

AI505  
Optimization

## Constrained Optimization

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# Constrained Optimization

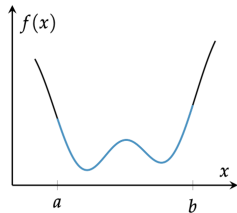
- Minimizing an objective subject to design point restrictions called **constraints**
- A variety of techniques transform constrained optimization problems into unconstrained problems
- New optimization problem statement

$$\begin{array}{ll}\text{minimize} & f(\mathbf{x}) \\ \text{subject to} & \mathbf{x} \in \mathcal{X}\end{array}$$

- The set  $\mathcal{X} \subset \mathbb{R}$  is called the **feasible set**

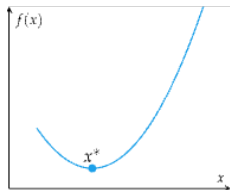
# Constrained Optimization

Constraints that bound feasible set can change the optimizer

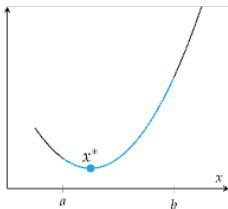


$$\begin{aligned} &\underset{x}{\text{minimize}} && f(x) \\ &\text{subject to} && x \in [a, b] \end{aligned}$$

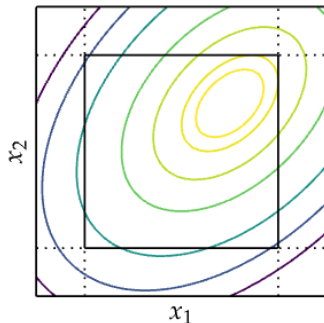
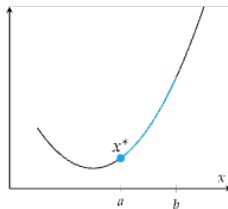
Unconstrained



Constrained, Same Solution



Constrained, New Solution



# Constraint Types

- Generally, constraints are formulated using two types:
  1. **Equality constraints:**  $h(\mathbf{x}) = 0$
  2. **Inequality constraints:**  $g(\mathbf{x}) \leq 0$
- Any optimization problem can be written as

$$\begin{aligned} & \underset{\mathbf{x}}{\text{minimize}} \quad f(\mathbf{x}) \\ & \text{subject to} \quad g_i(\mathbf{x}) \leq 0 \text{ for all } i \text{ in } \{1, \dots, m\} \\ & \quad \quad \quad h_j(\mathbf{x}) = 0 \text{ for all } j \text{ in } \{1, \dots, \ell\} \end{aligned}$$

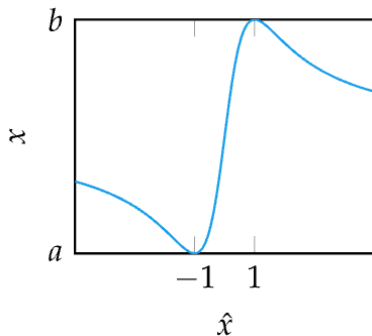
$$\begin{aligned} & \underset{\mathbf{x}}{\text{minimize}} \quad f(\mathbf{x}) \\ & \text{subject to} \quad \mathbf{g}(\mathbf{x}) \leq 0 \\ & \quad \quad \quad \mathbf{h}(\mathbf{x}) = 0 \end{aligned}$$

$f$  and the functions  $h$  and  $g$  are all smooth, real-valued functions on a subset of  $\mathbb{R}^n$

# Transformations to Remove Constraints

- If necessary, some problems can be reformulated to incorporate constraints into the objective function
- If  $x$  is constrained between  $a$  and  $b$

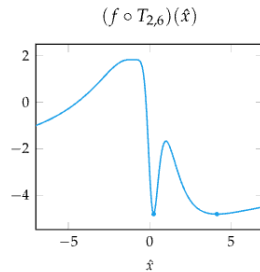
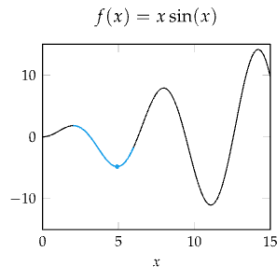
$$x = t_{a,b}(\hat{x}) = \frac{b+a}{2} + \frac{b-a}{2} \left( \frac{2\hat{x}}{1+\hat{x}^2} \right)$$



