# **Indian Institute of Technology Patna**



# Mechatronics, Instrumentation And Controls Laboratory Lab 7 Report

Topic: Introduction to Microprocessor

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# 1. Aim of the Experiments.

- 1.Write a program to ADD and subtract two 16 bit number and also store carry for addition and borrow for subtraction.
- 2. Write a program to compute factorial of a number.
- 3. Write a program to multiply to 16 bit no using repetitive addition.
- 4. Write a program to divide to 16 bit no using repetitive subtraction.

# 2. Pre-Requisites/Components Required

In the Simulation Implementation I have used the following Software:

- ➤ EMU8086 THE MICROPROCESSOR EMULATOR
- 1.1) Write a program to ADD and subtract two 16 bit number and also store carry for addition and borrow for subtraction.

Addition: -

Program for ADD

```
org 100h

; set location counter to 100h

; add your code here
MOU CX.0000h

B MOU AX.[3000h]

MOU BX.[3002h]

ADD AX.BX

JNC STORE

INC CX

STORE:

MOU [3004h],AX

MOU [3006h],CX

; Store the AX content into memory location 3004h

MOU [3006h],CX

; If Carry Flag bit=1,Increse Carry register by 1

; Store the AX content into memory location 3006h

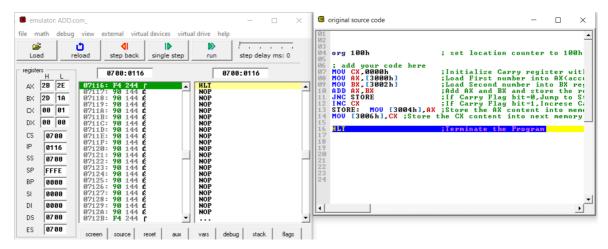
; Store the CX content into next memory location 3006h

; Terminate the Program
```

Data stored in memory location 3000h(First Number) and 3002h(Second Number) respectively

```
Random Access Memory
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                     update
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0700:3060
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                                                      00
```

#### Execution



#### Result



#### Verification

# Result

Hex value:

FE14 + 2D1A = 12B2E

Decimal value:

65044 + 11546 = **76590** 



Indeed my result stored in memory location 3004h(SUM) and 3006h(CARRY), matches with calculated values as shown above.

#### Subtraction: -

Program for Subtraction

```
org 100h

; set location counter to 100h

; add your code here

MOU CX.0000

MOU AX.[3000h]

BX.[3002h]

SUB AX.BX

JNC STORE

INC CX

MOU [3004h], AX

MOU [3006h], CX

Thitialize Borrow register CX with 0000h

;Load First number into AX(accumulator) from memory location 3000h

;Load First number into BX register from memory location 3000h

;Subtract BX from AX, and store the result to AX

;If Carry Flag bit=0, Jump to Store Label

;If Carry Flag bit=1, that means there is a borrow, Increse Borrow register by 1

STORE: MOU [3004h], AX

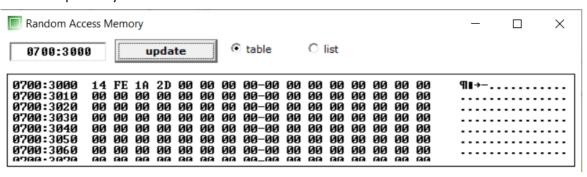
MOU [3006h], CX

;Store the AX content into memory location 3004h

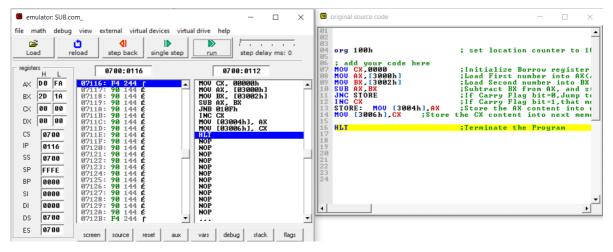
;Store the CX content into next memory location 3006h

;Terminate the Program
```

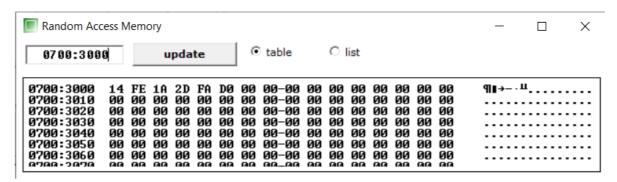
Data stored in memory location 3000h(First Number) and 3002h(Second Number) respectively



Execution



Result



Verification

# Result

Hex value:

FE14 - 2D1A = D0FA

Decimal value:

65044 - 11546 = **53498** 



- Indeed my result stored in memory location 3004h(Difference) and 3006h(Borrow), matches with calculated values as shown above.
- 1.2) Write a program to compute factorial of a number.

