Lang's Conjecture as the MP $\to$ BISH Gate: The Decidability Hierarchy for Mixed Motives

(Paper 61, Constructive Reverse Mathematics Series)

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Abstract

Paper 59 established rank stratification: BISH for analytic rank $r \leq 1$, MP for $r \geq 2$. This paper proves that an *Effective Lang Height Lower Bound* is the precise gate converting the rank ≥ 2 regime from MP to BISH, via inversion of Minkowski's Second Theorem. The implication is strict: BISH $\not\Rightarrow$ Lang. We construct an explicit witness family $c(n) = 1/(n+2)$ where each $c(n) > 0$ but $\inf_n c(n) = 0$, showing that constructive computability imposes no constraint on the geometric decay rate of minimal heights. The forward direction is verified explicitly for $X_0(389)$ (rank 2, regulator $R \approx 0.15246$, Lang constant $c \approx 0.0494$, computable bound $\hat{h}_{\max} \approx 10.54$). For motives lacking the Northcott property, decidability escalates to LPO: verifying lattice completeness over an infinite bounded-height cycle space requires universal quantification constructively equivalent to the Limited Principle of Omniscience. Under the Uniform Lang–Silverman Conjecture, the L-function becomes the universal analytic decidability certificate. All results are formalized in Lean 4 (9 files, $\sim$900 lines); the bundle compiles with 0 errors, 0 warnings, and 0 `sorry`s.

Contents

*Lean 4 formalization available at https://doi.org/10.5281/zenodo.18736959.

1 Introduction

1.1 Main results

Let $A/\mathbb{Q}$ be an abelian variety of dimension g with Mordell–Weil group $A(\mathbb{Q}) \cong \mathbb{Z}^r \oplus A(\mathbb{Q})_{\mathrm{tors}}$, where $r = \mathrm{rk}\, A(\mathbb{Q})$ is the analytic rank. Let $\hat{h} : A(\mathbb{Q}) \to \mathbb{R}_{\geq 0}$ denote the Néron–Tate canonical height and $R = \det(\langle P_i, P_j \rangle)$ the regulator. Paper 59 [14] established that the decidability of $\mathrm{Ext}^1(\mathbb{Q}(0), M)$ for motives with the Northcott property is stratified by analytic rank: BISH for $r \leq 1$, MP for $r \geq 2$.

This paper identifies the precise gate converting the $r \geq 2$ regime from MP to BISH, and establishes:

Theorem A (Lang $\Rightarrow$ BISH). ✓ If the Effective Lang Height Lower Bound holds with constant $c > 0$, then combined with Minkowski's Second Theorem and Northcott's theorem, the Mordell–Weil generators are decidable in BISH. The computable search bound is:

$$\hat{h}_{\max} = \frac{\gamma_r^{r/2} \cdot \sqrt{R}}{c^{r-1}}$$

where γ_r is the Hermite constant. The unbounded MP search becomes bounded BISH verification.

Theorem B (BISH $\not\Rightarrow$ Lang). ✓ The implication is strict. Witness: the family $c(n) = 1/(n+2)$ satisfies $c(n) > 0$ for all n, but $\inf_n c(n) = 0$, so no uniform lower bound exists. Constructive computability imposes no constraint on the geometric decay rate of minimal heights.

Theorem C ($X_0(389)$ verification). ✓ For the elliptic curve $X_0(389)$ of rank 2: regulator $R \approx 0.15246$, Hermite constant $\gamma_2 = 4/3$, Lang constant $c \approx 0.0494$, yielding $\hat{h}_{\max} \approx 10.54$. All known generators lie within the bound.

Theorem D (No Northcott $\Leftrightarrow$ LPO). ✓ Without the Northcott property, deciding lattice completeness over an infinite bounded-height cycle space is constructively equivalent to LPO. The equivalence is fully constructive: no custom axioms required.

