Solving two-person zero-sum sequential games via efficient computation of Nash equilibria: proof-of-concept on Kuhn's 3-card Poker

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### **Problem**

Given a sequential two-person zero-sum game (e.g Texas Hold'em Poker), construct an optimal player.

- There are 3 players: nature (player 0), Alice (player 1), and Bob (player 2).
- Players take turns in making moves
- Sets of moves of any two distinct players don't overlap (i.e  $A_p \cap A_q = \emptyset$  if  $p \neq q$ )

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## Strategy profiles, simplexes and complexes

#### Definition

• The *n-simplex*, denoted  $\Delta_n$ , is the subset of  $\mathbb{R}^{n+1}$  defined by

$$\Delta_n := \left\{ x \in \mathbb{R}^{n+1} \mid x_i \ge 0 \ \forall i, \text{ and } \sum_{i=1}^{n+1} x_i = 1 \right\}$$
 (1)

- Each kink  $\delta_i := (0, ..., 0, 1, 0, ..., 0)$  of  $\Delta_n$  represents a pure-strategy. Any non-kink point x of  $\Delta_n$  represents a mixed-strategy. Thus  $\Delta_n$  is simply the convex hall of (n+1)-dimensional pure strategies  $\delta_i$ .
- A complex is an abstract polyhedron whose kinks are themselves simplexes. Points on a complex correspond to behavioural strategies.

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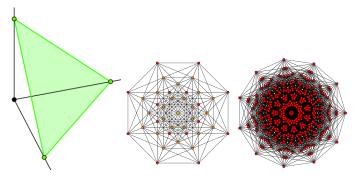
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# Simplexes and complexes (examples of)



Left: 2-simplex Middle: 7-simplex Right: 10-simplex

## NE minimax and best-response in sequential-games

Thanks to the sequence-form representation, we have:

NE problem for Alice and Bob is the the saddle-point problem

$$\max_{x \in Q_1} \min_{y \in Q_2} w^T A x, \tag{2}$$

where:

$$Q_1 := \{ x \in \mathbb{R}_+^n \mid Ex = e \} \text{ is Alice's complex, and }$$

$$Q_2 := \{ y \in \mathbb{R}_+^m \mid Fy = f \} \text{ is Bob's complex.}$$

$$(3)$$

A saddle-point  $(x^*, y^*)$  corresponds to a NE for the game.

• Given a fixed behavioural strategy  $y_0 \in Q_2$  for Bob, a best-response behavioural strategy  $x_0$  for Alice  $^1$  is a solution to the LCP

$$\max_{\mathbf{x} \in \mathcal{Q}_1} \max \mathbf{y}_0^T A \mathbf{x} \tag{4}$$

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The game tree T (from Alice's perspective) is defined as follows. T has a set V(T) of nodes (aka vertices) and a set E(T) of edges.

### Non-Linear Ledoit-Wolf: Why? How?

- (p/n large) or (Σ eigenvalues close to one another) ⇒ linear shrinkage OK.
   (p/n small) or (Σ eigenvalues dispersed) ⇒ non-linear shrinkage better.
- Marčenko-Pastur equation: relationship between  ${\bf S}$  and  ${\bf \Sigma}$  eigenvalues under large-dimensional asymptotics.

### Non-Linear Ledoit-Wolf: deeper into the "how?"

- $H_n$  empirical d.f. of sample  $(\boldsymbol{S}_n)$  eigenvalues  $\boldsymbol{\tau}_n$ .  $H_n(\boldsymbol{\tau}) \equiv \frac{1}{p} \sum_{i=1}^p \mathbb{1}_{[\tau_i, +\infty[}(\boldsymbol{\tau})$
- $F_n$  e.d.f. of population  $(\Sigma_n)$  eigenvalues  $\lambda_n$ .  $Y_n = X_n \Sigma_n^{1/2}$ ,  $X_n \sim \mathcal{N}(\mathbf{0}, \mathsf{Id})$
- $F_n$  (resp.  $H_n$ ) converges almost surely to F (resp. H).
- Stieltjes transform of a non-decreasing function G:  $\forall z \in \mathbb{C}^+$ ,  $m_G(z) \equiv \int_{-\infty}^{+\infty} \frac{1}{u-z} dG(u)$
- Marčenko-Pastur equation:  $\forall z \in \mathbb{C}^+, m_F(z) = \int_{-\infty}^{+\infty} \frac{1}{\tau[1-\frac{p}{2}-\frac{p}{n}zm_F(z)]-z} dH(\tau)$

### Non-Linear Ledoit-Wolf: Implementation...

$$\begin{aligned} & Q_{n,p} : [0, \infty[^p \to [0, \infty[^p \\ & \boldsymbol{t} \equiv (t_1, \dots, t_p)^\mathsf{T} \mapsto Q_{n,p}(\boldsymbol{t}) \equiv (q_{n,p}^1(t), \dots, q_{n,p}^p(t))^\mathsf{T} \end{aligned}$$

$$\hat{\boldsymbol{\tau}}_{\boldsymbol{n}} = \underset{\boldsymbol{t} \in [0, \infty[^p}{\operatorname{argmin}} \frac{1}{p} \sum_{i=1}^p [\boldsymbol{q}_{\boldsymbol{n}, \boldsymbol{p}}^i(\boldsymbol{t}) - \lambda_{\boldsymbol{n}, i}]^2$$

$$\overset{\leftarrow}{m}_{n,p}^{\hat{\tau}_n}(\lambda_i) \simeq \frac{n-p}{p\lambda_i} - \frac{n}{p} \frac{1}{\hat{\tau}_{n,i}\lambda_i}$$

$$\hat{m{d}}_i^* = rac{\lambda_i}{|1-rac{p}{c}-rac{p}{c}|^2 \widetilde{m{m}}_{n,n}^{\hat{m{n}}}(\lambda_i)};$$
 Covariance:  $S_n^* m{U}_n \hat{m{D}}_n^* m{U}_n^\mathsf{T}$ , with  $m{S}_n = m{U}_n m{D}_n m{U}_n^\mathsf{T}$ 

# About the $Q_{n,p}$ function

$$Q_{n,p}:[0,\infty[^p \to [0,\infty[^p$$

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$$\forall i=1,\ldots,p \quad q_{n,p}^i(t) \equiv p \int_{(i-1)/p}^{i/p} (F_{n,p}^t)^{-1}(u) du$$

$$(F_{n,p}^t)^{-1}(u) \equiv \sup\{x \in \mathbb{R} : F_{n,p}^t(x) \le u\}$$

$$\begin{split} F_{n,p}^{t}(x) &\equiv \max\left\{1-\frac{n}{p},\frac{1}{p}\sum_{i=1}^{p}\mathbb{1}_{\{t_i=0\}}\right\} \text{ if } x=0,\\ F_{n,p}^{t}(x) &\equiv \lim_{\eta\to 0^+}\frac{1}{\pi}\int_{-\infty}^{+\infty}\operatorname{Im}\left[\frac{\mathbf{m}_{n,p}^{t}(\zeta+i\,\eta)}{\mathbf{m}_{n,p}^{t}(\zeta+i\,\eta)}\right]d\zeta \text{ otherwise} \end{split}$$

$$m \equiv m_{n,p}^t(z)$$
 sol. in  $\left\{m \in \mathbb{C}: -\frac{n-p}{nz} + \frac{p}{n}\right\}$  of  $m = \frac{1}{p} \sum_{i=1}^p \frac{1}{t_i(1-\frac{p}{n}-\frac{p}{n}zm)-z}$