#### Factorial: -

Program

```
org 100h ; set location counter to 100h

add your code here

MOU CX, [0500h] ; Load the Number into CX register from memory location 0500h

[Initialize AX(accumulator) with 000h

Li: MUL CX ; Multiply AX with CX and store the result in AX

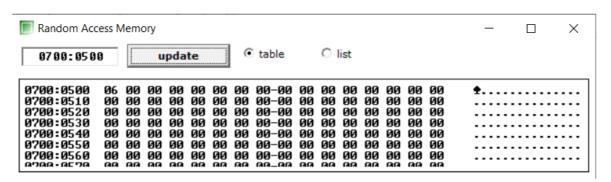
LOOP L1 ; Repeat a series of instructions. If CX is not 0, execution will jump to L1 Lable.

[If CX = 0 after the auto decrement, execution will simply go on to the next instruction after LOOP.

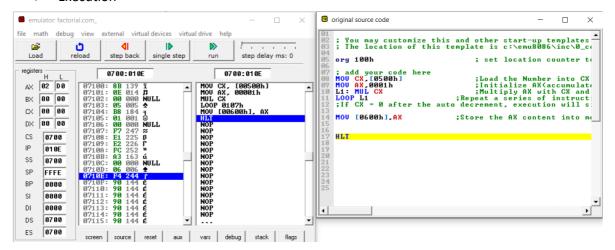
MOU [0600h],AX ; Store the AX content into memory location 0600h

HLT ; Terminate the Program
```

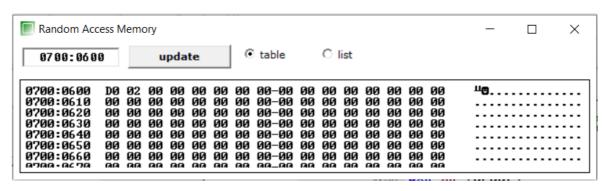
Data stored in memory location 0500h(Number)



#### Execution



#### Result



#### Verification

# Result

Hex value:

 $2d0 \times 1 = 2D0$ 

## Decimal value:

 $720 \times 1 = 720$ 



Indeed my result stored in memory location 0600h,matches with calculated values as shown above.

1.3) Write a program to multiply to 16 bit no using repetitive addition.

Multiplication using repetitive addition: -

Program

```
; You may customize this and other start-up templates;
The location of this template is c:\emu8086\inc\0_com_template.txt

org 100h

; set location counter to 100h
; add your code here
HOU CX, [8500h]

EC CX

; Load the First Number into CX register from memory location 0500h
; Because at first we are taking two number already for MUL,
; So Decreasing CX register by 1 initially. So it will itterate one less than CX.

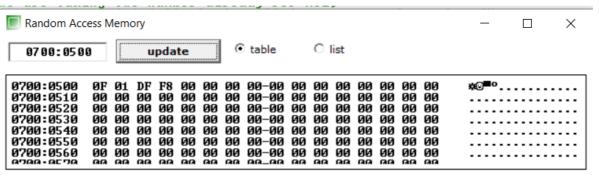
HOU BX, 0800h
HOU BX, 0800h
HOU BX, 0800h
HOU DX, 0800h
HOU DX, 0800h
HOU DX, 0800h
I ; Load the Second Number into AX register from memory location 0502h
; Load the Second Number into AX register from memory location 0502h
; Hou BX giff Carry Flag bit-0, Junp to CARRY Label
INC BX
I ; If Carry Flag bit-0, Junp to CARRY Label
INC BX
I ; If CARRY i ; If CARRY Flag bit-0, Junp to CARRY Label
INC BX
I ; If CARRY i ; If CARRY Flag bit-0, Junp to CARRY Label
I ; If CX = 0 after the auto decrement, execution will jump to L1 Lable.

HOU [0600h], AX
HOU [0600h], BX
; Store the AX content into memory location 0600h
HLT
; Terminate the Program

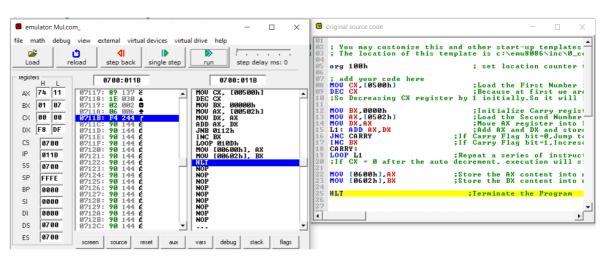
Terminate the Program

Terminate the Program
```

