#### AI505 Optimization

# Discrete Optimization Heuristic Methods, Randomized Optimization Algorithms

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#### Outline

1. Constraint Programming

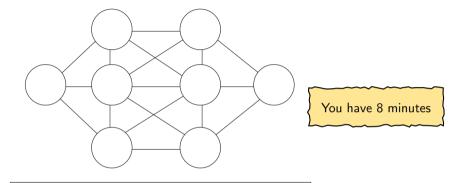
2. Randomized Optimization Heuristics

#### Outline

1. Constraint Programming

2. Randomized Optimization Heuristic

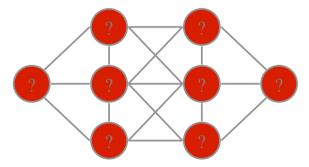
#### Number Circle Puzzle

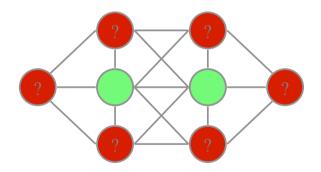


Put a different number (1 to 8) in each circle such that adjacent circles do not have consecutive numbers.

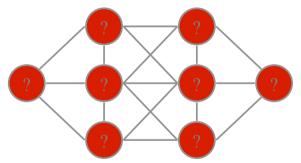
Example by Patrick Prosser with the help of Toby Walsh, Chris Beck, Barbara Smith, Peter van Beek,

Which nodes are hardest to number?

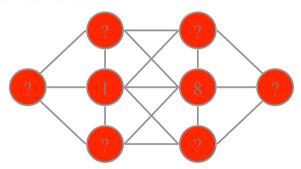




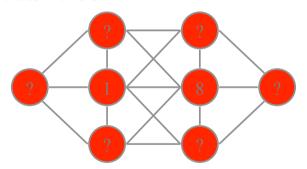
Which are the least constraining values to use?



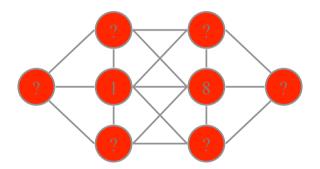
Values 1 and 8



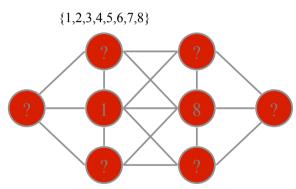
Values 1 and 8

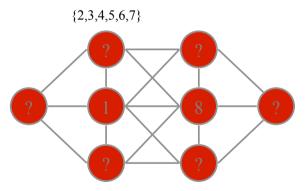


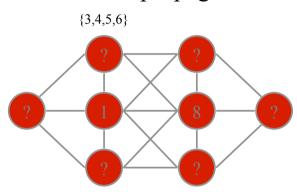
Symmetry means we don't need to consider: 8 1

