

# **02157 Functional Programming**

Collections: Sets and Maps

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### Overview



Sets and Maps as abstract data types

- Useful in the modelling and solution of many problems
- Many similarities with the list library

Recommendation: Use these libraries whenever it is appropriate.

### The set concept (1)



A set (in mathematics) is a collection of element like

$$\{Bob,Bill,Ben\},\{1,3,5,7,9\},\mathbb{N},$$
 and  $\mathbb{R}$ 

- the sequence in which elements are enumerated is of no concern, and
- repetitions among members of a set is of no concern either

$$Alice \not\in \{Bob, Bill, Ben\} \qquad \text{and} \qquad 7 \in \{1, 3, 5, 7, 9\}$$

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 and  $7 \in \{1, 3, 5, 7, 9\}$ 

The empty set containing no element is written  $\{\}$  or  $\emptyset$ .

### The sets concept (2)



A set A is a *subset* of a set B, written  $A \subseteq B$ , if all the elements of A are also elements of B, for example

$$\{Ben, Bob\} \subseteq \{Bob, Bill, Ben\}$$
 and  $\{1, 3, 5, 7, 9\} \subseteq \mathbb{N}$ 

Two sets A and B are equal, if they are both subsets of each other:

$$A = B$$
 if and only if  $A \subseteq B$  and  $B \subseteq A$ 

i.e. two sets are equal if they contain exactly the same elements.

The subset of a set *A* which consists of those elements satisfying a predicate *p* can be expressed using a *set-comprehension*:

$$\{x \in A \mid p(x)\}$$

For example

$$\{1,3,5,7,9\} = \{x \in \mathbb{N} \mid \text{odd}(x) \text{ and } x < 11\}$$

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### The set concept (3)



#### Some standard operations on sets:

```
A \cup B = \{x \mid x \in A \text{ or } x \in B\} union

A \cap B = \{x \mid x \in A \text{ and } x \in B\} intersection

A \setminus B = \{x \in A \mid x \notin B\} difference
```

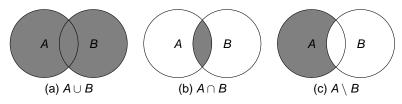


Figure: Venn diagrams for (a) union, (b) intersection and (c) difference

#### For example

```
\{Bob, Bill, Ben\} \cup \{Alice, Bill, Ann\} = \{Alice, Ann, Bob, Bill, Ben\} 
\{Bob, Bill, Ben\} \cap \{Alice, Bill, Ann\} = \{Bob, Ben\} 
\{Bob, Bill, Ben\} \setminus \{Alice, Bill, Ann\} = \{Bob, Ben\}
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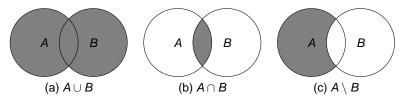


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# Abstract Data Types



# An abstract Data Type: A type together with a collection of operations, where

the representation of values is hidden.

An abstract data type for sets must have:

- Operations to generate sets from the elements. Why?
- Operations to extract the elements of a set. Why?
- Standard operations on sets.

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#### Sets in F#



The Set library of F# supports finite sets. An efficient implementation is based on a balanced binary tree.

### Examples:

```
set ["Bob"; "Bill"; "Ben"];;
val it : Set<string> = set ["Ben"; "Bill"; "Bob"]
set [3; 1; 9; 5; 7; 9; 1];;
val it : Set<int> = set [1; 3; 5; 7; 9]
```

Equality of two sets is tested in the usual manner:

```
set["Bob";"Bill";"Ben"] = set["Bill";"Ben";"Bill";"Bob"];;
val it : bool = true
```

Sets are order on the basis of a lexicographical ordering:

```
compare (set ["Ann";"Jane"]) (set ["Bill";"Ben";"Bob"]);;
val it : int = -1
```

### Selected operations (1)



- ofList: 'a list -> Set<'a>. where ofList  $[a_0; ...; a_{n-1}] = \{a_0; ...; a_{n-1}\}$
- toList: Set<'a> -> 'a list, where toList  $\{a_0, ..., a_{n-1}\} = [a_0; ...; a_{n-1}]$
- add: 'a -> Set<'a> -> Set<'a>. where add  $a A = \{a\} \cup A$
- remove: 'a -> Set<'a> -> Set<'a>. where remove  $aA = A \setminus \{a\}$
- contains: 'a -> Set<'a> -> bool, where contains  $a A = a \in A$
- minElement: Set<'a> -> 'a) where minElement  $\{a_0, a_1, \dots, a_{n-2}, a_{n-1}\} = a_0$  when n > 0

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### Notice that minElement is well-defined due to the ordering:

```
Set.minElement (Set.ofList ["Bob"; "Bill"; "Ben"]);;
val it : string = "Ben"
```

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### Selected operations (2)



- union: Set<'a> -> Set<'a> -> Set<'a>, where union  $AB = A \cup B$
- intersect: Set<'a> -> Set<'a> -> Set<'a>,
   where intersect A B = A ∩ B
- difference: Set<'a> -> Set<'a> -> Set<'a>,
   where difference A B = A \ B
- exists: ('a -> bool) -> Set<'a> -> bool, where exists  $pA = \exists x \in A.p(x)$
- forall: ('a -> bool) -> Set<'a> -> bool, where forall  $p A = \forall x \in A.p(x)$
- fold: ('a -> 'b -> 'a) -> 'a -> Set<'b> -> 'a,
   where

fold 
$$f$$
 a  $\{b_0, b_1, \dots, b_{n-2}, b_{n-1}\}$   
=  $f(f(f(\dots f(f(a, b_0), b_1), \dots), b_{n-2}), b_{n-1})$ 

These work similar to their List siblings, e.g.

Set.fold (-) 0 (set 
$$[1; 2; 3]$$
) =  $((0-1)-2)-3=-6$   
where the ordering is exploited.

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### Selected operations (2)



- union: Set<'a> -> Set<'a> -> Set<'a>. where union  $AB = A \cup B$
- intersect: Set<'a> -> Set<'a> -> Set<'a>. where intersect  $AB = A \cap B$
- difference: Set<'a> -> Set<'a> -> Set<'a>. where difference  $AB = A \setminus B$
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### Example: Map Coloring (1)



### Maps and colors are modelled in a more natural way using sets:

```
type country = string;;
type map = Set<country*country>;;
type color = Set<country>;;
type coloring = Set<color>;;
```

WHY?

```
Two countries c_1, c_2 are neighbors in a map m, if either (c_1, c_2) \in m or (c_2, c_1) \in m: let areNb c1 c2 m = Set.contains (c1,c2) m || Set.contains (c2,c1) m;
```

Color col and be extended by a country c given map m, if for every country c' in col: c and c' are not neighbours in m

```
let canBeExtBy m col c =
   Set.forall (fun c' -> not (areNb c' c m)) col;;
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### Example: Map Coloring (2)



#### The function

```
extColoring: map -> coloring -> country -> coloring
is declared as a recursive function over the coloring:
```

#### WHY not use a fold function?

```
let rec extColoring m cols c =
   if Set.isEmpty cols
   then Set.singleton (Set.singleton c)
   else let col = Set.minElement cols
        let cols' = Set.remove col cols
        if canBeExtBy m col c
        then Set.add (Set.add c col) cols'
        else Set.add col (extColoring m cols' c);;
```

#### Notice similarity to a list recursion:

- base case [] corresponds to the empty set
- for a recursive case x::xs, the head x corresponds to the minimal element col and the tail xs corresponds to the "rests" set cols'

The list-based version is more efficient (why?) and more readable

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### Example: Map Coloring (3)



A set of countries is obtained from a map by the function:

```
countries: map -> Set<country>
```

that is based on repeated insertion of the countries into a set:

```
let countries m =
   Set.fold
      (fun set (c1,c2) -> Set.add c1 (Set.add c2 set))
      Set.empty
      m;;
```

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    Set.empty
    m;;
```

The function

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```
colCntrs: map -> Set<country> -> coloring
```

is based on repeated insertion of countries in colorings using the extColoring function:

```
let colCntrs m cs = Set.fold (extColoring m) Set.empty cs;;
```

# Example: Map Coloring (4)



The function that creates a coloring from a map is declared using functional composition:

### The map concept



A map from a set A to a set B is a finite subset A' of A together with a function m defined on A':  $m: A' \rightarrow B$ .

The set A' is called the *domain* of m: dom m = A'.