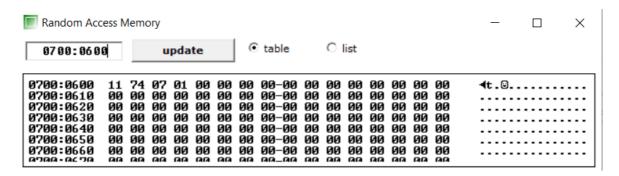
Data stored in memory location 0500h(First Number) and 0502h(Second Number) respectively



Execution



> Result



Verification

## Result

Hex value:

F8DF × 010F = 1077411

Decimal value:

63711 × 271 = **17265681** 



- Indeed my result stored in memory location 0600h(Lower 16 bit Result) and 0602h(Higher 16 bit Result), matches with calculated values as shown above.
- 1.4) Write a program to divide to 16 bit no using repetitive subtraction.

Division using repetitive subtraction: -

Program

```
org 100h

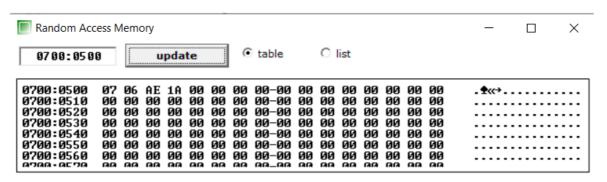
; add your code here

MOU CX, [0500h]

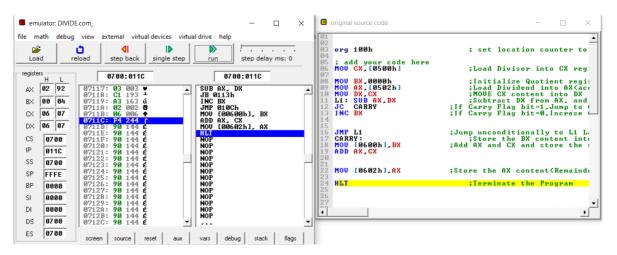
; Load Divisor into CX register from memory location 0500h

MOU BX, 0000h
MOU BX, 0000h
MOU AX, [0502h]
MOU BX, 0000h
MOU
```

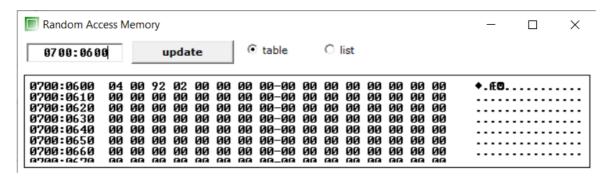
Data stored in memory location 0500h(Divisor) and 0502h(Dividend) respectively



#### Execution



#### Result



#### Verification

## Result

Hex value:

1AAE ÷ 0607 = 4 Remainder : 292

Decimal value:

6830 ÷ 1543 = 4 Remainder : 658



Indeed my result stored in memory location 0600h(Quotient) and 0602h(Remainder), matches with calculated values as shown above.

#### Question: -

1.) What you have understood of processors in short, one para?

Ans. A processor is the logic circuitry that responds to and processes the basic instructions that drive a computer. Thus it's basic job is to receive input and provide the appropriate output. Nowadays, Processors are found in all modern electronic devices, including PCs, smartphones, tablets, etc. Their purpose is to receive input in the form of program instructions and execute trillions of calculations to provide the output that the user will interface with.

A processor is made of four basic elements:

- The arithmetic logic unit (ALU),
- The floating point unit (FPU),
- Registers
- Cache memories.

The **ALU and FPU** carry basic and advanced arithmetic and logic operations on numbers, and then results are sent to the **registers**, which also store instructions. **Caches** are small and fast memories that store copies of data for frequent use, and act similarly to a random-access memory (RAM).

