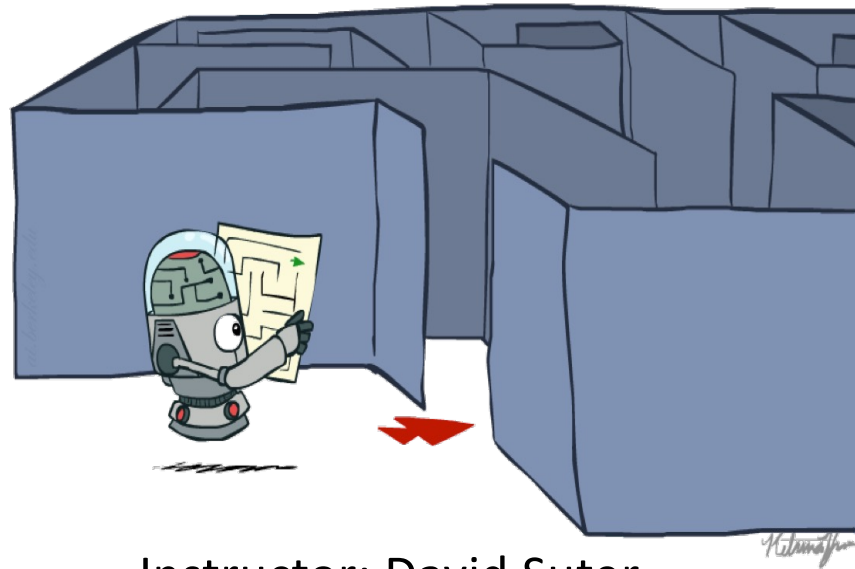


# Search



Instructor: David Suter

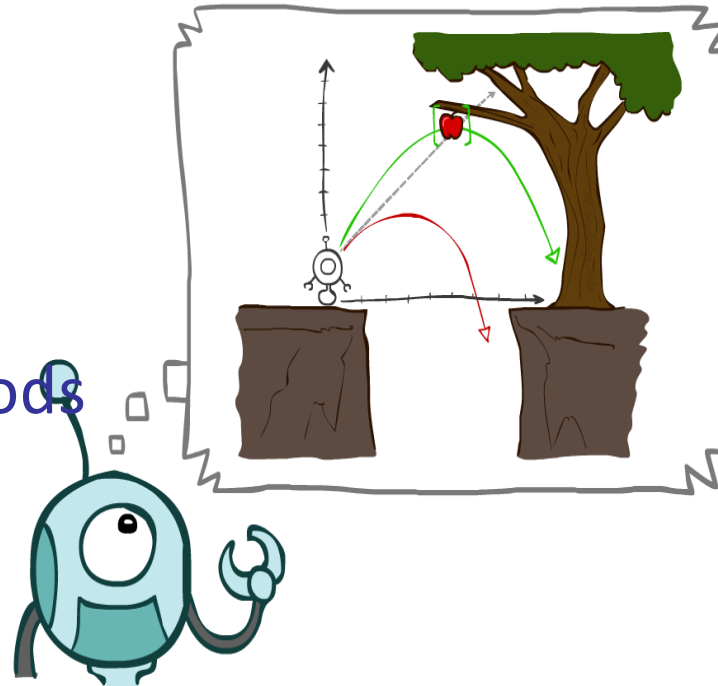
Course Delivered for Xidian University

[Many slides adapted from those created by Dan Klein and Pieter Abbeel for CS188 Intro to AI at UC Berkeley. Some others from colleagues at Adelaide University.]

# Topics

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- Agents that Plan Ahead
- Search Problems
- Uninformed Search Methods
  - Depth-First Search
  - Breadth-First Search
  - Uniform-Cost Search



# Search Problems

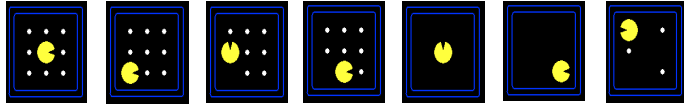
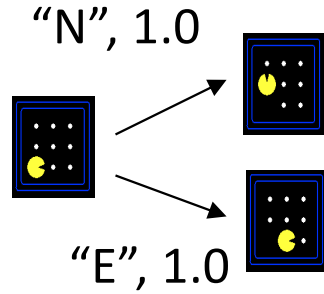
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# Search Problems

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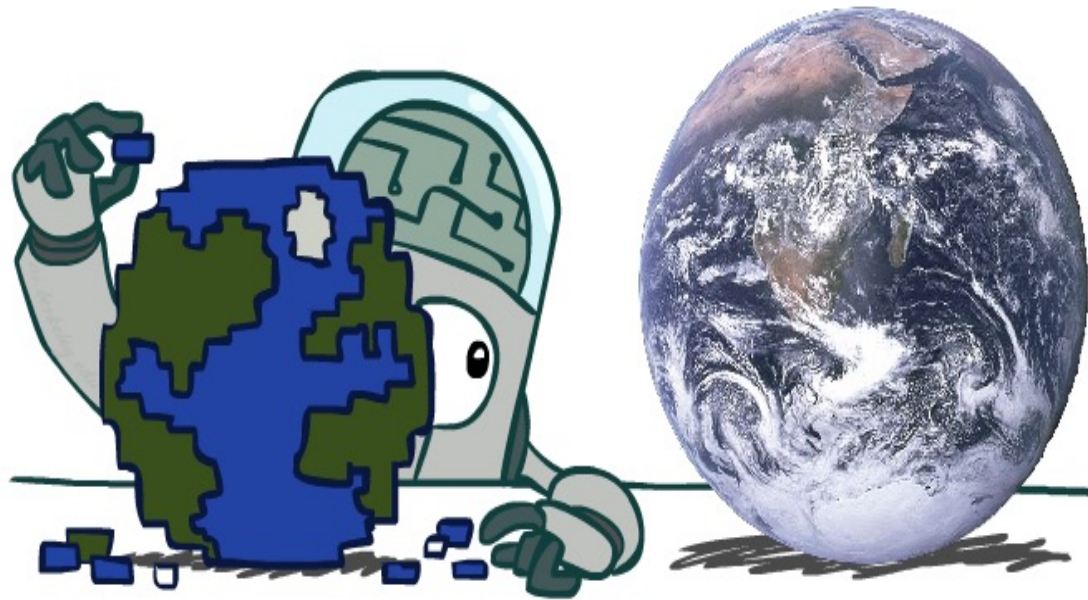
- A **search problem** consists of:

- A state space 
- A successor function (with actions, costs)  

- A start state and a goal test

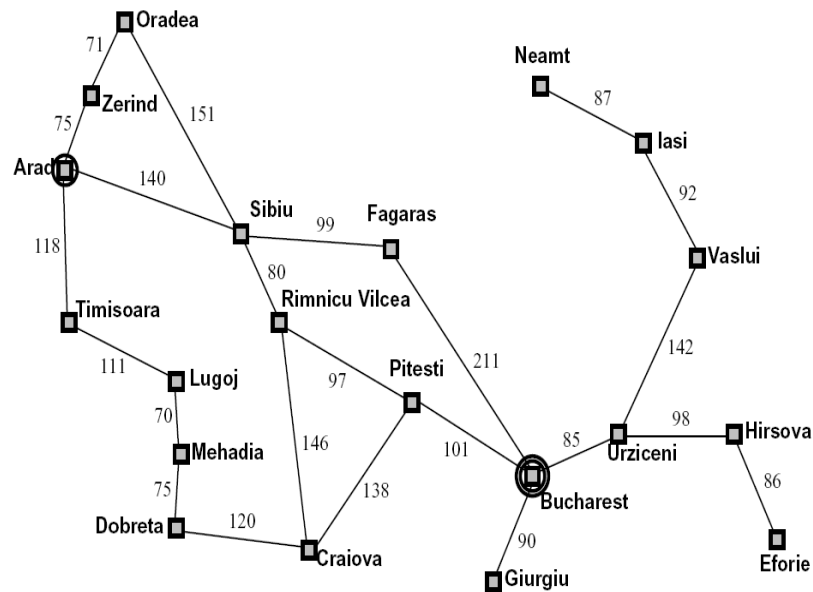
- A **solution** is a sequence of actions (a plan) which transforms the start state to a goal state

