Constraints Satisfaction Problems

Dra. Ma Dolores Rodríguez Moreno





Objectives

Specific Objectives

- Review what CSP is
- Main CSP techniques

Source

• Stuart Russell & Peter Norvig (2009). Chapter 6. Artificial Intelligence: A Modern Approach. (3rd Edition). Ed. Pearsons



Outline

- Introduction
- Example
- Techniques
 - Bactracking search in CSPs
 - Heuristics
 - Foward checking
 - Constraint Propagation
 - Local search in CSPs
- Conclusions





Introduction

- A CSP is defined by:
 - I. A set of variables {VI,...,Vn}
 - 2. A domain of values for each variable Dom [Vi]
 - 3. A set of constraints {C1,...,Cm}
- A goal test is a set of constraints specifying allowable combinations of values for subsets of variables
- An assignment that does not violate any restriction is called a consistent assignment
- Some problems require a solution that maximizes an objective function





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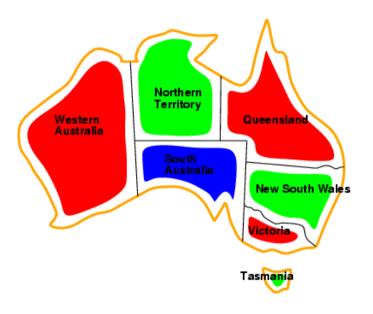


Example (I)

- Variables WA, NT, Q, NSW, V, SA, T
- Domains Di = {red,green,blue}
- Constraints: adjacent regions must have different colors
- e.g., WA ≠ NT, or (WA,NT) in
 {(red,green),(red,blue),(green,red), (green,blue),(blue,red),
 (blue,green)}



Example (II)

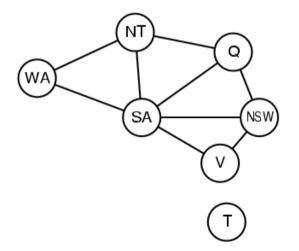


• Solutions are complete and consistent assignments, e.g., WA = red, NT = green, Q = red, NSW = green, V = red, SA = blue, T = green



Example (III)

- Binary CSP: each constraint relates two variables
- Constraint graph: nodes are variables, arcs are constraints



Varieties of CSPs

- There is a great variety of CSPs
- CSP with discrete variables
 - Finite domains
 - *n* variables, domain size $d \rightarrow O(d^n)$ complete assignments
 - eg. Boolean CSPs (var. can be T/F) although it can include satisfiability problemas where NP-complete (exponential time)
 - Infinite domains
 - Integers, strings, etc.
 - e.g., job scheduling, variables are start/end days for each job
 - Need a constraint language, e.g., $StartJob_x + 5 \le StartJob_3$
- CSP with continous variables
 - e.g., start/end times for Hubble Space Telescope observations
 - linear constraints solvable in polynomial time by linear programming





Type of constraints

- Unary constraints involve a single variable,
 - e.g., SA ≠ green
- Binary constraints involve pairs of variables,
 - e.g., $SA \neq WA$
- Higher-order constraints involve 3 or more variables,
 - e.g., cryptarithmetic column constraints

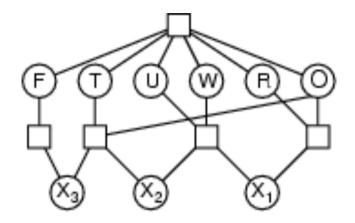


Type of constraints: Higher-order

• Each letter represents a different digit

T W O + T W O F O U R

- Variables: FTUWRO
 - Auxiliar: $X_1 X_2 X_3$
 - Represent o or 1 transfered to the next column
- Domain: {*o,1,2,3,4,5,6,7,8,9*}
- Constraints: Alldiff(F, T, U, W, R, O) or also $F \neq T$, $F \neq U$, ...
 - $O + O = R + IO \cdot X_I$
 - $X_1 + W + W = U + 10 \cdot X_2$
 - $X_2 + T + T = O + 10 \cdot X_3$
 - $X_3 = F$, $T \neq 0$, $F \neq 0$



Real-world CSPs

- Assignment problems
 - e.g., who teaches what class
- Timetabling problems
 - e.g., which class is offered when and where?
- Transportation scheduling
- Factory scheduling
- Notice that many real-world problems involve real-valued variables



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Introduction (I)

- States are defined by the values assigned so far
- Every solution appears at depth n with n variables
 use depth-first search
- The states are defined by the assigned values (incremental formulation):
 - Initial state: the empty assignment {}
 - Successor function: assign a value to an unassigned variable that does not conflict with current assignment → fail if there is no legal assignments
 - Goal test: the current assignment is complete
 - Cost path: a constant cost (eg 1) for each step
- Path is irrelevant, so we can also use **complete-state formulation** (each state is a complete assignment that may or may not meet the restrictions)