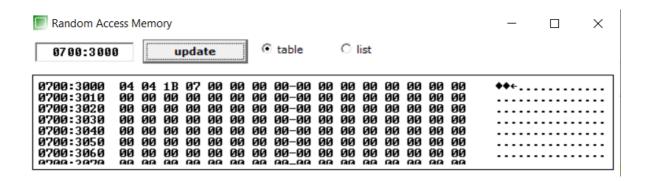
Basically, a processor carries out his operations through the three main steps of the instruction cycle: fetch, decode, and execute.

- Fetch: The processor retrieves instructions, usually from a RAM.
- > Decode: A decoder converts the instruction into signals to the other components of processing unit .

- Execute: The now decoded instructions are sent to each component so that the desired operation can be performed.
- 2.) What happens when you subtract a-b where b>a and a+bwhen (a+b)> 16bit ? What would be stored in the accumulator

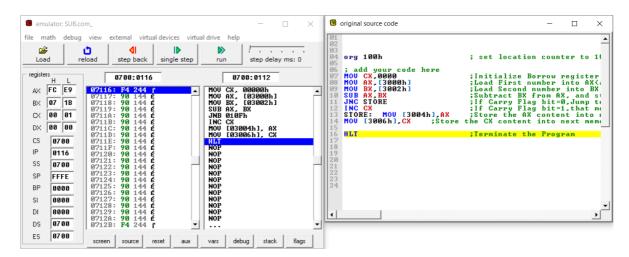
Ans)

I have taken a=0404h and b as 071Bh, where b>a.



#### Inference: -

➤ After Subtraction Operation, result being stored in Accumulator contains FCE9 i.e, AX= FCE9, which shows the difference between the two, taking the borrow and the borrow being saved in CX. We can see borrow=1/CX=1



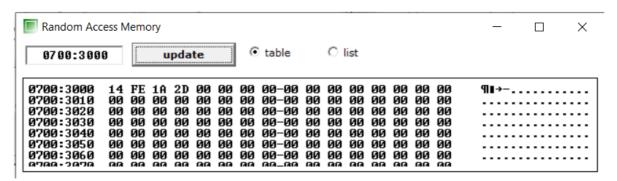
The same can be referred by seeing the FLAG Status. We will observe that the Carry Flag bit is set to 1. Concluding the 1<sup>st</sup> no saved in AX is smaller than 2<sup>nd</sup> number saved in BX, and thus an borrow has occurred.



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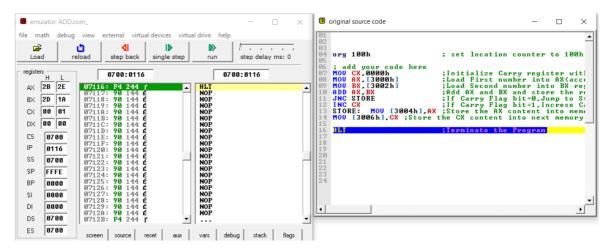
Name: Aditya Shah

I have taken a = FE14 and b = 2D1A, when (a+b) > 16bit,



#### Inference: -

➤ The result shows addition of two number ,being stored in AX with carry,and for carry we have to use another register to store higher bit results.



- Result stored in AX=2B2E and carry is saved in CX.
- The same can be inferred from carry flag, as CF bit is set to 1, which infer that there is overflow.



Name: Aditya Shah

Roll: 2011mt02

- 3.) What are instruction sets?
- •You can make perform a instruction like multiplication using two methods like in a single a go or using repeated addition?
- How would the size of instruction set affect the performance and price?
- •Look at https://www.computerhope.com/jargon/i/instset.htm

#### Ans)

An instruction set, is a list of all the commands (instructions), with all their variations, that a processor can execute. It is more or less machine language. The instruction set provides commands to the processor, to tell it what it needs to do.

#### Instructions include:

- Arithmetic such as add and subtract, etc.
- Logic instructions such as and, or, and not, etc.
- Data instructions such as move, input, output, load, and store, etc.
- Control flow instructions such as goto, if ... goto, call, and others.
- 3.2)Yes, It has been shown in exp-3 and also single go multiplication instruction is also used in factorial experiment.
- 3.3) The Size of instruction set greatly affect the performance and price: -
  - Like in our experiment, we have used 8086, which is enhanced version of 8085 microprocessor.
  - There are many difference between the two, starting from architectural difference's.
  - From Instruction set view point, there is also a difference.
    - o No multiplication and division instruction in 8085 Microprocessor.
    - Multiplication and Division operations are present
  - Also 8086 supports Instruction Queue, Pipelining, Memory Segmentation and many more advantages, which we don't have in 8085.
  - Thus, with enhanced Instruction set, it ease various task and performance is also better, but this comes at higher cost, which we see in the case of 8085 and 8086 too. (The cost of 8085 is low whereas that of 8086 is high)
- 4.) Can we add 32bit number in a 16bit processor, how, example ? Ans)

Yes, we can add very easily two 32bit number in a 16bit processor.