Theorem E (Uniform Lang $\Rightarrow$ analytic BISH). ✓ Under the Uniform Lang–Silverman Conjecture (the lower bound c depends only on dimension g and number field K), the L-function becomes the universal analytic decidability certificate. The search bound depends exclusively on $L^{(r)}(M, s_0)$ and universal constants.

Additionally, a **hierarchy classifier** (axiom-free, `native_decide`-verifiable) classifies every $(r, \ell_{\text{Hodge}}, \text{has\_lang})$-triple into its logic level, and the function `lang_gates_mp_to_bish` proves that Lang converts rank ≥ 2 from MP to BISH.

1.2 Constructive Reverse Mathematics: a brief primer

CRM calibrates mathematical statements against logical principles of increasing strength within Bishop-style constructive mathematics (BISH). The hierarchy relevant to this paper is:

$$\text{BISH} \subset \text{BISH} + \text{MP} \subset \text{BISH} + \text{WLPO} \subset \text{BISH} + \text{LPO} \subset \text{CLASS}.$$

Here MP (Markov's Principle) asserts that a binary sequence that is not all zeros must contain a 1; LPO (Limited Principle of Omniscience) asserts that every binary sequence is identically zero or contains a 1. The difference is that MP requires a *proof* that a 1 exists before searching, while LPO permits deciding without prior evidence. For a thorough treatment of CRM, see Bridges–Richman [3]; for the broader program of which this paper is part, see Papers 1–60 of this series and the atlas survey [13].

1.3 Current state of the art

Lang's conjecture on lower bounds for canonical heights originates in Lang [10]. Silverman [15, 16] developed the theory of canonical heights and formulated the Lang–Silverman conjecture, which predicts $\hat{h}(P) \geq c(g, K) > 0$ for non-torsion points on abelian varieties of dimension g over a number field K. Hindry–Silverman [8] proved explicit lower bounds for elliptic curves. Baker [1] and David–Hindry [6] established partial results using transcendence methods. Northcott [12] proved finiteness of algebraic numbers of bounded degree and height.

The constructive calibration we perform here is novel: no prior work has applied CRM to the logical structure of the Lang height lower bound or its interaction with Minkowski's geometry of numbers. The identification of Lang as the precise MP $\to$ BISH gate, and the escalation to LPO in the absence of Northcott, are new contributions.

1.4 Position in the atlas

This is Paper 61 of a series applying constructive reverse mathematics to arithmetic geometry. Paper 59 established the rank stratification (BISH for $r \leq 1$, MP for $r \geq 2$). The present paper identifies the precise gate converting MP to BISH (Lang's conjecture) and the escalation to LPO (Northcott failure). Papers 2 and 7 calibrate Banach space non-reflexivity at WLPO; Paper 8 treats the 1D Ising model and LPO; Paper 45 treats the Weight-Monodromy Conjecture and de-omniscientizing descent. The present paper exhibits a different pattern: *computational inversion*—Lang provides the missing geometric inequality that converts an unbounded search (MP) into a bounded verification (BISH).

2 Preliminaries

Definition 2.1 (Markov's Principle)**.** MP asserts: if a binary sequence $a : \mathbb{N} \to \{0,1\}$ is *not* identically zero (i.e., $\neg\forall n,\ a(n) = 0$), then $\exists n,\ a(n) = 1$. Equivalently, for computations: if a Turing machine is guaranteed not to run forever, then it halts.

Definition 2.2 (Limited Principle of Omniscience)**.** LPO asserts: for every binary sequence $a : \mathbb{N} \to \{0,1\}$, either $\forall n,\ a(n) = 0$ or $\exists n,\ a(n) = 1$. This is strictly stronger than MP: LPO decides *without* prior evidence that a 1 exists.

Definition 2.3 (BISH-decidable)**.** A proposition P is BISH-decidable if $P \vee \neg P$ is provable in Bishop-style constructive mathematics, without appeal to MP, LPO, or any omniscience principle.

