- Keys help you to identify any row of data in a table. In a real-world application, a table could contain thousands of records. Moreover, the records could be duplicated. Keys ensure that you can uniquely identify a table record despite these challenges.
- Allows you to establish a relationship between and identify the relation between tables
- Help you to enforce identity and integrity in the relationship.

## Types of Keys in Database Management System

There are mainly seven different types of Keys in DBMS and each key has its different functionality:

- **Super Key** - A super key is a group of single or multiple keys which identifies rows in a table.
- **Primary Key** - is a column or group of columns in a table that uniquely identify every row in that table.
- **Candidate Key** - is a set of attributes that uniquely identify tuples in a table. Candidate Key is a super key with no repeated attributes.
- **Alternate Key** - is a column or group of columns in a table that uniquely identify every row in that table.
- **Foreign Key** - is a column that creates a relationship between two tables. The purpose of Foreign keys is to maintain data integrity and allow navigation between two different instances of an entity.
- **Compound Key** - has two or more attributes that allow you to uniquely recognize a specific record. It is possible that each column may not be unique by itself within the database.
- **Composite Key** - An artificial key which aims to uniquely identify each record is called a surrogate key. These kind of key are unique because they are created when you don't have any natural primary key.
- **Surrogate Key** - An artificial key which aims to uniquely identify each record is called a surrogate key. These kind of key are unique because they are created when you don't have any natural primary key.

Primary key example:

```
CREATE TABLE Persons (  
    ID int NOT NULL,  
    LastName varchar(255) NOT NULL,  
    FirstName varchar(255),  
    Age int,  
    PRIMARY KEY (ID)  
);
```

## Create Primary Key (ALTER TABLE statement)

### Syntax

The syntax to create a primary key using the ALTER TABLE statement in SQL is:

```
ALTER TABLE table_name  
  ADD CONSTRAINT constraint_name  
  PRIMARY KEY (column1, column2, ... column_n);
```

### FOREIGN KEY on CREATE TABLE

The following SQL creates a FOREIGN KEY on the "PersonID" column when the "Orders" table is created:

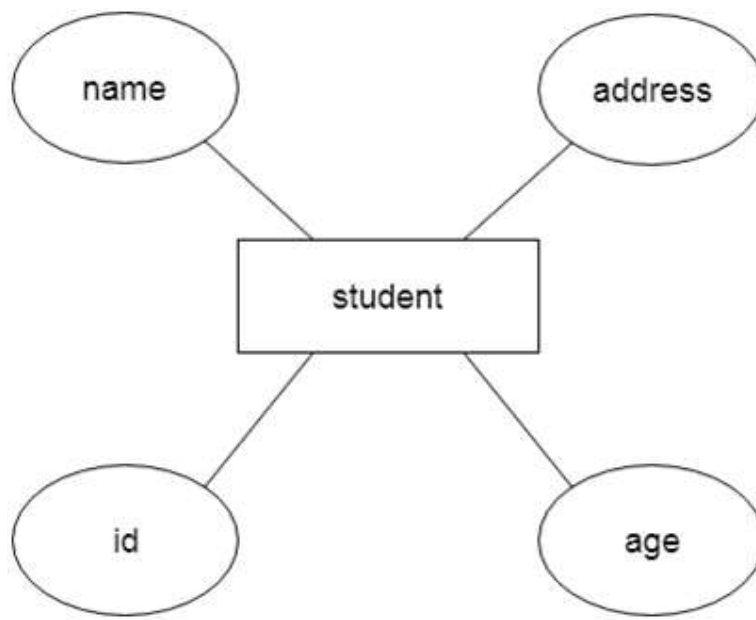
```
CREATE TABLE Orders (  
  OrderID int NOT NULL,  
  OrderNumber int NOT NULL,  
  PersonID int,  
  PRIMARY KEY (OrderID),  
  FOREIGN KEY (PersonID) REFERENCES Persons(PersonID)  
);
```

### ER model

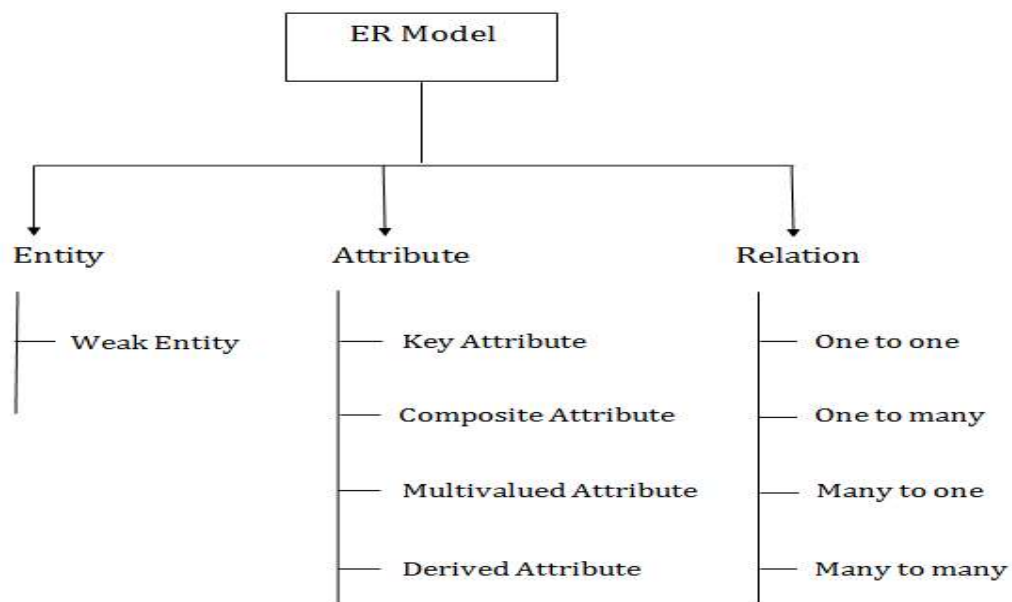
- ER model stands for an Entity-Relationship model. It is a high-level data model. This model is used to define the data elements and relationship for a specified system.
- It develops a conceptual design for the database. It also develops a very simple and easy to design view of data.
- In ER modeling, the database structure is portrayed as a diagram called an entity-relationship diagram.

**For example,** Suppose we design a school database. In this database, the student will be an entity with attributes like address, name, id, age, etc. The address can be another entity with attributes like city, street name, pin code, etc and there will be a relationship between them.





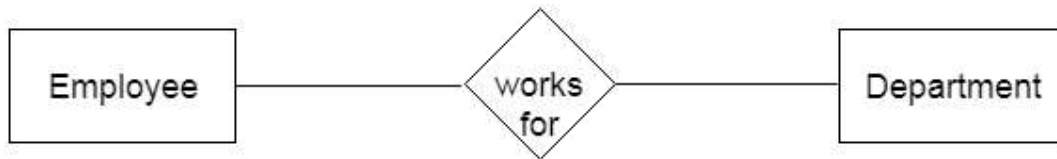
Component of ER Diagram



### 1. Entity:

An entity may be any object, class, person or place. In the ER diagram, an entity can be represented as rectangles.