A map *m* can be described in a tabular form:

$\boldsymbol{a}_0$	$b_0$
<b>a</b> <sub>1</sub>	<i>b</i> <sub>1</sub>
	h
$a_{n-1}$	$b_{n-1}$

- An element a<sub>i</sub> in the set A' is called a key
- A pair (a<sub>i</sub>, b<sub>i</sub>) is called an entry, and
- b<sub>i</sub> is called the *value* for the key a<sub>i</sub>.

We denote the sets of entries of a map as follows:

entriesOf(
$$m$$
) = {( $a_0, b_0$ ), . . . , ( $a_{n-1}, b_{n-1}$ )}

### Selected map operations in F#



- ofList: ('a\*'b) list -> Map<'a,'b> ofList  $[(a_0,b_0);...;(a_{n-1},b_{n-1})]=m$
- add: 'a -> 'b -> Map<'a,'b> -> Map<'a,'b> add a b m = m', where m' is obtained m by overriding m with the entry (a,b)
- find: 'a → Map<'a,'b> → 'b
  find a m = m(a), if a ∈ dom m;
  otherwise an exception is raised
- tryFind: 'a -> Map<'a,'b> -> 'b option tryFind a m = Some (m(a)), if  $a \in \text{dom } m$ ; None otherwise

foldBack: ('a->'b->'c->'c) -> Map<'a,'b> -> 'c -> 'c foldBack  $f \ m \ c = f \ a_0 \ b_0 \ (f \ a_1 \ b_1 \ (f \ \ldots \ (f \ a_{n-1} \ b_{n-1} \ c) \cdots))$ 

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### A few examples



An entry can be added to a map using add and the value for a key in a map is retrieved using either find or tryFind:

# An example using Map.foldBack



We can extract the list of article codes and prices for a given register using the fold functions for maps:

This and other higher-order functions are similar to their List and Set siblings.

### Example: Cash register (1)



```
type articleCode = string;;
type articleName = string;;
type noPieces = int;;
type price = int;;

type info = noPieces * articleName * price;;
type infoseq = info list;;
type bill = infoseq * price;;
```

The natural model of a register is using a map

```
type register = Map<articleCode, articleName*price>;;
```

since an article code is a arrigae racritinoation of a

#### First version

```
type item = noPieces * articleCode;;
type purchase = item list;;
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### The natural model of a register is using a map:

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type register = Map<articleCode, articleName*price>;;
```

since an article code is a unique identification of an article.

#### First version:

```
type item = noPieces * articleCode;;
type purchase = item list;;
```

### Example: Cash register (1) - a recursive program



```
exception FindArticle;;
(* makebill: register -> purchase -> bill *)
let rec makeBill reg = function
         -> ([],0)
     (np,ac)::pur ->
       match Map.tryFind ac reg with
        None
                       -> raise FindArticle
        Some(aname,aprice) ->
           let tprice = np*aprice
           let (infos, sumbill) = makeBill reg pur
           ((np,aname,tprice)::infos, tprice+sumbill);;
let pur = [(3, a2); (1, a1)];
makeBill reg1 pur;;
val it : (int * string * int) list * int =
  ([(3, "herring", 12); (1, "cheese", 25)], 37)
```

• the lookup in the register is managed by a Map.tryFind

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the lookup in the register is managed by a Map.tryFind

# Example: Cash register (2) - using List.foldBack



```
let makeBill' reg pur =
   let f (np,ac) (infos,billprice)
          = let (aname, aprice) = Map.find ac req
            let tprice = np*aprice
            ((np,aname,tprice)::infos, tprice+billprice)
   List.foldBack f pur ([],0);;
makeBill' reg1 pur;;
val it : (int * string * int) list * int =
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```

- the recursion is handled by List.foldBack
- the exception is handled by Map.find

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# Example: Cash register (2) - using List.foldBack



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# Example: Cash register (2) - using maps for purchases



The purchase: 3 herrings, one piece of cheese, and 2 herrings, is the same as a purchase of one piece of cheese and 5 herrings.

A purchase associated number of pieces with article codes:

```
type purchase = Map<articleCode,noPieces>;;
```

A bill is produced by folding a function over a map-purchase:

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### Summary



- The concepts of sets and maps.
- Fundamental operations on sets and maps.
- Applications of sets and maps.