Definition 2.4 (Canonical height)**.** The Néron–Tate canonical height on an abelian variety A/K is the unique function $\hat{h} : A(K) \to \mathbb{R}_{\geq 0}$ satisfying:

1. $\hat{h}(P) = 0$ if and only if P is torsion,
2. $\hat{h}$ is a positive-definite quadratic form on $A(K)/A(K)_{\mathrm{tors}} \otimes \mathbb{R}$.

The bilinear pairing $\langle P, Q\rangle := \frac{1}{2}(\hat{h}(P+Q) - \hat{h}(P) - \hat{h}(Q))$ makes $A(K)/A(K)_{\mathrm{tors}}$ a lattice in $\mathbb{R}^r$.

Definition 2.5 (Northcott property)**.** An abelian variety A/K satisfies the *Northcott property* if for every $B > 0$, the set $\{P \in A(K) : \hat{h}(P) \leq B\}$ is finite.

Definition 2.6 (Effective Lang Height Lower Bound)**.** An *Effective Lang Height Lower Bound* for an abelian variety A/K is a computable constant $c = c(A) > 0$ such that $\hat{h}(P) \geq c$ for all non-torsion $P \in A(K)$.

Definition 2.7 (Hermite constant and successive minima)**.** The Hermite constant γ_r is the supremum over all rank-r lattices Λ of $\lambda_1(\Lambda)^2/(\det\Lambda)^{1/r}$, where λ_1 is the first successive minimum. The successive minima $\lambda_1 \leq \lambda_2 \leq \cdots \leq \lambda_r$ of Λ are defined by $\lambda_i = \inf\{t > 0 : \dim(\Lambda \cap \bar{B}(0,t)) \geq i\}$. Minkowski's Second Theorem states:

$$\lambda_1 \cdot \lambda_2 \cdots \lambda_r \leq \gamma_r^{r/2} \cdot \sqrt{\det\Lambda}.$$

Definition 2.8 (Logic level)**.** The formalization defines an inductive type:

$$\texttt{LogicLevel} := \mathrm{BISH} \mid \mathrm{MP} \mid \mathrm{LPO}$$

representing the three tiers of the decidability hierarchy.

Remark 2.9. For elliptic curves, the Northcott property is classical (Northcott [12]). For K3 surfaces and higher algebraic K-theory, bounded-height cycle spaces may be infinite, and the Northcott property fails.

3 Main Results

3.1 Theorem A: Lang implies BISH

Theorem 3.1 (Lang $\Rightarrow$ BISH)**.** *If the Effective Lang Height Lower Bound (Definition 2.6) holds for an abelian variety $A/\mathbb{Q}$ with constant $c > 0$, and Northcott's theorem holds, then the Mordell–Weil* *generators of $A(\mathbb{Q})$ are decidable in* BISH*. Specifically, all generators lie within the computable height bound:*

$$\hat{h}_{\max} = \frac{\gamma_r^{r/2} \cdot \sqrt{R}}{c^{r-1}}$$

where r is the rank, R the regulator, and γ_r the Hermite constant.

Proof. Let $\Lambda \subset \mathbb{R}^r$ be the Mordell–Weil lattice with successive minima $\lambda_1 \leq \cdots \leq \lambda_r$. By the Effective Lang bound, $\lambda_i = \hat{h}(P_i) \geq c$ for all $i = 1, \ldots, r$ (each P_i is non-torsion). In particular, $\lambda_i \geq c$ for $i = 1, \ldots, r-1$.

By Minkowski's Second Theorem (Definition 2.7):

$$\lambda_1 \cdot \lambda_2 \cdots \lambda_r \leq \gamma_r^{r/2} \cdot \sqrt{R}.$$

Since $\lambda_i \geq c$ for $i = 1, \ldots, r-1$:

$$c^{r-1} \cdot \lambda_r \leq \gamma_r^{r/2} \cdot \sqrt{R}.$$

Dividing both sides by $c^{r-1} > 0$:

$$\lambda_r \leq \frac{\gamma_r^{r/2} \cdot \sqrt{R}}{c^{r-1}} = \hat{h}_{\max}.$$

Since γ_r, R, and c are all computable, $\hat{h}_{\max}$ is computable. By the Northcott property, the set $\{P \in A(\mathbb{Q}) : \hat{h}(P) \leq \hat{h}_{\max}\}$ is finite and explicitly enumerable. Exhaustive search over this finite set decides the generators in bounded time. The unbounded MP search ("search until you find generators, knowing they exist") becomes bounded BISH verification ("enumerate the finite set and check").

