Space and Congruence Compression of Proofs

Andreas Fellner





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Congruence

Knowledge

- **1** f(a) = a
- a = b
- **3** b = f(b)
- $f(a) \neq f(b)$

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Proof

Equality is transitive, therefore from f(a) = a, a = b and b = f(b) follows f(a) = f(b), which contradicts $f(a) \neq f(b)$

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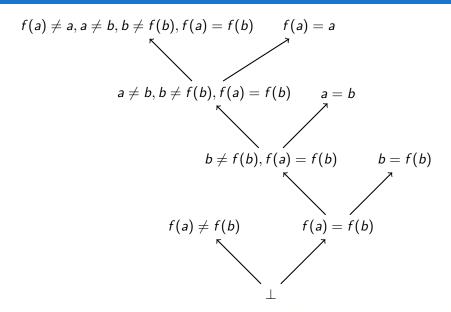
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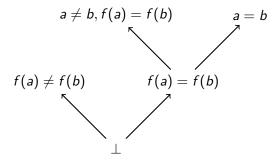
A different Proof

f(.) is a function, therefore from a = b follows f(a) = f(b), which contradicts $f(a) \neq f(b)$

A proof



A different proof



Ground Terms

- Constants a, b, c, \dots
- Compound Terms $f(t_1, \ldots, t_n)$

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Congruence Relation R

- Reflexive: $\forall t(t, t) \in R$
- Symmetric: $(s, t) \in R \Rightarrow (t, s) \in R$
- Transitive: $(t_1, t_2) \in R \dots (t_{m-1}, t_m) \in R \Rightarrow (t_1, t_m) \in R$
- Compatible: $\forall i(t_i, s_i) \in R \Rightarrow (f(t_1, \dots, t_n), f(s_1, \dots, s_n)) \in R$

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Congruence Closure E^* of set of equations E

- Smallest Congruence Relation containing E
- Computable in $O(n \log(n))$
- E is explanation for $(s, t) \in E^*$

Knowledge

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Explanation for f(a) = f(b)

$$\{ f(a) = a, a = b, b = f(b) \}$$

Knowledge

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- $\mathbf{a} = \mathbf{b}$
- **3** b = f(b)
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Explanation for f(a) = f(b)

$$\{ a = b, b = f(b) \}$$

Knowledge

- f(a) = a
- $\mathbf{a} = \mathbf{b}$
- **3** b = f(b)
- **4** $f(a) \neq f(b)$

Explanation for f(a) = f(b)

$$a = b$$

Knowledge

- f(a) = a
- $\mathbf{a} = \mathbf{b}$
- **3** b = f(b)
- $f(a) \neq f(b)$

Explanation for
$$f(a) = f(b)$$

{ $a = b$ }

Short explanation → short (sub)proof

Short Explanation Decision Problem

Given a set of input equations E, a target equation s = t and $k \in \mathbb{N}$, does there exist an explanation $E' \subseteq E$ of s = t with $|E'| \le k$?

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NP-complete

NP-completeness proof sketch

From a propositional logic formula Φ obtain ...

- a set of equations E_{Φ}
- a target equation $s_{\Phi} = t_{\Phi}$
- $k_{\Phi} \in \mathbb{N}$

such that ...

 Φ is satisfiable if and only if there is an explanation $E'\subseteq E_{\Phi}$ of $s_{\Phi}=t_{\Phi}$ with $|E'|\leq k_{\Phi}$

Formula

$$(x_1 \lor x_2 \lor \neg x_3) \land (\neg x_2 \lor x_3) \land (\neg x_1 \lor \neg x_2)$$

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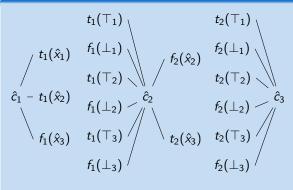
$$t_{1}(\top_{1})$$
 $t_{1}(\hat{x}_{1})$
 $f_{1}(\perp_{1})$
 $t_{1}(\top_{2})$
 $\hat{c}_{1} - t_{1}(\hat{x}_{2})$
 $f_{1}(\perp_{2})$
 $f_{1}(\hat{x}_{3})$
 $f_{1}(\perp_{3})$

Formula

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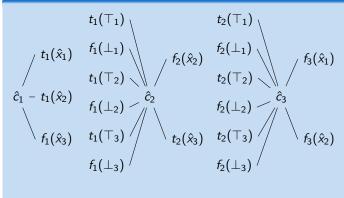
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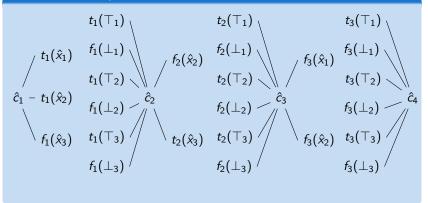
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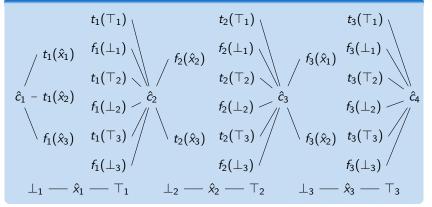
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Formula

$$(x_1 \lor x_2 \lor \neg x_3) \land (\neg x_2 \lor x_3) \land (\neg x_1 \lor \neg x_2)$$

Small subset corresponding to satisfying assignment

$$\hat{c}_1 - t_1(\hat{x}_1) \quad t_1(\top_1) - \hat{c}_2 - f_2(\hat{x}_2) \quad f_2(\bot_2) - \hat{c}_3 - f_3(\hat{x}_2) \quad f_3(\bot_2) - \hat{c}_4$$

$$\hat{x}_1 - \!\!\!\!- \top_1$$

$$\perp_2 - - \hat{x}_2$$

$$\hat{x}_1 - T_1$$
 $\perp_2 - \hat{x}_2$ $\hat{x}_3 - T_3$

Space

Is it possible to remove parts of the proof from memory during proof processing?

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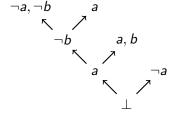
Which parts?

Is it possible to remove parts of the proof from memory during proof processing?

Space measure of a proof

Maximal amount of nodes that have to be kept in memory at once while processing the proof

Space Measure Example



Space Measure Example

Maximum number of nodes in memory: 0

- Not in memory
- In memory

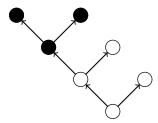
Space Measure Example

Maximum number of nodes in memory: 1

- Not in memory
- In memory

- Not in memory
- In memory

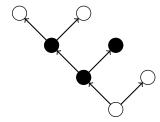
- Not in memory
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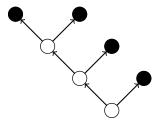
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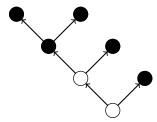
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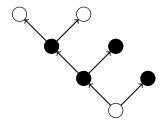


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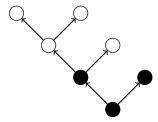
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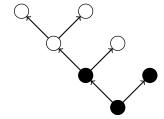
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Maximum number of nodes in memory: 5

- Not in memory
- In memory



Good Traversal Orders are key!

Construct Traversal Orders

Construct Optimal Order

- NP-complete
- Optimal strategy in some pebbling game

Construct Good Order

- Top-Down
- Bottom-Up
- Heuristics