Consider an organization as an example- manager, product, employee, department etc. can be taken as an entity.



#### a. Weak Entity

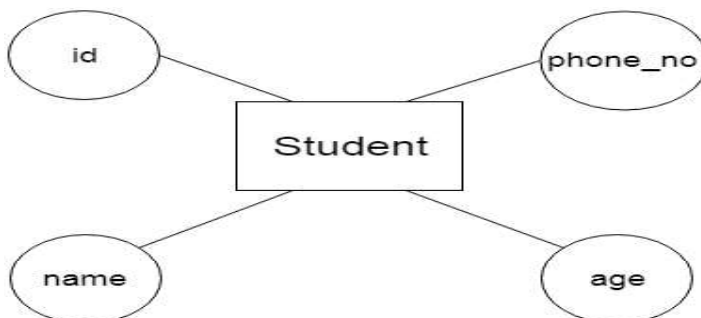
An entity that depends on another entity called a weak entity. The weak entity doesn't contain any key attribute of its own. The weak entity is represented by a double rectangle.



## 2. Attribute

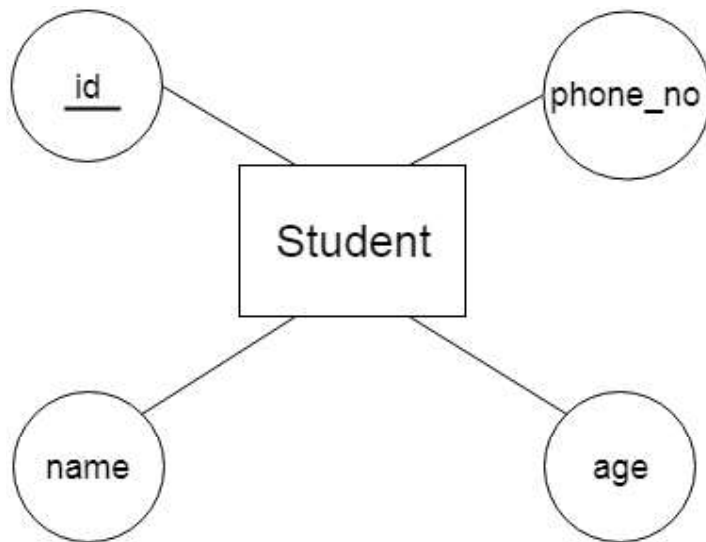
The attribute is used to describe the property of an entity. Eclipse is used to represent an attribute.

**For example,** id, age, contact number, name, etc. can be attributes of a student.



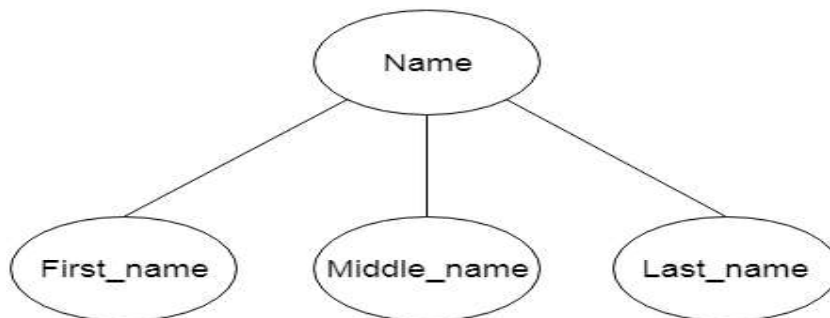
#### a. Key Attribute

The key attribute is used to represent the main characteristics of an entity. It represents a primary key. The key attribute is represented by an ellipse with the text underlined.



### b. Composite Attribute

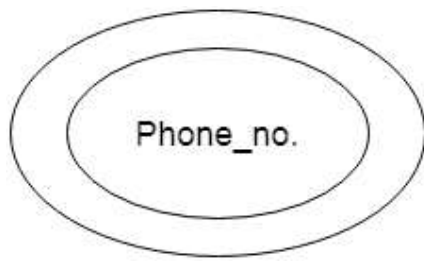
An attribute that composed of many other attributes is known as a composite attribute. The composite attribute is represented by an ellipse, and those ellipses are connected with an ellipse.



### c. Multivalued Attribute

An attribute can have more than one value. These attributes are known as a multivalued attribute. The double oval is used to represent multivalued attribute.

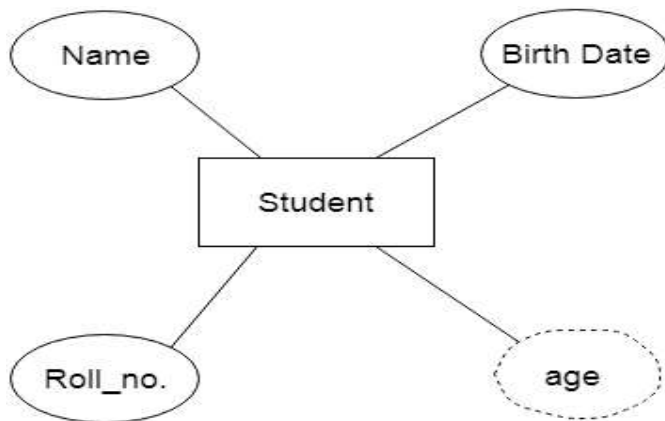
**For example,** a student can have more than one phone number.



#### d. Derived Attribute

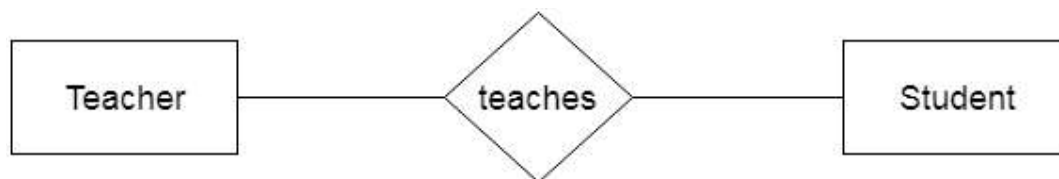
An attribute that can be derived from other attribute is known as a derived attribute. It can be represented by a dashed ellipse.

**For example,** A person's age changes over time and can be derived from another attribute like Date of birth.



#### 3. Relationship

A relationship is used to describe the relation between entities. Diamond or rhombus is used to represent the relationship



Types of relationship are as follows:

### a. One-to-One Relationship

When only one instance of an entity is associated with the relationship, then it is known as one to one relationship.

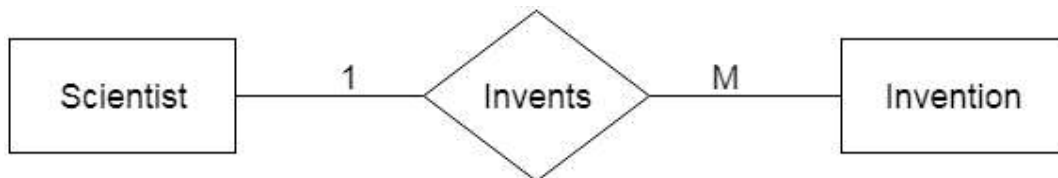
**For example,** A female can marry to one male, and a male can marry to one female.