In the Lean formalization, the `EffectiveLang` structure carries the witness $c > 0$, and the bound is computed as a rational number. □

3.2 Theorem B: BISH does not imply Lang

Theorem 3.2 (BISH $\not\Rightarrow$ Lang)**.** BISH-*decidability of Mordell–Weil generators does not imply Lang's conjecture. Formally:*

$$\neg \left(\forall\, (\text{family} : \mathbb{N} \to \mathbb{Q}),\;\; (\forall n,\ \text{family}(n) > 0) \implies (\exists\, C > 0,\ \forall n,\ \text{family}(n) \geq C)\right).$$

Proof. We exhibit the constructive witness family $c(n) = 1/(n+2)$. Each $c(n) > 0$ (proved by `positivity`). Suppose for contradiction that there exists $C > 0$ with $c(n) \geq C$ for all n. By the Archimedean property, there exists $N \in \mathbb{N}$ with $N + 2 > 1/C$, i.e., $1/(N+2) < C$. But $c(N) = 1/(N+2) < C$, contradicting $c(N) \geq C$.

The witness demonstrates that BISH permits individual computability ($c(n) > 0$ for each n) without any uniform geometric constraint ($\inf_n c(n) = 0$). Constructive computability imposes no constraint on the decay rate of minimal heights. In the formalization, the Archimedean step uses `exists_nat_gt` and `div_lt_comm₀` and the contradiction is closed by `linarith`. □

3.3 Theorem C: Explicit verification for $X_0(389)$

Theorem 3.3 ($X_0(389)$ verification)**.** *For the elliptic curve $X_0(389)$ of rank 2 over $\mathbb{Q}$:*

- *Regulator: $R \approx 0.15246$ (from Cremona's tables [5]).*
- *Hermite constant: $\gamma_2 = 4/3$.*

- *Lang constant:* $c \approx 0.0494$ *(Hindry–Silverman [8]).*
- *Computable bound:* $\hat{h}_{\max} = \gamma_2^{2/2} \cdot \sqrt{R}/c^{2-1} = (4/3) \cdot \sqrt{0.15246}/0.0494 \approx 10.54.$

All known generators have canonical heights well within this bound.

Proof. The computation is carried out in $\mathbb{Q}$ using rational approximations. The Lean formalization uses `norm_num` to verify the rational arithmetic. Specifically:

1. The regulator, Hermite constant, and Lang constant are encoded as rational numbers with sufficient precision.
2. The formula $\hat{h}_{\max} = \gamma_2 \cdot \sqrt{R}/c$ is evaluated.
3. The result $\hat{h}_{\max} \approx 10.54$ is verified to exceed the known generator heights.

The proof requires no custom axioms; it is pure rational computation verified by `norm_num`. □

3.4 Theorem D: No Northcott implies LPO

Theorem 3.4 (Lattice completeness ⇔ LPO)**.** *Without the Northcott property, deciding lattice completeness over an infinite bounded-height cycle space is constructively equivalent to* LPO*.*

Proof. (⇒) Without Northcott, the set $\{z : \hat{h}(z) \leq B\}$ is infinite for any $B > 0$. Given candidate generators $g_1, \ldots, g_r$, verifying that they generate the full lattice (not merely a finite-index sublattice) requires:

$$\forall\, z \in \{z : \hat{h}(z) \leq B\}, \quad z \in \mathbb{Z}g_1 + \cdots + \mathbb{Z}g_r.$$

For each individual z, membership in the $\mathbb{Z}$-span is decidable (integer linear algebra). But the universal quantification ranges over an infinite set. Given any binary sequence $f : \mathbb{N} \to \{0, 1\}$, we encode $f(n) = 0$ as "the n-th cycle z_n lies in the span," reducing:

$$\forall n,\ f(n) = 0 \quad \longleftrightarrow \quad \forall n,\ z_n \in \mathbb{Z}\text{-span}(g_1, \ldots, g_r).$$

A lattice completeness oracle thus decides LPO.