# Transformations to Remove Constraints

## Example

$$\begin{array}{ll}\underset{x}{\text{minimize}} & x \sin(x) \\ \text{subject to} & 2 \leq x \leq 6\end{array}$$



$$\underset{\hat{x}}{\text{minimize}} \quad t_{2,6}(\hat{x}) \sin(t_{2,6}(\hat{x}))$$

$$\underset{\hat{x}}{\text{minimize}} \quad \left( 4 + 2 \left( \frac{2\hat{x}}{1 + \hat{x}^2} \right) \right) \sin \left( 4 + 2 \left( \frac{2\hat{x}}{1 + \hat{x}^2} \right) \right)$$

# Transformations to Remove Constraints

## Example

$$\begin{aligned} & \underset{\mathbf{x}}{\text{minimize}} \quad f(\mathbf{x}) \\ & \text{subject to} \quad h(\mathbf{x}) = x_1^2 + x_2^2 + \dots + x_n^2 - 1 = 0 \end{aligned}$$

- Solve for one of the variables to eliminate constraint:

$$x_n = \pm \sqrt{1 - x_1^2 - x_2^2 - \dots - x_{n-1}^2}$$

- Transformed, unconstrained optimization problem:

$$\underset{\mathbf{x}}{\text{minimize}} \quad \left( \left[ x_1, x_2, \dots, x_{n-1}, \pm \sqrt{1 - x_1^2 - x_2^2 - \dots - x_{n-1}^2} \right] \right)$$

# Lagrangian Relaxation

- With only equality constraints, critical points (local minima, global minima, or saddle points optimal) where gradient of  $f$  and the gradient of  $h$  are aligned
- The method of Lagrangian relaxation is used to optimize a function subject to (equality) constraints
- Lagrangian multipliers refer to the variables introduced by the method denoted by  $\lambda$

## 1. Form **Lagrangian relaxation**

$$\begin{array}{ll}\underset{\mathbf{x}}{\text{minimize}} & f(\mathbf{x}) \\ \text{subject to} & h(\mathbf{x}) = 0\end{array}$$

$$\mathcal{L}(\mathbf{x}, \lambda) = f(\mathbf{x}) - \lambda h(\mathbf{x})$$

## 2. Set $\nabla_{\mathbf{x}}\mathcal{L}(\mathbf{x}, \lambda) = 0$ and $\nabla_{\lambda}\mathcal{L}(\mathbf{x}, \lambda) = 0$ to get

$$\nabla f(\mathbf{x}) = \lambda \nabla h(\mathbf{x}) \quad h(\mathbf{x}) = 0$$

## 3. solve for $\mathbf{x}$ and $\lambda$

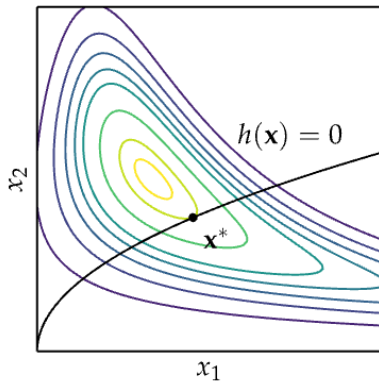


## Example

$$\begin{aligned} & \text{minimize} \quad -\exp\left(-\left(x_1x_2 - \frac{3}{2}\right)^2 - \left(x_2 - \frac{3}{2}\right)^2\right) \\ & \text{subject to} \quad x_1 - x_2^2 = 0 \end{aligned}$$

# Lagrangian Relaxation

Intuitively, the method of Lagrange multipliers finds the point  $\mathbf{x}^*$  where the constraint function is orthogonal to the gradient



# Lagrangian Relaxation with Inequality Constraints

$$\begin{array}{ll}\text{minimize} & f(\mathbf{x}) \\ \text{subject to} & g(\mathbf{x}) \leq 0\end{array}$$

- If solution lies at the constraint boundary, the constraint is called **active**, and the Lagrangian condition holds for a non-negative constant  $\mu$ :

$$\nabla f(\mathbf{x}) + \mu \nabla g(\mathbf{x}) = 0$$

- If the solution lies within the boundary, the constraint is called **inactive**, and the optimal solution simply lies where

$$\nabla f(\mathbf{x}) = 0$$

that is, the Lagrangian condition holds with  $\mu = 0$

# Lagrangian Relaxation with Inequality Constraints

$$\begin{aligned} & \underset{\mathbf{x}}{\text{minimize}} && f(\mathbf{x}) \\ & \text{subject to} && g(\mathbf{x}) \leq 0 \end{aligned}$$

- We create the Lagrangian relaxation such that it goes to  $\infty$  outside the feasibility set ( $g(\mathbf{x}) \not\leq 0$ ):

$$\mathcal{L}_{\infty}(\mathbf{x}) = f(\mathbf{x}) + \infty(g(\mathbf{x}) > 0)$$

impractical: discontinuous and nondifferentiable.