### b. One-to-many relationship

When only one instance of the entity on the left, and more than one instance of an entity on the right associates with the relationship then this is known as a one-to-many relationship.

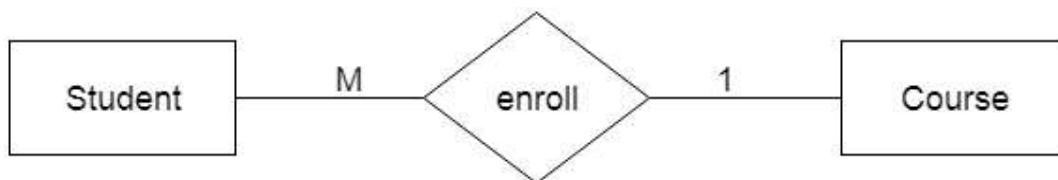
**For example,** Scientist can invent many inventions, but the invention is done by the only specific scientist.



### c. Many-to-one relationship

When more than one instance of the entity on the left, and only one instance of an entity on the right associates with the relationship then it is known as a many-to-one relationship.

**For example,** Student enrolls for only one course, but a course can have many students.



#### d. Many-to-many relationship

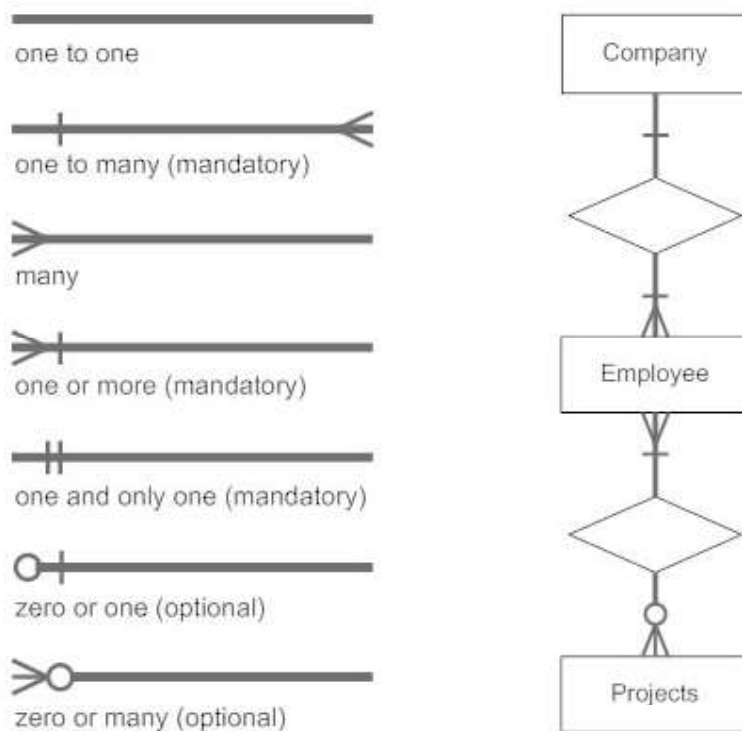
When more than one instance of the entity on the left, and more than one instance of an entity on the right associates with the relationship then it is known as a many-to-many relationship.

**For example,** Employee can assign by many projects and project can have many employees.



Notation of ER diagram

Database can be represented using the notations. In ER diagram, many notations are used to express the cardinality. These notations are as follows:

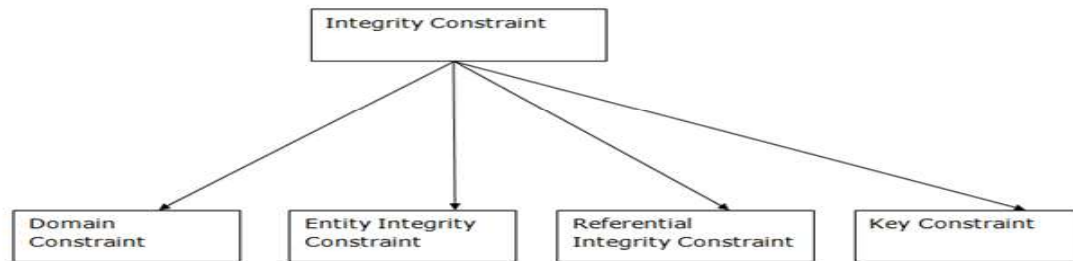


#### Integrity Constraints

- Integrity constraints are a set of rules. It is used to maintain the quality of information.
- Integrity constraints ensure that the data insertion, updating, and other processes have to be performed in such a way that data integrity is not affected.

- Thus, integrity constraint is used to guard against accidental damage to the database.

### Types of Integrity Constraint



#### 1. Domain constraints

- Domain constraints can be defined as the definition of a valid set of values for an attribute.
- The data type of domain includes string, character, integer, time, date, currency, etc. The value of the attribute must be available in the corresponding domain.

#### Example:

ID	NAME	SEMENSTER	AGE
1000	Tom	1 <sup>st</sup>	17
1001	Johnson	2 <sup>nd</sup>	24
1002	Leonardo	5 <sup>th</sup>	21
1003	Kate	3 <sup>rd</sup>	19
1004	Morgan	8 <sup>th</sup>	A

Not allowed. Because AGE is an integer attribute

#### 2. Entity integrity constraints

- The entity integrity constraint states that primary key value can't be null.
- This is because the primary key value is used to identify individual rows in relation and if the primary key has a null value, then we can't identify those rows.
- A table can contain a null value other than the primary key field.

## EMPLOYEE

EMP_ID	EMP_NAME	SALARY
123	Jack	30000
142	Harry	60000
164	John	20000
	Jackson	27000

Not allowed as primary key can't contain a NULL value

### 3. Referential

## Integrity Constraints

- A referential integrity constraint is specified between two tables.
- In the Referential integrity constraints, if a foreign key in Table 1 refers to the Primary Key of Table 2, then every value of the Foreign Key in Table 1 must be null or be available in Table 2.

(Table 1)

EMP_NAME	NAME	AGE	D_No
1	Jack	20	11
2	Harry	40	24
3	John	27	18
4	Devil	38	13

Foreign key

Not allowed as D\_No 18 is not defined as a Primary key of table 2 and In table 1, D\_No is a foreign key defined

Relationships

(Table 2)

Primary Key

<u>D_No</u>	D_Location
11	Mumbai
24	Delhi
13	Noida

## 4. Key constraints

- Keys are the entity set that is used to identify an entity within its entity set uniquely.
- An entity set can have multiple keys, but out of which one key will be the primary key. A primary key can contain a unique and null value in the relational table.



ID	NAME	SEMENSTER	AGE
1000	Tom	1 <sup>st</sup>	17
1001	Johnson	2 <sup>nd</sup>	24
1002	Leonardo	5 <sup>th</sup>	21
1003	Kate	3 <sup>rd</sup>	19
1002	Morgan	8 <sup>th</sup>	22

Not allowed. Because all row must be unique

## ER Design Issues

- ER design issues need to be discussed for better ER- design
- users often mislead the concept of the elements and the design process of the ER diagram. Thus, it leads to a complex structure of the ER diagram and certain issues that does not meet the characteristics of the real-world enterprise model.
- Here, we will discuss the basic design issues of an ER database schema in the following points:

### 1) Use of Entity Set vs Attributes

The use of an entity set or attribute depends on the structure of the real-world enterprise that is being modelled and the semantics associated with its attributes. It leads to a mistake when the user use the primary key of an entity set as an attribute of another entity set. Instead, he should use the relationship to do so. Also, the primary key attributes are implicit in the relationship set, but we designate it in the relationship sets.