# Search Problems Are Models

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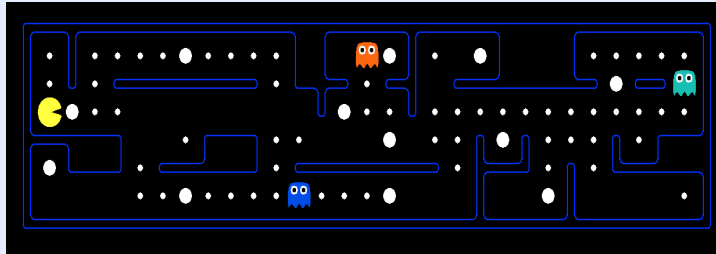
# Example: Traveling in Romania



- State space:
  - Cities
- Successor function:
  - Roads: Go to adjacent city with cost = distance
- Start state:
  - Arad
- Goal test:
  - Is state == Bucharest?
- Solution?

# What's in a State Space?

The **world state** includes every last detail of the environment



A **search state** keeps only the details needed for planning (abstraction)

## ■ Problem: Path Finding

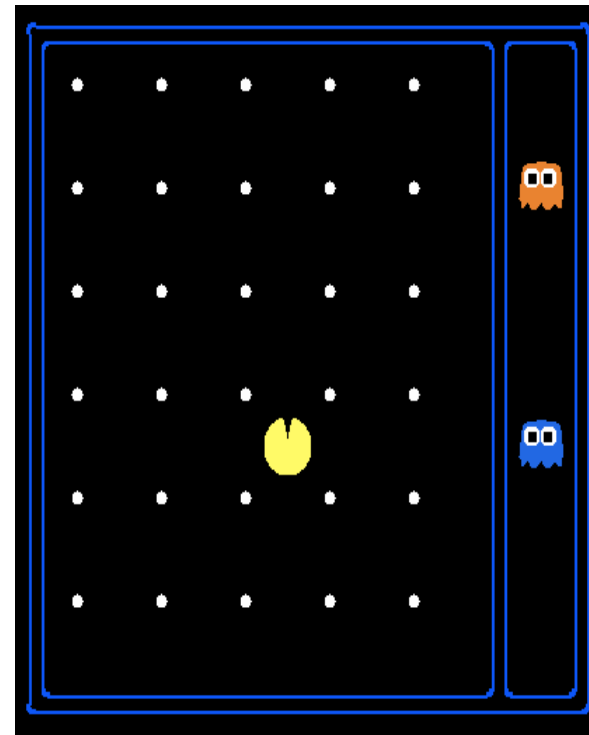
- States:  $(x,y)$  location
- Actions: NSEW
- Successor: update location only
- Goal test: is  $(x,y)=\text{END}$

## ■ Problem: Eat-All-Dots

- States:  $\{(x,y), \text{dot booleans}\}$
- Actions: NSEW
- Successor: update location and possibly a dot boolean
- Goal test: dots all

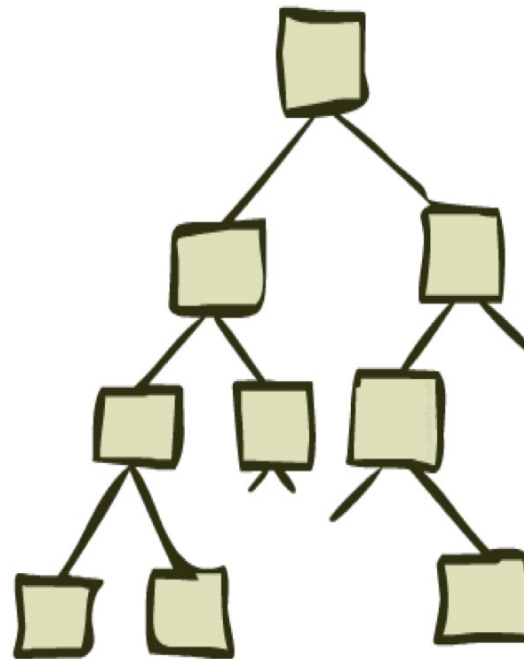
# State Space Sizes?

- World state:
  - Agent positions: 120
  - Food count: 30
  - Ghost positions: 12
  - Agent facing: NSEW
- How many
  - World states?  
 $120 \times (2^{30}) \times (12^2) \times 4$
  - States for path finding?  
120
  - States for eat-all-dots?  
 $120 \times (2^{30})$



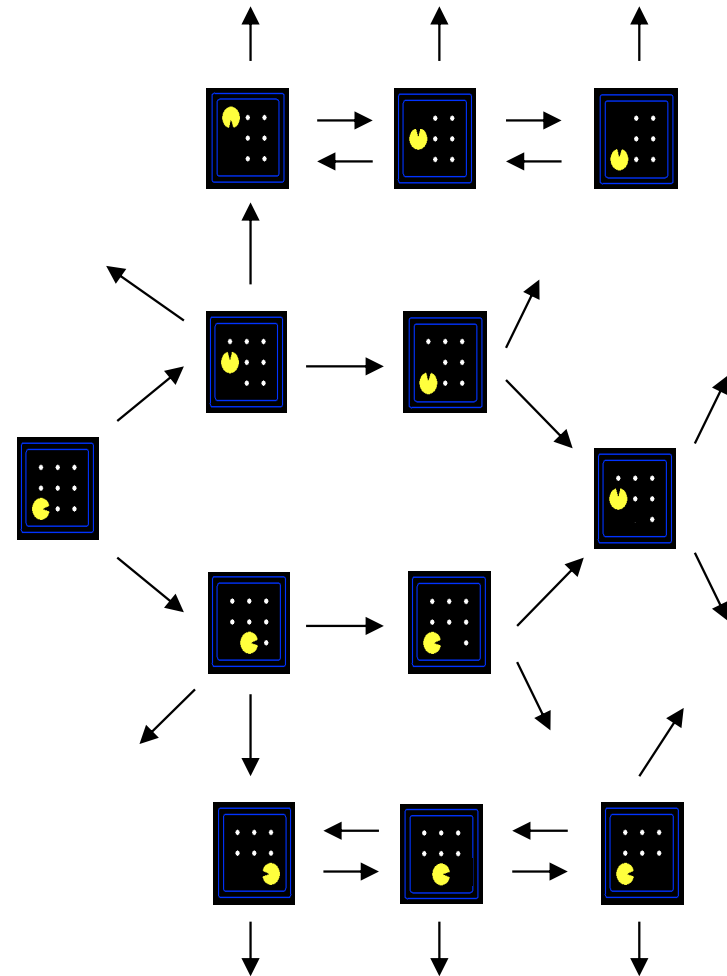


# State Space Graphs and Search Trees



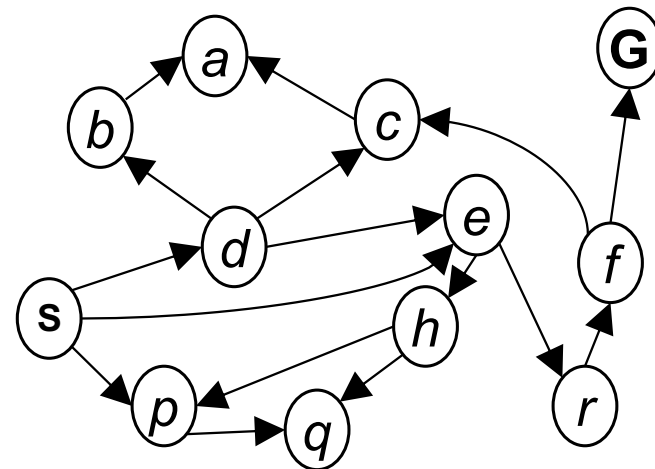
# State Space Graphs

- State space graph: A mathematical representation of a search problem
  - Nodes are (abstracted) world configurations
  - Arcs represent successors (action results)
  - The goal test is a set of goal nodes (maybe only one)
- In a state space graph, each state occurs only once!
- We can rarely build this full graph in memory (it's too big), but it's a useful idea



