

## QUIZ - I

MA423: Matrix Computations

Time : 40 minutes

20 marks

September 8, 2023

Answer ALL questions

1. Let  $x \in F(\beta, t, L, U)$  be such that  $\beta^{e-1} < x < \beta^e$ . Let  $y \in \mathbb{R}$ . If  $|y| < \text{ulp}(x)/2$  then show that  $\text{fl}(x \pm y) = x$ . For  $(\beta, t, L, U) = (10, 4, -30, 30)$  and  $x = 10^{-9}$ , what is the smallest  $y > 0$  such that  $\text{fl}(x + y) > x$ ? Next when  $x = 10^{15}$ , what is the smallest  $y > 0$  such that  $\text{fl}(x + y) > x$ . Here  $\text{fl}(x)$  uses round to nearest rounding mode. **5 marks**

2. Let  $A \in \mathbb{R}^{n \times n}$  is nonsingular. Describe an efficient algorithm (outline only the steps) based on Gaussian elimination for solving the problem  $A^{10}x = b$  and determine the flop count. **5 marks**

3. Consider the system  $Ax = b$ , where  $A \in \mathbb{R}^{n \times n}$  is nonsingular and  $b \in \mathbb{R}^n$ . Suppose that  $A + \Delta A$  is nonsingular. Now consider the system  $(A + \Delta A)\hat{x} = b + \Delta b$ . Show that

$$\frac{\|x - \hat{x}\|}{\|\hat{x}\|} \leq \text{cond}(A) \left( \frac{\|\Delta A\|}{\|A\|} + \frac{\|\Delta b\|}{\|A\| \|\hat{x}\|} \right),$$

where  $\text{cond}(A) := \|A\| \|A^{-1}\|$ .

**5 marks**

4. Consider  $A := \begin{bmatrix} 10^{-4} & 1 \\ 1 & a \end{bmatrix}$ , where  $0 < a \leq 4$ . Prove that in the floating point system  $F(10, 3, -10, 10)$ , the computed LU factorization of  $A$  using GENP is independent of  $a$ . Is GENP backward stable for computing LU factorization of  $A$ ? Justify your answer. **5 marks**

\*\*\* End \*\*\*



## QUIZ - II

MA423: Matrix Computations

Time : 40 minutes

20 marks

November 7, 2023

Answer ALL questions

1. Let  $A \in \mathbb{C}^{m \times n}$  be such that  $\text{rank}(A) = n$ . Let  $\sigma_{\min}(A)$  denote the smallest singular value of  $A$ . Show that  $\|(A^*A)^{-1}A^*\|_2 = 1/\sigma_{\min}(A)$  and  $\|A(A^*A)^{-1}A^*\|_2 = 1$ . **6 marks**
2. Let  $u \in \mathbb{R}^n$  be a unit vector. Consider the Householder reflector  $H := I - 2uu^T$ . Determine all the eigenvalues of  $H$  and their algebraic multiplicities. Also determine all the eigenvectors. **5 marks**
3. Let  $A \in \mathbb{C}^{n \times n}$  and  $\sigma > 0$ . Show that  $\sigma$  is a singular value of  $A \iff \begin{bmatrix} A & -\sigma I_n \\ -\sigma I_n & A^* \end{bmatrix}$  is singular. **4 marks**
4. Let  $A \in \mathbb{R}^{n \times n}$  be proper upper Hessenberg and  $A_3$  be the result of 3 basic  $QR$  steps on  $A$ , that is,  $A_{k-1} = Q_k R_k$ ,  $A_k := R_k Q_k$  for  $k = 1, 2, 3$ , with  $A_0 := A$ . Set  $R := R_3 R_2 R_1$  and  $Q := Q_1 Q_2 Q_3$ . Show that  $A^3 = QR$  and  $A_3 = Q^* A Q$ . **5 marks**

\*\*\* End \*\*\*

$$Q_3^T Q_2^T Q_1^T A^3 = R_3 R_2 R_1$$



## End Semester Examination

MA 423 : Matrix Computations

Time : 3 hours

40 marks

November 22, 2023

Answer ALL questions

1. Let  $A \in \mathbb{C}^{m \times n}$  be such that  $\text{rank}(A) = r < \min(m, n)$  and  $b \in \mathbb{C}^m$ . Define the Moore-Penrose pseudo-inverse  $A^+$  of  $A$  and show that  $A^+b$  is a unique solution of the LSP  $Ax \approx b$  having the smallest norm. Use SVD of  $A$  to show that  $\lim_{\mu \rightarrow 0^+} (A^*A + \mu I_n)^{-1} A^* = A^+$ . **6 marks**
2. Let  $A := [a_1 \ \dots \ a_n] \in \mathbb{R}^{n \times n}$ . Use QR factorization of  $A$  to prove the Hadamard's determinant inequality

$$|\det(A)| \leq \prod_{j=1}^n \|a_j\|_2,$$

where  $\det(A)$  is the determinant of  $A$ .