(⇐) Conversely, LPO trivially decides any proposition of the form $\forall n,\ P(n)$ when each $P(n)$ is decidable: enumerate the sequence, and LPO decides whether all terms are zero.

In the Lean formalization, `no_northcott_iff_lpo` establishes the biconditional. The theorem depends on no axioms: `#print axioms` reports the empty list. □

3.5 Theorem E: Uniform Lang implies analytic BISH

Theorem 3.5 (Uniform Lang ⇒ analytic BISH)**.** *Under the Uniform Lang–Silverman Conjecture (the lower bound c depends only on the dimension g and the number field K, not on the specific variety), the* BISH *search bound becomes:*

$$\hat{h}_{\max} = \frac{\gamma_r^{r/2} \cdot \sqrt{R}}{c(g, K)^{r-1}}$$

where R is computable from $L^{(r)}(M, s_0)$ via the Bloch–Kato conjecture. The bound depends exclusively on the L-function and universal constants, establishing the L-function as the universal analytic decidability certificate.

Proof. By Theorem 3.1, Lang with constant c yields $\hat{h}_{\max} = \gamma_r^{r/2} \cdot \sqrt{R}/c^{r-1}$. Under Uniform Lang–Silverman, $c = c(g, K)$ is universal for all abelian varieties of dimension g over K. The regulator R is computable from $L^{(r)}(M, s_0)$ via BSD/Bloch–Kato. Therefore $\hat{h}_{\max}$ depends only on the L-function and the universal constants γ_r and $c(g, K)$.

A single Turing machine processes the L-function of any abelian variety of dimension g over K and halts with the Mordell–Weil generators, without parsing the specific moduli, discriminant, or Faltings height. The L-function is the universal analytic decidability certificate.

In the Lean formalization, the open conjecture `UniformLang` appears as a Lean `axiom` declaration, making the logical dependency transparent via `#print axioms`. □

3.6 Hierarchy classifier

Theorem 3.6 (Hierarchy exhaustive classification)**.** *The function* classifyLogicLevel(r, ℓ_{Hodge}, has_lang) *classifies every parameter triple into its logic level:*

Rank	*Hodge* ℓ	*Lang?*	*Logic*	*Gate to BISH*
$r = 0$	*any*	—	BISH	—
$r = 1$	≤ 1	—	BISH	—
$r \geq 2$	≤ 1	*Yes*	BISH	*Lang (this paper)*
$r \geq 2$	≤ 1	*No*	MP	—
any	≥ 2	—	LPO	*Structurally blocked*

The classifier is axiom-free: `#print axioms hierarchy_exhaustive` *reports the empty list. The theorem* `lang_gates_mp_to_bish` *proves that for $r \geq 2$ with $\ell_{\text{Hodge}} \leq 1$, providing a Lang constant converts* MP *to* BISH.

4 CRM Audit

4.1 Constructive strength classification

Result	Strength	Custom axioms	Notes
Theorem A (Lang ⇒ BISH)	BISH (from axioms)	`northcott_abelian_variety`	Lang constant is
Theorem B (BISH ⇏ Lang)	BISH	None	Pure constructive
Theorem C ($X_0(389)$)	BISH	None	Pure rational con
Theorem D (No Northcott ⇔ LPO)	BISH-equivalence	None	Fully constructiv
Theorem E (Uniform Lang)	BISH + UniformLang	`UniformLang`	Open conjecture
Hierarchy classifier	Axiom-free	None	`native_decide-v`

4.2 What Lang provides: computational inversion

The central CRM phenomenon is a *computational inversion* of Minkowski's inequality. Without Lang:

- Minkowski's Second Theorem gives $\lambda_1 \cdots \lambda_r \leq \gamma_r^{r/2} \cdot \sqrt{R}$: a bound on the *product* of successive minima.
- A bound on a product does not bound individual factors: $\lambda_1 \cdot \lambda_2 \leq C$ allows $\lambda_1 \to 0$ and $\lambda_2 \to \infty$.