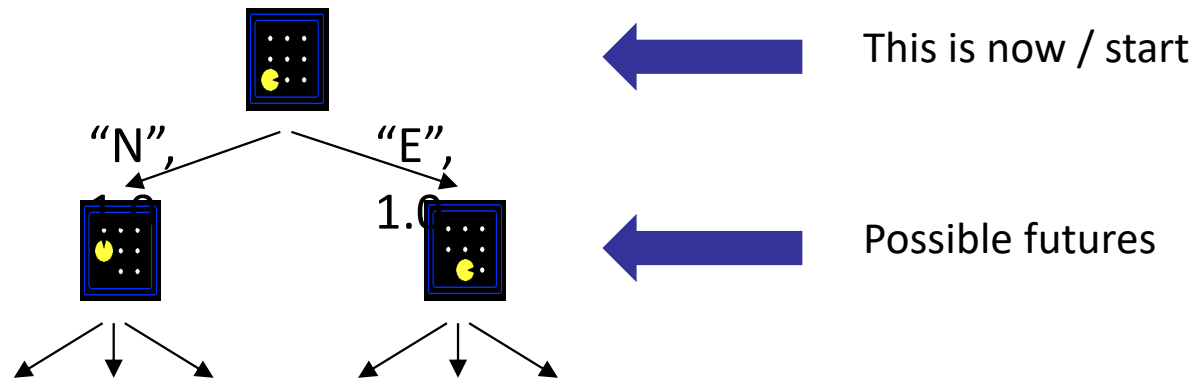
# State Space Graphs

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- In a search graph, each state occurs only once!
- We can rarely build this full graph in memory (it's too big), but it's a useful idea



*Tiny search graph  
for a tiny search  
problem*

# Search Trees

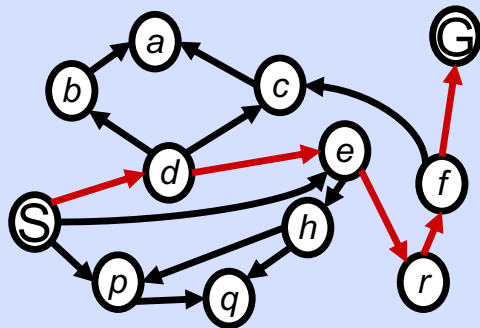


- A search tree:

- A “what if” tree of plans and their outcomes
- The start state is the root node
- Children correspond to successors
- Nodes show states, but correspond to PLANS that achieve those states
- For most problems, we can never actually build the whole tree

# State Space Graphs vs. Search Trees

State Space Graph

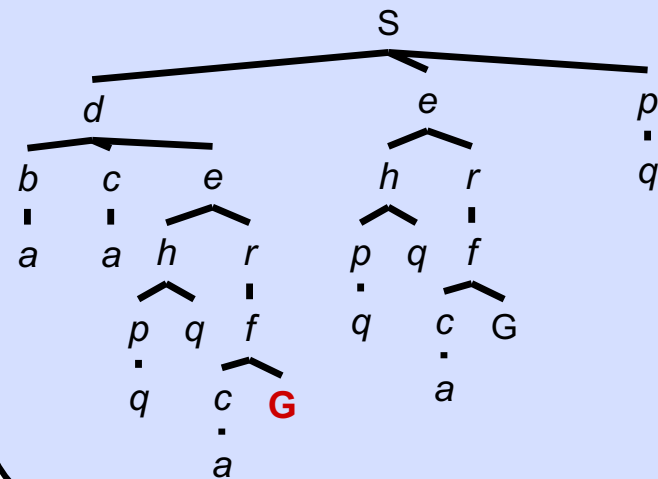


*Each NODE in the search tree is an entire PATH in the state space graph.*

*we construct both on demand – and we construct as little as*

*...*

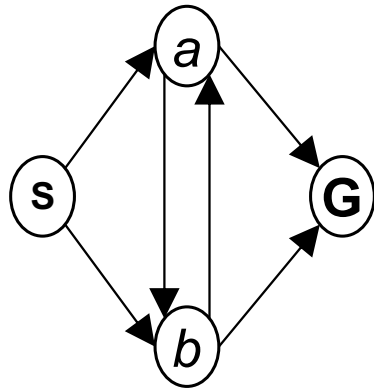
Search Tree



# Quiz: State Space Graphs vs. Search Trees

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Consider this 4-state graph:



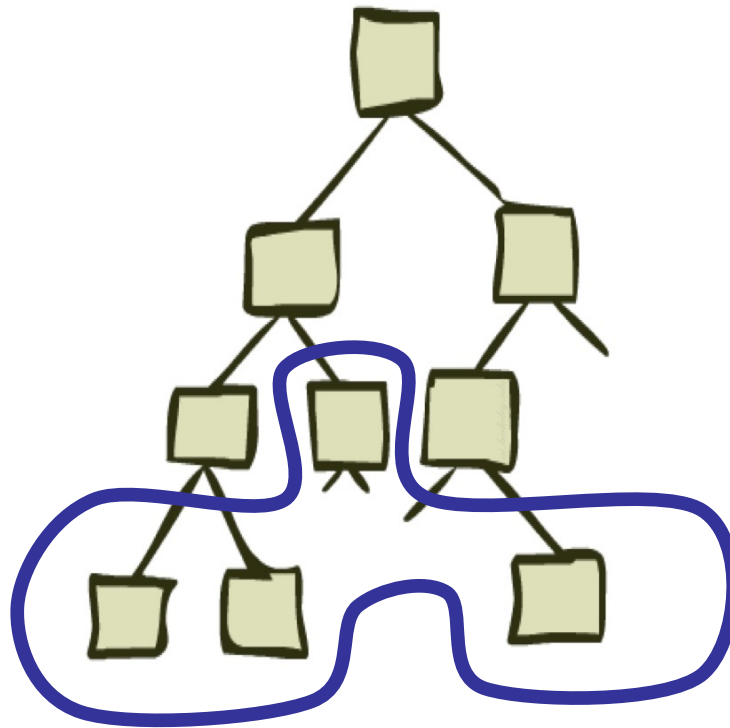
How big is its search tree (from S)?



