

# Category Theory

Projektarbejde i Datalogi 10ECTS (E24.520202U002.A)

AARHUS UNIVERSITET 

Abdul Haliq Abdul Latiff [202303466]

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This report is written for the course on Category Theory undertaken as a project module:  
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## §1 Version2

This section onwards departs from the previous due to new understanding of proof writing in category theory, particularly with UMPs.

EXERCISE 3, CHAPTER 4: regular mono stable under pullback

Given  $h', e'$  is a pullback of  $h, e$

$$\text{Pullback}_{h,e}(e', h') \Leftrightarrow he' = eh' \wedge \quad (1.1)$$

$$\forall f, g. hf = eg \Rightarrow \quad (1.2)$$

$$\exists! u. h'u = g \wedge e'u = f \quad (1.3)$$

Given  $e$  is a monomorphism

$$\text{isMono}(e) \Leftrightarrow \forall g, h. eg = eh \Rightarrow \quad (2.1)$$

$$g = h \quad (2.2)$$

We shown  $e'$  is a monomorphism

$$\text{isMono}(e') \Leftrightarrow \forall u_1, u_2. e'u_1 = e'u_2 \Rightarrow \quad (3.1)$$

$$u_1 = u_2 \quad (3.2)$$

By Equation 1.1 and Equation 3.1 above, we have

$$(eh')u_1 = h(e'u_1) = h(e'u_2) = eh'u_2$$

This satisfies Equation 2.1 by  $g = h'u_1, h = h'u_2$  thus quotienting them

$$h'u_1 = h'u_2$$

Both  $u_1, u_2$  now satisfies Equation 1.3 and by uniqueness we get the consequent of Equation 3.2

$$u_1 = u_2$$

$\therefore e'$  is a mono.  $\square$

Given  $e$  is an equalizer; we assume for some  $a, a'$

$$\text{Equalizer}_{a,a'}(e) \Leftrightarrow ae = a'e \wedge \quad (4.1)$$

$$\forall z. az = a'z \Rightarrow \quad (4.2)$$

$$\exists! u. eu = z \quad (4.3)$$

We show equalizers are stable under pullback;  $e'$  is an equalizer on  $ah$  and  $a'h$ .

$$\text{Equalizer}_{ah,a'h}(e') \Leftrightarrow ahe' = a'he' \wedge \quad (5.1)$$

$$\forall z. ahz = a'hz \Rightarrow \quad (5.2)$$

$$\exists! u. e'u = z \quad (5.3)$$

Equation 5.1 derived by Equation 1.1 and Equation 4.1

$$a(he') = (ae)h' = a'(eh') = a'he'$$

Supposing Equation 5.2 for  $z = f$

$$\forall f. ahf = a'hf$$

this satisfies Equation 4.3 giving us some unique  $g$  by Equation 4.3

$$\exists! g. eg = hf$$

this satisfies Equation 1.3 giving us some unique  $u_1$  by Equation 1.3

$$\exists! u_1. h'u_1 = g \wedge e'u_1 = f$$

this satisfies 3.3 but stronger, thus not necessarily unique, assuming another morphism  $u_2$  exists

$$\forall u_2. e'u_2 = f$$

both  $u_1, u_2$  satisfies Equation 3.2 thus quotienting them, specifically:

$$e'u_1 = e'u_2 \Rightarrow u_1 = u_2$$

Thus the consequent of Equation 5.3 holds

$$\exists! u. e'u = f$$

$\therefore e'$  is an equalizer on  $ah$  and  $a'h$ . ■

$$\begin{array}{c}
 \frac{\text{Pullback}_{h,e}(e', h')}{1.1} \quad \frac{\text{isMono}(e)}{3.1_{u_1, u_2}} \quad \frac{\text{Pullback}_{h,e}(e', h')}{1.3} \\
 \hline
 \frac{3.1_{u_1, u_2} \Rightarrow 3.2}{\text{isMono}(e')} \\
 \\
 \frac{\text{Pullback}_{h,e}(e', h')}{1.1} \quad \frac{\text{Equalizer}_{a,a'}(e)}{2.1} \quad \frac{\text{Equalizer}_{a,a'}(e)}{4.2_{hf} \Rightarrow 4.3_g} \quad \frac{\text{Pullback}_{h,e}(e', h')}{1.2_{f,g} \Rightarrow 1.3_{u_1}} \quad \frac{\text{isMono}(e')}{u_2} \quad \frac{\text{isMono}(e')}{3.1_{u_1, u_2} \Rightarrow 3.2} \\
 \hline
 \frac{5.1}{\text{Equalizer}_{ah, a'h}(e')} \quad \frac{5.2_f \Rightarrow 5.3_u}{5.2_f \Rightarrow 5.3_u}
 \end{array}$$