- Instead, for  $\mu > 0$ :

$$\mathcal{L}(\mathbf{x}, \mu \geq 0) = f(\mathbf{x}) + \mu g(\mathbf{x})$$

$$\mathcal{L}_{\infty}(\mathbf{x}) = \underset{\mu \geq 0}{\text{maximize}} \mathcal{L}(\mathbf{x}, \mu)$$

for  $\mathbf{x}$  infeasible,  $\mathcal{L}_{\infty}(\mathbf{x}) = \infty$ ; for  $\mathbf{x}$  feasible,  $\mathcal{L}_{\infty}(\mathbf{x}) = f(\mathbf{x})$

- The new optimization problem becomes

$$\underset{\mathbf{x}}{\text{minimize}} \underset{\mu \geq 0}{\text{maximize}} \mathcal{L}(\mathbf{x}, \mu)$$

This is called the **primal problem**

# Necessary Conditions – KKT Conditions

$$\begin{aligned} & \underset{\mathbf{x}}{\text{minimize}} && f(\mathbf{x}) \\ & \text{subject to} && \mathbf{g}(\mathbf{x}) \leq 0 \\ & && \mathbf{h}(\mathbf{x}) = 0 \end{aligned}$$

Any critical point  $\mathbf{x}^*$  must satisfy the **Karush-Kuhn-Tucker conditions**

1. primal feasibility:  $\mathbf{g}(\mathbf{x}^*) \leq 0$  and  $\mathbf{h}(\mathbf{x}^*) = 0$
2. dual feasibility: penalization is towards feasibility  $\mu \geq 0$
3. complementary slackness: either  $\mu_i$  or  $g_i(\mathbf{x}^*)$  is zero.

$$\mu_i g_i(\mathbf{x}^*) = 0, \text{ for } i = 1, \dots, m.$$

4. stationarity: objective function tangent to each active constraint

$$\nabla f(\mathbf{x}^*) + \sum_i \mu_i \nabla g_i(\mathbf{x}^*) + \sum_j \lambda_j \nabla h_j(\mathbf{x}^*) = 0$$

# Necessary Conditions – KKT Conditions

Particular cases

- $f$  concave,  $g$  convex: then KKT are also sufficient
- Patological cases

In vector form:

$$\begin{cases} \nabla f(\mathbf{x}^*) + \mu \nabla \cdot \mathbf{g}(\mathbf{x}^*) + \lambda \cdot \mathbf{h}(\mathbf{x}^*) = 0 \\ \mu \cdot \mathbf{g}(\mathbf{x}^*) = 0 \\ \mathbf{g}(\mathbf{x}^*) \leq 0, \mathbf{h}(\mathbf{x}^*) = 0 \\ \mu \geq 0 \end{cases}$$

# Duality

- Generalized Lagrangian Relaxation:

$$\mathcal{L}(\mathbf{x}, \boldsymbol{\mu}, \boldsymbol{\lambda}) = f(\mathbf{x}) + \sum_i \mu_i g_i(\mathbf{x}) + \sum_j \lambda_j h_j(\mathbf{x})$$

- the primal form is

$$\underset{\mathbf{x}}{\text{minimize}} \underset{\boldsymbol{\mu} \geq 0, \boldsymbol{\lambda}}{\text{maximize}} \mathcal{L}(\mathbf{x}, \boldsymbol{\mu}, \boldsymbol{\lambda})$$

- Reversing the order of operations leads to the **dual form**

$$\underset{\boldsymbol{\mu} \geq 0, \boldsymbol{\lambda}}{\text{maximize}} \underset{\mathbf{x}}{\text{minimize}} \mathcal{L}(\mathbf{x}, \boldsymbol{\mu}, \boldsymbol{\lambda})$$

- In some cases, the dual problem is easier to solve computationally than the original problem. In other cases, the dual can be used to obtain easily a lower bound on the optimal value of the objective for the primal problem. The dual has also been used to design algorithms for solving the primal problem.