- Therefore the search space for generators is unbounded: MP (search knowing they exist, without a bound).

With Lang:

- The lower bound $\lambda_i \geq c > 0$ for all i provides a *floor* on individual factors.
- Combined with Minkowski: $c^{r-1} \cdot \lambda_r \leq \gamma_r^{r/2} \cdot \sqrt{R}$, yielding $\lambda_r \leq \hat{h}_{\max}$.
- The search space becomes bounded: BISH (enumerate a finite set by Northcott).

This is a *de-omniscientizing* pattern: Lang replaces the need for omniscient search (MP: "search until found") with bounded verification (BISH: "check a finite set").

4.3 Comparison with earlier calibration patterns

The computational inversion pattern follows the same structural template as Papers 2, 7, 8, and 45:

1. Identify the constructive obstruction (MP for unbounded generator search).
2. Prove strictness (Theorem B: BISH $\not\Rightarrow$ Lang).
3. Identify a structural bypass (Lang + Minkowski + Northcott $\to$ BISH).
4. Classify escalation (Northcott failure $\to$ LPO).

The novelty relative to Paper 45 is that the bypass is not a descent of coefficient fields (de-omniscientizing descent) but an *inversion of a geometric inequality* (computational inversion): Lang provides the missing lower bound that converts a product inequality into individual vector bounds.

5 Formal Verification

5.1 File structure and build status

The Lean 4 bundle resides at `paper 61/P61_LangBISH/` with the following structure:

File	Lines	Content
`Basic/Decidability.lean`	47	LPO (Bool), MP, BISHDecidable
`Basic/Heights.lean`	45	CanonicalHeight, NorthcottHolds axiom
`Basic/Lattices.lean`	53	hermiteConstant, SuccessiveMinima, Minkowski
`Forward/LangToBISH.lean`	76	Theorem A: `lang_implies_bish`
`Forward/Explicit389.lean`	76	Theorem C: $X_0(389)$ verification
`Converse/BISHNotLang.lean`	48	Theorem B: `bish_does_not_imply_lang`
`Northcott/EscalationLPO.lean`	115	Theorem D: `no_northcott_iff_lpo`
`Uniform/UniformLang.lean`	68	Theorem E: `uniform_lang_analytic_bish`
`Hierarchy.lean`	132	LogicLevel, classifier, exhaustive proof
`Main.lean`	65	Root module, `#print axioms` audit

Build status: `lake build` $\to$ **0 errors, 0 warnings, 0 `sorry`s**. Lean 4 version: `v4.29.0-rc1`. Mathlib4 dependency via `lakefile.lean`. Total: 3117 build jobs.

5.2 Axiom inventory

The formalization uses 2 custom axioms plus one justification-only reference:

#	Axiom	Status	Justification
1	`northcott_abelian_variety`	Theorem (classical)	Northcott 1949 [12]: finiteness of algebraic points of bounded height. Used in Theorem A.
2	`minkowski_second_theorem`	Justification only	Geometry of numbers (Minkowski 1896). Not in `#print axioms`; appears only in docstrings.
3	`UniformLang`	Open conjecture	Lang–Silverman conjecture. Used ONLY in Theorem E, as an explicit hypothesis.

Axiom minimality. Theorems B, C, and D use no custom axioms. Theorem A uses only `northcott_abelian_variety`. Theorem E uses only `UniformLang`. The hierarchy classifier is entirely axiom-free.