### 2) Use of Entity Set vs. Relationship Sets

It is difficult to examine if an object can be best expressed by an entity set or relationship set.

### 3) Use of Binary vs n-ary Relationship Sets

Generally, the relationships described in the databases are binary relationships. However, non-binary relationships can be represented by several binary relationships.

### 4) Placing Relationship Attributes

The cardinality ratios can become an affective measure in the placement of the relationship attributes. So, it is better to associate the attributes of one-to-one or one-to-many relationship sets with any participating entity sets, instead of any relationship set.

## **Conceptual design**

Conceptual design is the first stage in the database design process. The goal at this stage is to design a database that is independent of database software and physical details. The output of this process is a conceptual data model that describes the main data entities, attributes, relationships, and constraints of a given problem domain.

Keep in mind the following minimal data rule:

"All that is needed is there, and all that is there is needed".

1. Data analysis and requirements
2. Entity relationship modeling and normalization
3. Data model verification
4. Distributed database design

## UNIT-2

### Relational Algebra

- Relational Algebra is procedural query language, which takes Relation as input and generates relation as output. Relational algebra mainly provides theoretical foundation for relational databases and SQL.
- Relational algebra is a procedural query language, it means that it tells what data to be retrieved and how to be retrieved.
- Relational Algebra works on the whole table at once, so we do not have to use loops etc to iterate over all the rows (tuples) of data one by one.
- All we have to do is specify the table name from which we need the data, and in a single line of command, relational algebra will traverse the entire given table to fetch data for you.

We can use Relational Algebra to fetch data from this Table(relation)

**Select Name students with age less than 17**

Output

Name
Ckon
Dkon

ID	Name	Age
1	Akon	17
2	Bkon	19
3	Ckon	15
4	Dkon	13

The output for query is also in form of a table(relation), with results in different columns

## Basic/Fundamental Operations:

1. Select ( $\sigma$ )
2. Project ( $\Pi$ )
3. Union ( $\cup$ )
4. Set Difference ( $-$ )
5. Cartesian product ( $\times$ )
6. Rename ( $\rho$ )

**1. Select Operation ( $\sigma$ ) :** This is used to fetch rows (tuples) from table(relation) which satisfies a given condition.

**Syntax:**  $\sigma_p(r)$

- $\sigma$  is the predicate
- $r$  stands for relation which is the name of the table
- $p$  is propositional logic

ex:  $\sigma_{age > 17}(\text{Student})$

This will fetch the tuples(rows) from table **Student**, for which **age** will be greater than **17**.

$\sigma_{age > 17 \text{ and gender} = \text{'Male'}}(\text{Student})$

This will return tuples(rows) from table **Student** with information of male students, of age more than 17.

BRANCH_NAME	LOAN_NO	AMOUNT
Downtown	L-17	1000
Redwood	L-23	2000
Perryride	L-15	1500
Downtown	L-14	1500
Mianus	L-13	500
Roundhill	L-11	900
Perryride	L-16	1300

**Input:**

$\sigma$  BRANCH\_NAME="perryride" (LOAN)

**Output:**

BRANCH_NAME	LOAN_NO	AMOUNT
Perryride	L-15	1500
Perryride	L-16	1300

**Project Operation ( $\Pi$ ):**

- Project operation is used to project only a certain set of attributes of a relation. In simple words, If you want to see only the **names** all of the students in the **Student** table, then you can use Project Operation.
- It will only project or show the columns or attributes asked for, and will also remove duplicate data from the columns.

**Syntax of Project Operator ( $\Pi$ )**

$\Pi$  column\_name1, column\_name2, ....., column\_nameN(table\_name)

Example:

$\Pi$  Name, Age(Student)

Above statement will show us only the **Name** and **Age** columns for all the rows of data in **Student** table.

**Example: CUSTOMER RELATION**

NAME	STREET	CITY
Jones	Main	Harrison
Smith	North	Rye
Hays	Main	Harrison

Curry	North	Rye
Johnson	Alma	Brooklyn
Brooks	Senator	Brooklyn

**Input:**

$\Pi$  NAME, CITY (CUSTOMER)

**Output:**

NAME	CITY
Jones	Harrison
Smith	Rye
Hays	Harrison
Curry	Rye
Johnson	Brooklyn
Brooks	Brooklyn

**Union Operation (U):**

- This operation is used to fetch data from two relations (tables) or temporary relation (result of another operation).
- For this operation to work, the relations (tables) specified should have same number of attributes (columns) and same attribute domain. Also the duplicate tuples are automatically eliminated from the result.

**Syntax:**  $A \cup B$

$\Pi_{\text{Student}}(\text{RegularClass}) \cup \Pi_{\text{Student}}(\text{ExtraClass})$

Example:

#### DEPOSITOR RELATION

CUSTOMER_NAME	ACCOUNT_NO
Johnson	A-101
Smith	A-121
Mayes	A-321
Turner	A-176
Johnson	A-273
Jones	A-472
Lindsay	A-284

#### BORROW RELATION

CUSTOMER_NAME	LOAN_NO
Jones	L-17
Smith	L-23
Hayes	L-15
Jackson	L-14
Curry	L-93
Smith	L-11

Williams	L-17
----------	------

**Input:**

$\Pi$  CUSTOMER\_NAME (BORROW)  $\cup$   $\Pi$  CUSTOMER\_NAME (DEPOSITOR)

**Output:**

CUSTOMER_NAME
Johnson
Smith
Hayes
Turner
Jones
Lindsay
Jackson
Curry
Williams
Mayes

**Set Difference (-):**

This operation is used to find data present in one relation and not present in the second relation. This operation is also applicable on two relations, just like Union operation.

**Syntax:** A - B

where A and B are relations.

For example, if we want to find name of students who attend the regular class but not the extra class, then, we can use the below operation:

$\Pi$  Student(RegularClass) -  $\Pi$  Student(ExtraClass)



**Input:**  $\{ \text{CUSTOMER\_NAME (BORROW)} \} \cap \{ \text{CUSTOMER\_NAME (DEPOSITOR)} \}$

CUSTOMER_NAME
Smith
Jones

### **Cartesian Product (X):**

This is used to combine data from two different relations (tables) into one and fetch data from the combined relation.

**Syntax:** A X B

For example, if we want to find the information for Regular Class and Extra Class which are conducted during morning, then, we can use the following operation:

$\sigma_{\text{time} = \text{'morning'}} (\text{RegularClass X ExtraClass})$

For the above query to work, both **RegularClass** and **ExtraClass** should have the attribute **time**.