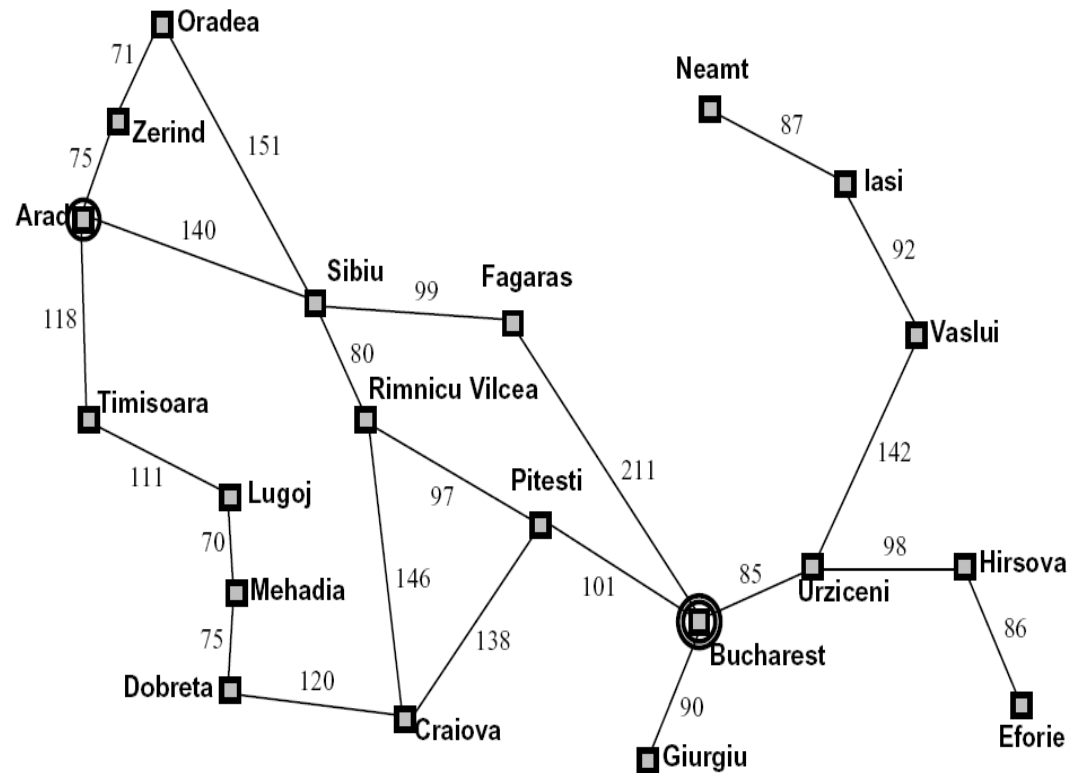
Important: Lots of repeated structure in the search tree!

# Tree Search

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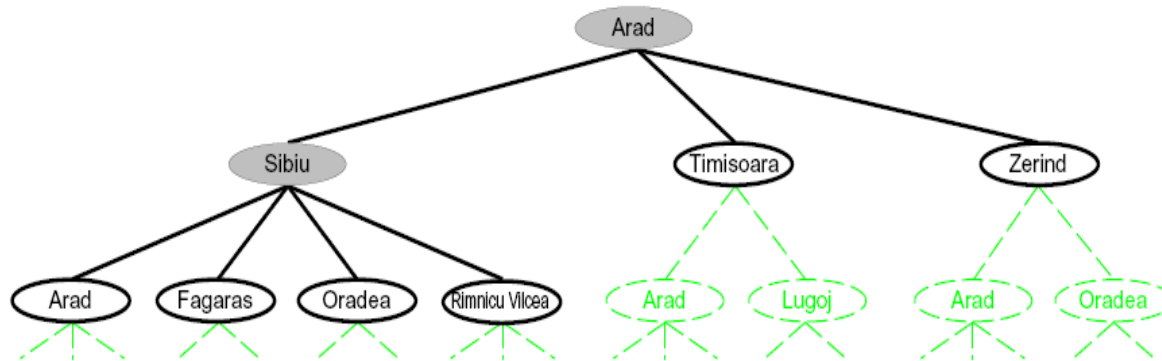
# Search Example: Romania





# Searching with a Search Tree

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- Search:

- Expand out potential plans (tree nodes)
- Maintain a **fringe** of partial plans under consideration
- Try to expand as few tree nodes as possible

