

Name:- Aditya Gupta.

Univ. Roll no:- 2215000094

Sub:- Computer network.

Sec:- B (03)

Answer 1:- Transmission speed = 20 bps

Time to travel 1 bit =  $1/20 = 0.05$  sec.

Propagation speed = 180 cm/ms.

Distance travelled = propagation speed  $\times$  time

$$1 \text{ bit} = 180 \times 0.5 \\ = 9 \text{ m.}$$

Therefore, the 1-bit delay in this network is equivalent to approx 9m of cable.

Answer 2:- Transmission time =  $2 \times$  propagation time.

$$= 2 \times \text{length} / 250000$$

$$= 8000 \times \frac{400}{250000} \Rightarrow 21 \text{ cm}$$

Answer 3:- Length = 12 km = 12000 m.

propagation speed  $v = 8 \text{ ns/m}$

No. of student in the ring = 40

frame size = 220 byte =  $220 \times 8 = 1760$  bits.

Data rate = 100 mbps.

Answer 4:- The cable needed  $n_{c_2} = \frac{n(n-1)}{2}$

$$n=6 \Rightarrow n_{c_2} = \frac{6(6-1)}{2} = 3 \times 5 = 15$$

No. of ports needed for each cable device

$$n-1 \Rightarrow 6-1 \Rightarrow 5 \text{ ports.}$$



Answer 5:- (a) If the connection fails, the other connections will be working.

- (b) If the other devices will still be able to send data through the hub. there will be no access to the device which has the failed connections to the hub.
- (c) All transmission stops if the failure is in the bus if the drop line fails, only the corresponding device cannot be operate.
- (d) The failed connections may disable the whole network unless it is a dual ring or there is a by pass mechanism.

Answer 6:- (a) E-mail is not an interactive application, Even if it is delivered immediately, it may stay in the mail box of the receiver for a while. It is not sensitive to delay.

- (b) we normally do not expect a file to be copied immediately by it is not very sensitive to delay.
- (c) Surfing the internet is an application very sensitive to delay, we expect to get access, to the site we are searching.

Answer 7:- In this case, the communication is only between a caller and the callee. A point to point connections provides a dedicated link between two devices & the other entire capacity of the link is reserved for transmission entire b/w those two devices, since only two devices are involved in a voice telephone call, it is a point-to-point connection.

Answer 8:-

- a) Network layer (layer-3)
- b) Data-link layer (layer-2)
- c) physical-layer (layer-1)
- d) Application layer (layer-7)



sw 4-9 → (a) Presentation Layer

(b) Session Layer

(c) Transport Layer

(d) Session-layer

(e) Presentation-layer.

Answer 10: → If the physical destination address of a frame is corrupted during transmission, the frame is either delivered to the memory station if the corrupted address matches one of the stations, or lost if it does not match any station. However, the error detection mechanism in most datalink protocols will find the error and discard the frame.

Answer 11: → If the logical destination address of a packet becomes corrupted, the packet may not be delivered to the correct destination. The destination computer can send an ICMP message back to the source computer to inform it of the error, the source computer can take corrective action.

Answer 12:-

- i) Packet loss
- ii) No Response
- iii) Error Handling.
- iv) Timeouts.

when no process with the specified destination port address is running, the transport layer cannot deliver the packet to any process, resulting in packet loss and communication failure.

Answer 13:-  $2 \times 1 \text{ MHz} \times \log_2 L = 4 \text{ mbps bit rate}$   
 $\text{signal level} = 2^{\text{bit rate}} = 2^4 = 16$   
bit rate 4 mbps & signal level is 16.



Answer 14:- Shannon capacity formula:  $C = B \times \log_2(1 + \text{SNR})$

Given,

$$\text{SNR} = 0$$

$$C = B \times \log_2(1)$$

Since,

$\log_2(1) = 0$  the channel capacity is this extremely noisy channel.

$$\boxed{C = 0}$$

Answer 15 :- Given,

Token ring LAN data rate = 3 mbps.

frame size = 200 bytes =  $200 \times 8 = 1600$  bits.

Ring latency = 150 ms

Now,

$$\begin{aligned} \text{Round trip time} &= \text{Ring latency} \times 2 \\ &= 150 \times 2 = 300 \text{ ms} \end{aligned}$$

$$\text{Effective data rate of LAN} = \frac{\text{frame size}}{\text{Round trip time}}$$

$$= \frac{1600 \text{ bits}}{300 \text{ ms}} = \frac{5.333}{10^6} \text{ mbps.}$$

$$= 0.005333 \text{ mbps} \text{ \textbf{Ans.}}$$