Notation: E X D

### **EMPLOYEE**

EMP_ID	EMP_NAME	EMP_DEPT
1	Smith	A
2	Harry	C
3	John	B

## DEPARTMENT

DEPT_NO	DEPT_NAME
A	Marketing
B	Sales
C	Legal

### Input:

EMPLOYEE X DEPARTMENT

### Output:

EMP_ID	EMP_NAME	EMP_DEPT	DEPT_NO	DEPT_NAME
1	Smith	A	A	Marketing
1	Smith	A	B	Sales
1	Smith	A	C	Legal
2	Harry	C	A	Marketing
2	Harry	C	B	Sales
2	Harry	C	C	Legal
3	John	B	A	Marketing
3	John	B	B	Sales

3	John	B	C	Legal
---	------	---	---	-------

### Rename Operation ( $\rho$ ):

This operation is used to rename the output relation for any query operation which returns result like Select, Project etc. Or to simply rename a relation(table)

**Syntax:**  $\rho(\text{RelationNew}, \text{RelationOld})$

The rename operation is used to rename the output relation. It is denoted by  **$\rho$**  ( $\rho$ ).

**Example:** We can use the rename operator to rename STUDENT relation to STUDENT1.  
 $\rho(\text{STUDENT1}, \text{STUDENT})$

### Join in DBMS:

- A JOIN clause is used to combine rows from two or more tables, based on a related column between them.
- **Join in DBMS** is a binary operation which allows you to combine join product and selection in one single statement.
- The goal of creating a join condition is that it helps you to combine the data from two or more DBMS tables.
- The tables in DBMS are associated using the primary key and foreign keys.

### Types of SQL JOIN

1. INNER JOIN
2. LEFT JOIN
3. RIGHT JOIN
4. FULL JOIN

**Table name: EMPLOYEE**

EMP_ID	EMP_NAME	CITY	SALARY	AGE
1	Angelina	Chicago	200000	30
2	Robert	Austin	300000	26
3	Christian	Denver	100000	42
4	Kristen	Washington	500000	29
5	Russell	Los angels	200000	36
6	Marry	Canada	600000	48

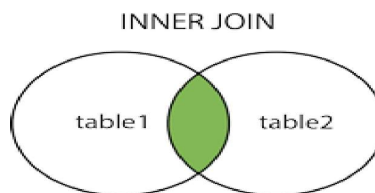
## PROJECT

PROJECT_NO	EMP_ID	DEPARTMENT
101	1	Testing
102	2	Development
103	3	Designing
104	4	Development

### 1. INNER JOIN

In SQL, INNER JOIN selects records that have matching values in both tables as long as the condition is satisfied.

It returns the combination of all rows from both the tables where the condition satisfies.



**Syntax**

```
SELECT table1.column1, table1.column2  
FROM table1 INNER JOIN table2  
ON table1.matching_column = table2.matching_column;
```

**Query**

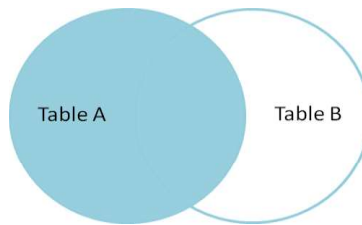
```
SELECT EMPLOYEE.EMP_NAME, PROJECT.DEPARTMENT  
FROM EMPLOYEE INNER JOIN PROJECT  
ON PROJECT.EMP_ID = EMPLOYEE.EMP_ID;
```

**Output**

EMP_NAME	DEPARTMENT
Angelina	Testing
Robert	Development
Christian	Designing
Kristen	Development

**2. LEFT JOIN**

The SQL left join returns all the values from left table and the matching values from the right table. If there is no matching join value, it will return NULL.



### Syntax

```
SELECT table1.column1, table1.column2 FROM table1
LEFT JOIN table2
ON table1.matching_column = table2.matching_column;
```

### Query

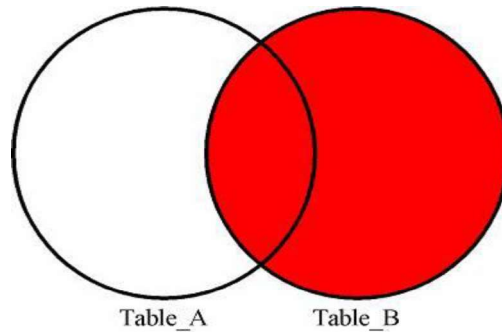
```
SELECT EMPLOYEE.EMP_NAME, PROJECT.DEPARTMENT
FROM EMPLOYEE LEFT JOIN PROJECT
ON PROJECT.EMP_ID = EMPLOYEE.EMP_ID;
```

### Output

EMP_NAME	DEPARTMENT
Angelina	Testing
Robert	Development
Christian	Designing
Kristen	Development
Russell	NULL
Marry	NULL

## 3. RIGHT JOIN

In SQL, RIGHT JOIN returns all the values from the values from the rows of right table and the matched values from the left table. If there is no matching in both tables, it will return NULL.



### Syntax

```
SELECT table1.column1, table1.column2
FROM table1 RIGHT JOIN table2
ON table1.matching_column = table2.matching_column;
```

### Query

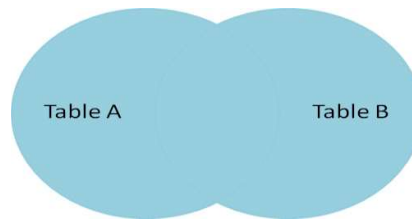
```
SELECT EMPLOYEE.EMP_NAME, PROJECT.DEPARTMENT
FROM EMPLOYEE RIGHT JOIN PROJECT
ON PROJECT.EMP_ID = EMPLOYEE.EMP_ID;
```

### Output

EMP_NAME	DEPARTMENT
Angelina	Testing
Robert	Development
Christian	Designing
Kristen	Development

#### 4. FULL JOIN

In SQL, FULL JOIN is the result of a combination of both left and right outer join. Join tables have all the records from both tables. It puts NULL on the place of matches not found.



##### Syntax

```
SELECT table1.column1, table1.column2  
FROM table1 FULL JOIN table2  
ON table1.matching_column = table2.matching_column;
```

##### Query

```
SELECT EMPLOYEE.EMP_NAME, PROJECT.DEPARTMENT  
FROM EMPLOYEE  
FULL JOIN PROJECT  
ON PROJECT.EMP_ID = EMPLOYEE.EMP_ID;
```

##### Output

EMP_NAME	DEPARTMENT
Angelina	Testing
Robert	Development
Christian	Designing
Kristen	Development
Russell	NULL



Marry	NULL
-------	------

### Division Operator in SQL

**Division Operator ( $\div$ ):** Division operator  $A \div B$  can be applied if and only if:

- Attributes of B is proper subset of Attributes of A.
- The relation returned by division operator will have attributes = (All attributes of A – All Attributes of B)
- The relation returned by division operator will return those tuples from relation A which are associated to every B's tuple.

The division operator is used when we have to evaluate queries which contain the keyword **ALL**.

