Concentration Review (Ref: Chap 5, Bandit Algorithms). {x1, x2, . - - , xn, ... } sequence of iid 1-05 independent, identically distributed Law of large numbers: $\hat{\mathcal{H}}_{n} = \frac{1}{n} \sum_{i=1}^{n} \chi_{i} \longrightarrow \mathbb{E}[\chi]$ a.s. Notation: E[x]=u, Var(x)=02. Question: With n samples [x, ..., xn] how "close" is the empirical estimate of

the mean to the true mean?

(i.e., $\widehat{M}_{n} = \frac{1}{n} \sum_{i=1}^{n} \chi_{i}$ a γ_{i} (empirical estimate)

M = E[X] a fixed number.

 $P\left(\frac{1}{M_{n}-M}>\epsilon\right)\leq\frac{2\pi}{2\pi}$ $P\left(\frac{1}{M_{n}-M}>\epsilon\right)\leq\frac{2\pi}{2\pi}$ $\frac{1}{M_{n}-M}>\epsilon$ $\frac{1}{M_{n}-M}>\epsilon$

why do we care: In bandit algorithms, past sample of rewards allows us to estimate "goodness" of actions. Concentrations permit us to grandify how much we trust these estimates.

Basic inequalities Markov Inequality: X > 0 a non-negative v.v. $P(X > \varepsilon) \leq E[X]$ Chebyshev's Inequality: X a r.o. with E[X)=M. Var(X) = 62 $\mathbb{P}\left(\left|\chi-\mu\right|>\varepsilon\right)\leq\frac{\sigma^{2}}{c^{2}}.$ Chernoff Bound: X a r.v. with MGF Mx(0) Mx(0) = E[e0x] Moment Generating Function emish atlent. for o " Mome " Then, $P(X>E) \leq C$ E>E[X] $T(\varepsilon) = \max_{\theta} \left[\frac{1}{2} \log - \ln M_{x}(\theta) \right].$

rate function. = Legendre-Fenchel Transform

	U	
of-	the	10g-MGF

Most often, we will simplify and assume
noise terms are Subgaussian
Recall: XNN (M 62) Ranssian 1.0- ata Normal 1.0
aka Normal 1.0
distributed as "
"has the paf of?"
$ (x-y)^2 _{20^2}$
fx(x) = 1 e fx(x)
$f_{\chi}(\pi) = \frac{1}{\sqrt{2\pi\sigma^2}}$
σ^2
. ! d
For $Z \sim N(0, \sigma^2)$
$\frac{\partial^2 \sigma^2}{\partial \sigma^2} = \frac{\nabla(x > x + \epsilon)}{\nabla(x > x + \epsilon)}$
$F \left[\begin{array}{c} 0 \\ \overline{z} \end{array} \right] = P \left[\begin{array}{c} 1 \\ \overline{z} \end{array} \right]$
$E\left[e^{0Z}\right] = e^{0^2\sigma^2/2} \qquad P(x>x+\epsilon)$ $\leq e^{-\epsilon^2/2\sigma^2}$
MGF = Moment Generating Function
Val 216 (-) (1012 - 12 - 12 - 12 - 12 - 12 - 12 -
YN SubGanssian (o) (also J-subgaussian) if

(b)
$$E[e^{0Y}] \leq e^{2\delta^2/2}$$

Properties IN Subgaussian (5)

$$\begin{array}{ccc}
\bullet & P(\forall \neq \epsilon) \leq e^{-c^2/26^2}
\end{array}$$

for any S>0

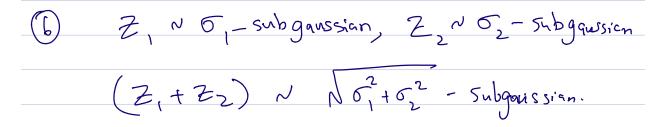
(3)
$$P\left(\left|\gamma\right| \geq \sqrt{2\sigma^2 \ln\left(2|\mathcal{S}\right)}\right) \leq S$$
 for

any 5>0.

$$(4)$$
 $Vor(1) \leq 6^2$

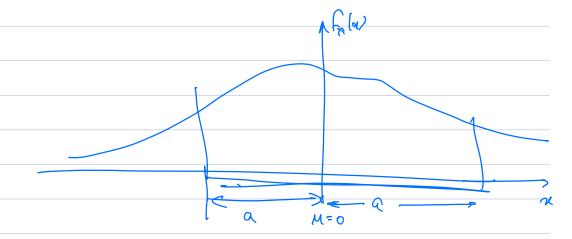
std-dev

(Proof: See Chapter 5, Bandit Algorithms).



Corrollary: $\chi \chi_i \chi_i = - subganssian, iid$

m = xi ~ I - subgaussian



 $P\left(\begin{array}{cc} \frac{1}{m} & \sum_{i=1}^{m} \chi_{i} > \varepsilon \\ \end{array}\right) \leq e^{\frac{2^{2}}{2\left(\frac{1}{\kappa_{m}}\right)}}$

 $\frac{-\varepsilon^2/2/m}{=}$

 $= e^{-\left(m\epsilon^2/2\right)}$