

A Proof of Dual Replacement Lemma

DING Chao

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Let φ be a formula in basic modal and we denote φ' as the formula obtained by replacing a “ \Diamond ” inside φ with “ $\neg\Box\neg$ ”. We claim that

$$\vdash_K \varphi \leftrightarrow \varphi'.$$

Proof. We shall prove this by structural induction on φ .

Case 1: $\varphi = \Diamond\psi$ and we are replacing the first \Diamond , then $\varphi' = \neg\Box\neg\psi$.

$$\begin{array}{ll} (1) \Diamond p \leftrightarrow \neg\Box\neg p & \text{DUAL} \\ (2) \Diamond\varphi \leftrightarrow \neg\Box\neg\varphi & \text{US:(1)} \end{array}$$

Case 2: $\varphi = \neg\psi$, $\varphi' = \neg\psi'$. We informally claim that the definition of ψ' is clear: For this case, if φ' replaces the i th \Diamond , ψ' does the i th.

$$\begin{array}{ll} (1) (p \leftrightarrow q) \rightarrow (\neg p \leftrightarrow \neg q) & \text{TAUT} \\ (2) (\psi \leftrightarrow \psi') \leftrightarrow (\neg\psi \leftrightarrow \neg\psi') & \text{US:(1)} \\ (3) \psi \leftrightarrow \psi' & \text{I.H.} \\ (4) \neg\psi \leftrightarrow \neg\psi' & \text{MP(2)(3)} \end{array}$$

Case 3: $\varphi = \psi * \rho$, $\varphi' = \psi' * \rho$, where $*$ $\in \{\wedge, \vee, \rightarrow, \leftrightarrow\}$.

$$\begin{array}{ll} (1) (p \leftrightarrow q) \rightarrow ((p * r) \leftrightarrow (q * r)) & \text{TAUT} \\ (2) (\psi \leftrightarrow \psi') \rightarrow ((\psi * \rho) \leftrightarrow (\psi' * \rho)) & \text{US:(1)} \\ (3) \psi \leftrightarrow \psi' & \text{I.H.} \\ (4) (\psi * \rho) \leftrightarrow (\psi' * \rho) & \text{MP(2)(3)} \end{array}$$

Case 4: $\varphi = \psi * \rho$, $\varphi' = \psi * \rho'$. A similar proof as above holds.

Case 5: $\varphi = \Box\psi$, $\varphi' = \Box\psi'$.

$$\begin{array}{ll} (1) \psi \rightarrow \psi' & \text{PL:I.H.} \\ (2) \Box(\psi \rightarrow \psi') & \text{N(1)} \\ (3) \Box(\psi \rightarrow \psi') \rightarrow (\Box\psi \rightarrow \Box\psi') & \text{US:K} \\ (4) \Box\psi \rightarrow \Box\psi' & \text{MP(3)(2)} \end{array}$$

Analogously, $\Box\psi' \rightarrow \Box\psi$. Therefore $\Box\psi \leftrightarrow \Box\psi'$.

Case 6: $\varphi = \Diamond\psi$ and we are replacing a \Diamond inside ψ , then $\varphi' = \Diamond\psi'$.

$$\begin{array}{ll} (1) \Box\neg\psi \leftrightarrow \Box\neg\psi' & \text{I.H.(Case 5)} \\ (2) \Diamond\psi \leftrightarrow \neg\Box\neg\psi & \text{I.H.(Case 1)} \\ (3) \Diamond\psi' \leftrightarrow \neg\Box\neg\psi' & \text{I.H.(Case 1)} \\ (4) \Diamond\psi \leftrightarrow \Diamond\psi' & \text{PL:(1),(2),(3)} \end{array}$$

Finally we have proved $\vdash_K \varphi \leftrightarrow \varphi'$ holds for any basic modal formula φ . \square