A	B1	B2	B3																													
<table><tr><th>sno</th><th>pno</th></tr><tr><td>s1</td><td>p1</td></tr><tr><td>s1</td><td>p2</td></tr><tr><td>s1</td><td>p3</td></tr><tr><td>s1</td><td>p4</td></tr><tr><td>s2</td><td>p1</td></tr><tr><td>s2</td><td>p2</td></tr><tr><td>s3</td><td>p2</td></tr><tr><td>s4</td><td>p2</td></tr><tr><td>s4</td><td>p4</td></tr></table>	sno	pno	s1	p1	s1	p2	s1	p3	s1	p4	s2	p1	s2	p2	s3	p2	s4	p2	s4	p4	<table><tr><th>pno</th></tr><tr><td>p2</td></tr></table>	pno	p2	<table><tr><th>pno</th></tr><tr><td>p2</td></tr><tr><td>p4</td></tr></table>	pno	p2	p4	<table><tr><th>pno</th></tr><tr><td>p1</td></tr><tr><td>p2</td></tr><tr><td>p4</td></tr></table>	pno	p1	p2	p4
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**Table 1: Course\_Taken** → It consists of the names of Students against the courses that they have taken.

Student_Name	Course
Robert	Databases

Robert	Programming Languages
David	Databases
David	Operating Systems
Hannah	Programming Languages
Hannah	Machine Learning
Tom	Operating Systems

**Table 2: Course\_Required** → It consists of the courses that one is required to take in order to graduate.

Course
Databases
Programming Languages

1. Find all the students

Create a set of all students that have taken courses. This can be done easily using the following command.

**CREATE TABLE AllStudents AS SELECT DISTINCT Student\_Name FROM Course\_Taken**

This command will return the table **AllStudents**, as the resultset:

Student_name
Robert
David
Hannah
Tom

## **2. Find all the students and the courses required to graduate**

Next, we will create a set of students and the courses they need to graduate. We can express this in the form of Cartesian Product of **AllStudents** and **Course\_Required** using the following command.

**CREATE table StudentsAndRequired AS  
SELECT AllStudents.Student\_Name, Course\_Required.Course  
FROM AllStudents, Course\_Required**

Now the new resultset - table **StudentsAndRequired** will be:

Student_Name	Course
Robert	Databases
Robert	Programming Languages
David	Databases
David	Programming Languages
Hannah	Databases
Hannah	Programming Languages
Tom	Databases
Tom	Programming Languages

### **Relational Calculus:**

Relational calculus is a non-procedural query language that tells the system what data to be retrieved but doesn't tell how to retrieve it. Relational Calculus exists in two forms:

1. Tuple Relational Calculus (TRC)
2. Domain Relational Calculus (DRC)

### **Tuple Relational Calculus (TRC)**

Tuple relational calculus is used for selecting those tuples that satisfy the given condition.

Table: Student

First_Name	Last_Name	Age
-----	-----	----
Ajeet	Singh	30
Chaitanya	Singh	31
Rajeev	Bhatia	27
Carl	Pratap	28

Lets write relational calculus queries.

Query to display the last name of those students where age is greater than 30

```
{ t.Last_Name | Student(t) AND t.age > 30 }
```

In the above query you can see two parts separated by | symbol. The second part is where we define the condition and in the first part we specify the fields which we want to display for the selected tuples.

The result of the above query would be:

Last_Name
-----
Singh

Query to display all the details of students where Last name is 'Singh'

```
{ t | Student(t) AND t.Last_Name = 'Singh' }
```

### Output:

First_Name	Last_Name	Age
-----	-----	----
Ajeet	Singh	30
Chaitanya	Singh	31

Ex:

**Table-1: Customer**

Customer name	Street	City
Saurabh	A7	Patiala
Mehak	B6	Jalandhar
Sumiti	D9	Ludhiana
Ria	A5	Patiala

**Table-2: Branch**

Branch name	Branch city
ABC	Patiala
DEF	Ludhiana
GHI	Jalandhar

**Table-3: Account**

Account number	Branch name	Balance
1111	ABC	50000
1112	DEF	10000
1113	GHI	9000

Account number	Branch name	Balance
1114	ABC	7000

**Table-4: Loan**

Loan number	Branch name	Amount
L33	ABC	10000
L35	DEF	15000
L49	GHI	9000
L98	DEF	65000

**Table-5: Borrower**

Customer name	Loan number
Saurabh	L33
Mehak	L49
Ria	L98

**Table-6: Depositor**

Customer name	Account number
Saurabh	1111

Customer name	Account number
Mehak	1113
Sumiti	1114

**Queries-1:** Find the loan number, branch, amount of loans of greater than or equal to 10000 amount.

$\{t \mid t \in \text{loan} \wedge t[\text{amount}] \geq 10000\}$

Resulting relation:

Loan number	Branch name	Amount
L33	ABC	10000
L35	DEF	15000
L98	DEF	65000

### Domain Relational Calculus (DRC)

In domain relational calculus the records are filtered based on the domains.

Again we take the same table to understand how DRC works.

Table: Student

First_Name	Last_Name	Age
-----	-----	----
Ajeet	Singh	30
Chaitanya	Singh	31
Rajeev	Bhatia	27
Carl	Pratap	28

Query to find the first name and age of students where student age is greater than 27

$\{ \langle \text{First\_Name}, \text{Age} \rangle \mid \in \text{Student} \wedge \text{Age} > 27 \}$



**Note:**

The symbols used for logical operators are:  $\wedge$  for AND,  $\vee$  for OR and  $\neg$  for NOT.

**Output:**

First_Name	Age
Ajeet	30
Chaitanya	31
Carl	28

**SQL Basic Structure**

1. Basic structure of an SQL expression consists of **select**, **from** and **where** clauses.
  - **select** clause lists attributes to be copied - corresponds to relational algebra **project**.
  - **from** clause corresponds to Cartesian product - lists relations to be used.
  - **where** clause corresponds to selection predicate in relational algebra.

The SELECT statement is used to select data from a database.

The data returned is stored in a result table, called the result-set.

To fetch the entire table or all the fields in the table:

```
SELECT * FROM table_name;
```

To fetch individual column data

```
SELECT column1,column2 FROM table_name
```

**WHERE SQL clause**

**WHERE** clause is used to specify/apply any condition while retrieving, updating or deleting data from a table. This clause is used mostly with **SELECT**, **UPDATE** and **DELETE** query.

The basic syntax of the SELECT statement with the WHERE clause is as shown below.

```
SELECT column1, column2, columnN
```

```
FROM table_name
```

WHERE [condition]

### Example

Consider the CUSTOMERS table having the following records –

ID	NAME	AGE	ADDRESS	SALARY
1	Ramesh	32	Ahmedabad	2000.00
2	Khilan	25	Delhi	1500.00
3	kaushik	23	Kota	2000.00
4	Chaitali	25	Mumbai	6500.00
5	Hardik	27	Bhopal	8500.00
6	Komal	22	MP	4500.00
7	Muffy	24	Indore	10000.00

The following code is an example which would fetch the ID, Name and Salary fields from the CUSTOMERS table, where the salary is greater than 2000 –

```
SQL> SELECT ID, NAME, SALARY
FROM CUSTOMERS
WHERE SALARY > 2000;
```

This would produce the following result –

ID	NAME	SALARY
4	Chaitali	6500.00
5	Hardik	8500.00
6	Komal	4500.00
7	Muffy	10000.00

**From clause:**

From clause can be used to specify a sub-query expression in SQL. The relation produced by the sub-query is then used as a new relation on which the outer query is applied.