EXERCISE DUAL: regular epi stable under pushout

## §2 Appendix

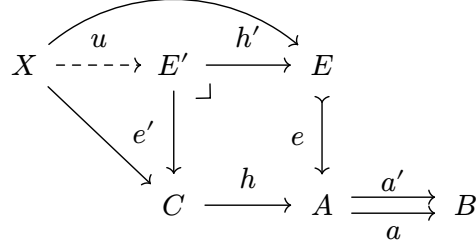
$$\begin{aligned}
\text{isCategory}(\mathcal{C}) = & \quad \forall A, B, C \in \mathbf{Ob}_{\mathcal{C}}, f \in \mathbf{Hom}_{\mathcal{C}}(A, B), g \in \mathbf{Hom}_{\mathcal{C}}(B, C). g \circ f \in \mathbf{Hom}_{\mathcal{C}}(A, C) \\
& \wedge \forall X \in \mathbf{Ob}_{\mathcal{C}}. \exists! 1_X \in \mathbf{Hom}(X, X) \\
& \wedge \forall h, g, f. h \circ (g \circ f) = (h \circ g) \circ f \\
& \wedge \forall A, B \in \mathbf{Ob}_{\mathcal{C}}, f \in \mathbf{Hom}_{\mathcal{C}}(A, B). f \circ 1_A = f = 1_B \circ f \\
\text{isFunctor}(F : C \rightarrow D) = & \quad \forall f \in \mathbf{Hom}_C(A, B). F(f) \in \mathbf{Hom}_D(F(A), F(B)) \\
& \wedge \forall 1_X \in \mathbf{Hom}_C(X, X). F(1_X) = 1_{F(X)} \\
& \wedge \forall g, f \in \mathbf{Hom}_C. F(g \circ f) = F(g) \circ F(f) \\
\text{isIso}(f : A \rightarrow B) = & \quad \exists! g. g \circ f = 1_A \wedge f \circ g = 1_B \\
A \cong B = & \quad \exists (f : A \rightarrow B). \text{isIso}(f) \\
C \times D = & \quad (\mathbf{Ob}_C \times \mathbf{Ob}_D, [(C_1, D_1), (C_2, D_2)] \mapsto \mathbf{Hom}_C(C_1, C_2) \times \mathbf{Hom}_D(D_1, D_2)) \\
C^{\text{op}} = & \quad (\mathbf{Ob}_C, [A, B \mapsto \{f^{\text{op}} : A \rightarrow B \mid f \in \mathbf{Hom}_C(B, A)\}]) \\
C^{\rightarrow} = & \quad (\mathbf{Hom}_C, [f, f' \mapsto \{(g, g') \mid g' \circ f = f' \circ g\}]) \\
C/A = & \quad (\{f \mid f \in \mathbf{Hom}_C(X, A)\}, [f, f' \mapsto \mathbf{Hom}_C(\text{dom}(f), \text{dom}(f'))]) \\
C \setminus A = & \quad (\{f \mid f \in \mathbf{Hom}_C(A, X)\}, [f, f' \mapsto \mathbf{Hom}_C(\text{cod}(f), \text{cod}(f'))]) \\
\text{isMono}(f) = & \quad \forall g, h. f \circ g = f \circ h \Rightarrow g = h \\
\text{isEpi}(f) = & \quad \forall g, h. g \circ f = h \circ f \Rightarrow g = h \\
\text{isSplitMono}(m, s) = & \quad s \circ m = 1_{\text{dom}(m)} \\
\text{isSplitEpi}(e, s) = & \quad e \circ s = 1_{\text{cod}(e)} \\
\text{areIso}(f, g) = & \quad \text{isSplitEpi}(f, g) \wedge \text{isSplitMono}(f, g) \\
\text{isProjective}(P) = & \quad \forall e : E \twoheadrightarrow X, f : P \rightarrow X. \exists \bar{f} : P \rightarrow E. e \circ \bar{f} = f \\
\text{UMP}_{\text{freemonoid}}(|\bar{f}|) = & \quad \forall i, f. \exists! \bar{f}. |\bar{f}| \circ i = f \\
\text{UMP}_{\text{terminal}}(0_X) = & \quad \forall X. \exists! 0_X \in \mathbf{Hom}(0, X) \\
\text{UMP}_{\text{initial}}(1_X) = & \quad \forall X. \exists! 1_X \in \mathbf{Hom}(X, 1) \\
\text{UMP}_{\text{product}}(p_1, p_2) = & \quad \forall x_1, x_2. \exists! u. x_1 = p_1 \circ u \wedge x_2 = p_2 \circ u \\
\text{UMP}_{\text{coproduct}}(q_1, q_2) = & \quad \forall x_1, x_2. \exists! u. x_1 = u \circ q_1 \wedge x_2 = u \circ q_2 \\
\text{UMP}_{\text{equalizer}}(e, f, g) = & \quad \forall z. (f \circ z = g \circ z) \wedge \exists! u. e \circ u = z \\
\text{UMP}_{\text{coequalizer}}(q, f, g) = & \quad \forall z. (z \circ f = z \circ g) \wedge \exists! u. u \circ q = z \\
\text{UMP}_{\text{pullback}}(p_1, p_2, f, g) = & \quad (f \circ p_1 = g \circ p_2) \\
& \wedge \forall z_1, z_2. \exists! u. z_1 = p_1 \circ u \wedge z_2 = p_2 \circ u \\
\text{UMP}_{\text{limit}}(\{p_i\}, D) = & \quad \forall \{c_j\}. \exists! u. \forall j. p_j \circ u = c_j \\
\text{isExponential}(C^B, \varepsilon) = & \quad \forall A, (f : A \times B \rightarrow C). \exists! (\tilde{f} : A \rightarrow C^B). \varepsilon \circ (\tilde{f} \times 1_B) = \tilde{f} = f \\
\text{isCCC}(\mathbf{Ob}, \mathbf{Hom}) = & \quad \text{isCategory}(\mathbf{Ob}, \mathbf{Hom}) \\
& \wedge \forall A, B. \exists! A \times B. \text{UMP}_{\text{product}}(p_1 : A \times B \rightarrow A, p_2 : A \times B \rightarrow B) \\
& \wedge \forall B, C. \text{isExponential}(C^B, \varepsilon)
\end{aligned}$$