Introduction (II)

- Variable assignments are commutative, i.e.,
 [WA = red then NT = green] same as [NT = green then WA = red]
- Only need to consider assignments to a single variable at each node
- Depth-first search for CSPs with single-variable assignments is called backtracking search
- Backtracking search is the basic uninformed algorithm for CSPs
- Can solve *n*-queens for $n \approx 25$

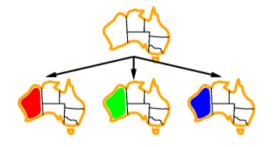






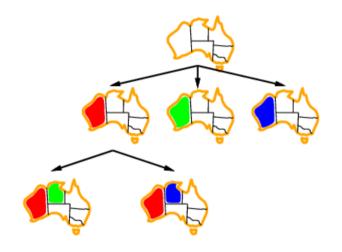






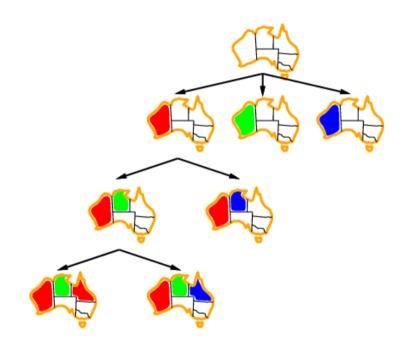
















Improve Backtracking search

- General-purpose methods (heuristics) can give huge gains in speed:
 - Which variable should be assigned next?
 - In what order should its values be tried?
 - Can we detect inevitable failure early?
- Let's see 3 heuristics





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Heuristics: MRV

- Most constrained variable:
 - Choose the variable with the fewest legal values
 - If X has o remaining legal values, selecting X will prevent assignments to other variables that fail when X is chosen



• a.k.a. minimum remaining values (MRV) heuristic



Heuristics (number of constraints)

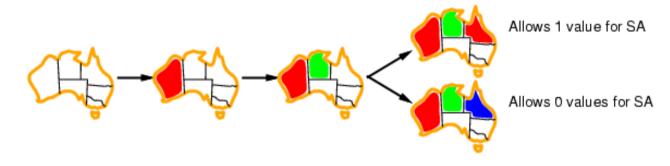
- Tie-breaker among most constrained variables
- Most constraining variable:
 - choose the variable with the most constraints on remaining variables





Heuristics (LCV)

- Given a variable, choose the least constraining value:
 - the one that rules out the fewest values in the remaining variables



• Combining these heuristics makes 1000 queens feasible



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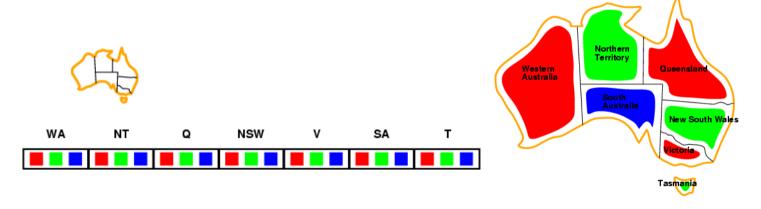


Introduction

- If we look at certain restrictions early or before the search starts, we can drastically reduce the search space
- There are 2 techniques:
 - Forward checking
 - Constraints Propagation

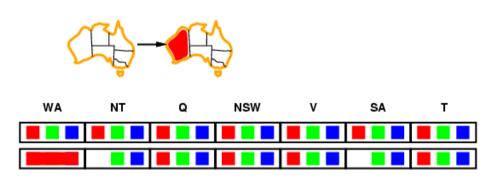


- Idea:
 - Keep track of remaining legal values for unassigned variables
 - Terminate search when any variable has no legal values



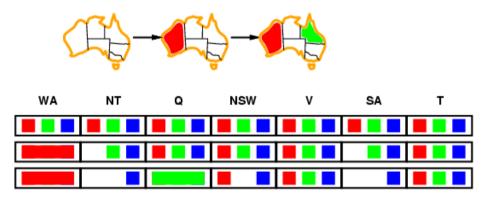


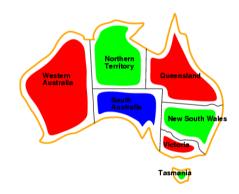
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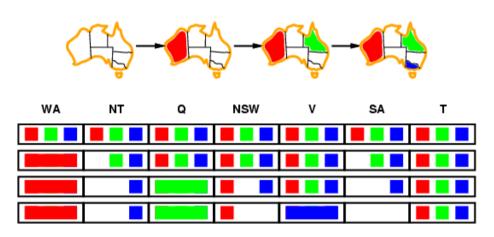
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• Instead of V, which would have been a better candidate according to the heuristic views? SA or NT