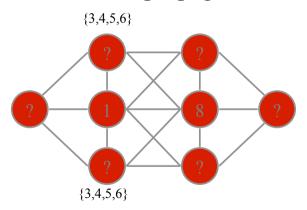


We can now eliminate many values for other nodes

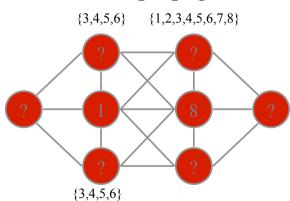


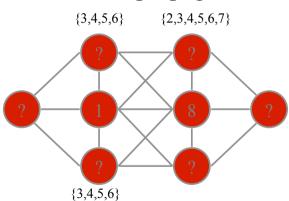


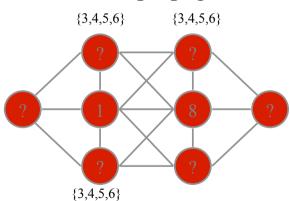


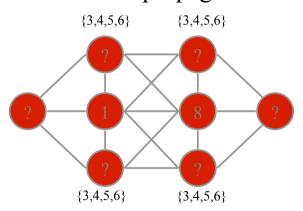


By symmetry

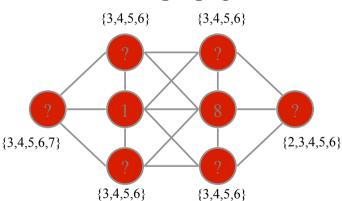


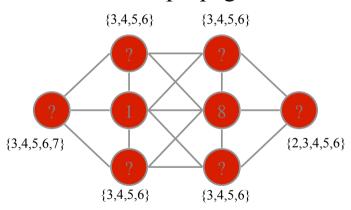




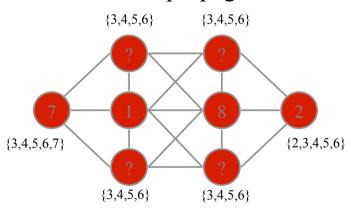


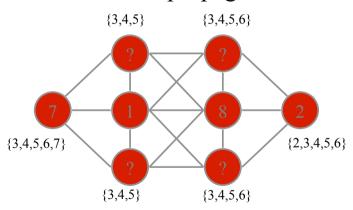
By symmetry

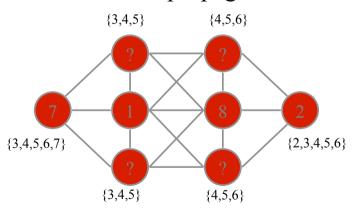


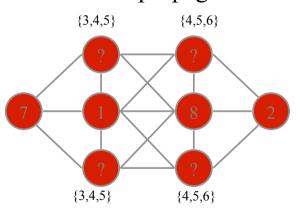


Value 2 and 7 are left in just one variable domain each

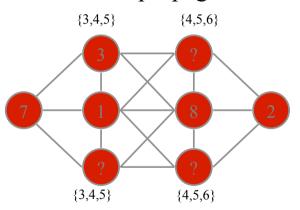




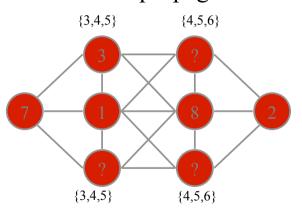


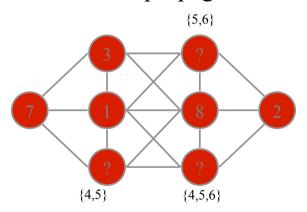


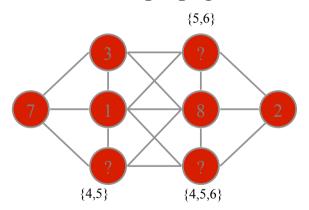
Guess a value, but be prepared to backtrack ...



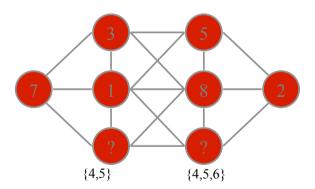
Guess a value, but be prepared to backtrack ...



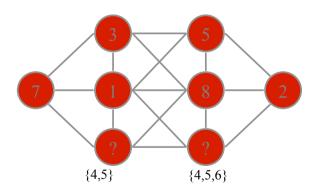


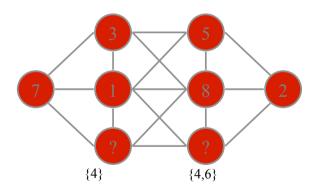


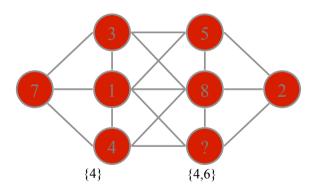
Guess another value ...



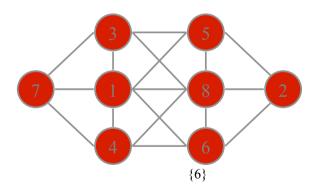
Guess another value ...