TODO LIST:

list of functors:  $| - |, (-)^A, (-)^{\text{op}}, \mathbf{Hom}(A, -), \mathbf{Hom}(-, A)$

todo in chapter4 notes

- $(f)^A = \widetilde{f \circ \varepsilon}$
- define Sub and Cone category



**Lemma 1:**  $e'$  is mono

$$\frac{\frac{\forall f g. e' f = e' g \quad \frac{\text{---} \forall X \rightarrow E}{u_1}}{e' u_1 = e' u_2} \quad \frac{\text{---} \forall X \rightarrow E}{u_2} \text{mono-antecedent}}{e' u_1 = e' u_2}$$

might be an error here, need to show  $h'$  is a mono, then  $\text{elim Mono}(eh')$ , quantifiers can't just be substituted by arbitrary composites

$$\frac{\frac{\text{Mono}(e) \quad h' u_1 \quad h' u_2 \text{mono-elim}}{eh' u_1 = eh' u_2 \Rightarrow h' u_1 = h' u_2} \quad \frac{\text{Pullback}(e', h')}{he' = eh'} \text{pullback-elim}^{\text{diagram}}}{he' u_1 = he' u_2 \Rightarrow h' u_1 = h' u_2} =$$

$$\frac{\frac{e' u_1 = e' u_2 \quad he' u_1 = he' u_2 \Rightarrow h' u_1 = h' u_2}{he' u_1 = he' u_1 \Rightarrow h' u_1 = h' u_2} \Rightarrow \text{-elim}}{h' u_1 = h' u_2}$$

$$\frac{\frac{\text{Pullback}(e', h') \quad u_1 \quad u_2 \text{pullback-elim}^{\text{unique}}}{e' u_1 = e' u_2 \wedge h' u_1 = h' u_2 \Rightarrow u_1 = u_2} \quad \frac{e' u_1 = e' u_2 \quad h' u_1 = h' u_2}{u_1 = u_2} \Rightarrow \text{-elim}}{u_1 = u_2}$$

$$\frac{\frac{e' u_1 = e' u_2 \quad u_1 = u_2}{e' u_1 = e' u_2 \Rightarrow u_1 = u_2} \Rightarrow \text{-intro}}{\text{Mono}(e')} \text{mono-intro}$$

**Lemma 2:**  $e'$  is regular / equalizer

$$\frac{\frac{\frac{\text{Equalizer}(e) \quad \frac{\text{---} \forall A \rightarrow B}{a}}{ae = a'e} \quad \frac{\text{---} \forall A \rightarrow B}{a'} \text{equalizer-elim}^{\text{diagram}}}{aeh' = a'eh'} = \frac{\frac{\text{Pullback}(e', h')}{he' = eh'} \text{pullback-elim}^{\text{diagram}}}{ahe' = a'he'}$$

$$\frac{\frac{\text{Mono}(e') \quad u_1 \quad u_2}{e'u_1 = e'u_2 \Rightarrow u_1 = u_2} \text{mono-elim}}{he'u_1 = he'u_2 \Rightarrow u_1 = u_2} =$$

$$\frac{he'u_1 = he'u_2 \Rightarrow u_1 = u_2 \quad ahe' = a'he'}{\text{Equalizer}(he')} \text{equalizer-intro}$$

**Proof:**  $he'$  is a regular mono

$$\frac{\frac{he'u_1 = he'u_2 \Rightarrow u_1 = u_2}{\text{Mono}(he')} \text{mono-intro} \quad \text{Equalizer}(he')}{\text{RegularMono}(he')} \text{regular-mono-intro}$$

■

assumptions

1. composing a morphism on both sides of an equality remains equal
  - i.e.  $f = g, h \vdash hf = hg$  and  $f = g, h \vdash fh = gh$
2. substituting a quantifier argument with another quantifier argument but composed is the same
  - i.e.  $\text{Mono}(e), ab, c \vdash eab = ec \Rightarrow ab = c$  where  $a$  is a quantified argument or fixed morphism and  $b$  otherwise
  - which should be true by definition of existence of composition morphism in a category
  - this might be wrong



## **Bibliography**