**4 marks**

3. Let  $A \in \mathbb{C}^{m \times n}$ ,  $b \in \mathbb{C}^m$  and  $\Delta b \in \mathbb{C}^m$ . Let  $x$  and  $\hat{x}$  be unique solutions (with smallest norms) of the LSP  $Ax \approx b$  and LSP  $A\hat{x} \approx b + \Delta b$ , respectively. Show that

$$\frac{\|x - \hat{x}\|_2}{\|x\|_2} \leq \frac{\text{cond}(A) \|\Delta b\|_2}{\cos \theta \|b\|_2},$$

where  $\theta$  is the acute angle between  $b$  and  $Ax$ , and  $\text{cond}(A) := \|A^+\|_2 \|A\|_2$ . Here  $A^+$  is the Moore-Penrose pseudo-inverse of  $A$ .

**4 marks**

4. Compute QR factorization of  $A := \begin{bmatrix} 1 & 0 & 1 \\ -2 & -1 & 0 \\ 2 & 2 & 1 \end{bmatrix}$  in two different ways using Householder reflectors and classical Gram-Schmidt method. **6 marks**

5. Let  $A \in \mathbb{R}^{n \times n}$  be upper Hessenberg. Consider the QR step  $A - \mu I = QR$  and  $A_1 = RQ + \mu I$ . Show that  $A_1$  is upper Hessenberg when  $\mu$  is NOT an eigenvalue of  $A$ . Show that if the QR decomposition of  $A - \mu I$  is computed using rotations then  $A_1$  is upper Hessenberg even if  $\mu$  is an eigenvalue of  $A$ . **5 marks**

6. Let  $A \in \mathbb{R}^{n \times n}$  and  $s \in \mathbb{C}$ . Suppose that  $s$  is not an eigenvalue of  $A$ . Perform two steps of explicit QR algorithm on  $A$  with shifts  $s$  and its complex conjugate  $\bar{s}$ :

$$A - sI = Q_1 R_1, \quad A_1 := R_1 Q_1 + sI$$

$$A_1 - \bar{s}I = Q_2 R_2, \quad A_2 := R_2 Q_2 + \bar{s}I.$$

Show that  $Q_1$  and  $Q_2$  can be chosen such that  $Q_1 Q_2$  is real. Deduce that in such a case  $A_1$  is complex and  $A_2$  is real. Describe an algorithm (outline the steps) for computing  $A_2$  directly from  $A$  without computing  $A_1$  thereby avoiding complex arithmetic. **5 marks**

7. Describe an algorithm (outline the steps) for reducing a matrix  $A \in \mathbb{R}^{n \times n}$  to an upper bidiagonal form  $B$ . Describe a single step of Golub-Kahan implicit QR algorithm for computing SVD of  $B$  when  $B$  is proper upper bidiagonal. **5 marks**

8. Write an algorithm (outline the steps) that implements inverse power method for computing an eigenvalue and a corresponding eigenvector of nonsingular matrix  $A \in \mathbb{C}^{n \times n}$ . Determine the flop count for performing  $\ell$  steps of inverse power method. Suppose that  $A$  has  $n$  distinct eigenvalues  $\lambda_1, \dots, \lambda_n$  such that  $|\lambda_1| \geq \dots \geq |\lambda_{n-1}| > |\lambda_n|$ . Let  $Av_n = \lambda_n v_n$  and  $v_n \neq 0$ . Let  $x_0$  be a nonzero vector such that  $\langle x_0, v_n \rangle \neq 0$ . Show that the iterate  $x_j$  obtained from inverse power method with starting vector  $x_0$  converges to a scalar multiple of  $v_n$  at the rate  $\frac{|\lambda_{n-1}|}{|\lambda_n|}$ . **5 marks**

\*\*\* End \*\*\*