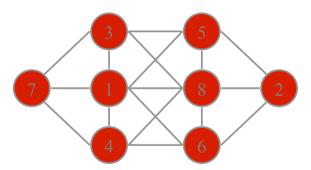




One node has only a single value left ...



# Solution

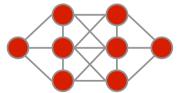


# The Core of Constraint Computation

- Modelling
  - Deciding on variables/domains/constraints
- Heuristic Search
- Inference/Propagation
- Symmetry
- Backtracking

#### Hardness

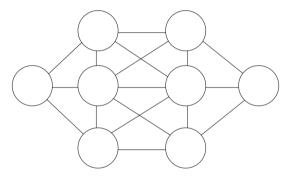
- The puzzle is actually a hard problem
  - NP-complete



### Constraint programming

- Model problem by specifying constraints on acceptable solutions:
  - define variables and domains
  - post constraints on these variables
- Solve model
  - choose algorithm
    - incremental assignment / backtracking search
    - complete assignments / stochastic search
  - design heuristics

#### Constraint Satisfaction Problem



- Variable  $x_i$  for each node i = 1, ..., 8
- Domain  $\{1, \ldots, 8\}$  for each variable  $x_i$
- Constraints:

allDifferent([
$$x_1, x_2, x_3, x_4, x_5, x_6, x_7, x_8$$
])  
 $|x_1 - x_2| > 1$   
 $|x_2 - x_3| > 1$   
 $|x_3 - x_4| > 1$   
 $\vdots$   
 $|x_7 - x_8| > 1$ 

# Modeling in Constraint Programming

The domain of a variable x, denoted D(x), is a finite set of elements that can be assigned to x.

A **constraint** C on X is a subset of the Cartesian product of the domains of the variables in X, i.e.,  $C \subseteq D(x_1) \times \cdots \times D(x_k)$ . A tuple  $(d_1, \ldots, d_k) \in C$  is called a **solution** to C.

Equivalently, we say that a solution  $(d_1, ..., d_k) \in C$  is an assignment of the value  $d_i$  to the variable  $x_i$  for all  $1 \le i \le k$ , and that this assignment satisfies C.

If  $C = \emptyset$ , we say that it is **inconsistent**.

# Modeling in Constraint Programming

#### Constraint Satisfaction Problem (CSP)

A CSP is a finite set of variables X with **domain extension**  $\mathcal{D} = D(x_1) \times \cdots \times D(x_n)$ , together with a finite set of constraints  $\mathcal{C}$ , each on a subset of X. A **solution** to a CSP is an assignment of a value  $d \in D(x)$  to each  $x \in X$ , such that all constraints are satisfied simultaneously.

#### Constraint Optimization Problem (COP)

A COP is a CSP  $\mathcal{P}$  defined on the variables  $x_1, \ldots, x_n$ , together with an objective function  $f: D(x_1) \times \cdots \times D(x_n) \to Q$  that assigns a value to each assignment of values to the variables. An **optimal solution** to a minimization (maximization) COP is a solution d to  $\mathcal{P}$  that minimizes (maximizes) the value of f(d).

### **Another Example: Social Golfers**

Example (Social Golfer Problem (Combinatorial Design))

- 9 golfers: 1, 2, 3, 4, 5, 6, 7, 8, 9
- wish to play in groups of 3 players in 4 days
- such that no golfer plays in the same group with any other golfer more than just once.

Is it possible?

	Group 1	Group 2	Group 3
Day 0	???	???	???
Day 1	???	???	???
Day 2	???	???	777
Day 3	???	???	???

This is an instance of a **constrained satisfaction problem**. Adding an optimizing criterion we get a **constrained optimization problem**.