# General Tree Search

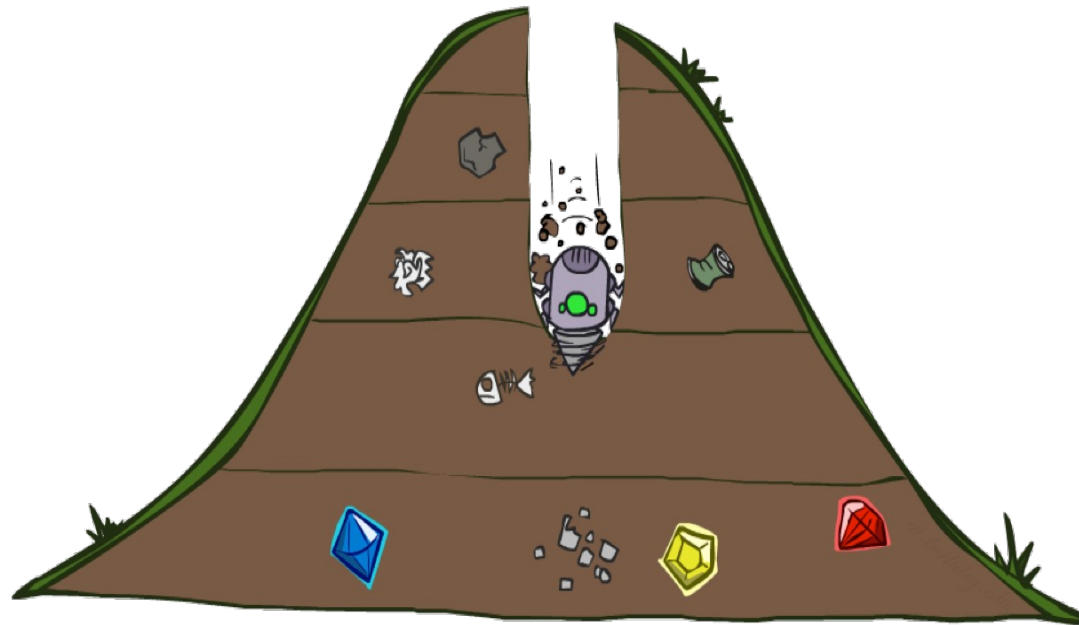
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```
function TREE-SEARCH(problem, strategy) returns a solution, or failure
  initialize the search tree using the initial state of problem
  loop do
    if there are no candidates for expansion then return failure
    choose a leaf node for expansion according to strategy
    if the node contains a goal state then return the corresponding solution
    else expand the node and add the resulting nodes to the search tree
  end
```

- Important ideas:
  - Fringe
  - Expansion
  - Exploration strategy
- Main question: which fringe nodes to explore?

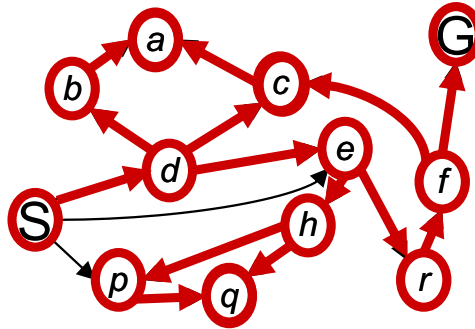
# Depth-First Search

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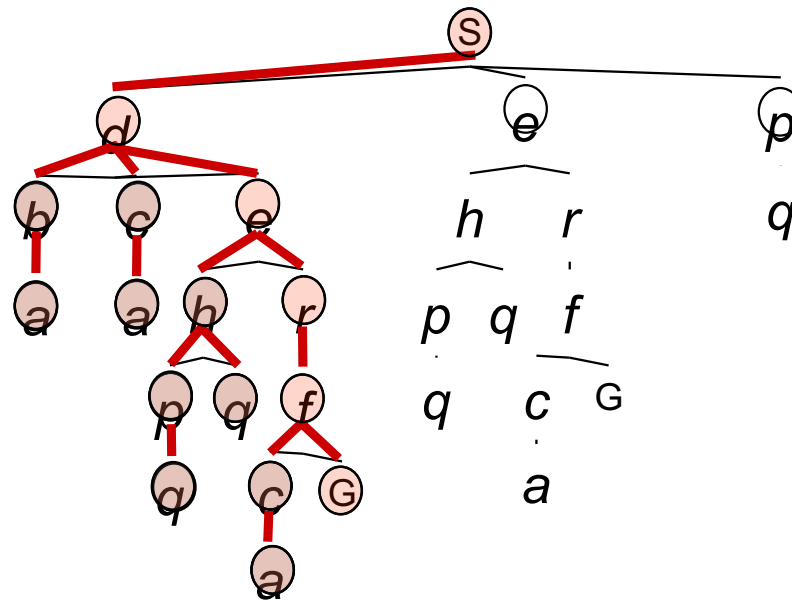


# Depth-First Search

*Strategy:*  
 expand a  
 deepest node  
 first

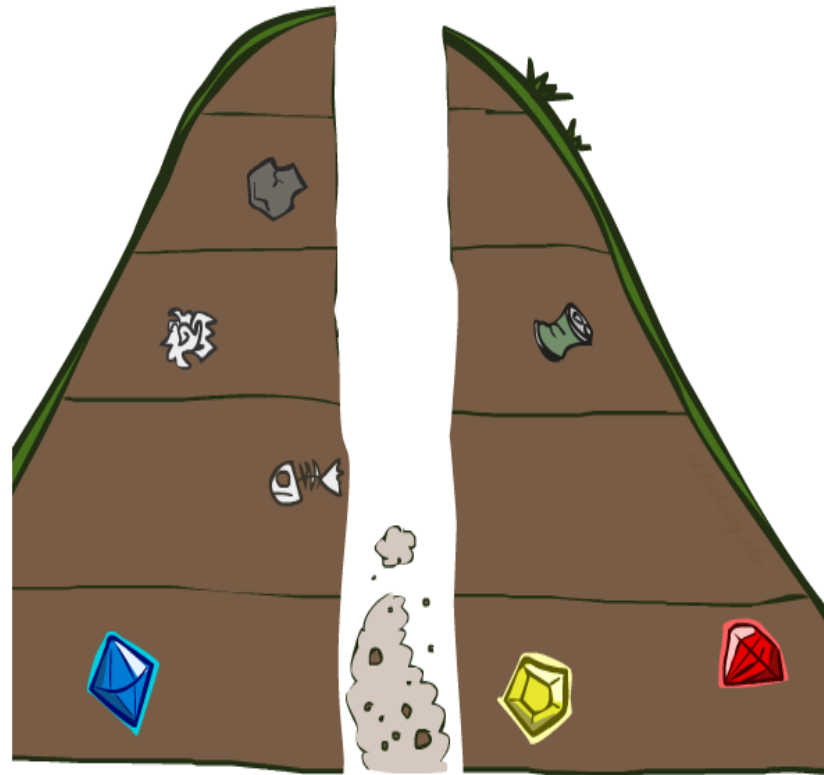


*Implementation:* Fringe is  
 a LIFO stack



# Search Algorithm Properties

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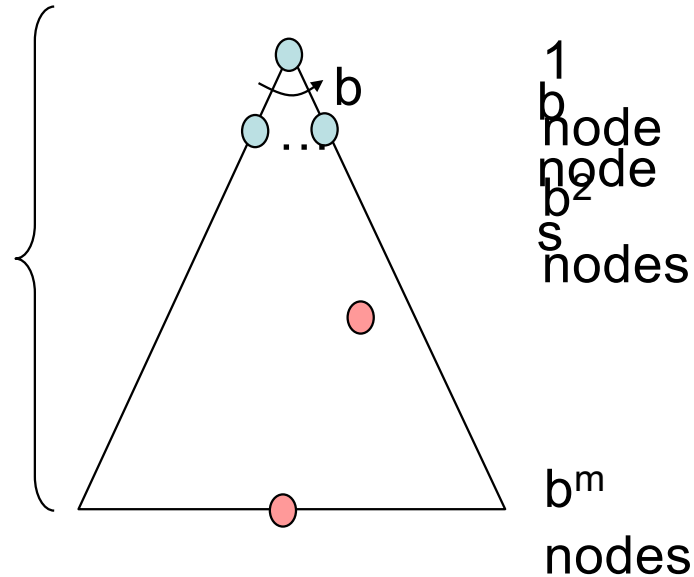


# Search Algorithm Properties

- Complete: Guaranteed to find a solution if one exists?
- Optimal: Guaranteed to find the least cost path?
- Time complexity?
- Space complexity?

- Cartoon of search tree:
  - $b$  is the branching factor
  - $m$  is the maximum depth
  - solutions at various depths

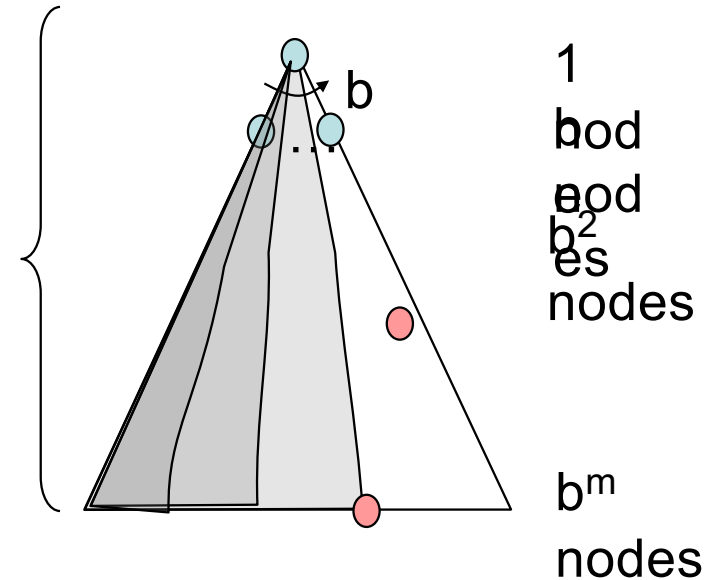
$m$   
tiers



- Number of nodes in entire tree?
  - $1 + b + b^2 + \dots + b^m = O(b^m)$

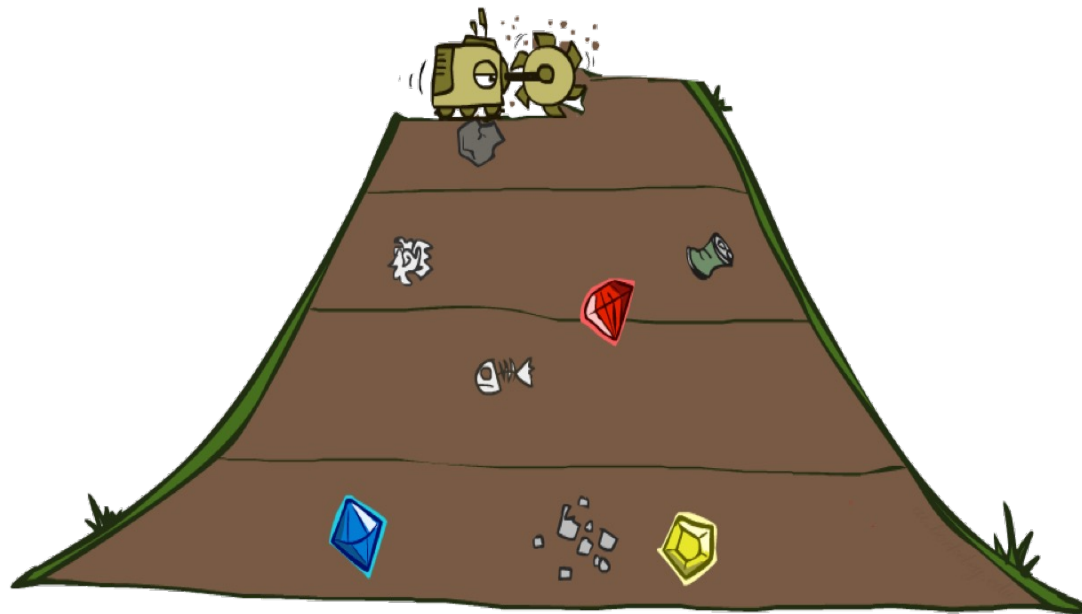
# Depth-First Search (DFS) Properties

- What nodes DFS expand?
  - Some left prefix of the tree.