- Sub queries in the from clause are supported by most of the SQL implementations.
- The correlation variables from the relations in from clause cannot be used in the sub-queries in the from clause.

**Syntax:**

**SELECT column1, column2 FROM**

**(SELECT column\_x as C1, column\_y FROM table WHERE PREDICATE\_X)**

**as table2**

**WHERE PREDICATE;**

**SET Operations**

SQL supports few Set operations which can be performed on the table data. These are used to get meaningful results from data stored in the table, under different special conditions.

In this tutorial, we will cover 4 different types of SET operations, along with example:

1. UNION
2. UNION ALL
3. INTERSECT
4. MINUS

**1. Union**

- The SQL Union operation is used to combine the result of two or more SQL SELECT queries.
- In the union operation, all the number of datatype and columns must be same in both the tables on which UNION operation is being applied.
- The union operation eliminates the duplicate rows from its resultset.

**Syntax**

```
SELECT column_name FROM table1
```

```
UNION
```

```
SELECT column_name FROM table2;
```

**The First table**

ID	NAME
1	Jack
2	Harry
3	Jackson

**The Second table**

ID	NAME
3	Jackson
4	Stephan
5	David

Union SQL query will be:

```
SELECT * FROM First
```

```
UNION
```

```
SELECT * FROM Second;
```

The resultset table will look like:

ID	NAME
1	Jack
2	Harry
3	Jackson
4	Stephan
5	David

## 2. Union All

Union All operation is equal to the Union operation. It returns the set without removing duplication and sorting the data.

### Syntax:

```
SELECT column_name FROM table1
UNION ALL
SELECT column_name FROM table2;
```

**Example:** Using the above First and Second table.

Union All query will be like:

```
SELECT * FROM First
UNION ALL
SELECT * FROM Second;
```

The resultset table will look like:

-ID	NAME
1	Jack
2	Harry
3	Jackson
3	Jackson
4	Stephan
5	David

### 3. Intersect

- It is used to combine two SELECT statements. The Intersect operation returns the common rows from both the SELECT statements.
- In the Intersect operation, the number of datatype and columns must be the same.
- It has no duplicates and it arranges the data in ascending order by default.

#### Syntax

```
SELECT column_name FROM table1
INTERSECT
SELECT column_name FROM table2;
```

#### Example:

**Using the above First and Second table.**

Intersect query will be:

```
SELECT * FROM First
INTERSECT
SELECT * FROM Second;
```

The resultset table will look like:

ID	NAME
3	Jackson

#### 4. Minus

- It combines the result of two SELECT statements. Minus operator is used to display the rows which are present in the first query but absent in the second query.
- It has no duplicates and data arranged in ascending order by default.

##### **Syntax:**

```
SELECT column_name FROM table1
MINUS
SELECT column_name FROM table2;
```

##### **Example**

**Using the above First and Second table.**

Minus query will be:

```
SELECT * FROM First
MINUS
SELECT * FROM Second;
```

The resultset table will look like:

ID	NAME
1	Jack
2	Harry

### Aggregate functions in SQL

- SQL aggregation function is used to perform the calculations on multiple rows of a single column of a table. It returns a single value.
- It is also used to summarize the data.

#### Aggregate Functions

- 1) Count()
- 2) Sum()
- 3) Avg()
- 4) Min()
- 5) Max()

### 1. COUNT FUNCTION

- COUNT function is used to Count the number of rows in a database table. It can work on both numeric and non-numeric data types.
- COUNT function uses the COUNT(\*) that returns the count of all the rows in a specified table. COUNT(\*) considers duplicate and Null.

**Count(\*):** Returns total number of records



### PRODUCT\_MAST

PRODUCT	COMPANY	QTY	RATE	COST
Item1	Com1	2	10	20
Item2	Com2	3	25	75
Item3	Com1	2	30	60
Item4	Com3	5	10	50
Item5	Com2	2	20	40
Item6	Cpm1	3	25	75
Item7	Com1	5	30	150
Item8	Com1	3	10	30
Item9	Com2	2	25	50
Item10	Com3	4	30	120

#### Example: COUNT()

```
SELECT COUNT(*) FROM PRODUCT_MAST;
```

#### Output:

```
10
```

#### Example: COUNT with WHERE

```
SELECT COUNT(*)  
FROM PRODUCT_MAST;  
WHERE RATE>=20;
```

#### Output:7

#### Example: COUNT() with DISTINCT

```
SELECT COUNT(DISTINCT COMPANY)
FROM PRODUCT_MAST;
```

**Output:**

```
3
```

## 2. SUM Function

Sum function is used to calculate the sum of all selected columns. It works on numeric fields only.

**Syntax**

```
SUM()
```

or

```
SUM( [ALL|DISTINCT] expression )
```

**Example: SUM()**

```
SELECT SUM(COST)
FROM PRODUCT_MAST;
```

**Output:**

```
670
```

**Example: SUM() with WHERE**

```
SELECT SUM(COST)
FROM PRODUCT_MAST
WHERE QTY>3;
```

**Output:**

```
320
```

## 3. AVG function

The AVG function is used to calculate the average value of the numeric type. AVG function returns the average of all non-Null values.

**Syntax**

AVG()

**Example:**

```
SELECT AVG(COST)
FROM PRODUCT_MAST;
```

**Output:**

67.00

**4. MAX Function**

MAX function is used to find the maximum value of a certain column. This function determines the largest value of all selected values of a column.

**Syntax:** MAX()**Example:**

```
SELECT MAX(RATE)
FROM PRODUCT_MAST;
```

30

**5. MIN Function**

MIN function is used to find the minimum value of a certain column. This function determines the smallest value of all selected values of a column.

**Syntax:** MIN() )**Example:** SELECT MIN(RATE)

```
FROM PRODUCT_MAST;
```

**Output:**10**GROUP BY Statement**

The GROUP BY statement groups rows that have the same values into summary rows, like "find the number of customers in each country".

The GROUP BY statement is often used with aggregate functions (COUNT, MAX, MIN, SUM, AVG) to group the result-set by one or more columns.

#### GROUP BY Syntax

```
SELECT column_name(s)
FROM table_name
WHERE condition
GROUP BY column_name(s)
ORDER BY column_name(s);
```

SI NO	NAME	SALARY	AGE	SUBJECT	YEAR	NAME
1	Harsh	2000	19	English	1	Harsh
2	Dhanraj	3000	20	English	1	Pratik
3	Ashish	1500	19	English	1	Ramesh
4	Harsh	3500	19	English	2	Ashish
5	Ashish	1500	19	English	2	Suresh
				Mathematics	1	Deepak
				Mathematics	1	Sayan

#### Example:

- **Group By single column:** Group By single column means, to place all the rows with same value of only that particular column in one group. Consider the query as shown below:
- SELECT NAME, SUM(SALARY) FROM Employee
- GROUP BY NAME;

The above query will produce the below output:

NAME	SALARY
Ashish	3000
Dhanraj	3000
Harsh	5500

**Group By multiple columns:** Group by multiple column is say for example, **GROUP BY column1, column2**. This means to place all the rows with same values of both the columns **column1** and **column2** in one group. Consider the below query:

```
SELECT SUBJECT, YEAR, Count(*)
FROM Student
GROUP BY SUBJECT, YEAR;
```

SUBJECT	YEAR	Count
English	1	3
English	2	2
Mathematics	1	2

### HAVING Clause:

We know that WHERE clause is used to place conditions on columns but what if we want to place conditions on groups?