### **Solution Paradigms**

- Dedicated algorithms
   (eg.: enumeration, branch and bound, dynamic programming)
- 2. Constraint Programming
- 3. Integer Linear Programming
- 4. Other modeling (SAT, SMT, etc.)
- 5. Randomized Search/Optimization Heuristics

Common to 2-5: Representation (modeling) + reasoning (search + inference)

# Constraint Programming: Representation

Golfers

	Group 1	Group 2	Group 3
Day 0	???	???	???
Day 1	???	???	???
Day 2	???	???	???
Day 3	???	???	???

#### Integer variables:

assign[i,j] variables whose value is from domain  $\{1,2,3\}$ 

#### Constraints:

C1: each group has exactly groupSize players

C2: each pair of players only meets once

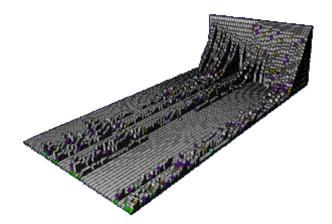
#### Groups

	Day 0	Day 1	Day 2	Day 3
Golfer 0	1	{1,2,3}	{1,2,3}	{1,2,3}
Golfer 1	1	{1,2,3}	{1,2,3}	{1,2,3}
Golfer 2	1	{1,2,3}	{1,2,3}	{1,2,3}
Golfer 3	{2,3}	{1,2,3}	{1,2,3}	{1,2,3}
Golfer 4	{2,3}	{1,2,3}	{1,2,3}	{1,2,3}
Golfer 5	{2,3}	{1,2,3}	{1,2,3}	{1,2,3}
Golfer 6	{2,3}	{1,2,3}	{1,2,3}	{1,2,3}
Golfer 7	{2,3}	{1,2,3}	{1,2,3}	{1,2,3}
Golfer 8	{2,3}	{1,2,3}	{1,2,3}	{1,2,3}

### Constraint Programming: Representation

```
int: golfers = 9:
int: groupSize = 3;
int: days = 4:
int: groups = golfers/groupSize:
set of int: Golfer = 1..golfers:
set of int: Day = 1..days;
set of int: Group = 1..groups;
array[Golfer, Dav] of var Group: assign; % Variables
constraint
 % C1: Each group has exactly groupSize players
 forall (gr in Group, d in Day) ( % c1
   sum (g in Golfer) (bool2int(assign[g,d] = gr)) = groupSize
 ) /\
 % C2: Each pair of players only meets at most once
 forall (q1, q2 in Golfer, d1, d2 in Day where q1 != q2 / d1 != d2) (
   (bool2int(assign[q1,d1] = assign[q2,d1]) + bool2int(assign[q1,d2] = assign[q2,d2])) <=1);
solve :: int search([assign[i, j] | i in Golfer, j in Day ],
                   first fail, indomain min, complete) satisfy;
```

### Constraint Programming: Reasoning



The solution process proceeds by propagating the constraints on the domanins of the variables (ie, removing values) and tentatively assigning variables until only feasible values are left or backtracking.

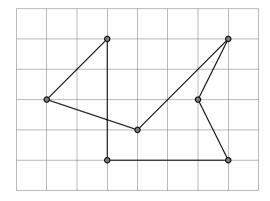
### Outline

1. Constraint Programming

2. Randomized Optimization Heuristics

### Yet Another Example: TSP

#### Example (Traveling Salesman Problem)



Can you find a better solution?

#### Heuristics

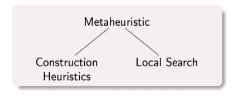
- Get inspired by approach to problem solving in human mind
   [A. Newell and H.A. Simon. "Computer science as empirical inquiry: symbols and search."
  - Communications of the ACM, ACM, 1976, 19(3)]
    effective rules without theoretical support
    - trial and error
- Applications:
  - Optimization
  - But also in Psychology, Economics, Management [Tversky, A.; Kahneman, D. (1974).
     "Judgment under uncertainty: Heuristics and biases". Science 185]
- Basis on empirical evidence rather than mathematical logic. Getting things done in the given time.