  - Could process the whole tree!
  - If  $m$  is finite, takes time  $O(b^m)$
- How much space does the fringe take?
  - Only has siblings on path to root, so  $O(bm)$  tiers
- Is it complete?
  - $m$  could be infinite, so only if we prevent cycles (more later)
- Is it optimal?
  - No, it finds the “leftmost” solution, regardless of depth or cost



# Breadth-First Search

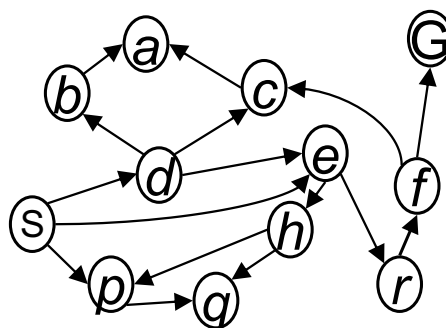
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# Breadth-First Search

*Strategy:  
expand a  
shallowest  
node first*

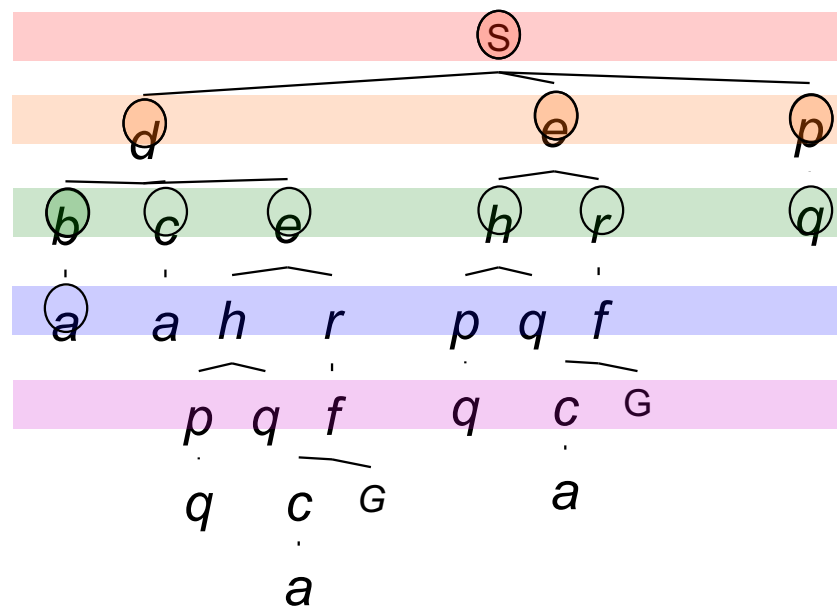


## Implementati

on: Fringe is a  
FIFO queue

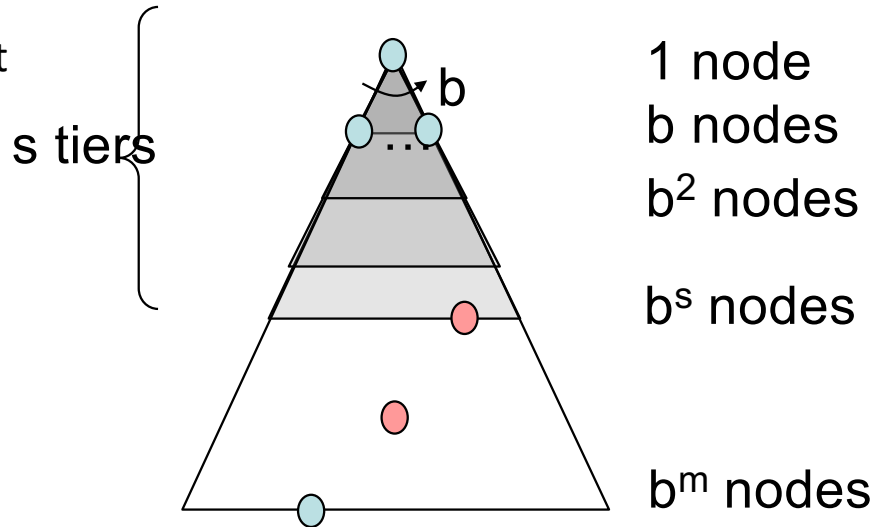
# Search

# Tiers



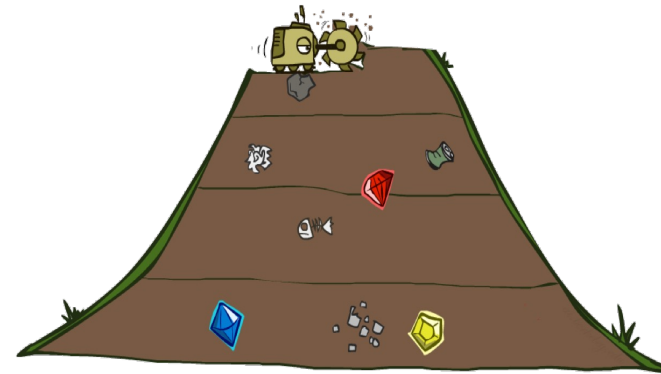
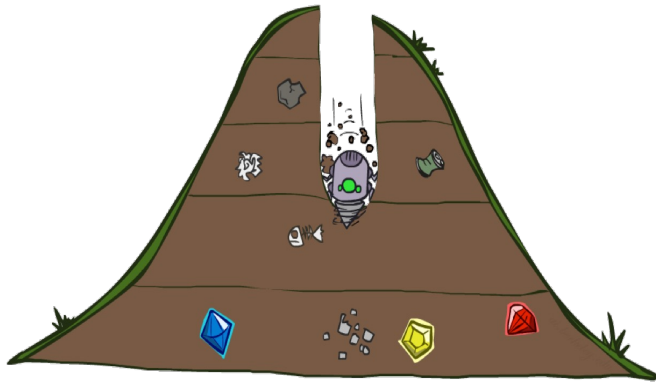
# Breadth-First Search (BFS) Properties

- What nodes does BFS expand?
  - Processes all nodes above shallowest solution
  - Let depth of shallowest solution be  $s$
  - Search takes time  $O(b^s)$
- How much space does the fringe take?
  - Has roughly the last tier, so  $O(b^s)$
- Is it complete?
  - $s$  must be finite if a solution exists, so yes!
- Is it optimal?
  - Only if costs are all 1 (more on costs later)



# Quiz: DFS vs BFS

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# Quiz: DFS vs BFS

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- When will BFS outperform DFS?
- When will DFS outperform BFS?

[Demo: dfs/bfs maze

# Video of Demo Maze Water DFS/BFS (part 1)

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# Video of Demo Maze Water DFS/BFS (part 2)

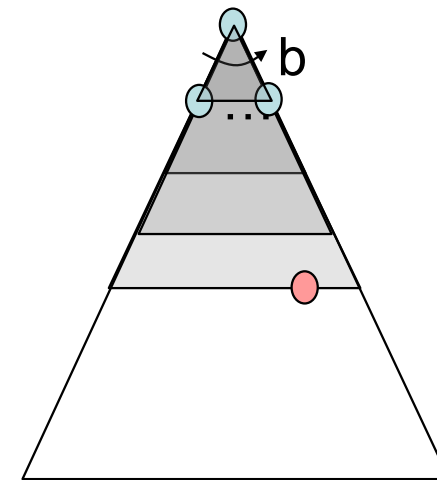
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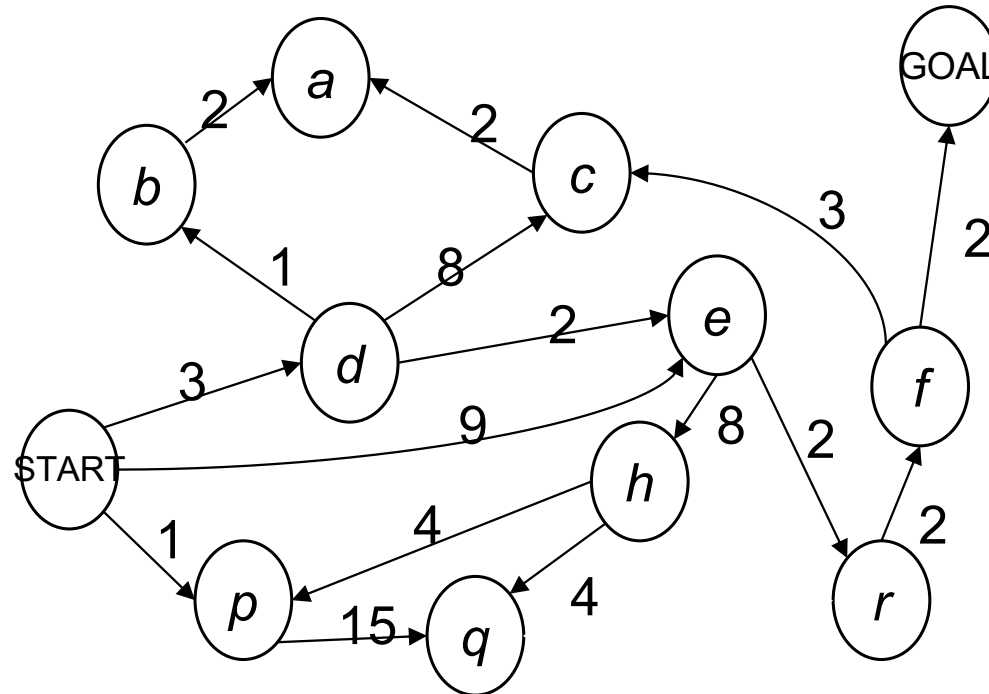
