

A Proof of Dual Replacement Lemma

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July 4, 2018

Let φ be a formula in basic modal and we denote φ' as the formula obtained by replacing a “ \Diamond ” inside φ with “ $\neg\Box\neg$ ”. We claim that

$$\vdash_K \varphi \leftrightarrow \varphi'.$$

Proof. We shall prove this by structural induction on φ .

Case 1: $\varphi = \Diamond\psi$ and we are replacing the first \Diamond , then $\varphi' = \neg\Box\neg\psi$.

$$\begin{aligned} (1) \quad & \Diamond p \leftrightarrow \neg\Box\neg p && \text{DUAL} \\ (2) \quad & \Diamond\varphi \leftrightarrow \neg\Box\neg\varphi && \text{US:(1)} \end{aligned}$$

Case 2: $\varphi = \neg\psi$, $\varphi' = \neg\psi'$. We informally claim that the definition of ψ' is clear: For this case, if φ' replaces the i th \Diamond , ψ' does the i th.

$$\begin{aligned} (1) \quad & (p \leftrightarrow q) \rightarrow (\neg p \leftrightarrow \neg q) && \text{TAUT} \\ (2) \quad & (\psi \leftrightarrow \psi') \leftrightarrow (\neg\psi \leftrightarrow \neg\psi') && \text{US:(1)} \\ (3) \quad & \psi \leftrightarrow \psi' && \text{I.H.} \\ (4) \quad & \neg\psi \leftrightarrow \neg\psi' && \text{MP(2)(3)} \end{aligned}$$

Case 3: $\varphi = \psi * \rho$, $\varphi' = \psi' * \rho$, where $*$ $\in \{\wedge, \vee, \rightarrow, \leftrightarrow\}$.

$$\begin{aligned} (1) \quad & (p \leftrightarrow q) \rightarrow ((p * r) \leftrightarrow (q * r)) && \text{TAUT} \\ (2) \quad & (\psi \leftrightarrow \psi') \rightarrow ((\psi * \rho) \leftrightarrow (\psi' * \rho)) && \text{US:(1)} \\ (3) \quad & \psi \leftrightarrow \psi' && \text{I.H.} \\ (4) \quad & (\psi * \rho) \leftrightarrow (\psi' * \rho) && \text{MP(2)(3)} \end{aligned}$$

Case 4: $\varphi = \psi * \rho$, $\varphi' = \psi * \rho'$. A similar proof as above holds.

Case 5: $\varphi = \Box\psi$, $\varphi' = \Box\psi'$.

$$\begin{aligned} (1) \quad & \psi \rightarrow \psi' && \text{PL:I.H.} \\ (2) \quad & \Box(\psi \rightarrow \psi') && \text{N(1)} \\ (3) \quad & \Box(\psi \rightarrow \psi') \rightarrow (\Box\psi \rightarrow \Box\psi') && \text{US:K} \\ (4) \quad & \Box\psi \rightarrow \Box\psi' && \text{MP(3)(2)} \end{aligned}$$

Analogously, $\Box\psi' \rightarrow \Box\psi$. Therefore $\Box\psi \leftrightarrow \Box\psi'$.

Case 6: $\varphi = \Diamond\psi$ and we are replacing a \Diamond inside ψ , then $\varphi' = \Diamond\psi'$.

$$\begin{aligned} (1) \quad & \psi \leftrightarrow \psi' && \text{I.H.} \\ (2) \quad & \neg\psi \leftrightarrow \neg\psi' && \text{PL:(1)} \\ (3) \quad & \Box\neg\psi \leftrightarrow \Box\neg\psi' && \text{I.H.(Case 5):(2)} \\ (4) \quad & \Diamond\psi \leftrightarrow \neg\Box\neg\psi && \text{I.H.(Case 1)} \\ (5) \quad & \Diamond\psi' \leftrightarrow \neg\Box\neg\psi' && \text{I.H.(Case 1)} \\ (6) \quad & \Diamond\psi \leftrightarrow \Diamond\psi' && \text{PL:(3),(4),(5)} \end{aligned}$$

Finally we have proved $\vdash_K \varphi \leftrightarrow \varphi'$ holds for any basic modal formula φ . \square