# Randomized Optimization Heuristics (ROHs)

#### Two main search paradigms:

- Constructive search
- Local search

plus high level guiding heuristics (ie, metaheuristics), eg, evolutionary algorithms.



# **ROHs: Representation**

	Group 1	Group 2	Group 3
Day 0	012	3 4 5	6 7 8
Day 1	<b>0 4</b> 6	1 <b>3 7</b>	2 5 8
Day 2	0 4 8	156	2 <b>3 7</b>
Day 3	0 5 7	1 3 8	2 4 6

- Variables = solution representation, tentative solution
- Constraints: relaxed = soft
- Evaluation function to guide the search

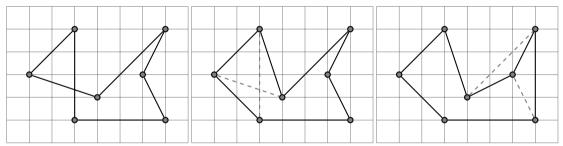
# ROHs: Reasoning, Local Search Solution: Trial and Error

	Group 1	Group 2	Group 3
Day 0	012	3 4 5	6 7 8
Day 1	<b>0 4</b> 6	1 3 7	2 5 8
Day 2	0 4 8	156	2 <b>3 7</b>
Day 3	0 5 7	138	2 4 6

**Heuristic algorithms**: compute, efficiently, **good** solutions to a problem (without caring for theoretical guarantees on running time and approximation quality).

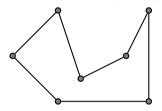
# ROHs: Reasoning, Local Search

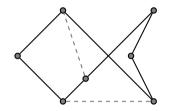
#### Example on Traveling Salesman Problem:

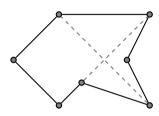


# ROHs: Reasoning, Metaheuristics

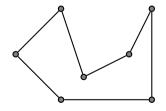
Accepting worsening changes

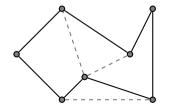


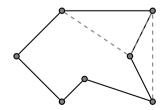




Trying different changes







### ROHs: Reasoning, Metaheuristics

- Stochastic Local Ssearch
- Simulated Annealing
- Iterated Local Search
- Tabu Search
- Variable Neighborhood Search
- Adaptive Large Neighborhood Search
- Evolutionary Algorithms
- Ant Colony Optimization
- Estimation-of-Distribution Algorithms
- Artificial Immune Systems
- ...
- Evolutionary Computation Bestiary http://fcampelo.github.io/EC-Bestiary/
- Supernatural inspired [Maturana, Fouhey, 2013]

#### **A** Classification

- White box optimization: models can be expressed mathematically
- Grey box optimization: internal information about objective function computation is often available models that have a mathematical expression but may need data to determine them (eg, neural networks)
- Black box optimization: no mathematical expression is available

### Approaches to ROHs

- White/Grey box: representation (modelling) + reasoning (search)
   constraint based local search, comet, local solver (Hexaly)
- Black Box: a different approach, framework separating problem from solvers and defining the interface specification
   EasyLocal, ...,

⇒ Cost Action: ROAR-NET



https://github.com/roar-net/roar-net-api-spec

#### A Search Problem

#### Definition (Problem statement)

Assume we want to solve a **constrained optimization problem**:  $\min f(x) \mid x \in F$  where is a set of **feasible solutions** and f an objective function. All parameters of the problem are known and deterministic.

#### Definition (Search or Optimization Algorithm)

**Goal formulation**: we want to find the minimum with respect to some criteria from a set of candidate elements.

**Problem formulation**: Given a description of the states, an initial state and actions necessary to reach the goal, find a sequence of actions to reach the goal.

**Search**: the algorithm simulates sequences of actions in the model of the goal, searching until it finds a sequence of actions that reaches the goal. The algorithm might have to simulate multiple tentative answers that do not meet the goal, but eventually it reach a solution, or it will find that no solution is possible.