# Iterative Deepening

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- Idea: get DFS's space advantage with BFS's time / shallow-solution advantages
  - Run a DFS with depth limit 1. If no solution...
  - Run a DFS with depth limit 2. If no solution...
  - Run a DFS with depth limit 3. ....
- Isn't that wastefully redundant?
  - Generally most work happens in the lowest level searched, so not so bad!



# Cost-Sensitive Search

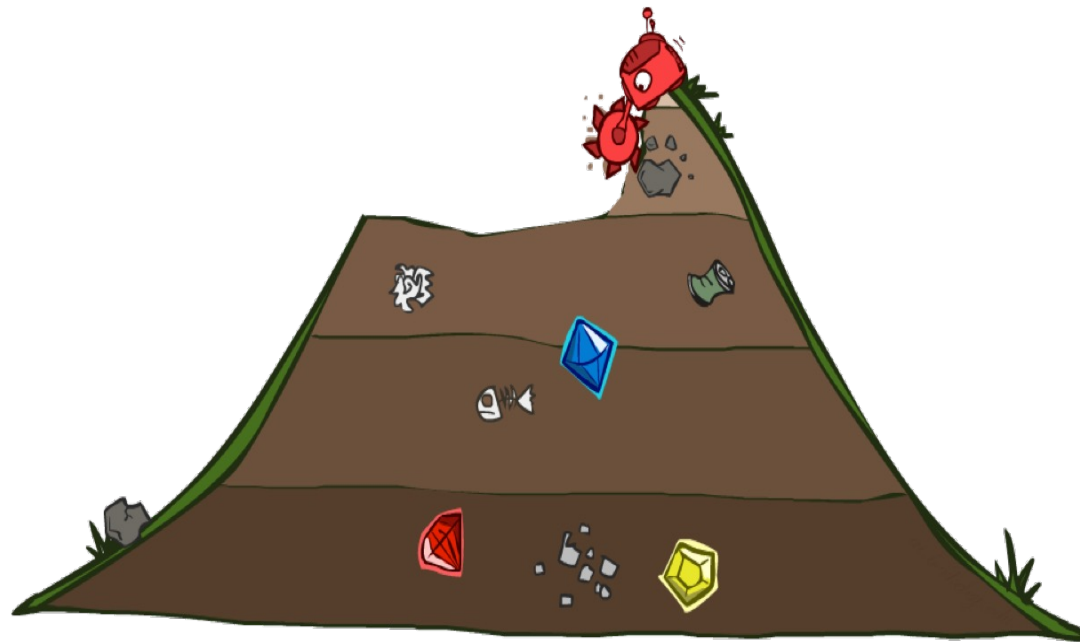


BFS finds the shortest path in terms of number of actions.  
It does not find the least-cost path. We will now cover  
a similar algorithm which does find the least-cost path.



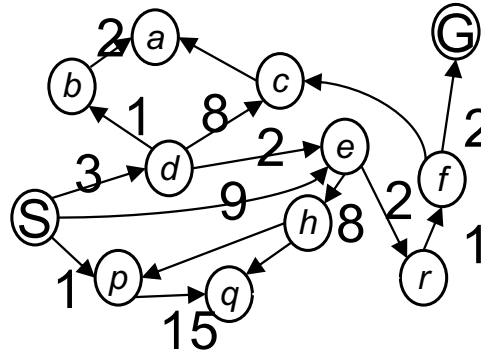
# Uniform Cost Search

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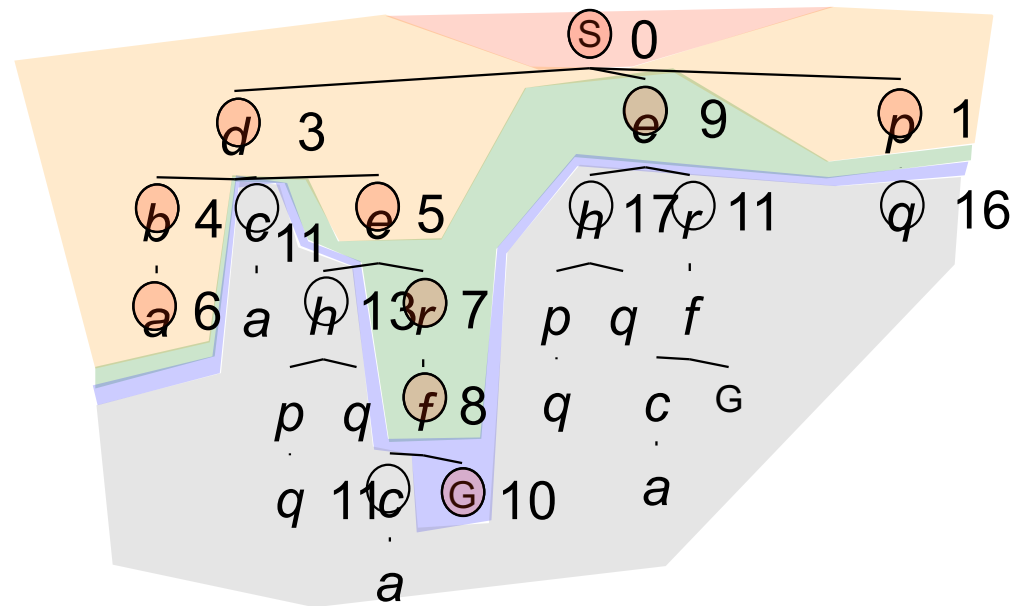
# Uniform Cost Search

*Strategy:*  
*expand a*  
*cheapest node*  
*first:*



*Fringe is a*  
*priority queue*  
*(priority:*  
*cumulative*  
*cost)*

Cost  
 contours



# Uniform Cost Search (UCS) Properties

- What nodes does UCS expand?

- Processes all nodes with cost less than cheapest solution!
- If that solution costs  $C^*$  and arcs cost at least  $\varepsilon$ , then the “effective depth” is roughly  $C^*/\varepsilon$
- Takes time  $O(b^{C^*/\varepsilon})$  (exponential in effective depth)

- How much space does the fringe take?

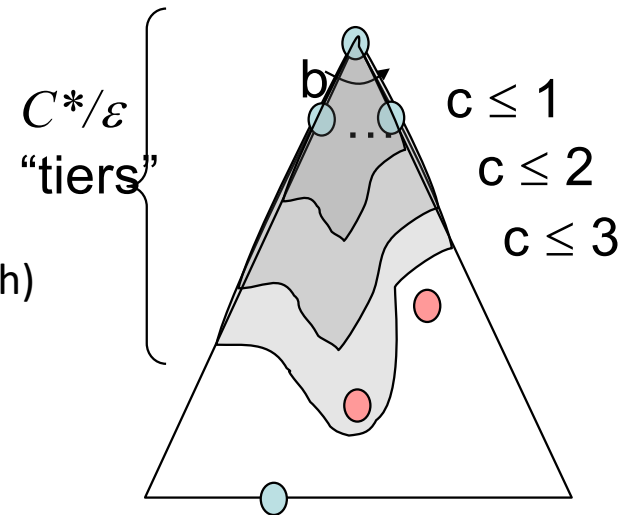
- Has roughly the last tier, so  $O(b^{C^*/\varepsilon})$

- Is it complete?

- Assuming best solution has a finite cost and minimum arc cost is positive, yes!

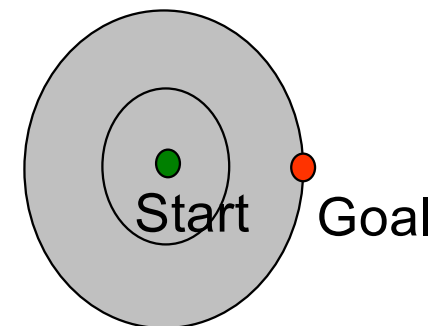
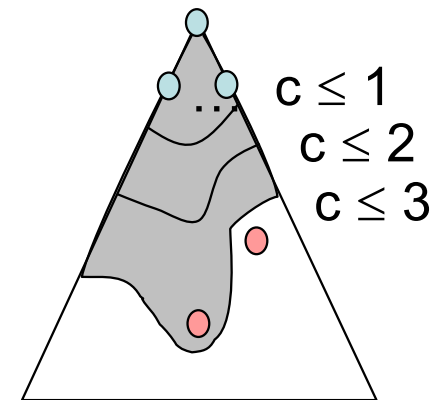