# Duality

## Theorem (Max-min inequality)

For any function  $f : Z \times W \rightarrow \mathbb{R}$ ,

$$\sup_{z \in Z} \inf_{w \in W} f(z, w) \leq \inf_{w \in W} \sup_{z \in Z} f(z, w) .$$

Proof: see wikipedia



- When  $f$ ,  $W$ , and  $Z$  are convex the inequality becomes equality and we have a strong max-min property (or a saddle-point property).
- For us:

$$\underset{\mu \geq 0, \lambda}{\text{maximize}} \underset{\mathbf{x}}{\text{minimize}} \mathcal{L}(\mathbf{x}, \mu, \lambda) \leq \underset{\mathbf{x}}{\text{minimize}} \underset{\mu \geq 0, \lambda}{\text{maximize}} \mathcal{L}(\mathbf{x}, \mu, \lambda)$$

- Therefore, the solution to the dual problem  $d^*$  is a lower bound to the primal solution  $p^*$
- The inner part of the dual problem can be used to define the **dual function** or **dual objective**

$$\mathcal{D}(\mu \geq 0, \lambda) = \underset{\mathbf{x}}{\text{minimize}} \mathcal{L}(\mathbf{x}, \mu, \lambda)$$



# Duality

- The dual function is concave. Gradient ascent on a concave function always converges to the global maximum.
- **Dual Problem:**  $\max \mathcal{D}(\lambda)$  subject to  $\lambda \geq 0$
- Optimizing the dual problem is easy whenever minimizing the Lagrangian with respect to  $x$  is easy.
- For any  $\mu \geq 0$  and any  $\lambda$ , we have

$$\mathcal{D}(\mu \geq 0, \lambda) \leq p^*$$

- The difference between dual and primal solutions  $d^*$  and  $p^*$  is called the **duality gap**
- Showing zero-duality gap is a “certificate” of optimality

# Penalty methods

- Penalty methods are a way of reformulating a constrained optimization problem as an unconstrained problem by penalizing the objective function value when constraints are violated

## Example

$$\begin{aligned} & \underset{\mathbf{x}}{\text{minimize}} && f(\mathbf{x}) \\ & \text{subject to} && \mathbf{g}(\mathbf{x}) \leq 0 \\ & && \mathbf{h}(\mathbf{x}) = 0 \end{aligned}$$

$$\begin{aligned} & \min_{\mathbf{x}} && f(\mathbf{x}) + \rho \cdot p_{count}(\mathbf{x}) \\ & \text{s.t.} && p_{count}(\mathbf{x}) = \sum_i (g_i(\mathbf{x}) > 0) + \sum_j (h_j(\mathbf{x}) \neq 0) \end{aligned}$$

# Penalty Methods

**Procedure** penalty\_method;

**Input:**  $f, p, \mathbf{x}, k_{max}; \rho = 1, \gamma = 2$

**Output:**  $\mathbf{x}$  solution

**for**  $k$  in  $1, \dots, k_{max}$  **do**

$\mathbf{x} \leftarrow \text{minimize}_{\mathbf{x}} \{f(\mathbf{x}) + \rho \cdot p(\mathbf{x})\};$

$\rho \leftarrow \rho \cdot \gamma;$

**if**  $p(\mathbf{x}) = 0$  **then**

        return  $\mathbf{x};$

return  $\mathbf{x};$

# Penalty methods

- Count penalty:

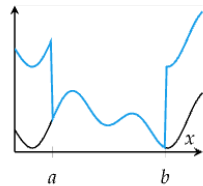
$$p_{count}(\mathbf{x}) = \sum_i (g_i(\mathbf{x}) > 0) + \sum_j (h_j(\mathbf{x}) \neq 0)$$

- Quadratic penalty:

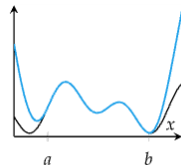
$$p_{quadratic}(\mathbf{x}) = \sum_i \max(g_i(\mathbf{x}), 0)^2 + \sum_j h_j(\mathbf{x})^2$$

- Mixed Penalty:

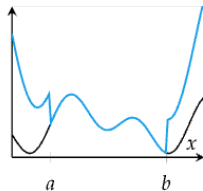
$$p_{mixed}(\mathbf{x}) = \rho_1 p_{count}(\mathbf{x}) + \rho_2 p_{quadratic}(\mathbf{x})$$



—  $f(x)$   
—  $f(x) + \rho p_{count}(x)$



—  $f(x)$   
—  $f(x) + \rho p_{quadratic}(x)$



—  $f(x)$   
—  $f(x) + \rho p_{mixed}(x)$

# Augmented Lagrange Method

- Adaptation of penalty method for equality constraints

$$p_{\text{Lagrangian}}(\mathbf{x}) \stackrel{\text{def}}{=} \frac{1}{2}\rho \sum_i (h_i(\mathbf{x}))^2 - \sum_i \lambda_i h_i(\mathbf{x})$$