The following algorithm shows ,how it can be done: -

#### Algorithm to Add Two 32 Bit Numbers: -

- Initialize the data segment.
- Load the LSB of first number into AX register.
- Load the MSB of first number into BX register.
- Load the LSB of the second number into CX register.
- Load the MSB of the second number into DX register.
- Add the LSBs of two number.

- Add the MSBs of two numbers along with carry.
- Display the result.
- Stop.
- 5.) PAE is useful when we have more than one application running, and consuming memory. Like if we have 16GB of RAM on a 32bit machine, without PAE, we'd be able to use only 4GB for all applications, and the OS itself. But with PAE, we could use all the 16GB of RAM, or more.

Physical Address Extension (PAE) is a processor feature that enables x86 processors to access more than 4 GB of physical memory on capable versions of Windows. Certain 32-bit versions of Windows Server running on x86-based systems can use PAE to access up to 64 GB or 128 GB of physical memory, depending on the physical address size of the processor.

With PAE, the operating system moves from **two-level linear address translation to three-level address translation**.

Instead of a linear address, it is being split into three separate fields for indexing into memory tables, it is split into four separate fields: a 2-bit bitfield, two 9-bit bitfields, and a 12-bit bitfield that corresponds to the page size implemented by Intel architecture (4 KB). The size of page table entries (PTEs) and page directory entries (PDEs) in PAE mode is increased from 32 to 64 bits. The additional bits allow an operating system PTE or PDE to reference physical memory above 4 GB.

Thus, x86 processor(32bit processor) hardware-architecture being augmented with additional address lines, which are used to select the additional memory, so physical address size increases from 32 bits to higher bits. This, theoretically, increases maximum physical memory size above 4 GB.

- 6.) Following two questions gives a hint how would you make a compiler for a high level language like "C" from assembly language.
- •Can you make a custom "For" loop and "while" loop show any example if possible
- Can you make a function using assembly language from using jump statements and other operations.

Ans.) 6.1)

Yes we can make a custom "For" loop and "While" loop: -

```
org 100h ; set location counter to 100h

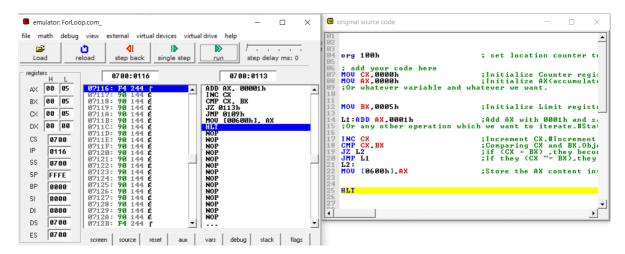
; add your code here
HOU CX,0000h ;Initialize Counter register CX with 0000h. #Initialization
HOU AX,0000h ;Initialize AX(accumulator) register CX with 0000h
;Or whatever variable and whatever we want.

HOU BX,0005h ;Initialize Limit register BX with 0005h.how many times to run.#Condition

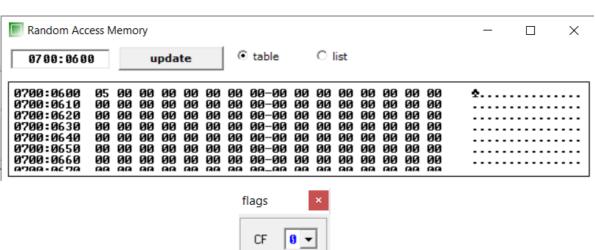
L1:
ADD AX,0001h ;Add AX with 0001h and save the result to AX.
;Or any other operation which we want to iterate.#Statement

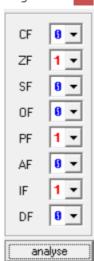
INC CX ;Increment CX.#Increment
CX #Increment CX.#Increment
CX #Increment
CX #
```

#### Result



#### Verification





Thus, we can infer from above, that when the condition is met,i.e., when ZF bit is set to high, the loops ends. Through the above program, we added 1 to AX iteratively 5 times and then when condition is met, the program counter went out of loop and then consequent instruction are executed.