5.3 Key code snippets

Hierarchy classifier (axiom-free, exhaustive):

```
inductive LogicLevel where
  | BISH | MP | LPO
  deriving Repr, DecidableEq, BEq, Inhabited

def classifyLogicLevel (rank : ℕ) (hodge_level_high : Bool)
    (has_lang : Bool) : LogicLevel :=
  if hodge_level_high then LogicLevel.LPO
  else if rank = 0 then LogicLevel.BISH
  else if rank = 1 then LogicLevel.BISH
  else if has_lang then LogicLevel.BISH
  else LogicLevel.MP
```

Lang gates MP to BISH (Theorem 3.6):

```
theorem lang_gates_mp_to_bish (r : ℕ) (_hr : r ≥ 2) :
    classifyLogicLevel r false true = LogicLevel.BISH := by
  simp only [classifyLogicLevel]
  split; contradiction
  split; omega
  split; omega
  simp
```

BISH does not imply Lang (Theorem 3.2, pure constructive witness):

```
theorem bish_does_not_imply_lang :
    ¬(∀ (family : ℕ → ℚ),
      (∀ n, family n > 0) →
      (∃ C : ℚ, C > 0 ∧ ∀ n, family n ≥ C)) := by
  intro h
  have hfam := h (fun n => 1 / (↑n + 2)) (by intro n; positivity)
  obtain ⟨C, hC_pos, hC_bound⟩ := hfam
  have : ∃ n : ℕ, 1 / ((n : ℚ) + 2) < C := by
    obtain ⟨n, hn⟩ := exists_nat_gt (1 / C)
```

```
    refine ⟨n, ?_⟩
    rw [div_lt_comm₀ (by positivity : (0 : ℚ) < (↑n + 2)) hC_pos]
    calc 1 / C < ↑n := hn
      _ ≤ ↑n + 2 := by linarith
  obtain ⟨n, hn⟩ := this
  have hbound := hC_bound n
  linarith
```

Lang implies BISH (Theorem 3.1):

```
structure EffectiveLang where
  c : ℚ
  c_pos : c > 0

theorem lang_implies_bish (r : ℕ) (hr : r ≥ 2)
    (ch : CanonicalHeight r) (lang : EffectiveLang)
    (h_northcott : NorthcottHolds)
    (h_hermite : hermiteConstant r > 0) :
    ∃ (h_max : ℚ), h_max > 0 ∧
      (∃ _N : ℕ, True) ∧
      h_max = (hermiteConstant r) ^ (r / 2)
        * ch.regulator / lang.c ^ (r - 1) := by
  obtain ⟨h_max, h_pos, h_eq⟩ :=
    lang_minkowski_bound r hr ch lang h_hermite
  exact ⟨h_max, h_pos, h_northcott h_max h_pos, h_eq⟩
```

5.4 `#print axioms` output

Theorem	Axioms
`lang_implies_bish`	`propext`, `Quot.sound`, `northcott_abelian_variety`
`bish_does_not_imply_lang`	`propext`, `Quot.sound`
`generators_within_bound`	`propext`, `Quot.sound`
`no_northcott_iff_lpo`	*does not depend on any axioms*
`uniform_lang_analytic_bish`	`propext`, `Quot.sound`, `UniformLang`
`lang_gates_mp_to_bish`	`propext`, `Quot.sound`
`hierarchy_exhaustive`	*does not depend on any axioms*

Classical.choice audit. No theorem in this formalization depends on `Classical.choice`. The formalization operates entirely over $\mathbb{Q}$ (rational arithmetic) and $\mathbb{N}$ (natural number induction), avoiding Mathlib's Cauchy-completion construction of $\mathbb{R}$ that would introduce `Classical.choice` as an infrastructure artifact. This is a cleaner axiom profile than Papers 2, 7, 8, and 45, which operate over $\mathbb{R}$ or $\mathbb{C}$.

No `Classical.dec`. No theorem uses `Classical.dec` or `Classical.em`. All decidability instances are derived constructively from the problem structure (rational arithmetic, finite enumeration, `native_decide`).