**Procedure** augmented\_lagrange\_method;

**Input:**  $f, h, \mathbf{x}, k_{\max}; \rho = 1, \gamma = 2)$

$\lambda \leftarrow 0;$

**for**  $k$  in  $1, \dots, k_{\max}$  **do**

$\rho \leftarrow (x \mapsto \rho/2 \cdot \sum_i (h_i(\mathbf{x}))^2) - \lambda \cdot h(\mathbf{x});$   
 $\mathbf{x} \leftarrow \text{minimize}_{\mathbf{x}} \{f(\mathbf{x}) + \rho(\mathbf{x})\};$   
 $\lambda \leftarrow \lambda - \rho \cdot h(\mathbf{x});$   
 $\rho \leftarrow \rho \cdot \gamma;$

**return**  $\mathbf{x};$

- $\lambda$  converges towards the Lagrangian multiplier

# Interior Point Methods

- Also called **barrier methods**, interior point methods ensure that each step is feasible
- This allows premature termination to return a nearly optimal, feasible point
- Barrier functions are implemented similar to penalties but must meet the following conditions:
  1. Continuous
  2. Non-negative
  3. Approach infinity as  $x$  approaches boundary

# Interior Point Methods

- Inverse Barrier:

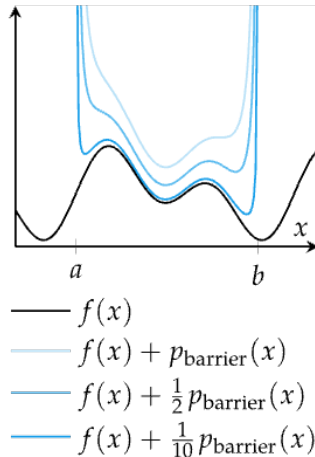
$$p_{\text{barrier}}(\mathbf{x}) = -\sum_i \frac{1}{g_i(\mathbf{x})}$$

- Log Barrier:

$$p_{\text{barrier}}(\mathbf{x}) = -\sum_i \begin{cases} \log(-g_i(\mathbf{x})) & \text{if } g_i(\mathbf{x}) \geq -1 \\ 0 & \text{otherwise} \end{cases}$$

New optimization problem:

$$\underset{\mathbf{x}}{\text{minimize}} \quad f(\mathbf{x}) + \frac{1}{\rho} p_{\text{barrier}}(\mathbf{x})$$



# Interior Point Methods

**Procedure** interior\_point\_method;

**Input:**  $f, p, \mathbf{x}; \rho = 1, \gamma = 2, \epsilon = 0.001$

$\Delta \leftarrow \infty;$

**while**  $\Delta > \epsilon$  **do**

$\mathbf{x}' \leftarrow \underset{\mathbf{x}}{\text{minimize}} \{f(\mathbf{x}) + p(\mathbf{x})/\rho\};$   
 $\Delta \leftarrow \|\mathbf{x}' - \mathbf{x}\|;$   
 $\mathbf{x} \leftarrow \mathbf{x}';$   
 $\rho \leftarrow \rho \cdot \gamma;$

**return**  $\mathbf{x};$

- Line searches  $f(\mathbf{x} + \alpha \mathbf{d})$  are constrained to the interval  $\alpha = [0, \alpha_u]$ , where  $\alpha_u$  is the step to the nearest boundary.  
In practice,  $\alpha_u$  is chosen such that  $\mathbf{x} + \alpha \mathbf{d}$  is just inside the boundary to avoid the boundary singularity.
- Needs an initial **feasible** solutions. Typically, found by solving:

$$\underset{\mathbf{x}}{\text{minimize}} \ p_{\text{quadratic}}(\mathbf{x})$$



# Summary

- Constraints are requirements on the design points that a solution must satisfy
- Some constraints can be transformed or substituted into the problem to result in an unconstrained optimization problem
- Analytical methods using Lagrange multipliers yield the generalized Lagrangian and the necessary conditions for optimality under constraints
- A constrained optimization problem has a dual problem formulation that is easier to solve and whose solution is a lower bound of the solution to the original problem
- Penalty methods penalize infeasible solutions and often provide gradient information to the optimizer to guide infeasible points toward feasibility
- Interior point methods maintain feasibility but use barrier functions to avoid leaving the feasible set