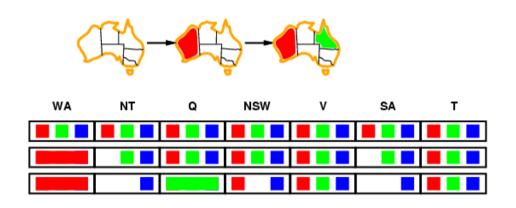
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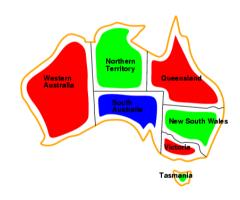
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Constraint propagation

• Forward checking propagates information from assigned to unassigned variables, but doesn't provide early detection for all failures:





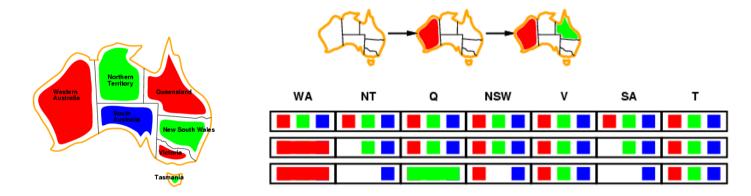
- NT and SA cannot both be blue!
- Constraint propagation repeatedly enforces constraints locally





Constraint propagation

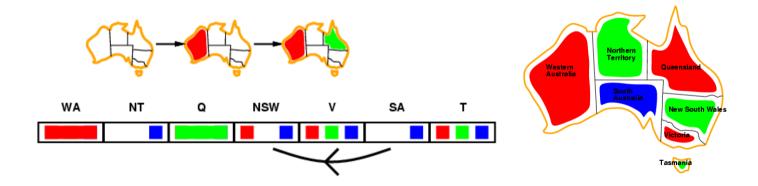
• We need to propagate from WA and Q to NT and SA (as we did in forward checking) and then, in the constraints between NT and SA, to detect the inconsistency



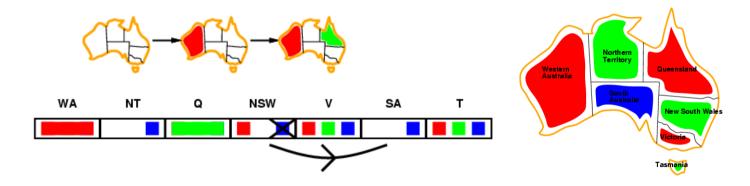
- The way to do it has to be fast: the propagation time cannot be longer than the simple search
- Arc consistency



- Simplest form of propagation makes each arc consistent
- $X \rightarrow Y$ is consistent iff for every value x of X there is some allowed y

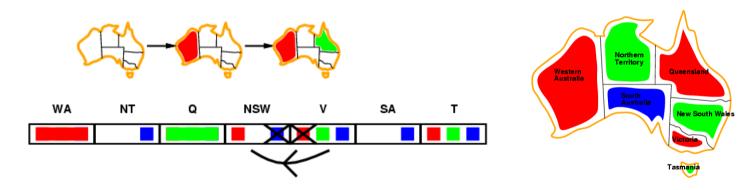


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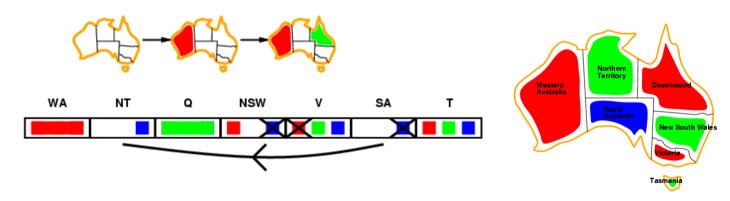
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• If X loses a value, neighbors of X need to be rechecked



- Simplest form of propagation makes each arc consistent
- $X \rightarrow Y$ is consistent iff for every value x of X there is some allowed y



- If Xloses a value, neighbors of Xneed to be rechecked
- Arc consistency detects failure earlier than forward checking
- Can be run as a preprocessor or after each assignment





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Local search in CSPs

- Hill-climbing, simulated annealing typically work with "complete" states, i.e., all variables assigned
- To apply to CSPs:
 - Allow states with unsatisfied constraints
 - Operators reassign variable values
- Variable selection: randomly select any conflicted variable
- Value selection by min-conflicts heuristic:
 - Choose value that violates the fewest constraints
 - i.e., hill-climb with h(n) = total number of violated constraints



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Conclusions

- CSPs are a special kind of problem:
 - States defined by values of a fixed set of variables
 - Goal test defined by constraints on variable values
- Backtracking = depth-first search with one variable assigned per node
- Variable ordering and value selection heuristics help significantly
- Forward checking prevents assignments that guarantee later failure
- Constraint propagation (e.g., arc consistency) does additional work to constrain values and detect inconsistencies