### Search Algorithms

#### Components of a Search Algorithm (1):

- State or Search Space A set of possible states that the search can be in.
- State or (Candidate) solution: a definition of the states of the search,
- Initial State that the search starts in. For example: an empty set of actions or a complete set
  of actions.
- Goal A set of one or more goal states. Sometimes there is one goal state sometimes there is a small set of alternative goal states
- Evaluation function f(s) assess the distance from a potential goal. It can also include relaxed constraints.

### Search Algorithms

#### Components of a Search Algorithm (2):

• Action Type t available to the algorithm.

- Neighborhood Structure
- For a given Action Type t and a State/Solution s, Actions(t,s) returns a finite set of actions of type t that can be executed in s. We say that each of these actions is applicable in s. A transition model, which describes what each action does.

  Neighborhood
- Result(s, a) returns the state that results from doing action a in state s. Apply Move
- Action-Cost(s, a, s') or c(s, a, s') action cost function gives the numeric cost of applying action a in state s to reach state s'. It reflects the evaluation of the state. Increment

### **Constraint Handling**

In the Constraint Based Local Search community, constraints in heuristic methods are handled:

- implicitly in the definition of the search space and of the actions
- as one way constraints
- as soft constraints
   ie, relaxed in the evaluation function as objectives with large weights or as lexicographically
   more important objectives

# Application Programming Interface (API)

(Here: not meant as Web API, network-based API, or REST API.)

The ROAR-NET API Specification is the definition of an interface or protocol between optimization problems seen as black box and their solvers in order to facilitate understanding, reusing and scaling of solution approaches.

#### We look for a model which

- ... allows one to use off the shelf components to solve it.
- ... assumes a separation between **problem specifics** and **solver**.
- .. is designed as a **software interface** offering a service to other pieces of software and is implemented by the user.
- ... promotes **reusability** of software components and minimizes the user's effort to deploy a solution for the specific optimization problem at hand.
- ... aims at maximizing code extensability, reusability, and simplicity.

### **Types**

- Problem the problem instance
- Solution implements the representation of a tentative solution
- Value represents points in objective space
- Neighborhood a function that given a solution, gives another solution neighbor: S → S, based on a neighborhood, compute a Move
- Move applied to a solution to get the novel solution
- Increment represents points in objective space

Simplification in single objective cases: Value and Increment are Reals (ie, Double or Integer)

### Operations: Problem Representation

#### 

• heuristic solution(P): Solution[0..1]

#### ... on Solution S:

```
copy_solution(S): Solutionlower_bound(S): Value[0..1]evaluation_value(): Value[0..1]
```

### **Operations: Search**

#### ... on Neighborhood (N):

- construction\_neighbourhood(Problem): Neighbourhood
- destruction\_neighbourhood(Problem): Neighbourhood
- local\_neighbourhood(Problem): Neighbourhood
- moves(N, Solution): Move[0..\*]
- random\_move(N, Solution): Move[0..1]
- random\_move\_without\_replacement(N, Solution): Move[0..\*]

### **Operations: Search**

#### ... on Move:

- lower\_bound\_increment(Move, Solution): double[0..1]
- objective\_value\_increment(Move, Solution): double[0..1] apply\_move(M, Solution): Solution applies the move to a solution, to get a novel solution
- invert\_move (M): Move computes the inverse of a move to revert it

#### The Full API

# Types Problem Solution Value Neighborhood Move # Operations on Problem empty\_solution random\_solution heuristic\_solution

# Operations on Solution objective\_value lower\_bound copy\_solution

# Operations on Neighborhood local\_neighbourhood construction\_neighbourhood destruction\_neighbourhood # Operations on Move
moves
random\_move
random\_moves\_without\_replacement
lower\_bound\_increment
objective\_value\_increment

apply\_move invert move