This is where HAVING clause comes into use. We can use HAVING clause to place conditions to decide which group will be the part of final result-set. Also we can not use the aggregate functions like SUM(), COUNT() etc. with WHERE clause. So we have to use HAVING clause if we want to use any of these functions in the conditions.

### Syntax:

```
SELECT column1, function_name(column2)
FROM table_name
WHERE condition
GROUP BY column1, column2
HAVING condition
ORDER BY column1, column2;
```

**function\_name:** Name of the function used for example, SUM() , AVG().

**table\_name:** Name of the table.

**condition:** Condition used.

**Example:**

```
SELECT NAME, SUM(SALARY) FROM Employee
GROUP BY NAME
HAVING SUM(SALARY)>3000;
```

Example

Consider the CUSTOMERS table having the following records.

ID	NAME	AGE	ADDRESS	SALARY
1	Ramesh	32	Ahmedabad	2000.00
2	Khilan	25	Delhi	1500.00
3	kaushik	23	Kota	2000.00
4	Chaitali	25	Mumbai	6500.00
5	Hardik	27	Bhopal	8500.00
6	Komal	22	MP	4500.00
7	Muffy	24	Indore	10000.00

Following is an example, which would display a record for a similar age count that would be more than or equal to 2.

```
SQL > SELECT ID, NAME, AGE, ADDRESS, SALARY
FROM CUSTOMERS
GROUP BY age
HAVING COUNT(age) >= 2;
```

This would produce the following result –

ID	NAME	AGE	ADDRESS	SALARY
2	Khilan	25	Delhi	1500.00

+---+-----+---+-----+-----+

### Nested Queries

In nested queries, a query is written inside a query. The result of inner query is used in execution of outer query. We will use **STUDENT**, **COURSE**, **STUDENT\_COURSE** tables for understanding nested queries.

#### STUDENT

S_ID	S_NAME	S_ADDRESS	S_PHONE	S_AGE
S1	RAM	DELHI	9455123451	18
S2	RAMESH	GURGAON	9652431543	18
S3	SUJIT	ROHTAK	9156253131	20
S4	SURESH	DELHI	9156768971	18

#### COURSE

C_ID	C_NAME
C1	DSA
C2	Programming
C3	DBMS

#### STUDENT\_COURSE

S_ID	C_ID
------	------

S1	C1
S1	C3
S2	C1
S3	C2
S4	C2
S4	C3

### Example

Consider the CUSTOMERS table having the following records –

```
+---+-----+---+-----+-----+
| ID | NAME   | AGE | ADDRESS | SALARY |
+---+-----+---+-----+-----+
| 1 | Ramesh | 35 | Ahmedabad | 2000.00 |
| 2 | Khilan | 25 | Delhi    | 1500.00 |
| 3 | kaushik | 23 | Kota     | 2000.00 |
| 4 | Chaitali | 25 | Mumbai   | 6500.00 |
| 5 | Hardik | 27 | Bhopal   | 8500.00 |
| 6 | Komal | 22 | MP       | 4500.00 |
| 7 | Muffy | 24 | Indore   | 10000.00 |
+---+-----+---+-----+-----+
```

Now, let us check the following subquery with a SELECT statement.

```
SQL> SELECT *
      FROM CUSTOMERS
      WHERE ID IN (SELECT ID
```



FROM CUSTOMERS

WHERE SALARY > 4500);

This would produce the following result.

```
+---+-----+---+-----+-----+
| ID | NAME   | AGE | ADDRESS | SALARY |
+---+-----+---+-----+-----+
| 4 | Chaitali | 25 | Mumbai | 6500.00 |
| 5 | Hardik   | 27 | Bhopal  | 8500.00 |
| 7 | Muffy    | 24 | Indore  | 10000.00 |
+---+-----+---+-----+-----+
```

### Students

id	name	class_id	GPA
1	Jack Black	3	3.45
2	Daniel White	1	3.15
3	Kathrine Star	1	3.85
4	Helen Bright	2	3.10
5	Steve May	2	2.40

### Teachers

id	name	subject	class_id	monthly_salary
1	Elisabeth Grey	History	3	2,500
2	Robert Sun	Literature	[NULL]	2,000
3	John Churchill	English	1	2,350
4	Sara Parker	Math	2	3,000

### Classes

id	grade	teacher_id	number_of_students
1	10	3	21
2	11	4	25
3	12	1	28

```
SELECT *  
FROM students  
WHERE GPA > (  
    SELECT AVG(GPA)  
    FROM students);
```

result:

id	name	class_id	GPA
1	Jack Black	3	3.45
3	Kathrine Star	1	3.85

```
SELECT AVG(number_of_students)
FROM classes
WHERE teacher_id IN (
    SELECT id
    FROM teachers
    WHERE subject = 'English' OR subject = 'History');
```

### Views in SQL

- Views in SQL are considered as a virtual table. A view also contains rows and columns.
- To create the view, we can select the fields from one or more tables present in the database.
- A view can either have specific rows based on certain condition or all the rows of a table.

### Sample table:

#### Student\_Detail

STU_ID	NAME	ADDRESS
1	Stephan	Delhi

2	Kathrin	Noida
3	David	Ghaziabad
4	Alina	Gurugram

### Student\_Marks

STU_ID	NAME	MARKS	AGE
1	Stephan	97	19
2	Kathrin	86	21
3	David	74	18
4	Alina	90	20
5	John	96	18

### 1. Creating view

A view can be created using the **CREATE VIEW** statement. We can create a view from a single table or multiple tables.

#### Syntax:

```
CREATE VIEW view_name AS
SELECT column1, column2.....
FROM table_name
WHERE condition;
```

## 2. Creating View from a single table

### Query:

```
CREATE VIEW DetailsView AS  
SELECT NAME, ADDRESS  
FROM Student_Details  
WHERE STU_ID < 4;
```

Just like table query, we can query the view to view the data.

```
SELECT * FROM DetailsView;
```

### Output:

NAME	ADDRESS
Stephan	Delhi
Kathrin	Noida
David	Ghaziabad

## 3. Creating View from multiple tables

View from multiple tables can be created by simply include multiple tables in the SELECT statement.

In the given example, a view is created named MarksView from two tables Student\_Detail and Student\_Marks.

### Query:

```
CREATE VIEW MarksView AS  
SELECT Student_Detail.NAME, Student_Detail.ADDRESS, Student_Marks.MARKS  
FROM Student_Detail, Student_Mark  
WHERE Student_Detail.NAME = Student_Marks.NAME;
```