- Is it optimal?

- Yes! (Proof via A\*)



# Uniform Cost Issues

- Remember: UCS explores increasing cost contours
- The good: UCS is complete and optimal!
- The bad:
  - Explores options in every “direction”
  - No information about goal location
- We’ll fix that soon!



[Demo: empty grid UCS (L2D5)]  
[Demo: maze with deep/shallow water DFS/BFS/UCS (L2D7)]

# Video of Demo Empty UCS

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## Video of Demo Maze with Deep/Shallow Water --- DFS, BFS, or UCS? (part 1)

---



## Video of Demo Maze with Deep/Shallow Water --- DFS, BFS, or UCS? (part 2)

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# Video of Demo Maze with Deep/Shallow Water --- DFS, BFS, or UCS? (part 3)

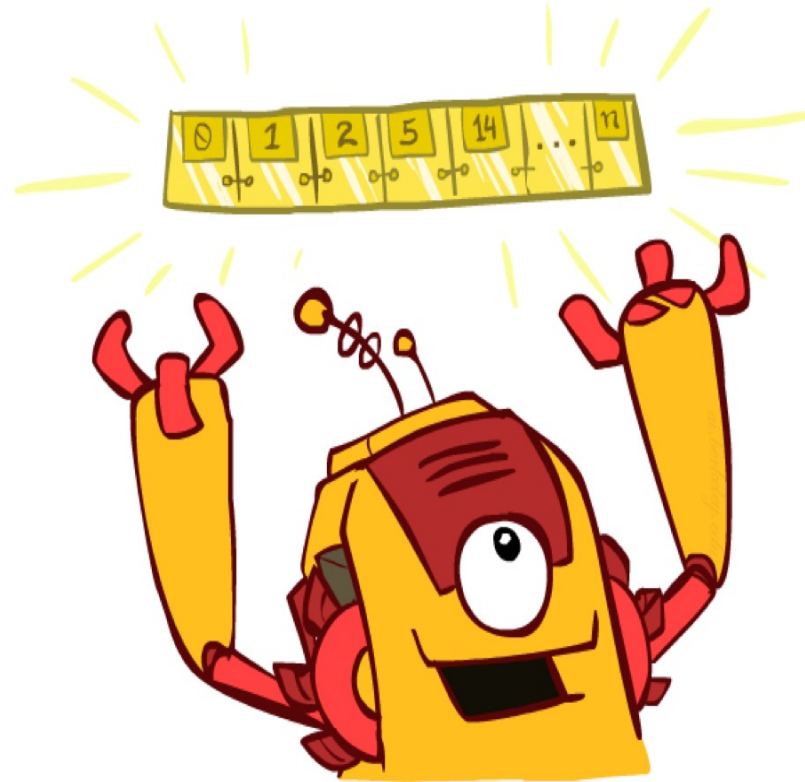
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# The One Queue

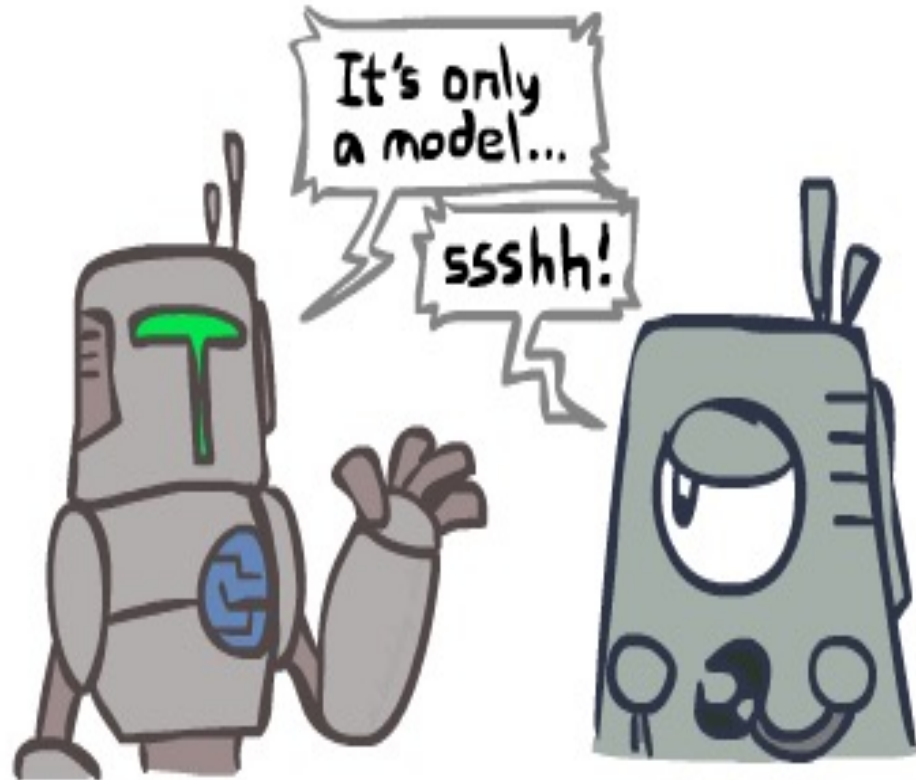
- All these search algorithms are the same except for fringe strategies
  - Conceptually, all fringes are priority queues (i.e. collections of nodes with attached priorities)
  - Practically, for DFS and BFS, you can avoid the  $\log(n)$  overhead from an actual priority queue, using stacks and queues
  - Can even code one implementation that takes a variable queuing object



# Search and Models

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- Search operates over models of the world
  - The agent doesn't actually try all the plans out in the real world!
  - Planning is all “in simulation”
  - Your search is only as good as your models...



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# Search

Prof. Abbeel steps through a few  
search examples