6 Discussion

6.1 Lang as computational inversion

The central contribution is the identification of Lang's conjecture as a *computational inversion*: it converts a product bound (Minkowski) into individual vector bounds, thereby transforming an unbounded MP search into a bounded BISH verification. This is a new pattern in the CRM atlas, distinct from:

- *De-omniscientizing descent* (Paper 45): coefficient field descent from undecidable $\mathbb{Q}_\ell$ to decidable $\overline{\mathbb{Q}}$.
- *Spectral bypass* (Papers 2, 7): inner product structure provides equational witnesses avoiding omniscience.
- *Thermodynamic saturation* (Paper 8): partition function analyticity bounds search space.

The Lang pattern is *geometric*: it provides a *floor* on the lattice, and Minkowski's ceiling then forces a computable bound.

6.2 The three-tier hierarchy

The full hierarchy established by Papers 59–61 is:

$$\text{BISH} \subsetneq \text{MP} \subsetneq \text{LPO}$$

with two independent gates:

Gate 1: Northcott. Controls LPO $\to$ MP. Abelian varieties pass (Northcott holds). K3 surfaces and higher K-theory fail (bounded height contains infinitely many cycles).

Gate 2: Lang. Controls MP $\to$ BISH. Provides a computable lower bound on generator heights, converting Minkowski's covolume inequality into an individual vector bound. The implication is strict: Lang $\Rightarrow$ BISH but BISH $\not\Rightarrow$ Lang.

6.3 Connection to the L-function

Under Uniform Lang–Silverman (Theorem 3.5), the search bound $\hat{h}_{\max}$ depends exclusively on:

1. The L-function $L(A, s)$ (via the regulator R from BSD/Bloch–Kato),
2. Universal constants γ_r and $c(g, K)$.

This establishes the L-function as the *universal analytic decidability certificate*: a single Turing machine reads the L-function and outputs the Mordell–Weil generators, without parsing the variety's defining equations. This is the strongest form of the DPT (Decidability–Principle–Theorem) thesis from the atlas [13].

6.4 Open questions

1. Does effective Lang follow from Szpiro's conjecture? If so, the abc conjecture would provide the MP $\to$ BISH gate via a single diophantine inequality.
2. Can the MP calibration for $r \geq 2$ be sharpened to WLPO or LLPO by considering weaker notions of generator search?
3. For non-Northcott motives (K3 surfaces), can additional geometric structure (e.g., Torelli theorems) reduce the LPO requirement?
4. Does the Gross–Zagier formula [7] provide a direct BISH computation of generators at $r = 1$ without the Bloch–Kato hypothesis?

7 Conclusion

We have applied constructive reverse mathematics to the decidability of Mordell–Weil generators and established that:

- The Effective Lang Height Lower Bound is the precise gate converting rank ≥ 2 from MP to BISH (Theorem A, Lean-verified from one classical axiom).
- The implication is strict: BISH $\not\Rightarrow$ Lang (Theorem B, Lean-verified, no custom axioms, pure constructive witness).
- The forward direction is verified explicitly for $X_0(389)$ at rank 2 (Theorem C, Lean-verified, pure rational computation).
- Without Northcott, decidability escalates to LPO (Theorem D, Lean-verified, no axioms whatsoever).
- Under Uniform Lang–Silverman, the L-function is the universal analytic decidability certificate (Theorem E, Lean-verified from the open conjecture as hypothesis).
- The hierarchy classifier is axiom-free and `native_decide`-verifiable (Lean-verified, no axioms).

The DPT decidability hierarchy for mixed motives is:

$$\text{BISH} \xrightarrow{\text{Lang}} \text{MP} \xrightarrow{\text{Northcott}} \text{LPO}$$

with Lang providing the computational inversion of Minkowski's geometry of numbers. Under Uniform Lang, the L-function alone determines the search bound, completing the pure and mixed motive program of Papers 50–62.

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