# While Loop Syntaxwhile(condition) {statement(s);}Program

```
org 100h

org 100h

; add your code here

MOU CX,0000h

MOU CX,0000h

; Initialize Counter register CX with 0000h. #Initialization
; Initialize AX(accumulator) register CX with 0000h

; or whatever variable and whatever we want.

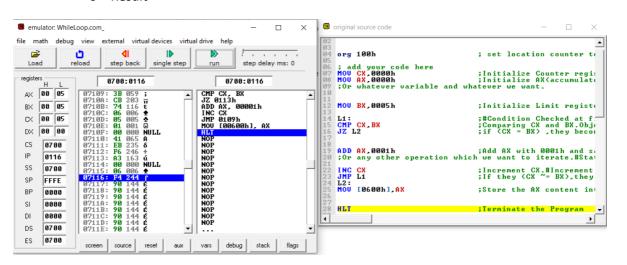
MOU BX,0005h

; Initialize Limit register BX with 0005h.how many times to run.#Condition

11.

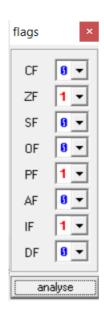
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```

#### o Result



#### Verification





Thus, we can infer from above, that the condition is checked at entry of loop and when the condition is met,i.e., when ZF bit is set to high, the loops ends. Through the above program, we added 1 to AX iteratively 5 times and then when condition is met, the program counter went out of loop and then consequent instruction are executed.

6.2) Yes we can make a function in assembly language, as shown below.

Program(where we call MainFunction.asm) : -

```
org 100h

org 100h

; set location counter to 100h

add your code here

MOU CX,0000h

MOU AX,0000h

; Initialize Counter register CX with 0000h. #Initialization
; Initialize AX(accumulator) register CX with 0000h
; or whatever variable and whatever we want.

MOU BX,0005h
include MainFunction.asm

; Initialize Limit register BX with 0005h.how many times to run.#Condition
; Compiler automatically searches for MainFunction.asm

; Terminate the Program
```

Program(MainFunction.asm): -

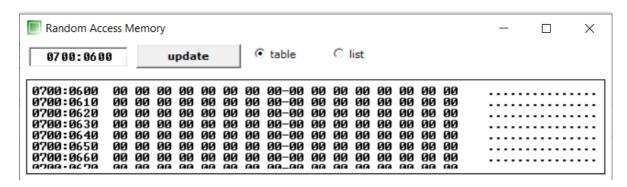
```
org 100h

; set location counter to 100h

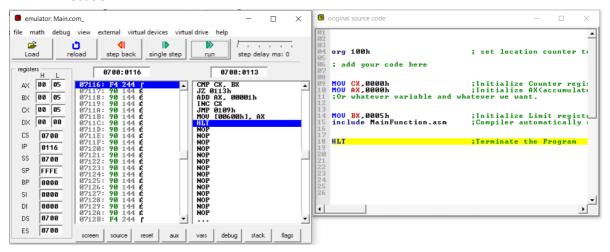
the add your code here

the comparing CX and BX.Objective to find when they become equal form of CX and CX
```

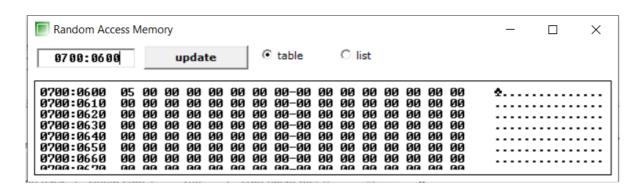
Before running Main.asm, content of Memory location 0600h

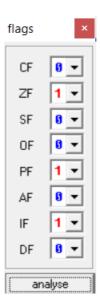


Execution



> Result





#### ❖ Inference: -

We run the Main.asm file, where we just initialized the three registers, namely AX, BX, CX and after initialization, we called the MainFunction.asm file, where we wrote the FOR-loop operation, using jump statements and other operations and thus we see the changes in memory location 0600h and other registers and the no of times the loop has been run is inferred from CX content. After executing this file and it then returned back to Main.asm file, where it halted the processor.

- ✓ To view the program, click on below link:
  - o **Program Files**